CSE 421 Algorithms

Richard Anderson

Autumn 2019

Lecture 8 – Greedy Algorithms II

1

Announcements

- Today's lecture
  - Kleinberg-Tardos, 4.2, 4.3
- Next week
  - Kleinberg-Tardos, 4.4, 4.5

2

Scheduling Intervals

- Given a set of intervals
  - What is the largest set of non-overlapping intervals
  - What is the minimum number of processors required to schedule all intervals
- Suppose the n intervals are "random"
  - What is the expected number of independent intervals
  - What is the expected depth

3

Generating a random set of intervals

- Method 1:
  - Each interval assigned random start position in [0.0, 1.0]
  - Each interval assigned a random length in [0.0, 1.0]
- Method 2:
  - Start with the array [1, 1, 2, 2, 3, 3, 4, 4, 5, 5]
  - Randomly permute it [2, 1, 4, 2, 3, 4, 5, 1, 3, 5]
  - Index of the first j is the start of interval j, and the index of the second j is the end of interval j

4

Greedy Algorithms

- Solve problems with the simplest possible algorithm
- The hard part: showing that something simple actually works
- Today's problems (Sections 4.2, 4.3)
  - Homework Scheduling
  - Optimal Caching
  - Subsequence testing

5

Homework Scheduling

- Tasks to perform
- Deadlines on the tasks
- Freedom to schedule tasks in any order

- Can I get all my work turned in on time?
- If I can't get everything in, I want to minimize the maximum lateness

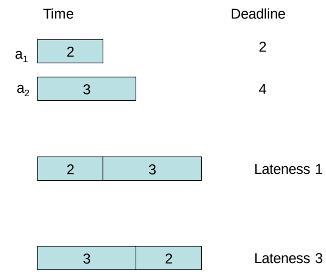
6

Scheduling tasks

- Each task has a length t_i and a deadline d_i
- All tasks are available at the start
- One task may be worked on at a time
- All tasks must be completed
- Goal minimize maximum lateness
  - Lateness: $L_i = f_i – d_i$ if $f_i >= d_i$

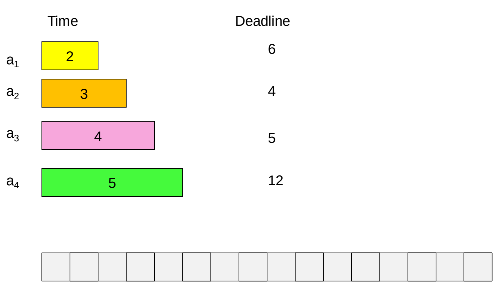
7

Example



8

Determine the minimum lateness



9

Greedy Algorithm

- Earliest deadline first
- Order jobs by deadline
- This algorithm is optimal

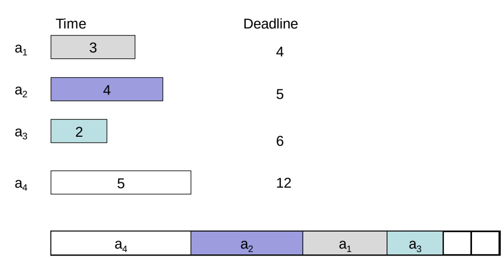
10

Analysis

- Suppose the jobs are ordered by deadlines, $d_1 <= d_2 <= . . . <= d_n$
- A schedule has an *inversion* if job j is scheduled before i where j > i
- The schedule A computed by the greedy algorithm has no inversions.
- Let O be the optimal schedule, we want to show that A has the same maximum lateness as O

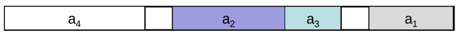
11

List the inversions



12

Lemma: There is an optimal schedule with no idle time

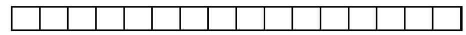


- It doesn't hurt to start your homework early!
- Note on proof techniques
  - This type of can be important for keeping proofs clean
  - It allows us to make a simplifying assumption for the remainder of the proof

13

Lemma

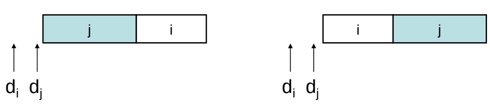
- If there is an inversion i, j, there is a pair of adjacent jobs i', j' which form an inversion



14

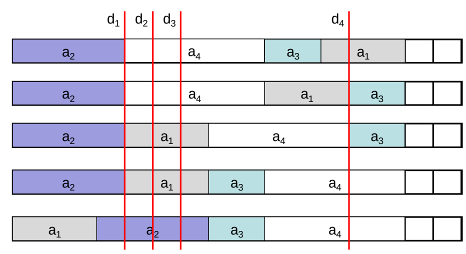
Interchange argument

- Suppose there is a pair of jobs i and j, with $d_i <= d_j$, and j scheduled immediately before i. Interchanging i and j does not increase the maximum lateness.



15

Proof by Bubble Sort



Determine maximum lateness

16

Real Proof

- There is an optimal schedule with no inversions and no idle time.
- Let O be an optimal schedule k inversions, we construct a new optimal schedule with k-1 inversions
- Repeat until we have an optimal schedule with 0 inversions
- This is the solution found by the earliest deadline first algorithm

17

Result

- Earliest Deadline First algorithm constructs a schedule that minimizes the maximum lateness

18

Homework Scheduling

- How is the model unrealistic?

19

Extensions

- What if the objective is to minimize the sum of the lateness?
  - EDF does not work
- If the tasks have release times and deadlines, and are non-preemptable, the problem is NP-complete
- What about the case with release times and deadlines where tasks are preemptable?

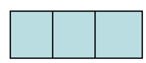
20

Optimal Caching

- Caching problem:
  - Maintain collection of items in local memory
  - Minimize number of items fetched

21

Caching example



A, B, C, D, A, E, B, A, D, A, C, B, D, A

22

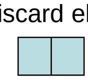
Optimal Caching

- If you know the sequence of requests, what is the optimal replacement pattern?
- Note – it is rare to know what the requests are in advance – but we still might want to do this:
  - Some specific applications, the sequence is known
    - Register allocation in code generation
  - Competitive analysis, compare performance on an online algorithm with an optimal offline algorithm

23

Farthest in the future algorithm

- Discard element used farthest in the future



A, B, C, A, C, D, C, B, C, A, D

24

Correctness Proof

- Sketch
- Start with Optimal Solution O
- Convert to Farthest in the Future Solution F-F
- Look at the first place where they differ
- Convert O to evict F-F element
  - There are some technicalities here to ensure the caches have the same configuration . . .

25

Subsequence Testing

- Is $a_1a_2\ldots a_m$ a subsequence of $b_1b_2\ldots b_n$?
  - e.g. is T,R,E,E a subsequence of S,T,U,A,R,T,R,E,G,E,S

T	R	E	E

S	T	U	A	R	T	R	E	G	E	S

26

Next week

27