

Lemma

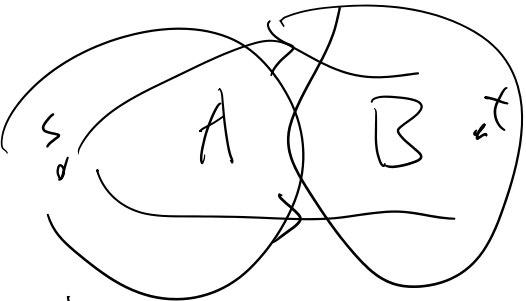
Let f be any st flow, let (A, B) be st cut

Then

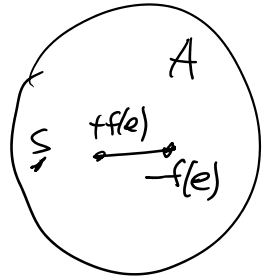
$$\underbrace{\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ into } A} f(e)}_{\text{net flow across } (A, B)} = \underbrace{v(f)}_{\text{flow value}}$$

Proof

$$v(f) \stackrel{\text{def}}{=} \sum_{e \text{ out of } s} f(e) - \sum_{e \text{ into } s} f(e)$$



$$\begin{aligned} &\rightarrow \text{flow conservation} \quad \left(\sum_{e \text{ out of } s} f(e) - \sum_{e \text{ into } s} f(e) \right) = 0 \\ &+ \sum_{\substack{v \text{ fs} \\ v \in A}} \left(\sum_{e \text{ out of } v} f(e) - \sum_{e \text{ into } v} f(e) \right) \\ &= \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ into } A} f(e) \end{aligned}$$



Def (A, B) is st cut
is A, B partition V

(A, A^c)

and $s \in A, t \in B$.

Weak duality

For any st flow f and st cut (A, B)

$$v(f) \leq c(A, B)$$

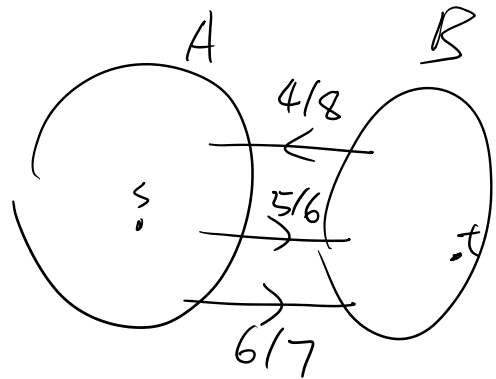
for any s, t ...

$$v(f) \leq \text{cap}(A, B)$$

Proof

$$v(f) = \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ into } A} f(e)$$

$$\leq \sum_{e \text{ out of } A} f(e)$$



$$\text{Cap}(A, B) = 6 + 7$$

$$\leq \sum_{e \text{ out of } A} c(e) = \text{Cap}(A, B) \quad v(f) = 5 + 6 - 4 = 7$$

Corollary If there is $s \overset{\text{cut}}{t}(A, B)$
 and s - t flow f
 $s.t.$ $v(f) = \text{cap}(A, B)$

Then (A, B) is min s - t cut
 f is max s - t flow

Proof If there A', B' $s.t.$ $\text{cap}(A', B') < \text{cap}(A, B)$

Then $v(f) \leq \text{cap}(A', B') < \text{cap}(A, B) = v(f)$
 it is contradiction.

$$T(n) = T(n-1) + T\left(\frac{n}{2}\right) + O(n^{2018})$$

$$= T(n-2) + T\left(\frac{n}{2}\right) + T\left(\frac{n-1}{2}\right) + O(n^{2018})$$

$$\vdots$$

$$= \underbrace{T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + \dots + T\left(\frac{n}{4}\right)}_{\geq \frac{n}{2} T\left(\frac{n}{2}\right)} + O(n^{2019})$$

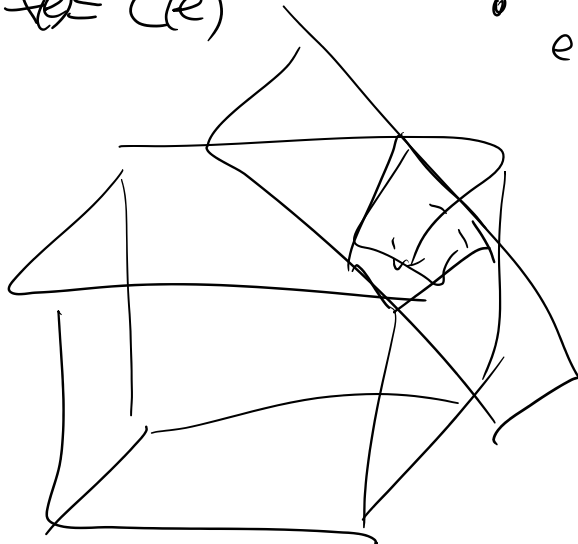
$$\begin{cases} \geq \frac{n}{2} T\left(\frac{n}{2}\right) + O(n^{2019}) \\ \leq \frac{n}{2} T\left(\frac{n}{4}\right) + O(n^{2019}) \end{cases}$$

How to draw s-t flow

- $0 \leq f(e) \leq c(e)$

- $\sum_{e \text{ out } v} f(e) = \sum_{e \text{ in } v} f(e)$

$$\forall v \notin \{s, t\}$$



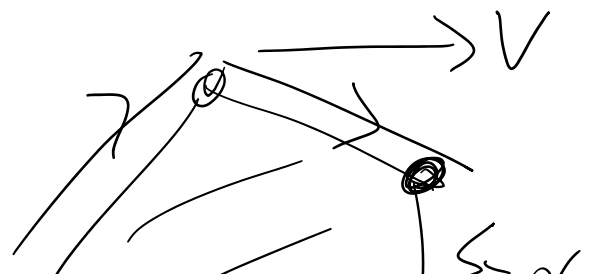
$$f_1 = f_2$$

$$\begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}^T (f) = 0$$

$$f = \begin{pmatrix} \\ \\ \end{pmatrix} \rightarrow |E| \text{ many coordinates}$$

What is max flow?

$$\max v^T f$$



$U \cup V \cup W \cup X \cup Y$
 $f \in$ 

