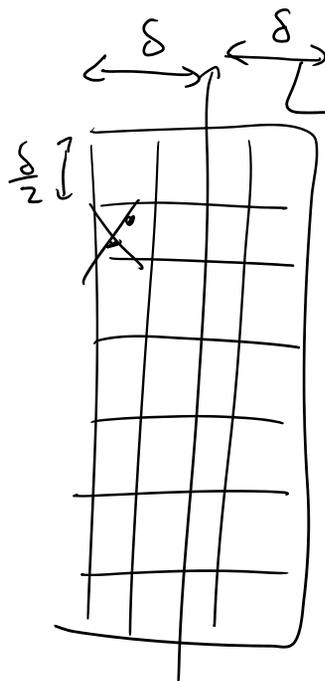


Claim: No 2 points lie in the same box

PF: If 2 points in the same box,

$$\text{then distance} \leq \sqrt{\left(\frac{\delta}{2}\right)^2 + \left(\frac{\delta}{2}\right)^2} < \delta$$

(contradicts to the def of δ).



ALG(P_1, \dots, P_n)

find vertical line cut $\{P_1, \dots, P_{\frac{n}{2}}\}$ on left

$\{P_{\frac{n}{2}+1}, \dots, P_n\}$ on right

$$\delta_1 = \text{ALG}(P_1, \dots, P_{\frac{n}{2}})$$

$$\delta_2 = \text{ALG}(P_{\frac{n}{2}+1}, \dots, P_n)$$

$$\delta = \min(\delta_1, \delta_2)$$

← we know there is a pair with dist exactly δ

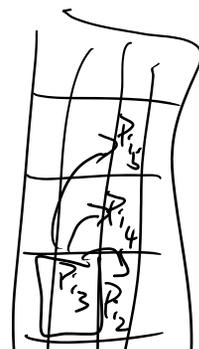
Let L be strip.

Let $P_{i_1}, P_{i_2}, \dots, P_{i_m}$ be the points in L .

Sort the P_{i_j} according to the y-axis

for $j = 1, \dots, m$

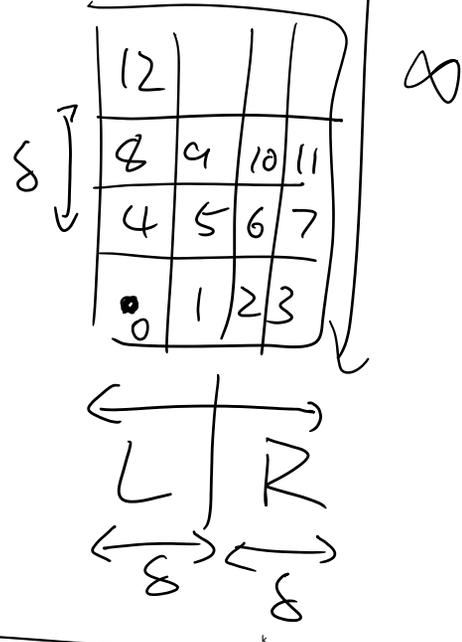
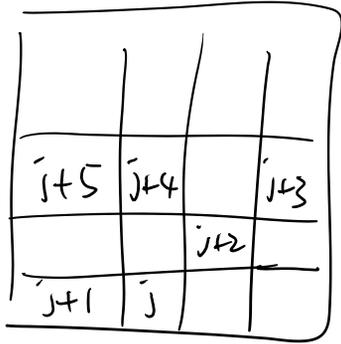
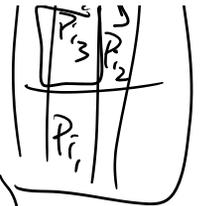
for $k = 1, 2, \dots, 11$



for $k = 1, 2, \dots, 11$

if $j+k \leq m$

$$\delta = \min(\delta, \text{dist}(P_{i,j}, P_{i,j+k}))$$



Runtime :

$$T(n) \leq \begin{cases} 1 & \text{if } n=1 \\ 2T(\frac{n}{2}) + O(n \log n) + O(n) & \text{else} \end{cases}$$

$$\begin{aligned} T(n) &= O(n \log n) + 2T(\frac{n}{2}) \\ &= O(n \log n) + 2 \left[O(\frac{n}{2} \log \frac{n}{2}) + 2T(\frac{n}{4}) \right] \\ &= O(n \log n) + O(n \log n) + 4T(\frac{n}{4}) \end{aligned}$$

$\log n$ level

\vdots
 $n, \dots, 1$

$$L = O(n \log^2 n)$$

Correctness:

Let $P_i, P_{i'}$ be the closest pair in the strip L with $\text{dist} \leq \delta$.

WLOG, $i < i'$

Claim $i' \leq i+11$

$P_{i'}$ must appear in the same / next / next next row as P_i .

• There are 12 boxes in those 3 rows.

• Each box contains 1 point

$\Rightarrow i' \leq i+11$