

CSE 421 Algorithms

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Lecture 29
Complexity Theory

Announcements

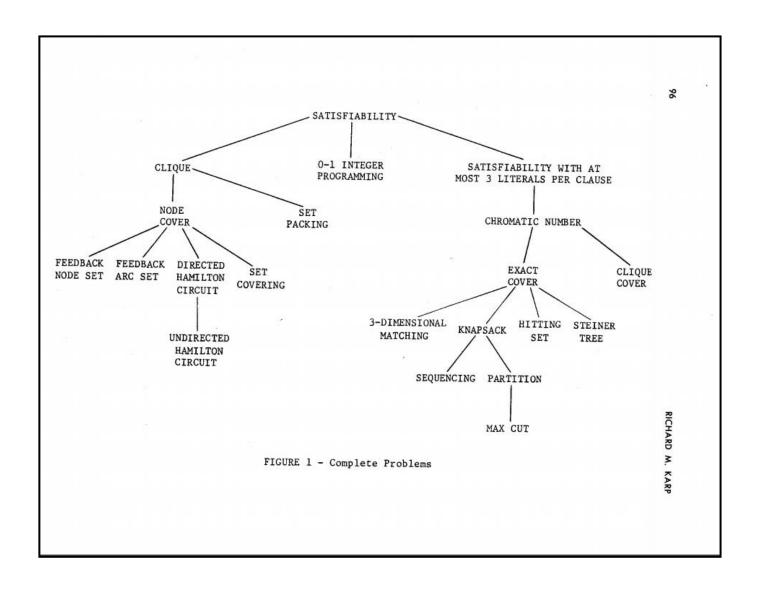
- Final exam,
 - Monday, December 14, 2:30-4:20 pm
 - Comprehensive (2/3 post midterm, 1/3 pre midterm)
- Review session
 - Friday, 3:30 5:00 pm. More 220
- Online course evaluations available

NP Complete Problems

- 1. Circuit Satisfiability
- 2. Formula Satisfiability
 - a. 3-SAT
- 3. Graph Problems
 - a. Independent Set
 - b. Vertex Cover
 - c. Clique
- 4. Path Problems
 - a. Hamiltonian cycle
 - b. Hamiltonian path

- 5. Partition Problems
 - a. Three dimensional matching
 - b. Exact cover
- 6. Graph Coloring
- 7. Number problems
 - a. Subset sum
- 8. Integer linear programming
- 9. Scheduling with release times and deadlines

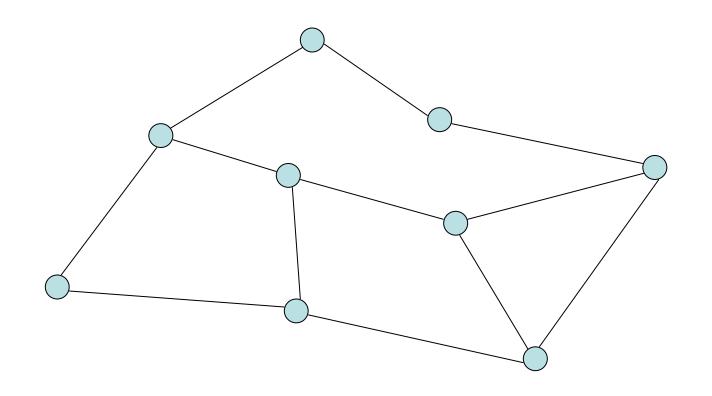
Karp's 21 NP Complete Problems



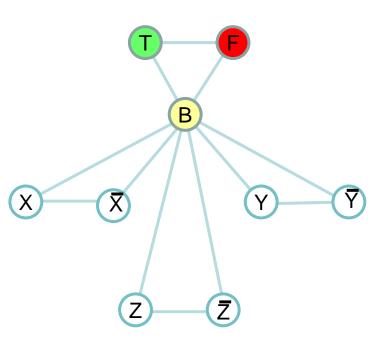
A final NP completeness result: Graph Coloring

- NP-Complete
 - Graph K-coloring
 - Graph 3-coloring

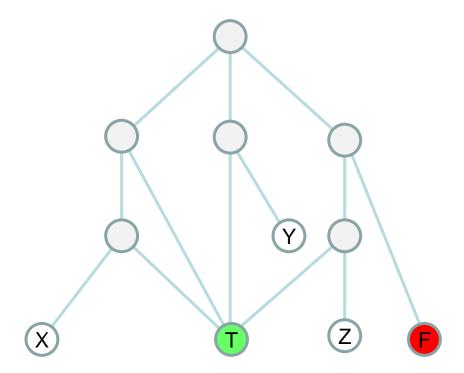
- Polynomial
 - Graph 2-Coloring



3-SAT <_P 3 Colorability



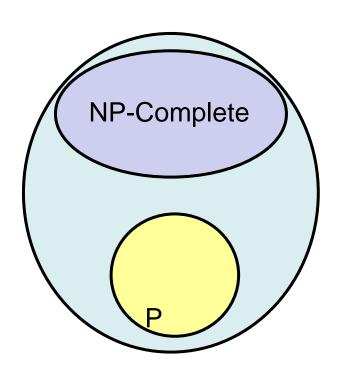
Truth Setting Gadget

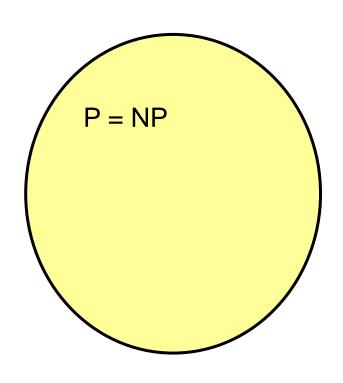


Clause Testing Gadget

What we don't know

• P vs. NP



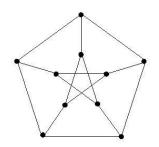


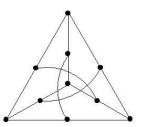
If P!= NP, is there anything in between

- Yes, Ladner [1975]
- Problems not known to be in P or NP Complete
 - Factorization
 - Discrete Log

Solve $g^k = b$ over a finite group

Graph Isomorphism



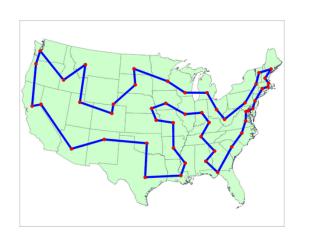


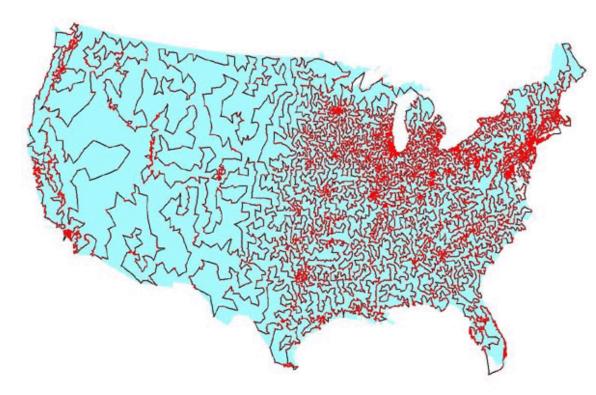
Coping with NP Completeness

- Approximation Algorithms
 - Christofides algorithm for TSP (Undirected graphs satisfying triangle inequality)
- Solution guarantees on greedy algorithms
 - Bin packing

Coping with NP-Completeness

- Branch and Bound
 - Euclidean TSP





Coping with NP-Completeness

- Local Search
 - Modify solution until a local minimum is reached
 - Interchange algorithm for TSP
 - Recoloring algorithms
 - Simulated annealing

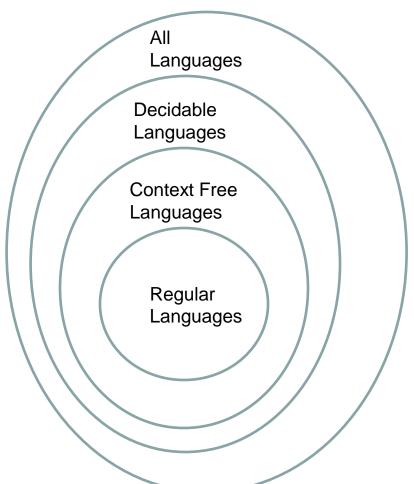
Complexity Theory

Computational requirements to recognize

languages

Models of Computation

- Resources
- Hierarchies



Time complexity

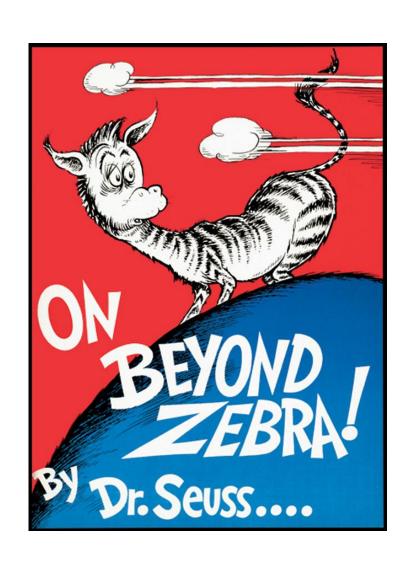
- P: (Deterministic) Polynomial Time
- NP: Non-deterministic Polynomial Time
- EXP: Exponential Time

Space Complexity

- Amount of Space (Exclusive of Input)
- L: Logspace, problems that can be solved in O(log n) space for input of size n

 PSPACE, problems that can be required in a polynomial amount of space

So what is beyond NP?



NP vs. Co-NP

 Given a Boolean formula, is it true for some choice of inputs

 Given a Boolean formula, is it true for all choices of inputs

Problems beyond NP

 Exact TSP, Given a graph with edge lengths and an integer K, does the minimum tour have length K

 Minimum circuit, Given a circuit C, is it true that there is no smaller circuit that computes the same function a C

Polynomial Hierarchy

- Level 1
 - $-\exists X_1 \Phi(X_1), \forall X_1 \Phi(X_1)$
- Level 2
 - $\forall X_1 \exists X_2 \Phi(X_1, X_2), \exists X_1 \forall X_2 \Phi(X_1, X_2)$
- Level 3
 - $\forall X_{1} \exists X_{2} \forall X_{3} \Phi(X_{1}, X_{2}, X_{3}), \exists X_{1} \forall X_{2} \exists X_{3} \Phi(X_{1}, X_{2}, X_{3})$

Polynomial Space

- Quantified Boolean Expressions
 - $-\exists X_{1} \forall X_{2} \exists X_{3} ... \exists X_{n-1} \forall X_{n} \Phi(X_{1}, X_{2}, X_{3} ... X_{n-1} X_{n})$
- Space bounded games
 - Competitive Facility Location Problem

- Counting problems
 - The number of Hamiltonian Circuits in a graph