

# CSE 421

## Algorithms

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Lecture 28

Survey of NP Complete Problems

# Announcements

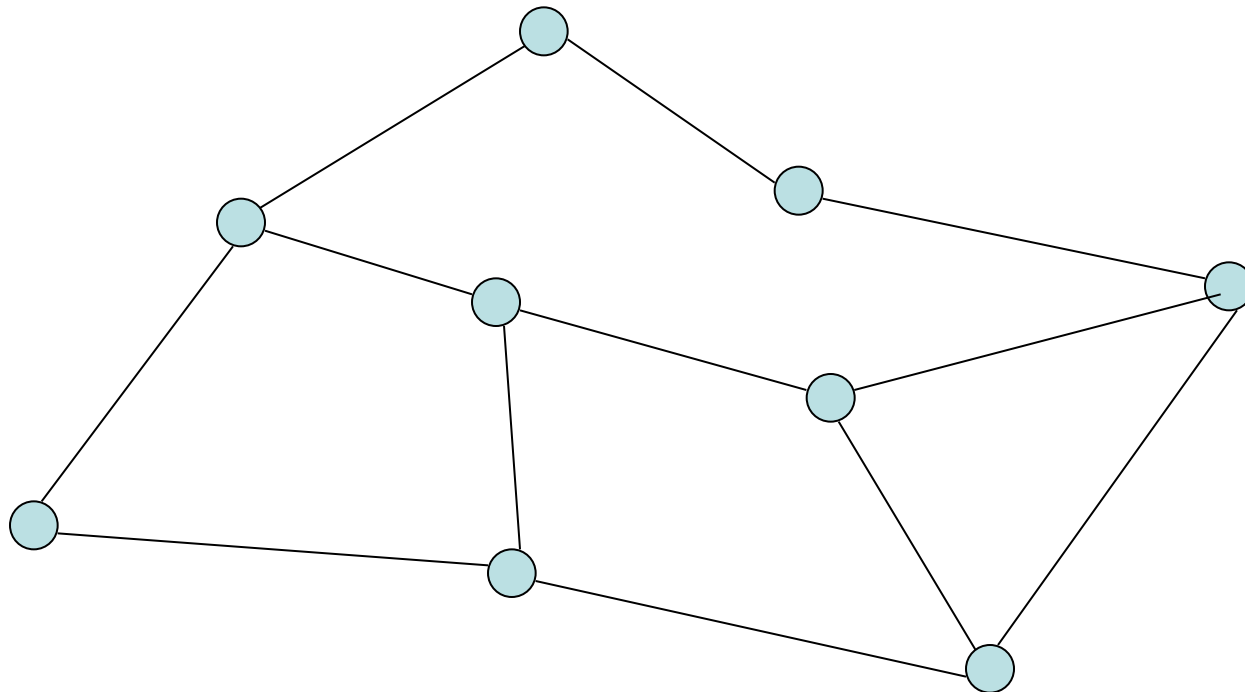
- Final exam,
  - Monday, December 14, 2:30-4:20 pm
  - Comprehensive (2/3 post midterm, 1/3 pre midterm)
- Review session
  - Friday, 3:30 – 5:00 pm. More 220
- Online course evaluations available

# NP Complete Problems

1. Circuit Satisfiability
2. Formula Satisfiability
  - a. 3-SAT
3. Graph Problems
  - a. Independent Set
  - b. Vertex Cover
  - c. Clique
4. Path Problems
  - a. Hamiltonian cycle
  - b. Hamiltonian path
5. Partition Problems
  - a. Three dimensional matching
  - b. Exact cover
6. Graph Coloring
7. Number problems
  - a. Subset sum
8. Integer linear programming
9. Scheduling with release times and deadlines

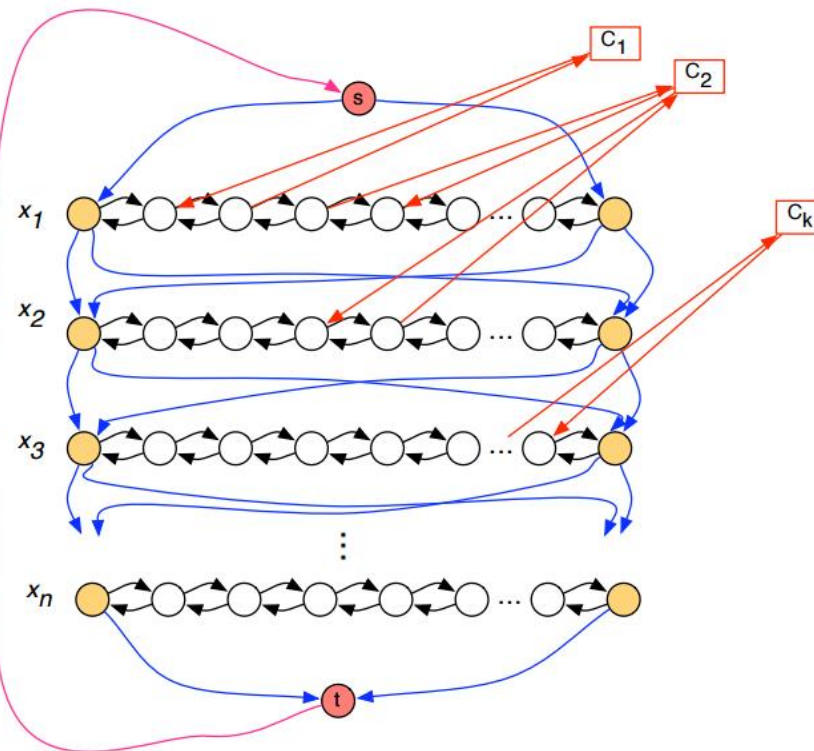
# Hamiltonian Circuit Problem

- Hamiltonian Circuit – a simple cycle including all the vertices of the graph



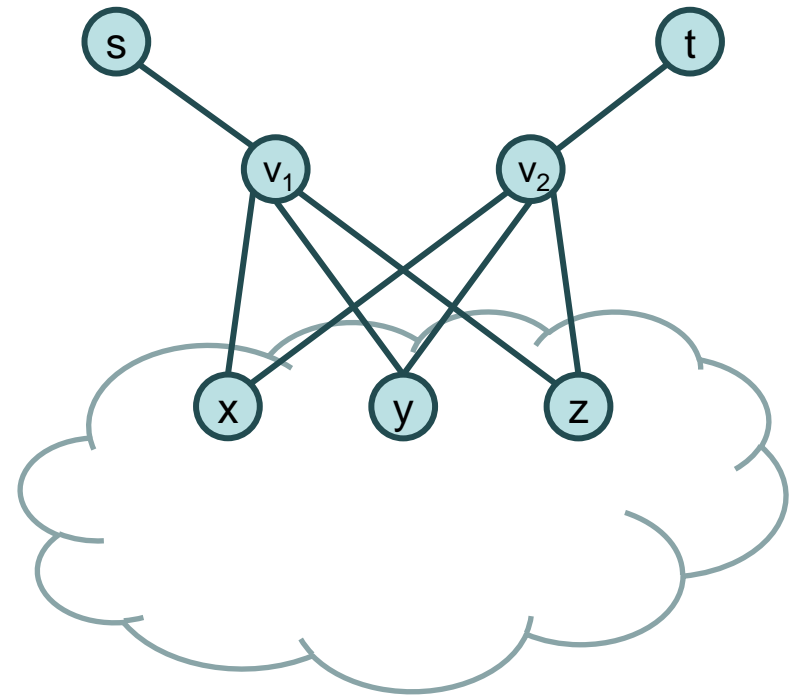
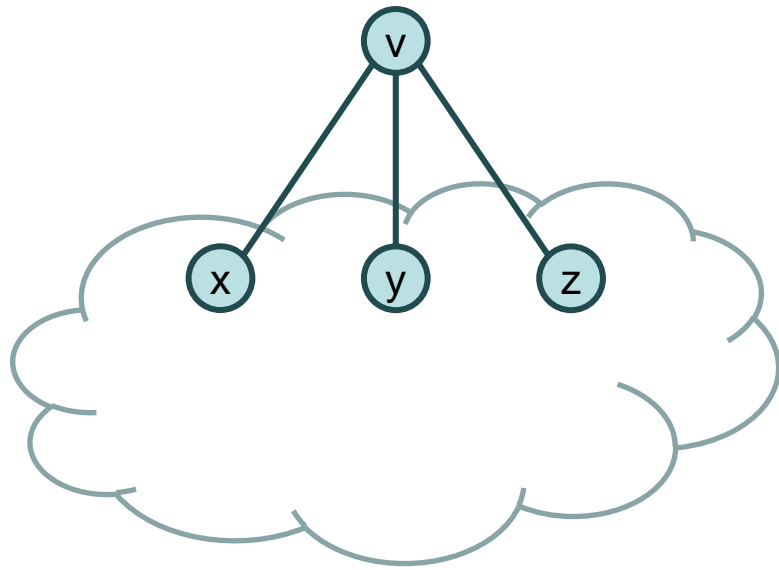
# Thm: Hamiltonian Circuit is NP Complete

- Reduction from 3-SAT



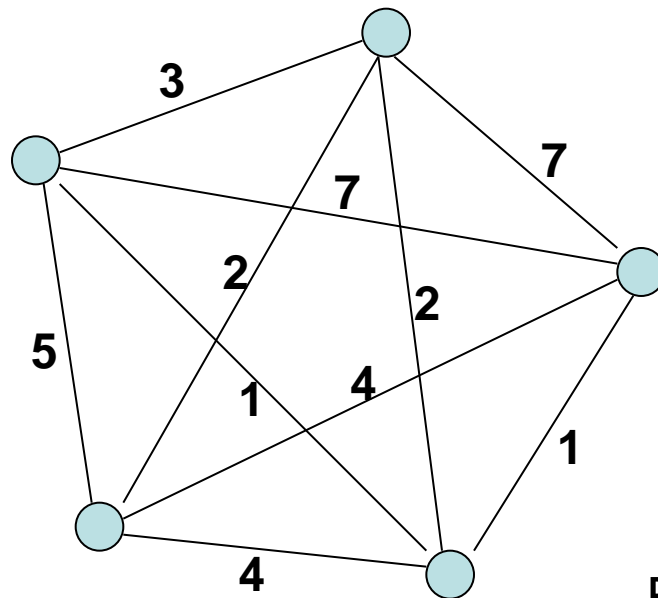
# Reduce Hamiltonian Circuit to Hamiltonian Path

$G_2$  has a Hamiltonian Path iff  $G_1$  has a Hamiltonian Circuit



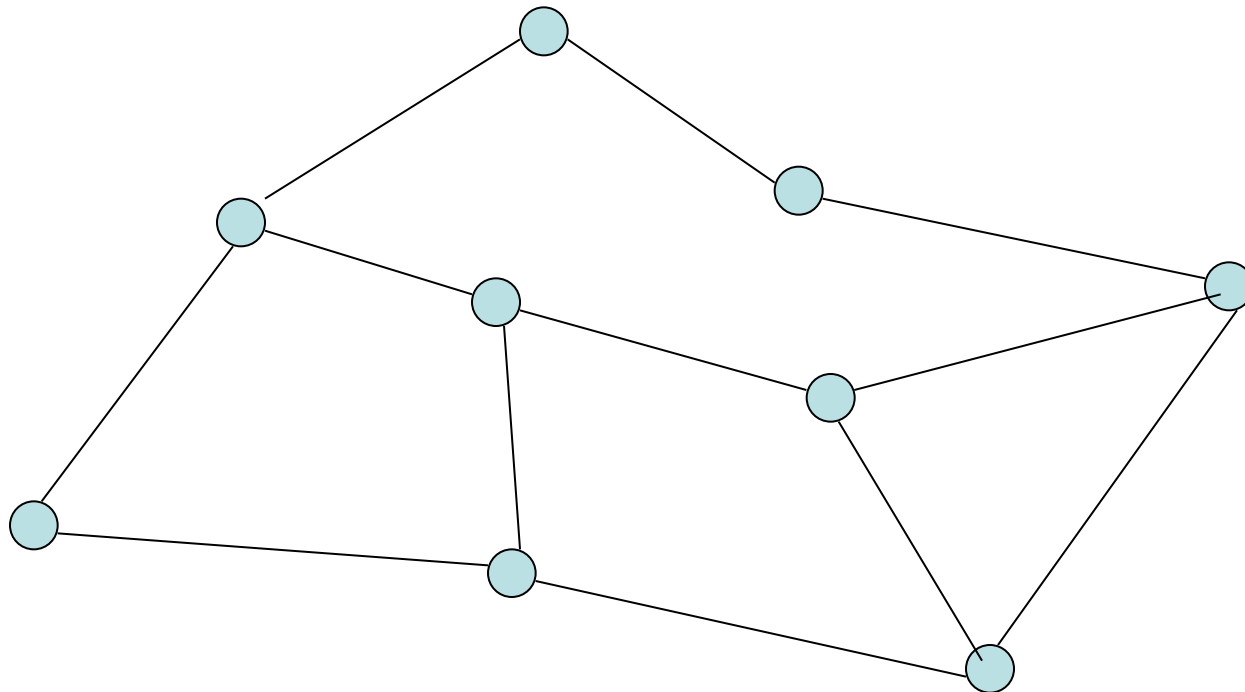
# Traveling Salesman Problem

- Given a complete graph with edge weights, determine the shortest tour that includes all of the vertices (visit each vertex exactly once, and get back to the starting point)



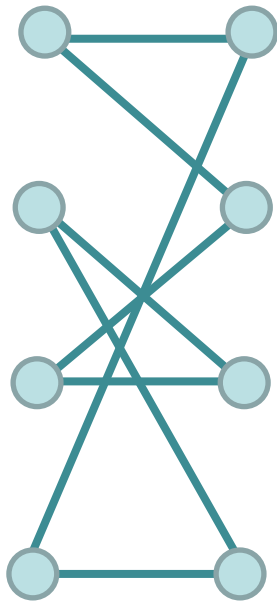
Find the minimum cost tour

Thm:  $HC \leq_p TSP$

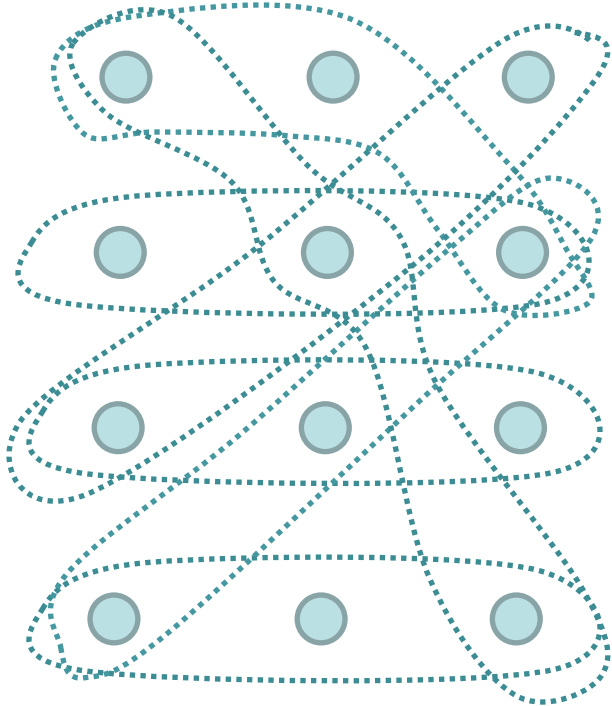




# Matching

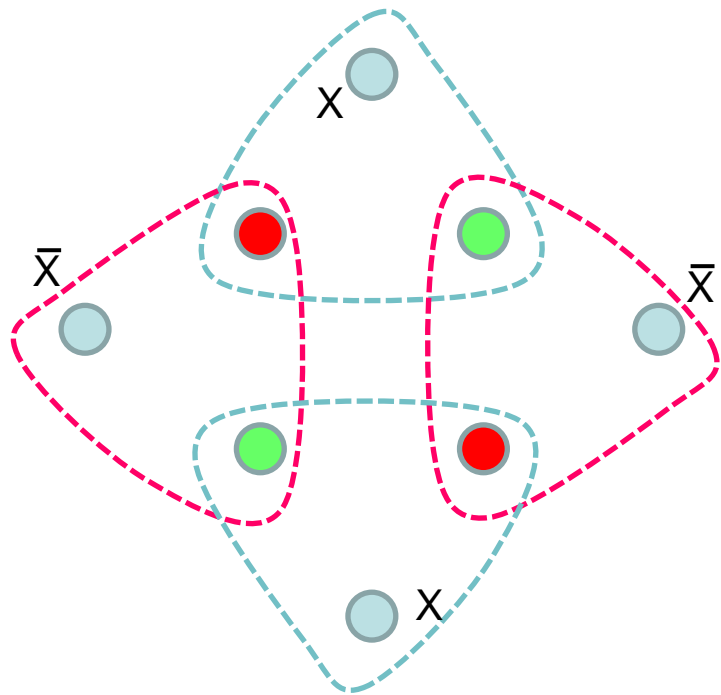


Two dimensional matching

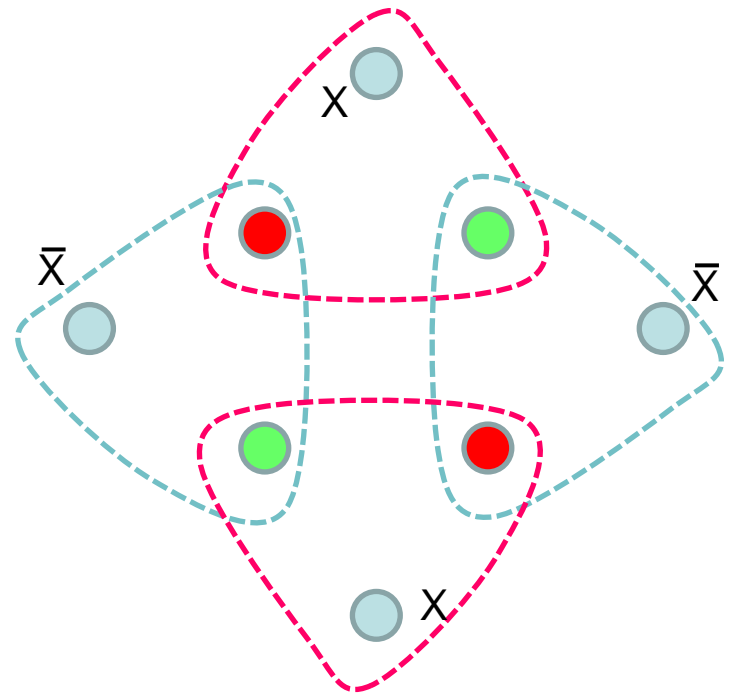


Three dimensional matching (3DM)

# 3-SAT $\leq_P$ 3DM



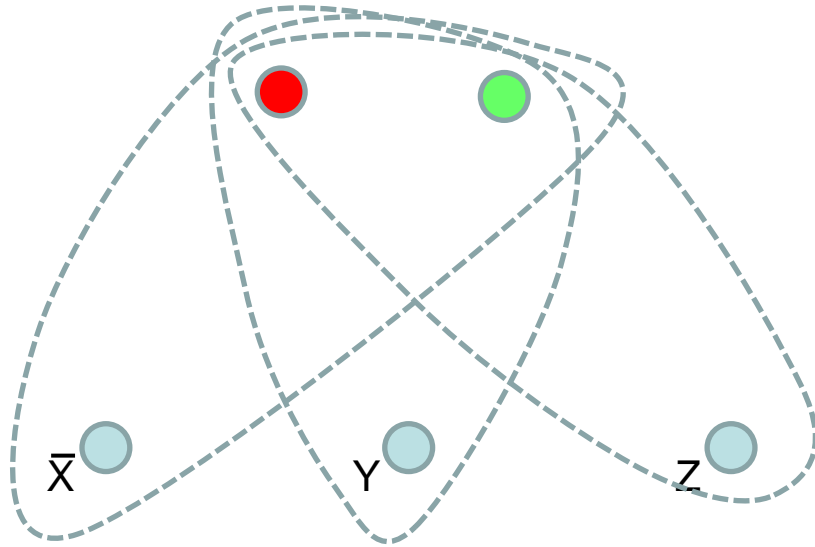
$X$  True



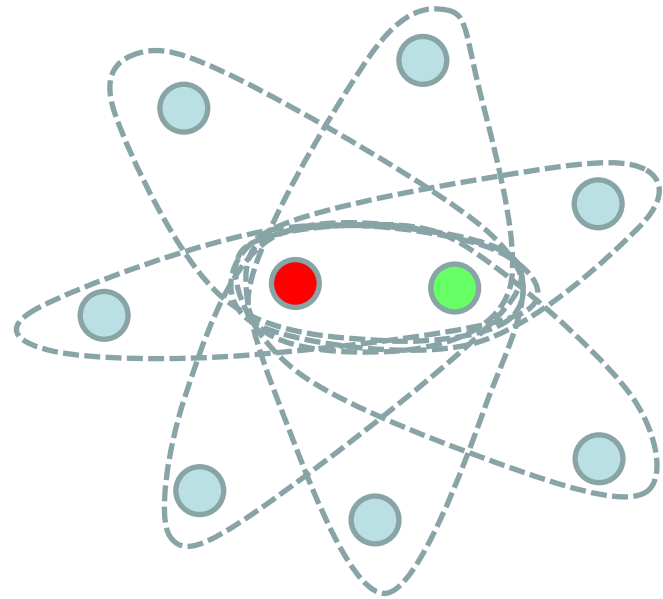
$X$  False

Truth Setting Gadget

# 3-SAT $\leq_P$ 3DM



Clause gadget for ( $\bar{X}$  OR Y OR Z)



Garbage Collection Gadget  
(Many copies)

# Exact Cover (sets of size 3) XC3

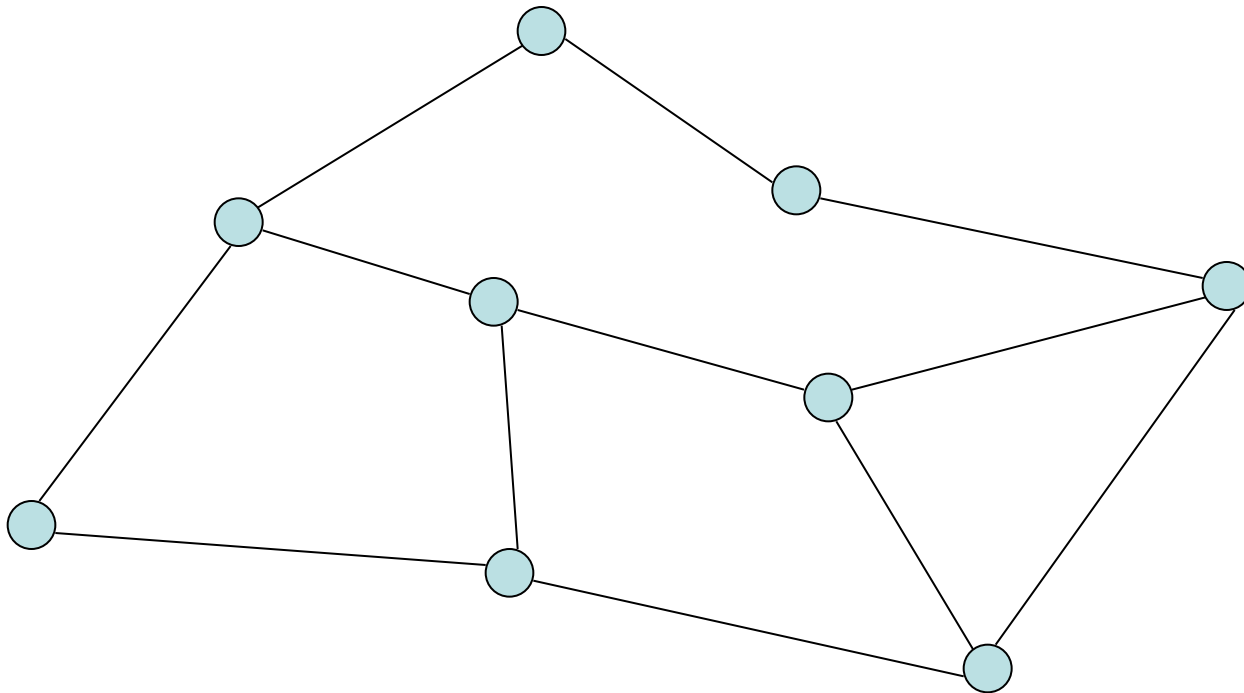
Given a collection of sets of size 3 of a domain of size  $3N$ , is there a sub-collection of  $N$  sets that cover the sets

(A, B, C), (D, E, F), (A, B, G),  
(A, C, I), (B, E, G), (A, G, I),  
(B, D, F), (C, E, I), (C, D, H),  
(D, G, I), (D, F, H), (E, H, I),  
(F, G, H), (F, H, I)

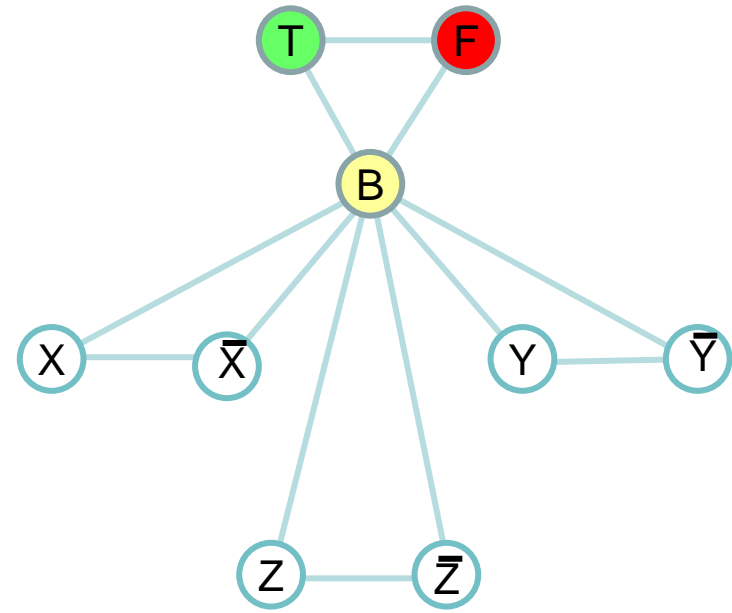
$$3DM \leq_P XC3$$

# Graph Coloring

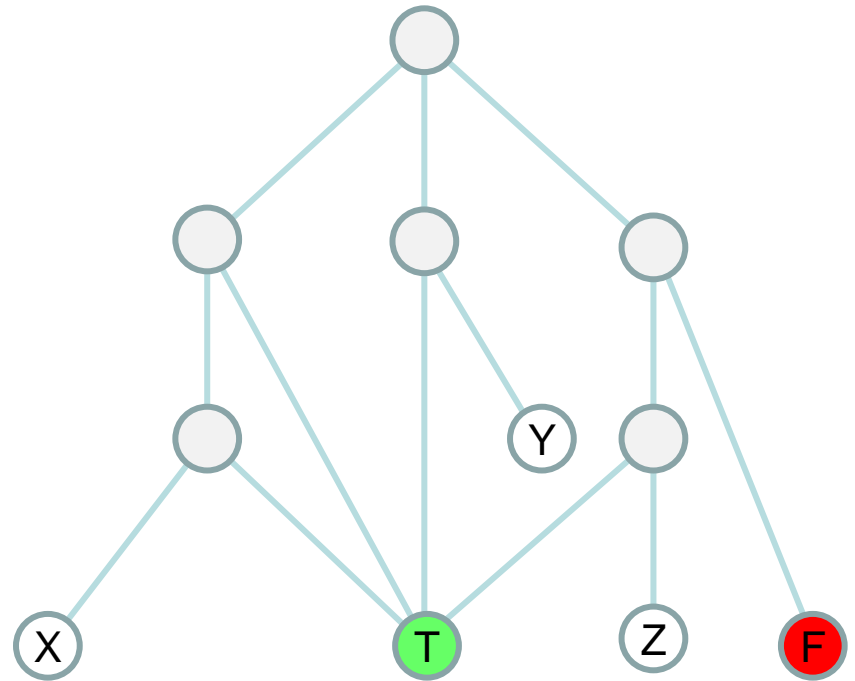
- NP-Complete
  - Graph K-coloring
  - Graph 3-coloring
- Polynomial
  - Graph 2-Coloring



# 3-SAT $\leq_P$ 3 Colorability



Truth Setting Gadget



Clause Testing Gadget

# Number Problems

- Subset sum problem
  - Given natural numbers  $w_1, \dots, w_n$  and a target number  $W$ , is there a subset that adds up to exactly  $W$ ?
- Subset sum problem is NP-Complete
- Subset Sum problem can be solved in  $O(nW)$  time

# $XC3 \leq_p$ SUBSET SUM

Idea: Represent each set as a bit vector, then interpret the bit vectors as integers. Add them up to get the all one's vector.

$\{x_3, x_5, x_9\} \Rightarrow 001010001000$

Does there exist a subset that sums to exactly 111111111111?

Annoying detail: What about the carries?



# Integer Linear Programming

# Scheduling with release times and deadlines