

CSE 42 I: Introduction to Algorithms

I: Overview

Larry Ruzzo



University of Washington

Computer Science & Engineering

CSE 421, Su '12: Introduction to Algorithms

▷ CSE Home

Administrative

[FAQ](#)
[Schedule & Reading](#)

Course Email/BBoard

[Subscription Options](#)
[Class List Archive](#)
[E-mail Course Staff](#)
[GoPost BBoard](#)

Lecture Notes

[1: Overview & Example](#)
[2: Analysis](#)

Lecture: [LOW 206](#) (schematic)

MW 10:50- 12:20

Instructor: [Larry Ruzzo](#), ruzzo@cs

Office Hours

W 1:00- 2:00

Location

6298

TA: Tamara Bonaci, tbonaci@uw.edu

<http://www.cs.washington.edu/421>

Course Email

and general interest student/staff Q&A

..., but probably should [change their default subscrip](#)

... use [Catalyst GoPost](#) to discuss homework, etc.

... techniques for design of efficient algorithms. Methods for showing lower bo
..., set manipulation, arithmetic, graph problems, pattern matching.

Prerequisites: either [CSE 312](#) or [CSE 322](#); either [CSE 326](#) or [CSE 332](#).

Credits: 3

Grading: Homework, Final. Homework will be a mix of paper & pencil exercises and programing. O

Late Policy: Unless otherwise announced, weekly homeworks will be due by 4:00PM on Thursdays;
electronically, including scanned versions of handwritten papers, via the Catalyst drop box link at left

Extra Credit: Assignments may include "extra credit" sections. These will enrich your understanding
glory, not the points, and don't start extra credit until the basics are complete.

Textbook: [Algorithm Design](#) by [Jon Kleinberg](#) and [Eva Tardos](#). Addison Wesley, 2006. (A₂ailable fi

What you have to do

Homework (~60% of grade)

Programming?

perhaps some small projects

Written homework assignments

English exposition and pseudo-code

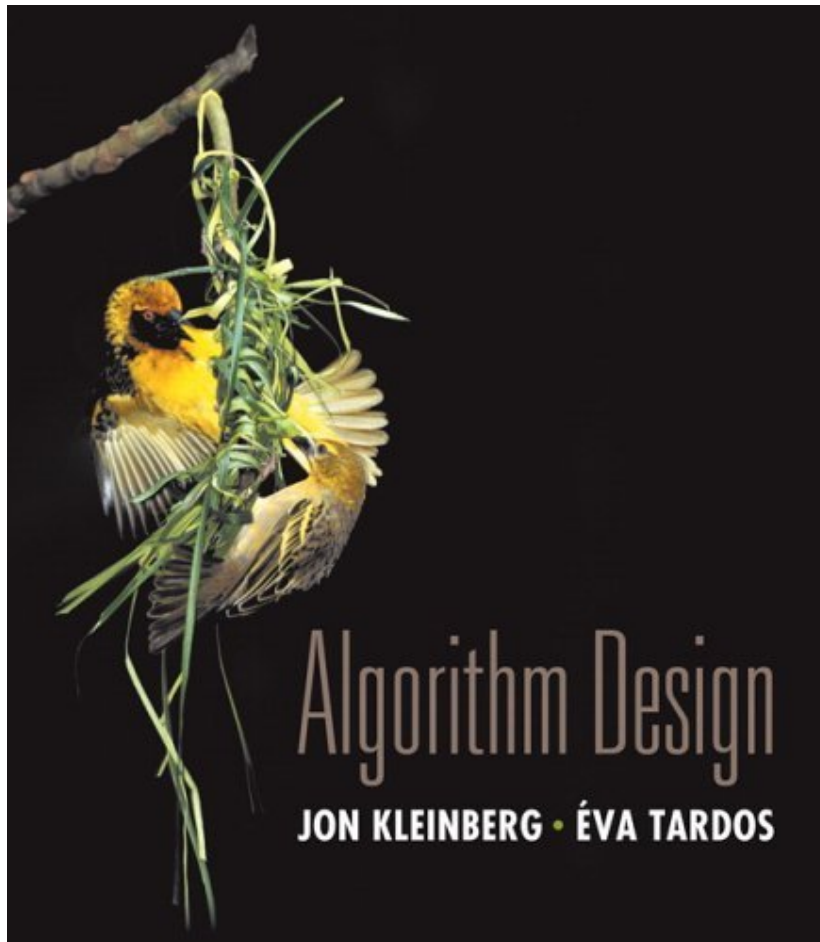
Analysis and argument as well as design

Final Exam (~40%)

Late Policy:

Papers and/or electronic turnins generally due Thursdays by 4:00 pm; minus 20% per day thereafter

Textbook



Algorithm Design by
Jon Kleinberg and
Eva Tardos. Addison
Wesley, 2006.

What the course is about

Design of Algorithms

design methods

common or important types of problems

analysis of algorithms - efficiency

correctness proofs

What the course is about

Complexity, NP-completeness and intractability

solving problems in principle is not enough

algorithms must be *efficient*

some problems have *no efficient solution*

NP-complete problems

important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

Very Rough Division of Time

Algorithms (6-7 weeks)

Analysis

Techniques: greedy, divide&conquer, dynamic programming

Applications: graph algorithms, flows & matchings

Complexity & NP-completeness (2-3 weeks)

Check online
schedule page for
(evolving) details



University of Washington
Computer Science & Engineering

CSE 417, Wi '06: *Approximate* Schedule

[CSE Home](#)

[About Us](#) [S](#)

		Due	Lecture Topic	Reading
Week 1 1/2-1/6	M		Holiday	
	W		Intro, Examples & Complexity	Ch. 1; Ch. 2
	F		Intro, Examples & Complexity	
Week 2 1/9-1/13	M		Intro, Examples & Complexity	
	W		Graph Algorithms	Ch. 3
	F		Graph Algorithms	

Complexity Example

Cryptography (e.g., RSA, SSL in browsers)

Secret: p, q prime, say 512 bits each

Public: n which equals $p \times q$, 1024 bits

In principle

there is an algorithm that given n will find p and q :
try all $2^{512} > 1.3 \times 10^{154}$ possible p 's: kinda slow...

In practice

no fast algorithm known for this problem (on non-quantum computers)

security of RSA depends on this fact

(“quantum computing”: strongly driven by possibility of changing this)

Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

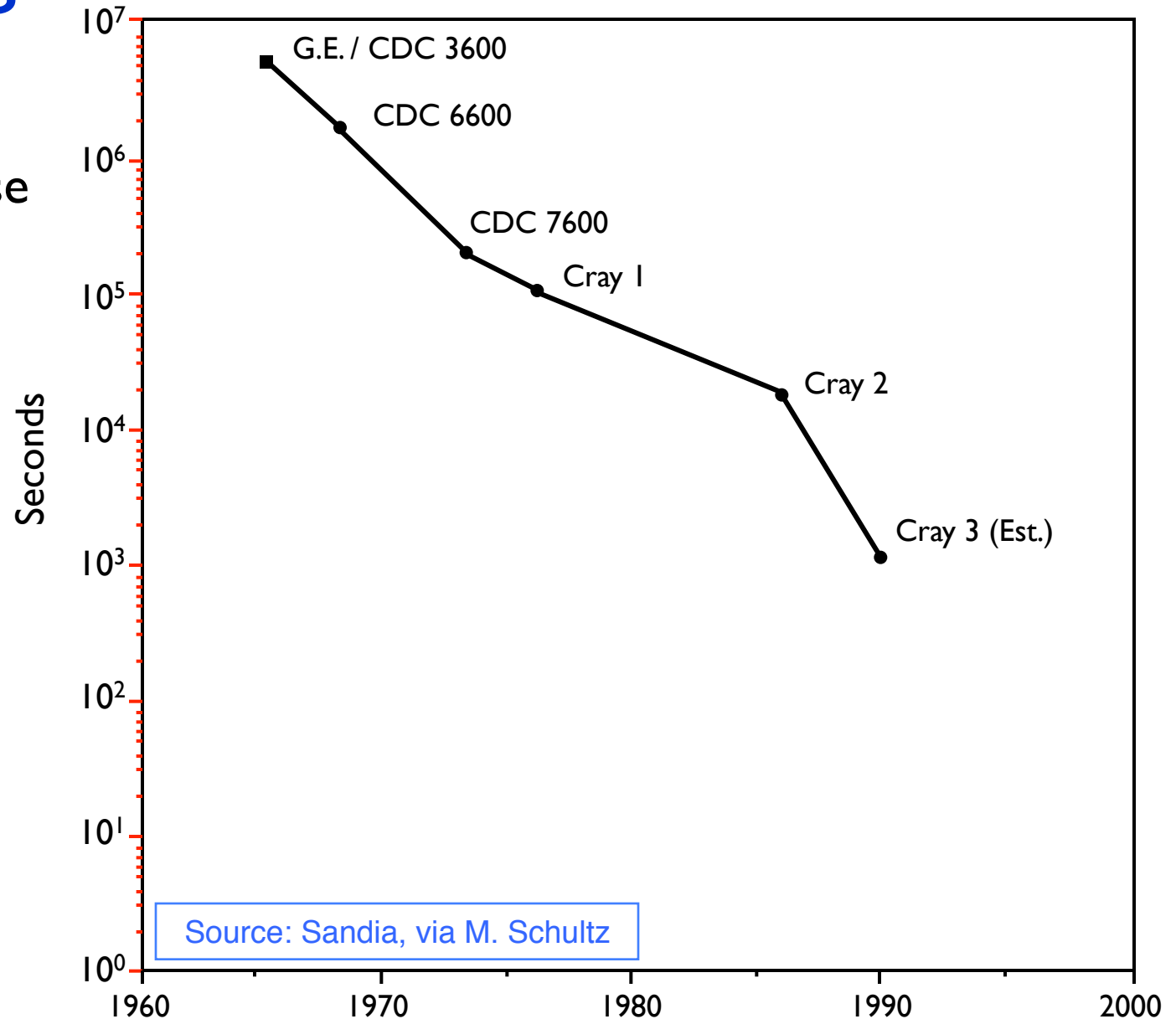
Ex: sparse linear equations over 25 years

10 orders of magnitude improvement!

Algorithms or Hardware?

25 years
progress
solving sparse
linear
systems

hardware: 4
orders of
magnitude

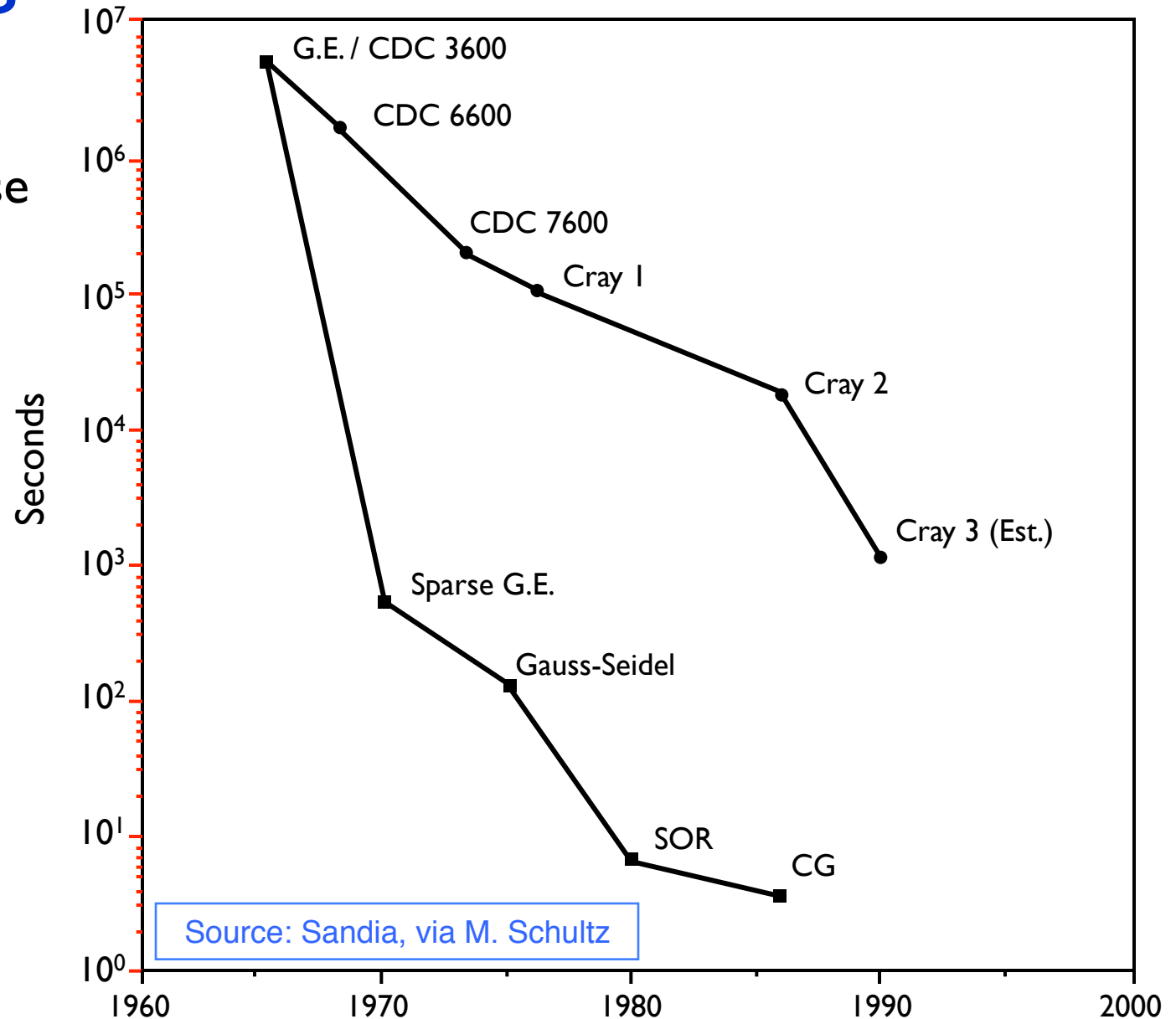


Algorithms or Hardware?

25 years
progress
solving sparse
linear
systems

hardware: 4
orders of
magnitude

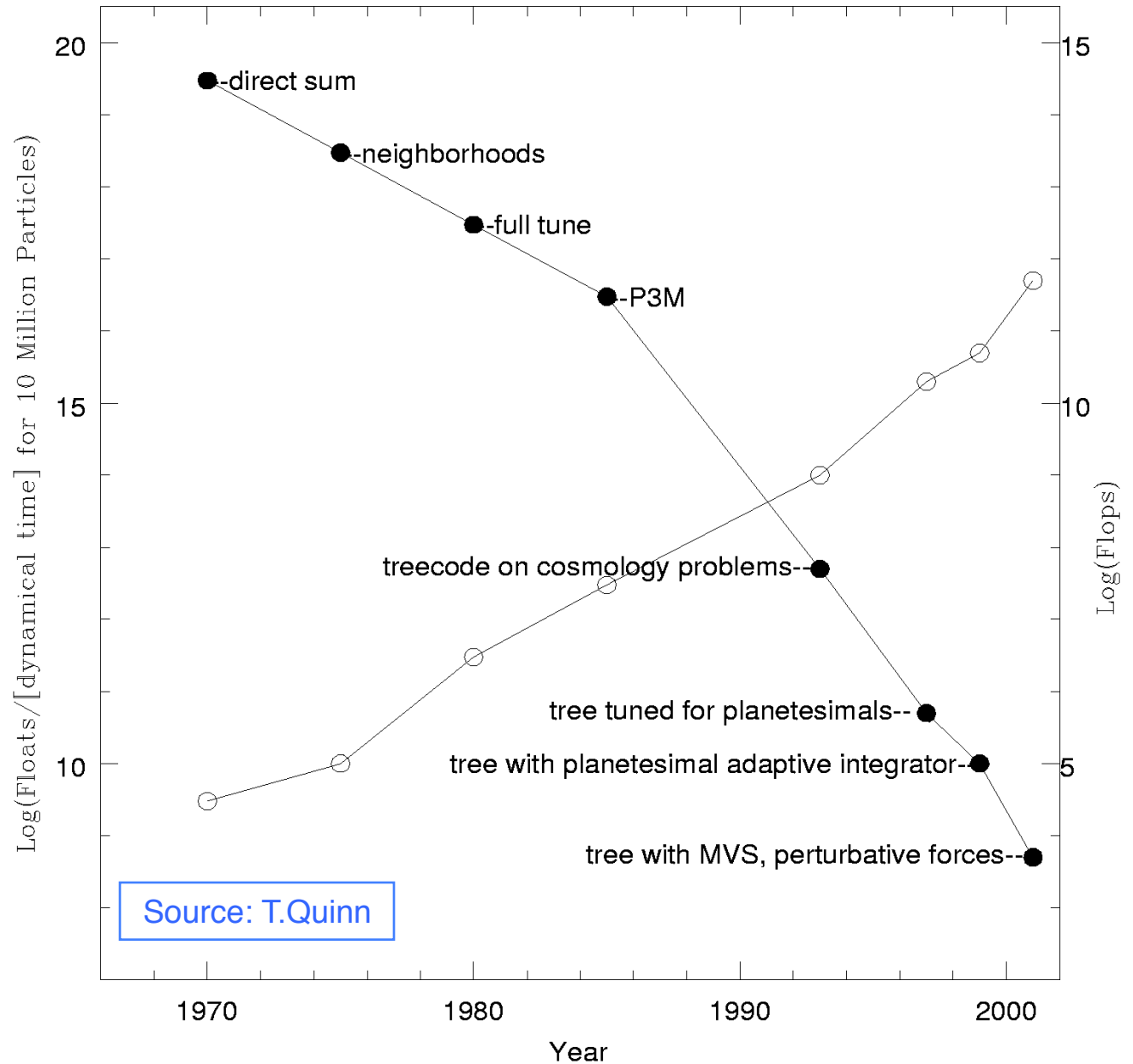
software: 6
orders of
magnitude



Algorithms or Hardware?

The
N-Body
Problem:

in 30 years
 10^7 hardware
 10^{10} software



Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them

“accomplish” via simple, well-defined steps

Ex: sorting names (via comparison)

Ex: checking for primality (via $+$, $-$, $*$, $/$, \leq)

Goals

Correctness

often subtle

Analysis

often subtle

Generality, Simplicity, 'Elegance'

Efficiency

time, memory, network bandwidth, ...

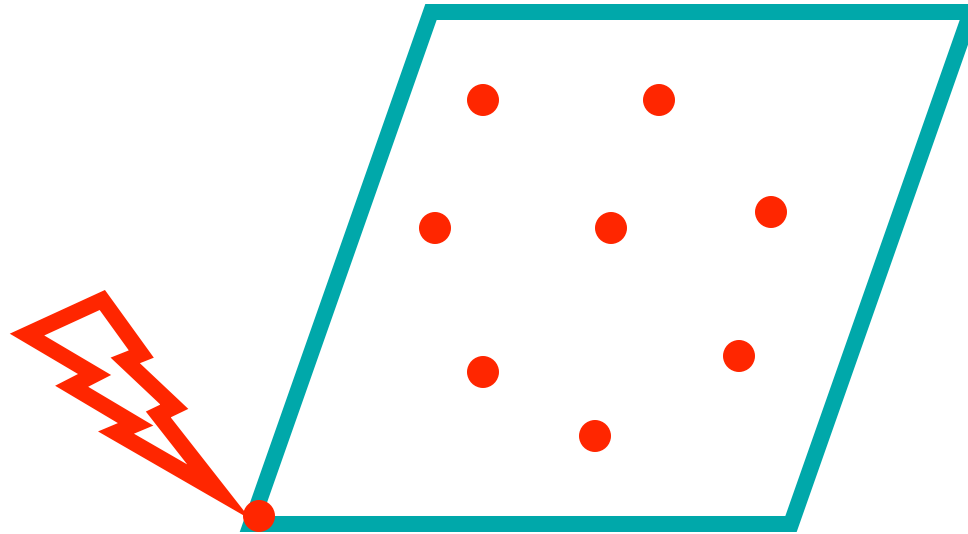
Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

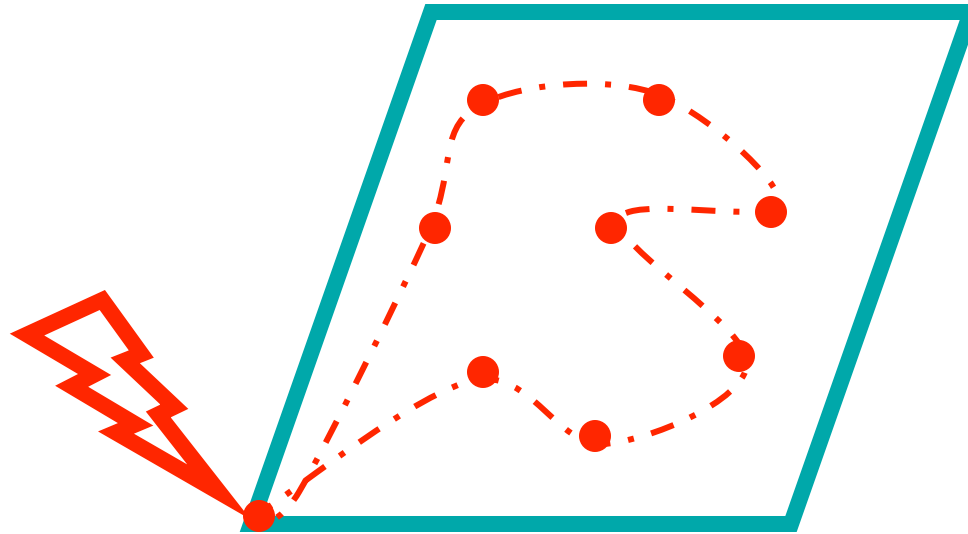
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

Printed Circuit Board



Printed Circuit Board



A Well-defined Problem

Input: Given a set S of n points in the plane

Output: The shortest cycle tour that visits each point in the set S .

Better known as “TSP”

How might you solve it?

Nearest Neighbor Heuristic

Start at some point p_0

Walk first to its
nearest neighbor p_1

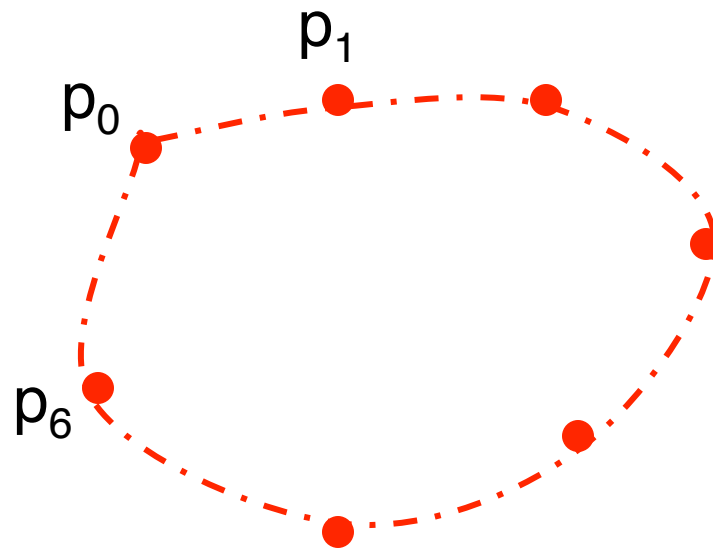
Repeatedly walk to the nearest unvisited neighbor
 p_2 , then p_3, \dots until all points have been visited

Then walk back to p_0

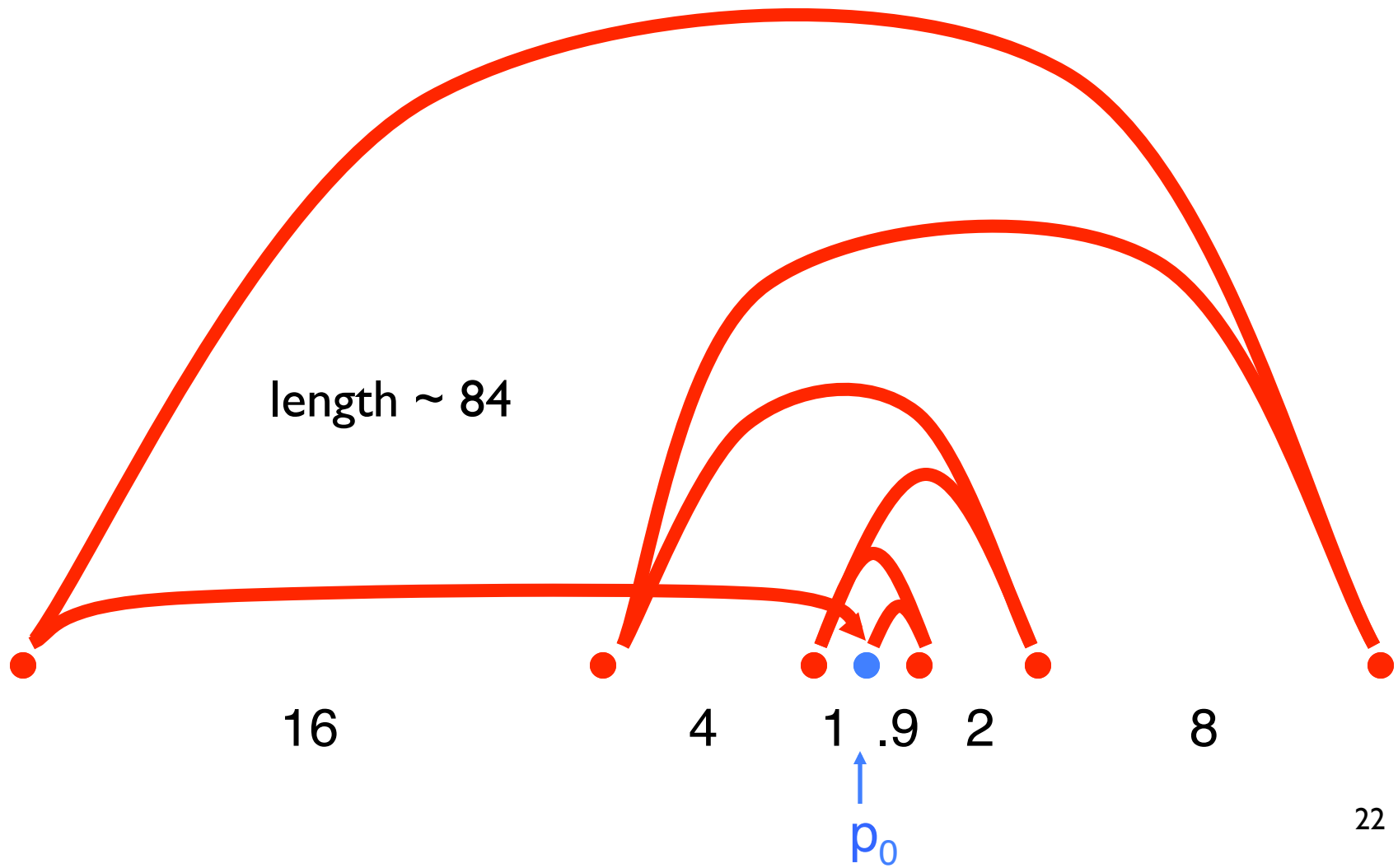
heuristic:

A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually *not* guaranteed to give the best or fastest solution.

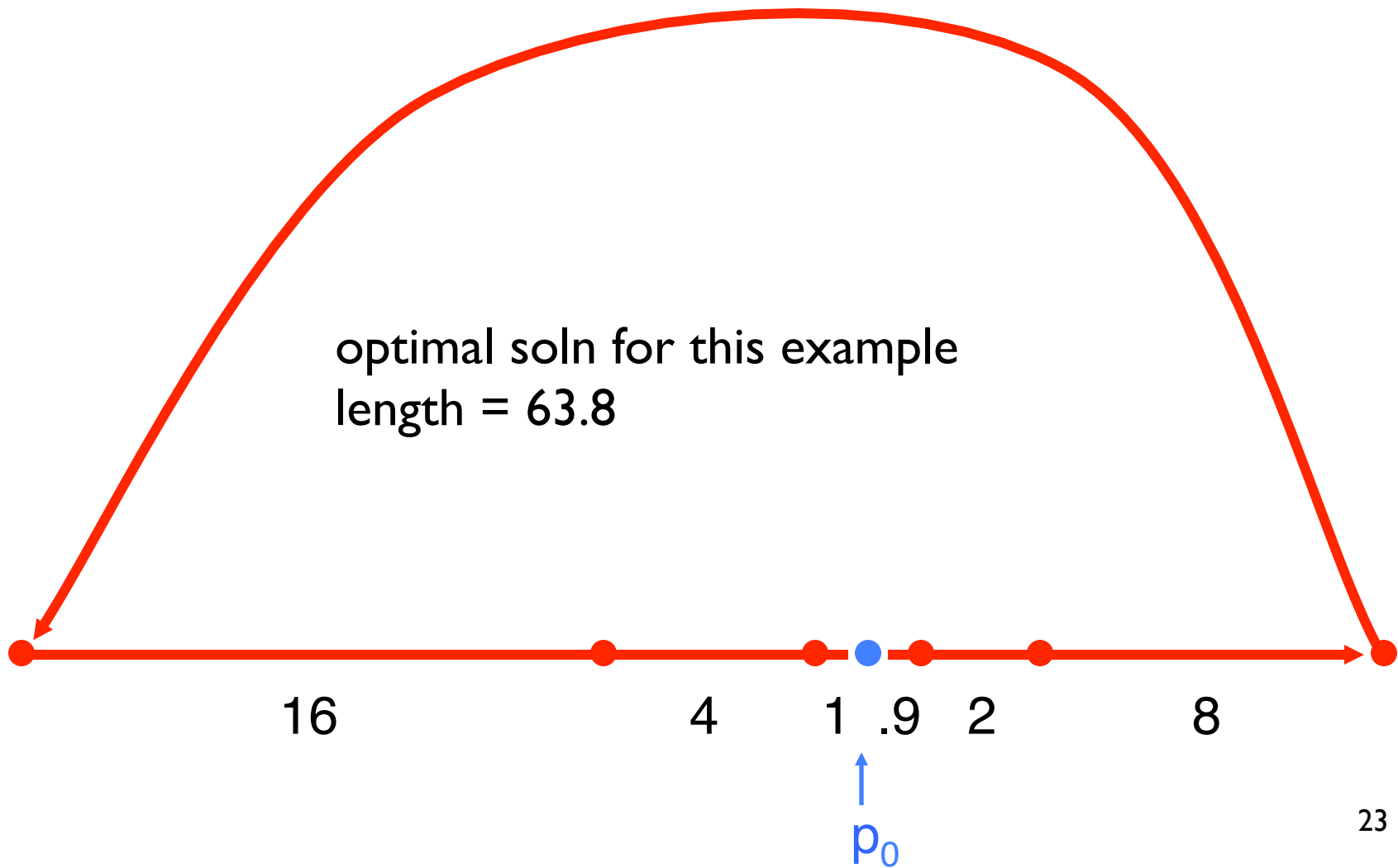
Nearest Neighbor Heuristic



An input where it works badly



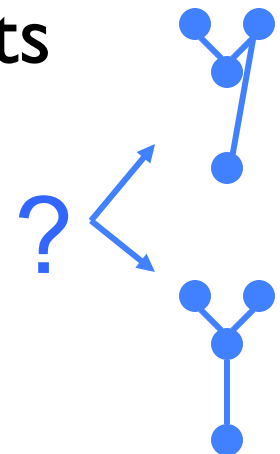
An input where it works badly



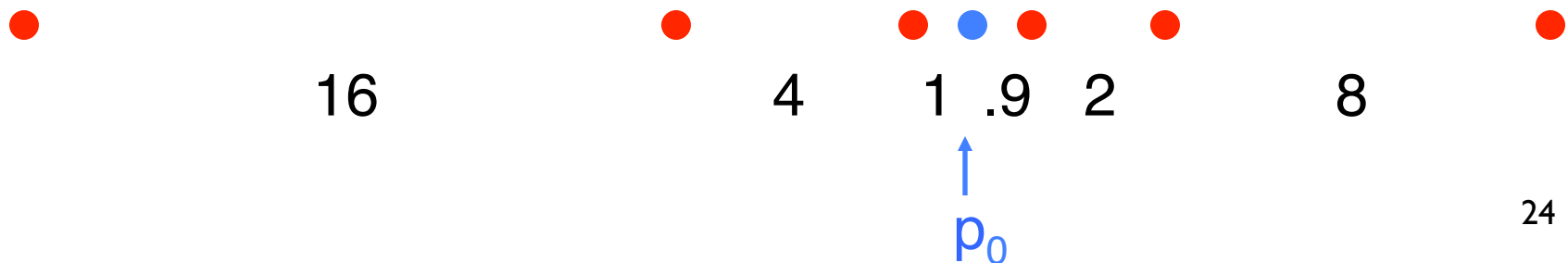
Revised idea - Closest pairs first

Repeatedly join the closest pair of points

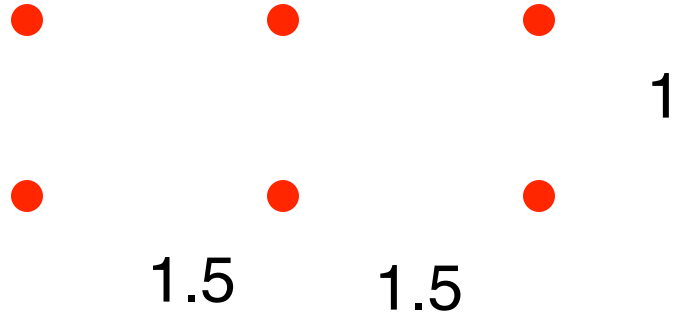
(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)



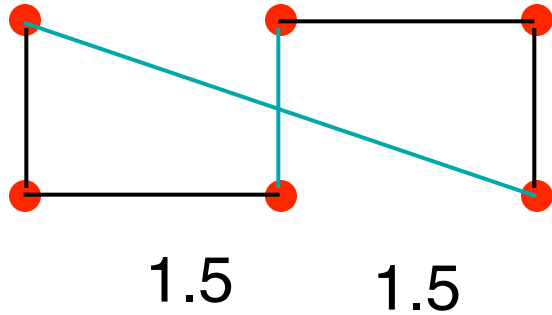
How does this work on our bad example?



Another bad example



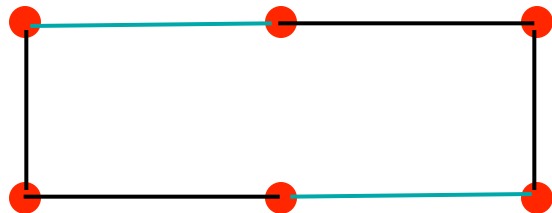
Another bad example



1

$$6 + \sqrt{10} = 9.16$$

VS



8

Something that works

“Brute Force Search”:

For each of the $n! = n(n-1)(n-2)\dots 1$ orderings of the points, check the length of the cycle you get

Keep the best one

Two Notes

The two *incorrect* algorithms were greedy

Often very natural & tempting ideas

They make choices that look great “locally” (and never reconsider them)

When greed works, the algorithms are typically efficient

BUT: often does not work - you get boxed in

Our correct alg avoids this, but is incredibly slow

20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)

And *growing*: $n! \sim \sqrt{2\pi n} \cdot (n/e)^n \sim 2^{O(n \log n)}$

The Morals of the Story

Algorithms are important

Many performance gains outstrip Moore's law

Simple problems can be hard

Factoring, TSP

Simple ideas don't always work

Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

Brute-force factoring, TSP

Changing your objective can be good

Guaranteed approximation for TSP

And: for some problems, even the *best* algorithms are slow