# CSE 421 Algorithms

Richard Anderson Lecture 22 Network Flow

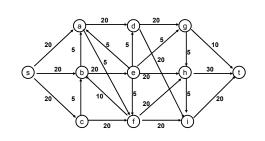
#### Outline

- · Network flow definitions
- · Flow examples
- · Augmenting Paths
- · Residual Graph
- · Ford Fulkerson Algorithm
- Cuts
- · Maxflow-MinCut Theorem

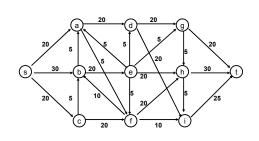
#### **Network Flow Definitions**

- Flowgraph: Directed graph with distinguished vertices s (source) and t (sink)
- Capacities on the edges, c(e) >= 0
- Problem, assign flows f(e) to the edges such that:
  - $0 \le f(e) \le c(e)$
  - Flow is conserved at vertices other than s and t
    - Flow conservation: flow going into a vertex equals the flow going out
  - The flow leaving the source is a large as possible

# Flow Example

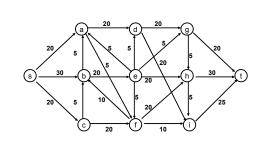


#### Find a maximum flow



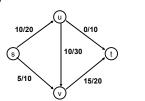
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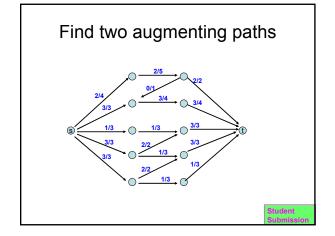
#### Find a maximum flow



### Augmenting Path Algorithm

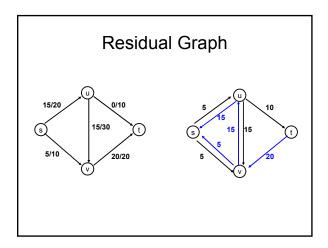
- Augmenting path
  - Vertices v<sub>1</sub>,v<sub>2</sub>,...,v<sub>k</sub>
    - $v_1 = s, v_k = t$
    - Possible to add b units of flow between v<sub>j</sub> and v<sub>j+1</sub> for j = 1 ... k-1



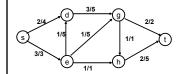


## Residual Graph

- Flow graph showing the remaining capacity
- Flow graph G, Residual Graph G<sub>R</sub>
  - G: edge e from u to v with capacity c and flow f
  - $-G_R$ : edge e' from u to v with capacity c-f
  - $-G_R$ : edge e" from v to u with capacity f



## Build the residual graph



Student

# Augmenting Path Lemma

- Let P = v<sub>1</sub>, v<sub>2</sub>, ..., v<sub>k</sub> be a path from s to t with minimum capacity b in the residual graph.
- b units of flow can be added along the path P in the flow graph.

