CSE 421
Algorithms
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Lecture 22
Network Flow

## Outline

- Network flow definitions
- Flow examples
- Augmenting Paths
- Residual Graph
- Ford Fulkerson Algorithm
- Cuts
- Maxflow-MinCut Theorem


## Network Flow Definitions

- Flowgraph: Directed graph with distinguished vertices s (source) and t (sink)
- Capacities on the edges, $c(e)>=0$
- Problem, assign flows $f(e)$ to the edges such that:
- $0<=\mathrm{f}(\mathrm{e})<=\mathrm{c}(\mathrm{e})$
- Flow is conserved at vertices other than $s$ and $t$
- Flow conservation: flow going into a vertex equals the flow going out
- The flow leaving the source is a large as possible

Find a maximum flow


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Flow Example


Find a maximum flow


## Augmenting Path Algorithm

- Augmenting path
- Vertices $\mathrm{v}_{1}, \mathrm{v}_{2}, \ldots, \mathrm{v}_{\mathrm{k}}$
- $v_{1}=s, v_{k}=t$
- Possible to add $b$ units of flow between $v_{j}$ and $v_{j+1}$ for $\mathrm{j}=1$... k-1



## Residual Graph

- Flow graph showing the remaining capacity
- Flow graph G, Residual Graph $G_{R}$
- G: edge e from $u$ to $v$ with capacity $c$ and flow $f$
$-G_{R}$ : edge e' from $u$ to $v$ with capacity $c-f$
$-G_{R}$ : edge e" from $v$ to $u$ with capacity $f$


Find two augmenting paths


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## Augmenting Path Lemma

- Let $P=v_{1}, v_{2}, \ldots, v_{k}$ be a path from $s$ to $t$ with minimum capacity $b$ in the residual graph.
- b units of flow can be added along the path $P$ in the flow graph.


