

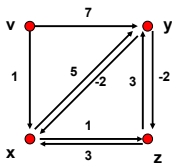
CSE 421 Algorithms

Richard Anderson
Lecture 21
Shortest Path
Network Flow Introduction

Announcements

- Friday, 11/18, Class will meet in CSE 305
- Reading 7.1-7.3, 7.5-7.6
 - Section 7.4 will not be covered

Find the shortest paths from v with exactly k edges



Express as a recurrence

- $Opt_k(w) = \min_x [Opt_{k-1}(x) + c_{xw}]$
- $Opt_0(w) = 0$ if $v=w$ and infinity otherwise

Algorithm, Version 1

```
foreach w
  M[0, w] = infinity;
M[0, v] = 0;
for i = 1 to n-1
  foreach w
    M[i, w] = min_x(M[i-1, x] + cost[x, w]);
```

Algorithm, Version 2

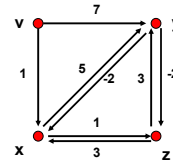
```
foreach w
  M[0, w] = infinity;
M[0, v] = 0;
for i = 1 to n-1
  foreach w
    M[i, w] = min(M[i-1, w], min_x(M[i-1, x] + cost[x, w]));
```

Algorithm, Version 3

```

foreach w
    M[w] = infinity;
M[v] = 0;
for i = 1 to n-1
    foreach w
        M[w] = min(M[w], min_x(M[x] + cost[x,w]))
    
```

Algorithm 2 vs Algorithm 3

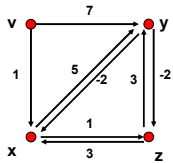


| i | v | x | y | z |
|---|---|---|---|---|
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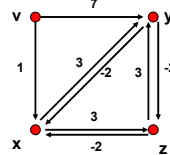
| i | v | x | y | z |
|---|---|---|---|---|
| | | | | |
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| | | | | |

Correctness Proof for Algorithm 3

- Key lemma – at the end of iteration i , for all w , $M[w] \leq M[i, w]$;
- Reconstructing the path:
 - Set $P[w] = x$, whenever $M[w]$ is updated from vertex x



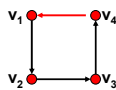
Negative Cost Cycle example



| i | v | x | y | z |
|---|---|---|---|---|
| | | | | |
| | | | | |
| | | | | |

If the pointer graph has a cycle, then the graph has a negative cost cycle

- If $P[w] = x$ then $M[w] \geq M[x] + \text{cost}(x,w)$
 - Equal after update, then $M[x]$ could be reduced
- Let v_1, v_2, \dots, v_k be a cycle in the pointer graph with (v_k, v_1) the last edge added
 - Just before the update
 - $M[v_j] \geq M[v_{j-1}] + \text{cost}(v_{j-1}, v_j)$ for $j < k$
 - $M[v_k] \geq M[v_1] + \text{cost}(v_1, v_k)$
 - Adding everything up
 - $0 > \text{cost}(v_1, v_2) + \text{cost}(v_2, v_3) + \dots + \text{cost}(v_k, v_1)$

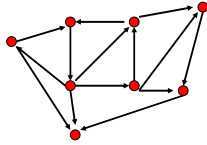


Negative Cycles

- If the pointer graph has a cycle, then the graph has a negative cycle
- Therefore: if the graph has no negative cycles, then the pointer graph has no negative cycles

Finding negative cost cycles

- What if you want to find negative cost cycles?



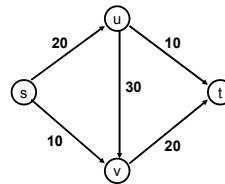
Network Flow



Network Flow Definitions

- Capacity
- Source, Sink
- Capacity Condition
- Conservation Condition
- Value of a flow

Flow Example



Residual Graph

