

# CSE 421 Algorithms

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Lecture 9  
Dijkstra's algorithm

## Who was Dijkstra?

- What were his major contributions?

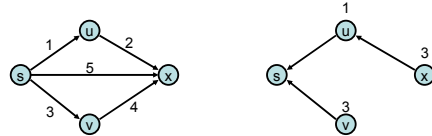
## Edsger Wybe Dijkstra

<http://www.cs.utexas.edu/users/EWD/>

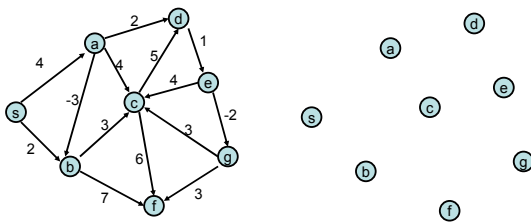


## Single Source Shortest Path Problem

- Given a graph and a start vertex  $s$ 
  - Determine distance of every vertex from  $s$
  - Identify shortest paths to each vertex
    - Express concisely as a "shortest paths tree"
    - Each vertex has a pointer to a predecessor on shortest path

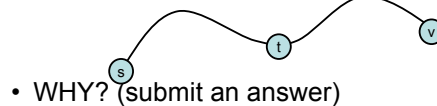


## Construct Shortest Path Tree from $s$



## Warmup

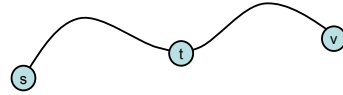
- If  $P$  is a shortest path from  $s$  to  $v$ , and if  $t$  is on the path  $P$ , the segment from  $s$  to  $t$  is a shortest path between  $s$  and  $t$



## Careful Proof

- Suppose s-v is a shortest path
- Suppose s-t is not a shortest path
- Therefore s-v is not a shortest path
- Therefore s-t is a shortest path

Prove if s-t not a shortest path then s-v is not a shortest path



## Dijkstra's Algorithm

$S = \{ \}$ ;  $d[s] = 0$ ;  $d[v] = \text{inf}$  for  $v \neq s$

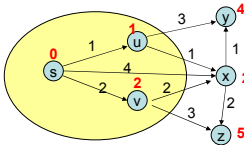
While  $S \neq V$

    Choose  $v$  in  $V-S$  with minimum  $d[v]$

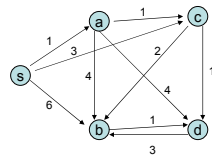
    Add  $v$  to  $S$

    For each  $w$  in the neighborhood of  $v$

$d[w] = \min(d[w], d[v] + c(v, w))$



Simulate Dijkstra's algorithm (starting from s) on the graph



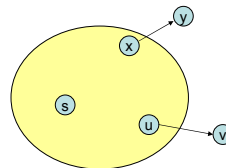
Round	Vertex Added	s	a	b	c	d
1						
2						
3						
4						
5						

## Dijkstra's Algorithm as a greedy algorithm

- Elements committed to the solution by order of minimum distance

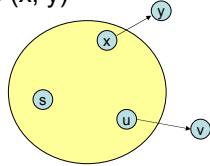
## Correctness Proof

- Elements in  $S$  have the correct label
- Key to proof: when  $v$  is added to  $S$ , it has the correct distance label.



## Proof

- Let  $P_v$  be the path of length  $d[v]$ , with an edge  $(u,v)$
- Let  $P$  be some other path to  $v$ . Suppose  $P$  first leaves  $S$  on the edge  $(x, y)$ 
  - $P = P_{sx} + c(x,y) + P_{yv}$
  - $\text{Len}(P_{sx}) + c(x,y) \geq d[x]$
  - $\text{Len}(P_{yv}) \geq 0$
  - $\text{Len}(P) \geq d[x] + 0 \geq d[v]$

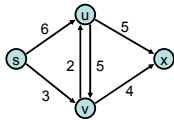


## Negative Cost Edges

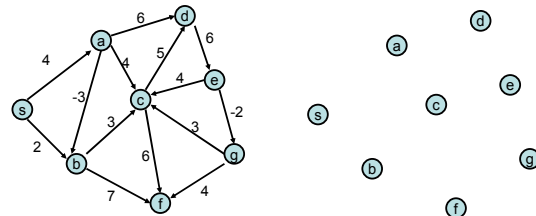
- Draw a small example a negative cost edge and show that Dijkstra's algorithm fails on this example

## Bottleneck Shortest Path

- Define the bottleneck distance for a path to be the maximum cost edge along the path



## Compute the bottleneck shortest paths



How do you adapt Dijkstra's algorithm to handle bottleneck distances

- Does the correctness proof still apply?