

# CSE 421 Algorithms

Richard Anderson  
Lecture 1

## Course Introduction

- Instructor
  - Richard Anderson, [anderson@cs.washington.edu](mailto:anderson@cs.washington.edu)
- Teaching Assistant
  - Yiannis Giotas, [giotas@cs.washington.edu](mailto:giotas@cs.washington.edu)

All of Computer Science is the  
Study of Algorithms

## Mechanics

- It's on the web
- Weekly homework
- Midterm
- Final exam
  
- Subscribe to the mailing list

## Text book

- Algorithm Design
- Jon Kleinberg, Eva Tardos
  
- Read Chapters 1 & 2

## How to study algorithms

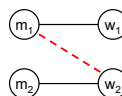
- Zoology
- Mine is faster than yours is
- Algorithmic ideas
  - Where algorithms apply
  - What makes an algorithm work
  - Algorithmic thinking

## Introductory Problem: Stable Matching

- Setting:
  - Assign TAs to Instructors
  - Avoid having TAs and Instructors wanting changes
    - E.g., Prof A. would rather have student X than her current TA, and student X would rather work for Prof A. than his current instructor.

## Formal notions

- Perfect matching
- Ranked preference lists
- Stability



## Examples

- |                  |                  |
|------------------|------------------|
| • $m_1: w_1 w_2$ | • $m_1: w_1 w_2$ |
| • $m_2: w_2 w_1$ | • $m_2: w_1 w_2$ |
| • $w_1: m_1 m_2$ | • $w_1: m_1 m_2$ |
| • $w_2: m_2 m_1$ | • $w_2: m_1 m_2$ |

## Examples

- $m_1: w_1 w_2$
- $m_2: w_2 w_1$
- $w_1: m_2 m_1$
- $w_2: m_1 m_2$

## Intuitive Idea for an Algorithm

- $m$  proposes to  $w$ 
  - If  $w$  is unmatched,  $w$  accepts
  - If  $w$  is matched to  $m_2$ 
    - If  $w$  prefers  $m$  to  $m_2$ ,  $w$  accepts
    - If  $w$  prefers  $m_2$  to  $m$ ,  $w$  rejects
- Unmatched  $m$  proposes to highest  $w$  on its preference list

## Algorithm

```

Initially all  $m$  in  $M$  and  $w$  in  $W$  are free
While there is a free  $m$ 
   $w$  highest on  $m$ 's list that  $m$  has not proposed to
  if  $w$  is free, then match ( $m$ ,  $w$ )
  else
    suppose ( $m_2$ ,  $w$ ) is matched
    if  $w$  prefers  $m$  to  $m_2$ 
      unmatched ( $m_2$ ,  $w$ )
      match ( $m$ ,  $w$ )
  
```

### Does this work?

- Does it terminate?
- Is the result a stable matching?
- Begin by identifying invariants and measures of progress
  - m's proposals get worse
  - Once w is matched, w stays matched
  - w's partners get better

Claim: The algorithm stops in at most  $n^2$  steps

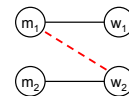
- Why?

The algorithm terminates with a perfect matching

- Why?

The resulting matching is stable

- Suppose
  - $m_1$  prefers  $w_2$  to  $w_1$
  - $w_2$  prefers  $m_1$  to  $m_2$



- How could this happen?