CSE 421
Introduction to Algorithms

Depth First Search and
Strongly Connected Components

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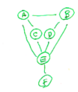
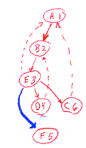
Undirected Depth-First Search



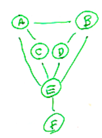


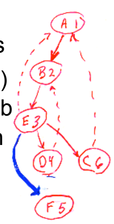
- It's not just for trees

```
DFS(v)
back edge  { if v marked then return;
             mark v; #v := ++count;
tree edge  { for all edges (v,w) do DFS(w);
Main()
    count := 0;
    for all unmarked v do DFS(v);
```

NB: book uses decreasing

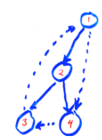
Undirected Depth-First Search





- Key Properties:
  1. No "cross-edges"; only tree- or back-edges
  2. Before returning, DFS(v) visits all vertices reachab from v via paths through previously unvisited vertices

Directed Depth First Search



- Algorithm: Unchanged
- Key Properties:
  2. Unchanged
  1'. Edge (v,w) is:

As before	Tree-edge	if w unvisited
	Back-edge	if w visited, #w<#v, on stack
New	Cross-edge	if w visited, #w<#v, not on stack
	Forward-edge	if w visited, #w>#v

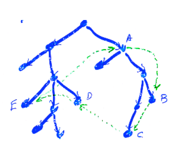
Note: Cross edges *only* go "Right" to "Left"

An Application:

G has a cycle ⇔ DFS finds a back edge

⇐ Easy - back edge (x,y) plus tree edges y, …, x form a cycle.

⇒ Why can't we have something like this?:



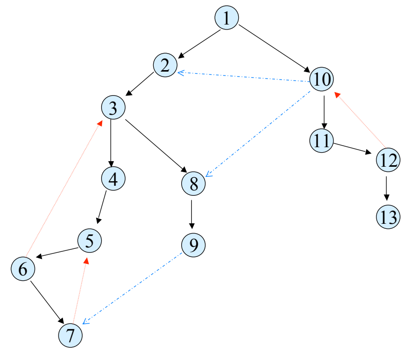
Lemma 1

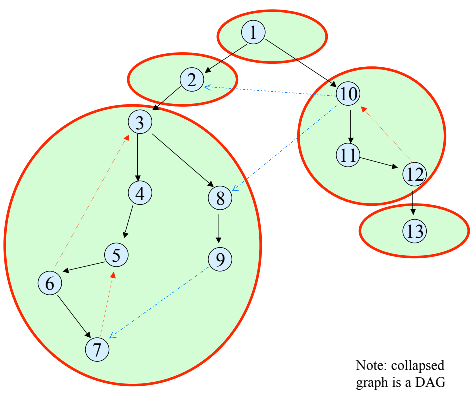
Before returning, dfs(v) visits w iff
- w is unvisited
- w is reachable from v via a path through unvisited vertices

Proof:
- dfs follows all direct out-edges
- call dfs recursively at each unvisited one
- by induction on path length, visits all

Strongly Connected Components

- <u>Defn:</u> G is *strongly connected* if for all u,v there is a (directed) path from u to v and from v to u.
  [Equivalently:
    there is a circuit through u and v.]
- <u>Defn:</u> a *strongly connected component* of G is a maximal strongly connected (vertex-induced) subgraph.





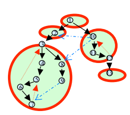
Note: collapsed graph is a DAG

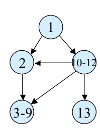
Uses for SCC's

- Optimizing compilers:
  - SCC's in program flow graph = loops
  - SCC's in call graph = mutual recursion
- Operating Systems: If (u,v) means process u is waiting for process v, SCC's show deadlocks.
- Econometrics: SCC's might show highly interdependent sectors of the economy.
- Etc.

Directed Acyclic Graphs

- If we collapse each SCC to a single vertex we get a directed graph with no cycles
  - a **directed acyclic graph** or **DAG**
- Many problems on directed graphs can be solved as follows:
  - Compute SCC's and resulting DAG
  - Do one computation on each SCC
  - Do another on the overall DAG
  - Example: Spreadsheet evaluation





Two Simple SCC Algorithms

- u,v in same SCC iff there are paths u → v & v → u
- Transitive closure: O(n^3)
- DFS from every u, v: O(ne) = O(n^3)

Goal:

- Find all Strongly Connected Components in linear time, i.e., time O(n+e)

(Tarjan, 1972)

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Definition

The *root* of an SCC is the first vertex in it visited by DFS.

Equivalently, the root is the vertex in the SCC with the smallest DFS number.

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Lemma 2

Exercise: show that each SCC is a *contiguous* subtree.

All members of an SCC are descendants of its root.

Proof:
- all members are reachable from all others
- so, all are reachable from its root
- all are unvisited when root is visited
- so, all are descendants of its root (Lemma 1)

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Subgoal

- Can we identify some root?
- How about the root of the first SCC completely explored (returned from) by DFS?
- Key idea: no exit from first SCC
  (first SCC is leftmost "leaf" in collapsed DAG)

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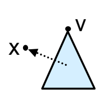
Definition



x is an *exit* from v (from v's subtree) if
- x is not a descendant of v, but
- x is the head of a (cross- or back-) edge from a descendant of v (including v itself)

NOTE: #x < #v

Ex: node #1 cannot have an exit.

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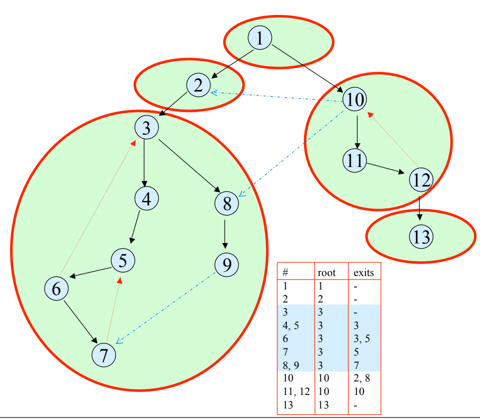

#	root	exits
1	1	-
2	2	-
3	3	-
4, 5	3	3
6	3	3, 5
7	3	5
8, 9	3	7
10	10	2, 8
11, 12	10	10
13	13	-

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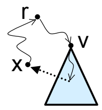
Lemma 3: Nonroots have exits



Idea: Follow cycle to root

If v is not a root, then v has an exit.

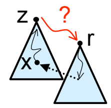
Proof:

- let r be root of v's SCC
- r is a proper ancestor of v (Lemma 2)
- let x be the first vertex that is not a descendant of v on a path v → r .
- x is an exit

Cor (contrapositive): If v has no exit, then v is a root.

NB: converse not true; some roots do have exits

Lemma 4: No Escaping 1st Root



Idea: Exit ⇒ Bigger Cycle

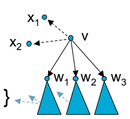
If r is the first root from which dfs returns, then r has no exit

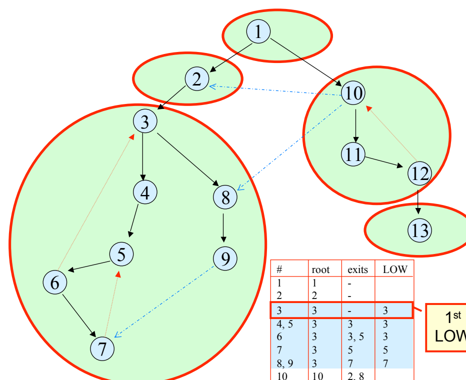
Proof (by contradiction):

- Suppose x is an exit
- let z be root of x's SCC
- r not reachable from z, else in same SCC
- #z ≤ #x (z ancestor of x; Lemma 2)
- #x < #r (x is an exit from r)
- #z < #r, no z → r path, so return from z first
- Contradiction

How to Find Exits (in 1st component)

- All exits x from v have #x < #v
- Suffices to find any of them, e.g. min #
- Defn:
  LOW(v) = min({ #x | x an exit from v} ∪ {#v})
- Calculate inductively:
  LOW(v) = min of:
  - #v
  - { LOW(w) | w a child of v }
  - { #x | (v,x) is a back- or cross-edge }
- 1st root : LOW(v)=v





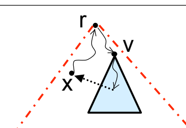
#	root	exits	LOW
1	1	-	
2	2	-	
3	3	-	3
4, 5	3	3	3
6	3	3, 5	3
7	3	5	5
8, 9	3	7	7
10	10	2, 8	
11, 12	10	10	
13	13	-	

1st root: LOW(v)=v

Finding Other Components

- Key idea: No exit from
  - 1st SCC
  - 2nd SCC, except maybe to 1st
  - 3rd SCC, except maybe to 1st and/or 2nd
  - ...

Lemma 3'



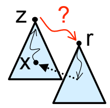
If v is not a root, then v has an exit in v's SCC.

Proof:

- let r be root of v's SCC
- r is a proper ancestor of v (Lemma 2)
- let x be the first vertex that is not a descendant of v on a path v → r .
- x is an exit in v's SCC

Cor: If v has no exit in v's SCC, then v is a root.

Lemma 4’



If r is the ~~first~~ kth root from which dfs returns, then r has no exit except possibly to the first (k-1) components

Proof:

– Suppose x is an exit
– let z be root of x’s SCC
– r not reachable from z, else in same SCC
– #z ≤ #x (z ancestor of x; Lemma 2)
– #x < #r (x is an exit from r)
– #z < #r, no z → r path, so return from z first
– ~~Contradiction~~ i.e., x in first (k-1)

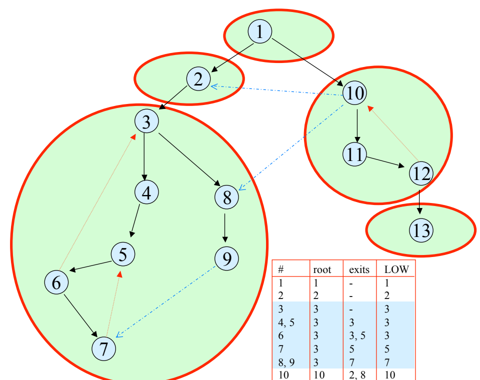
How to Find Exits (in ~~1st~~ kth component)

- All exits x from v have #x < #v
- Suffices to find any of them, e.g. min #
- Defn:
  LOW(v) = min({ #x | x an exit from v and x not in first (k-1) components } ∪ {#v})
- Calculate inductively:
  LOW(v) = min of:
  – #v
  – { LOW(w) | w a child of v }
  – { #x | (v,x) is a back- or cross-edge }

SCC Algorithm

#v = DFS number
v.low = LOW(v)
v.scc = component #

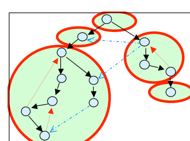
```
SCC(v)
   #v = vertex_number++; v.low = #v; push(v)
   for all edges (v,w)
     if #w == 0 then
        SCC(w); v.low = min(v.low, w.low) // tree edge
     else if #w < #v && w.scc == 0 then
        v.low = min(v.low, #w)       // cross- or back-edge
   if #v == v.low then               // v is root of new scc
     scc#++;
     repeat
        w = pop(); w.scc = scc#;     // mark SCC members
     until w==v
```

#	root	exits	LOW
1	1	-	1
2	2	-	2
3	3	-	3
4, 5	3	3	3
6	3	3, 5	3
7	3	5	5
8, 9	3	7	7
10	10	2, 8	10
11, 12	10	10	10
13	13	-	13

Complexity

- Look at every edge once
- Look at every vertex (except via in-edge) at most once
- Time = O(n+e)



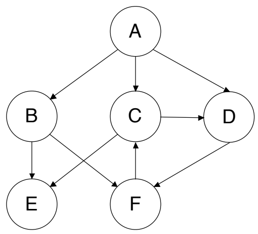
Where to start

- Unlike undirected DFS, start vertex matters
- Add “outer loop”:

  mark all vertices unvisited
  while there is unvisited vertex v do
       scc(v)

- Exercise: redo example starting from another vertex, e.g. #11 or #13 (which become #1)

Example



dfs#	v	root	exits	low(v)
1				
2				
3				
4				
5				
6				

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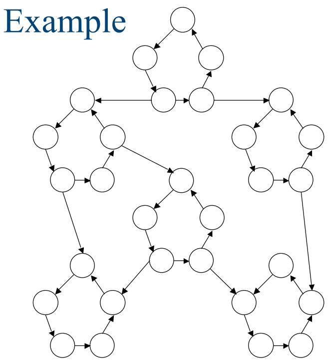
Example



v	Low(v)	v	Low(v)

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