

NP-Hard
 NP-Complete
 NP
 P
 $P = NP$

NP-Hard
 $P = NP = NP\text{-Complete}$

COMPUTERS AND INTRACTABILITY
 A GUIDE TO THE THEORY OF NP-COMPLETENESS
 Ronald L. Rivest / David S. Johnson

CSE 417

Algorithms and Complexity

Autumn 2024

Lecture 27

NP-Completeness I

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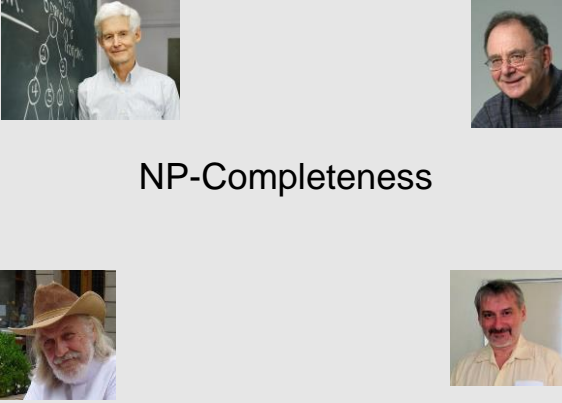
Announcements

- Homework 9
- Exam practice problems on course homepage
- Final Exam: Monday, December 9, 8:30 AM
 - One Hour Fifty Minutes
 - Closed book, no notes

Mon, Dec 2	NP-Completeness
Wed, Dec 4	NP-Completeness
Fri, Dec 6	Last Lecture: NP-Completeness and Beyond
Mon, Dec 9	Final Exam

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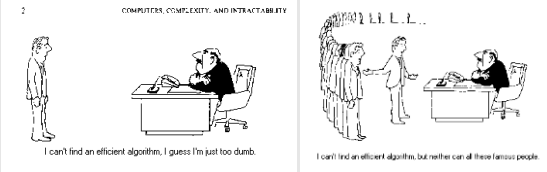
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NP-Completeness

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NP Completeness



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 COMPUTERS, COMPLEXITY, AND INTRACTABILITY
 I can't find an efficient algorithm, I guess I'm just too dumb.

I can't find an efficient algorithm, but neither can all these famous people.

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Algorithms vs. Lower bounds

- Algorithmic Theory
 - What we can compute
 - I can solve problem X with resources R
 - Proofs are almost always to give an algorithm that meets the resource bounds
- Lower bounds
 - How do we show that something can't be done?

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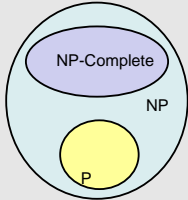
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Theory of NP Completeness

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The Universe



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Polynomial Time

- P: Class of problems that can be solved in polynomial time
 - Corresponds with problems that can be solved efficiently in practice
 - Right class to work with “theoretically”

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Decision Problems

- Theory developed in terms of yes/no problems
 - Independent set
 - Given a graph G and an integer K, does G have an independent set of size at least K
 - Shortest Path
 - Given a graph G with edge lengths, a start vertex s, and end vertex t, and an integer K, does the graph have a path between s and t of length at most K

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What is NP?

- Problems solvable in non-deterministic polynomial time . . .
- Problems where “yes” instances have polynomial time checkable certificates

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Certificate examples

- Independent set of size K
 - The Independent Set
- Satisfiable formula
 - Truth assignment to the variables
- Hamiltonian Circuit Problem
 - A cycle including all of the vertices
- K-coloring a graph
 - Assignment of colors to the vertices

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Certifiers and Certificates: 3-Satisfiability

SAT: Does a given CNF formula have a satisfying formula

Certificate: An assignment of truth values to the n boolean variables

Certifier: Check that each clause has at least one true literal,

instance s

$$(\bar{x}_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \bar{x}_2 \vee x_3) \wedge (x_1 \vee x_2 \vee x_4) \wedge (\bar{x}_1 \vee \bar{x}_3 \vee \bar{x}_4)$$

certificate t

$$x_1=1, x_2=1, x_3=0, x_4=1$$

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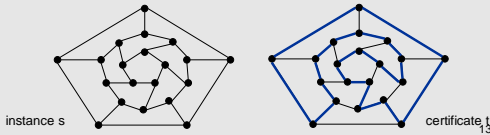
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Certifiers and Certificates: Hamiltonian Cycle

HAM-CYCLE. Given an undirected graph $G = (V, E)$, does there exist a simple cycle C that visits every node?

Certificate. A permutation of the n nodes.

Certifier. Check that the permutation contains each node in V exactly once, and that there is an edge between each pair of adjacent nodes in the permutation.



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Polynomial time reductions

- Y is Polynomial Time Reducible to X
 - Solve problem Y with a polynomial number of computation steps and a polynomial number of calls to a black box that solves X
 - Notations: $Y <_P X$
- Usually, this is converting an input of Y to an input for X , solving X , and then converting the answer back

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Composability Lemma

- If $X <_P Y$ and $Y <_P Z$ then $X <_P Z$

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Lemmas

- Suppose $Y <_P X$. If X can be solved in polynomial time, then Y can be solved in polynomial time.
- Suppose $Y <_P X$. If Y cannot be solved in polynomial time, then X cannot be solved in polynomial time.

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NP-Completeness

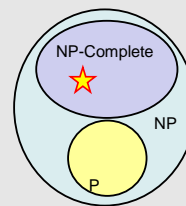
- A problem X is NP-complete if
 - X is in NP
 - For every Y in NP, $Y <_P X$
- X is a “hardest” problem in NP
- If X is NP-Complete, Z is in NP and $X <_P Z$
 - Then Z is NP-Complete

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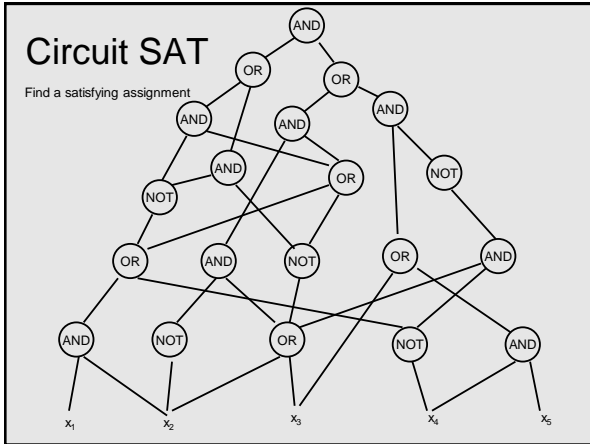
Cook’s Theorem

- There is an NP Complete problem
 - The Circuit Satisfiability Problem



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Populating the NP-Completeness Universe

- Circuit Sat \leq_p 3-SAT
- 3-SAT \leq_p Independent Set
- 3-SAT \leq_p Vertex Cover
- Independent Set \leq_p Clique
- 3-SAT \leq_p Hamiltonian Circuit
- Hamiltonian Circuit \leq_p Traveling Salesman
- 3-SAT \leq_p Integer Linear Programming
- 3-SAT \leq_p Graph Coloring
- 3-SAT \leq_p Subset Sum
- Subset Sum \leq_p Scheduling with Release times and deadlines

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Graph Coloring

- NP-Complete
 - Graph 3-coloring
- Polynomial
 - Graph 2-Coloring

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Graph 4-Coloring

- Given a graph G, can G be colored with 4 colors?
- Prove 4-Coloring is NP Complete
- Proof: 3-Coloring \leq_p 4-Coloring
- Show that you can 3-Color a graph if you have an algorithm to 4-Color a graph

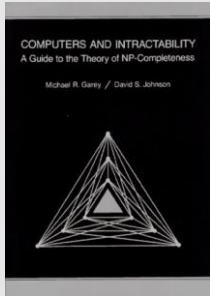
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3-Coloring \leq_p 4-Coloring

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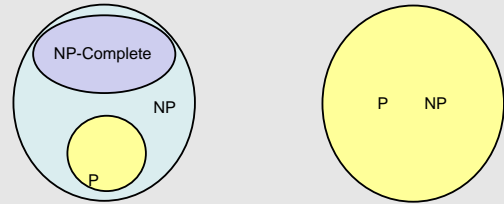
Garey and Johnson



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P vs. NP Question



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