CSE 417 Algorithms and Complexity

Lecture 19, Autumn 2024

Dynamic Programming

Announcements

- No class Monday, November 11
- Homework 7 – Due Wednesday, November 20
- Dynamic Programming Reading:
  - 6.1-6.2, Weighted Interval Scheduling
  - 6.4 Knapsack and Subset Sum
  - 6.6 String Alignment
    - 6.7* String Alignment in linear space
  - 6.8 Shortest Paths (again)
  - 6.9 Negative cost cycles
    - How to make an infinite amount of money

Dynamic Programming

- The most important algorithmic technique covered in CSE 417
- Key ideas
  - Express solution in terms of a polynomial number of sub problems
  - Order sub problems to avoid recomputation

Recursion vs Iteration

```
Factorial(n){
   if (n <= 1)
       return 1;
   else
       return n*Factorial(n-1);
}
```

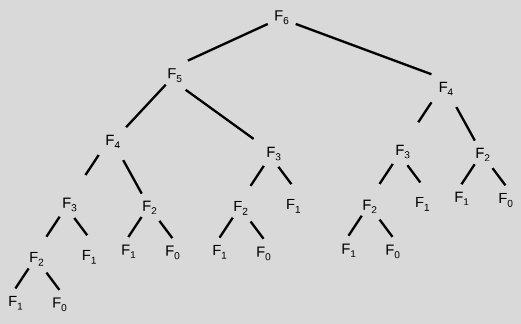
```
Factorial(n){
   v = 1;
   for (i = 2; i <= n; i++)
       v = v*I;
   return v;
}
```

Counting Rabbits

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, . . .

$F_0 = 0$; $F_1 = 1$; $F_n = F_{n-1} + F_{n-2}$

```
Fib(n){
   if (n ==  0)
       return 0;
   else if (n == 1)
       return 1;
   else
       return Fib(n-1) + Fib(n-2);
}
```



Fibonacci with Memoization

```
Fib(n){
   if (n ==  0)
       return 0;
   else if (n == 1)
       return 1;
   else
       return Fib(n-1) + Fib(n-2);
}
```

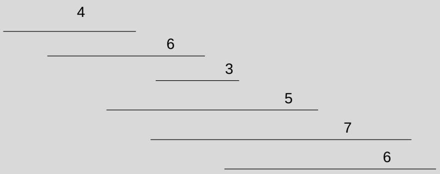
Reordering computation

```
Fib(n){
   int[ ] F = new int[n+1]

   F[0] = 0;
   F[1] = 1;
   for (i = 2; i <= n; i++)
       F[i] = F[i-1] + F[i-2];
   return F[n];
}
```

Intervals sorted by end time

Dynamic Programming

- Weighted Interval Scheduling
- Given a collection of intervals $I_1,\ldots,I_n$ with weights $w_1,\ldots,w_n$, choose a maximum weight set of non-overlapping intervals



Intervals sorted by end time

Optimality Condition

- Opt[j] is the maximum weight independent set of intervals $I_1, I_2, \ldots, I_j$
- Opt[j] = max(Opt[j – 1], w_j + Opt[p[j]])
  - Where p[j] is the index of the last interval which finishes before I_j starts

Algorithm

```
MaxValue(j) =
   if j = 0 return 0
   else
       return max( MaxValue(j-1),
                   $w_j$ + MaxValue(p[ j ]))
```

Worst case run time: 2^n

A better algorithm

M[j] initialized to -1 before the first recursive call for all j

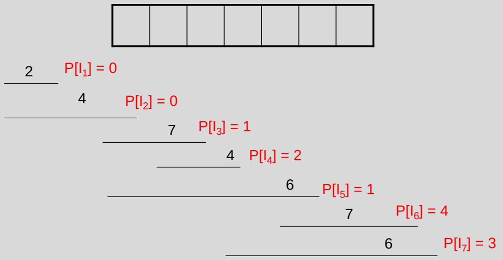
```
MaxValue(j) =
   if j = 0 return 0;
   else if M[ j ] != -1 return M[ j ];
   else {
       M[ j ] = max(MaxValue(j-1), $w_j$ + MaxValue(p[ j ]));
       return M[ j ];
   }
```

Iterative Algorithm

```
MaxValue(n){
    int[ ] M = new int[n+1];

    M[0] = 0;

    for (int i = 1; i <= n; i++){
          M[ j ] = max(M[j-1], $w_j$ + M[p[ j ]]);
    }

    return M[n];
}
```

Fill in the array with the Opt values

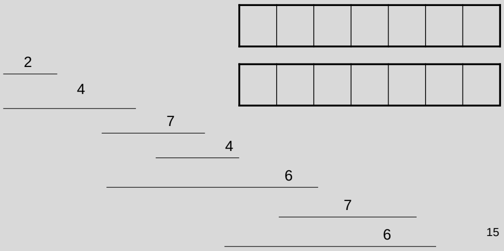
Opt[j] = max (Opt[j – 1], w_j + Opt[p[j]])



Computing the solution

Opt[j] = max (Opt[j – 1], w_j + Opt[p[j]])

Record which case is used in Opt computation



Iterative Algorithm

```
int[] M = new int[n+1];
char[] R = new char[n+1];

M[0] = 0;
for (int j = 1; j < n+1; j++){
        v1 = M[j-1];
        v2 = W[j] + M[P[j]];
        if (v1 > v2) {
                M[j] = v1;
                R[j] = ‘A’;
        }
        else {
                M[j] = v2;
                R[j] = ‘B’;
        }
}
```

Algorithm Summary

- O(n) time algorithm for finding maximum weight independent set of intervals
- Key idea: Creating an Opt function to express optimal set of $I_1, I_2, . . . I_k$ in terms of optimal set of $I_1, I_2, . . . I_{k-1}$
- Organize computation to avoid recomputation