

CSE 417: Algorithms and Computational Complexity

W. L. Ruzzo

Dynamic Programming, II
RNA Folding

Outline

A few (well, ~25) slides on *applications* of dynamic programming in biology
(You might enjoy a slightly deeper look at the use of some of the algorithms we study)

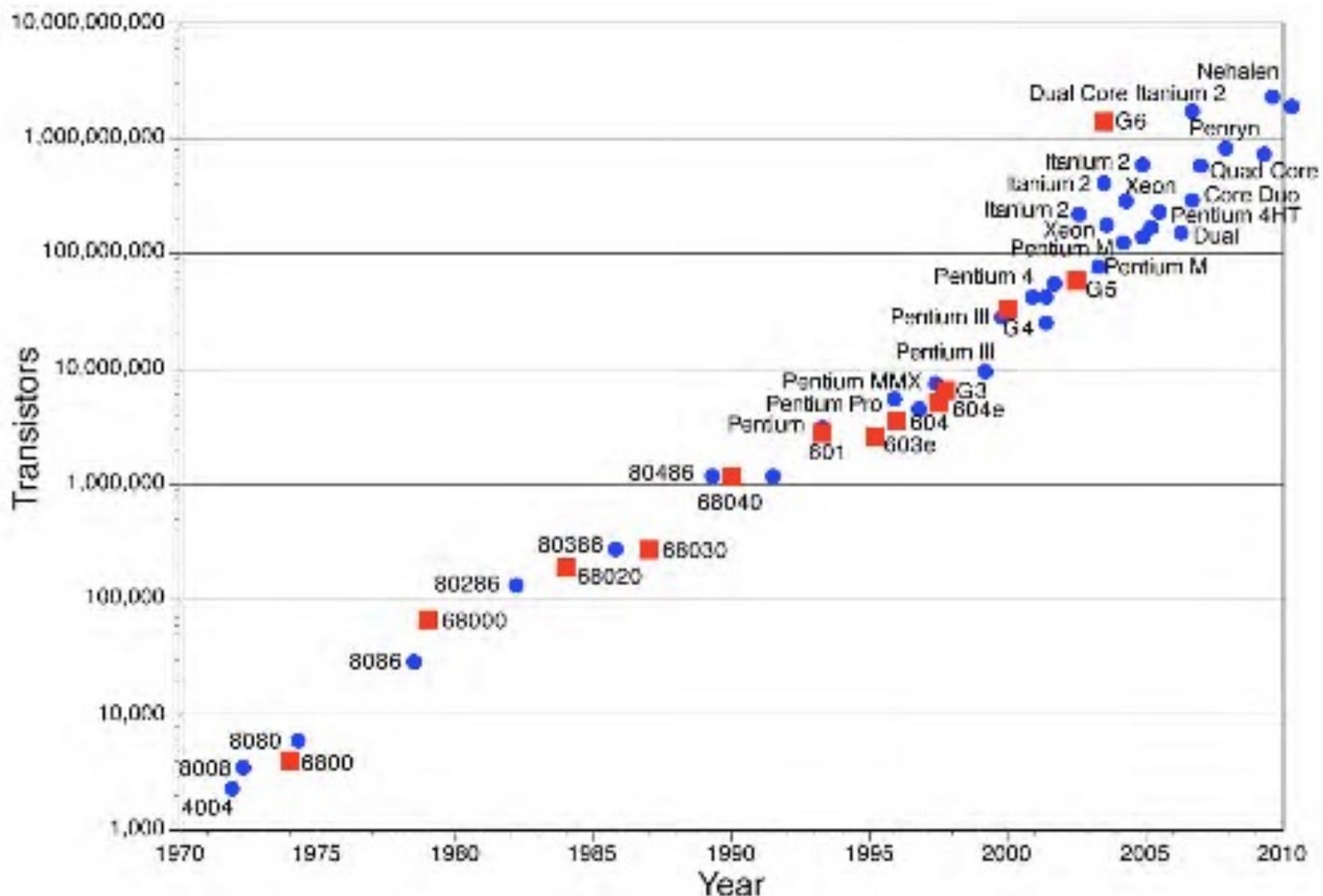
Sequence alignment

RNA structure

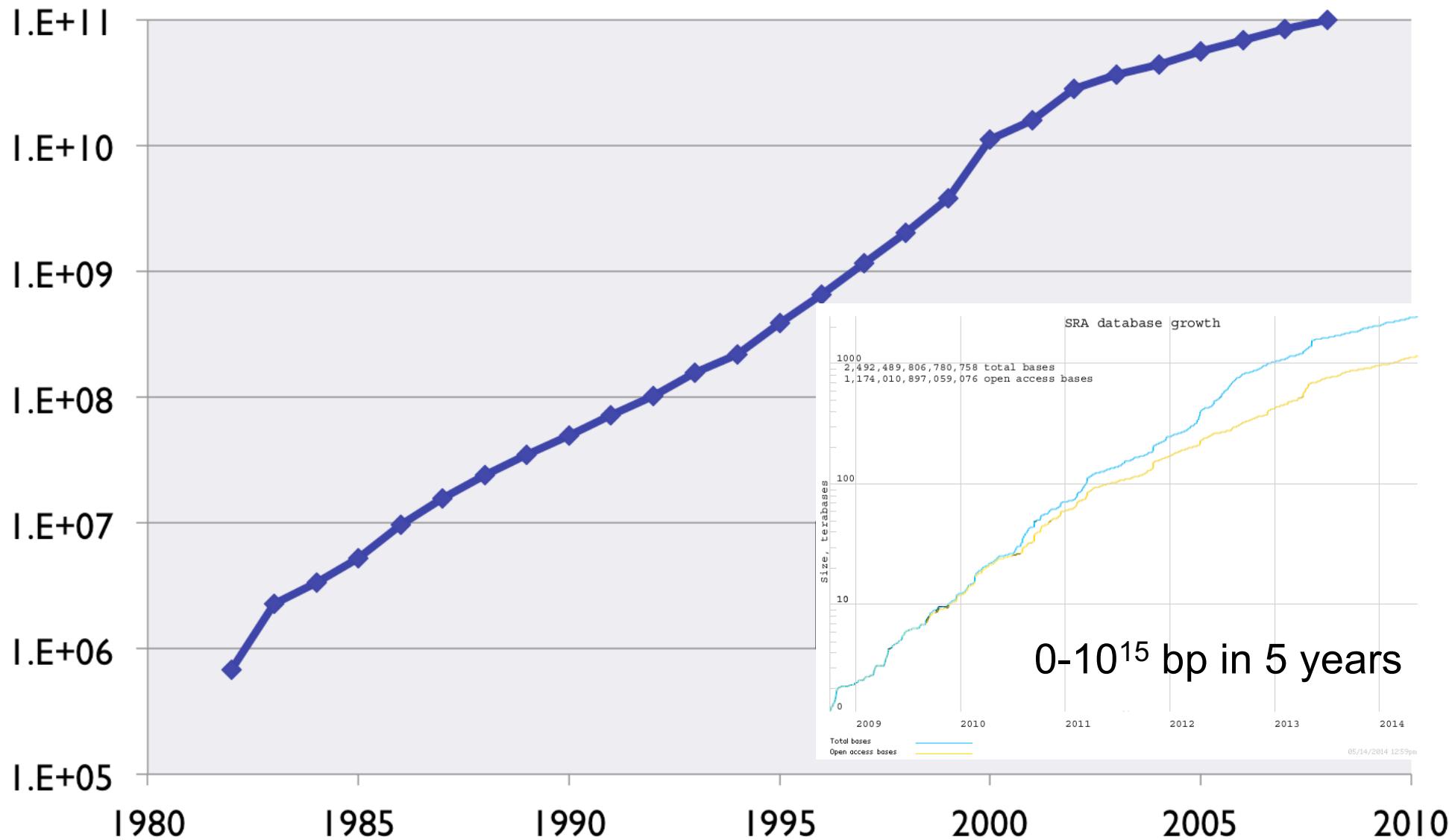
Algorithms for RNA structure (probable HW)

Application I: Sequence Search

Moore's Law

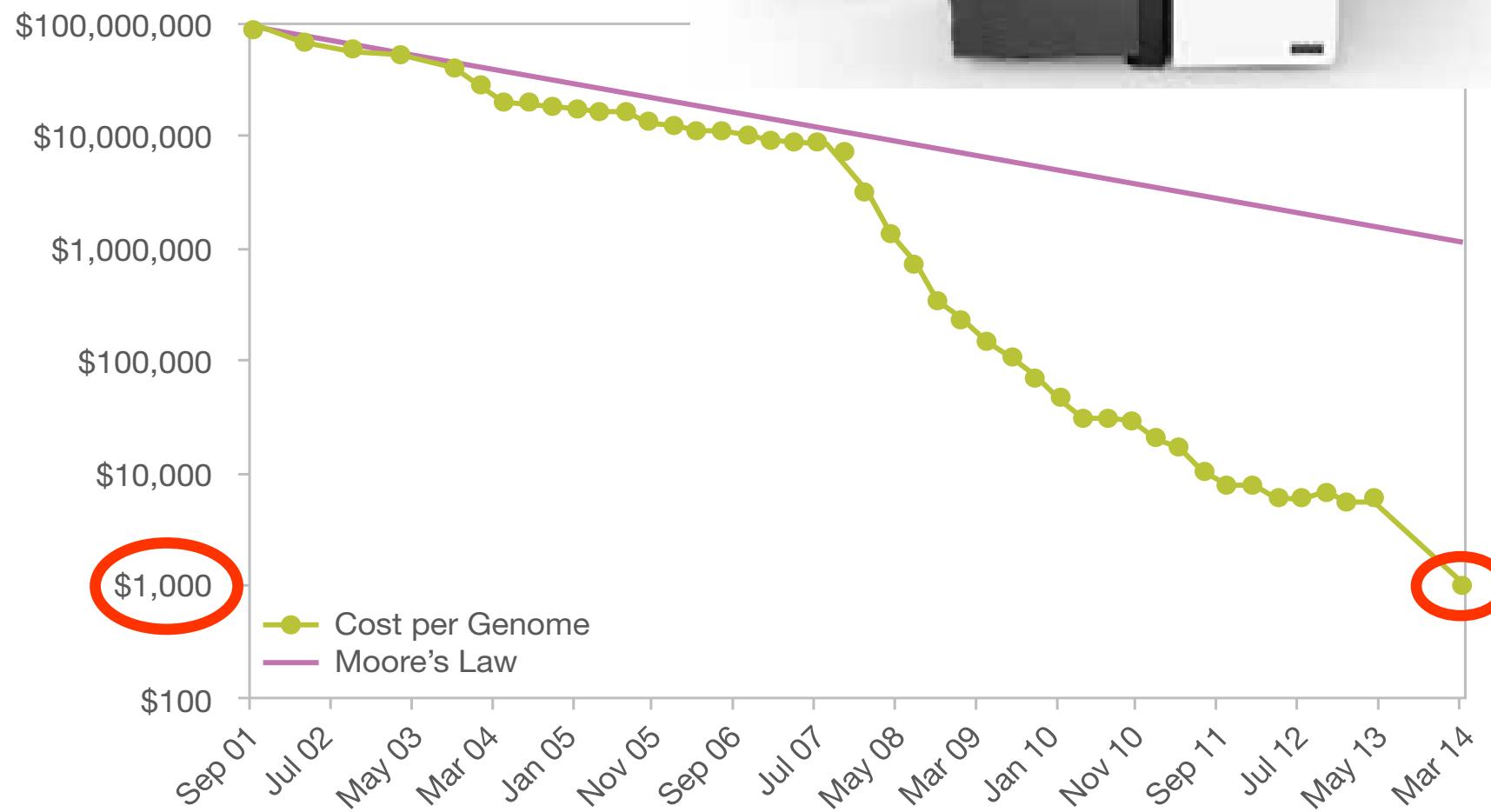


Growth of GenBank (Base Pairs)



Source: <http://www.ncbi.nlm.nih.gov/Genbank/genbankstats.html>

Sequencing Costs Outpace Moore's Law



A Database Search

go to, e.g., <http://www.uniprot.org/>, “blast” tab, and paste in this:

```
>sp|P15172|MYOD1_HUMAN Myoblast determination protein 1 OS=Homo sapiens GN=MYOD1 PE=1 SV=3
```

```
MELLSPPLRDVDLTAPDGSLCSFATTDDFYDDPCFDSPDLRFFEDLDPRLMHVGALLKPE  
EHSHFPAAVHPAPGAREDEHVRAPSGHHQAGRCLLWACKACKRKTTNADRRKAATMRERR  
RLSKVNEAFETLKRCTSSNPQNQLPKVEILRNAIRYIEGLQALLRDQDAAPPGAAAAYA  
PGPLPPGRGGEHYSGDSDASSPRSNCSDGMMDYSGPPSGARRRNCYEGAYYNEAPSEPRP  
GKSAAVSSLDCISSLVERISTESPAAPALLADVPSESPPRQEAAAPSEGESSIONGDPTQS  
PDAAPQCPAGANPNPIYQVL
```

A Few seconds Later...

Graphical overview

Color code for identity 0-
100% =



Accession	Entry name	0Query hit320	0Match hit (sqrt scale)17392	Name (Organism)
P15172	MYOD1_HUMAN			human
B2RC72	B2RC72_HUMAN			human
E2RT59	E2RT59_CANFA			dog
P49811	MYOD1_PIG			pig
D2KPI9	D2KPI9_PIG			pig
F1S9A9	F1S9A9_PIG			pig
D2I0V4	D2I0V4_AILME			panda
P29331	MYOD1_SHEEP			sheep
D2SP11	D2SP11_BUBBU			water buffalo
Q0VBX9	Q0VBX9_BOVIN			cow
Q7YS82	MYOD1_BOVIN			cow
Q8C6B1	Q8C6B1_MOUSE			mouse
A0JPK9	A0JPK9_RAT			rat
Q02346	MYOD1_RAT			rat
P10085	MYOD1_MOUSE			mouse
Q6DTY5	Q6DTY5_PIG			pig
P21572	MYOD1_COTJA			quail
Q6DV59	Q6DV59_MELGA			turkey
P16075	MYOD1_CHICK			chicken
C5J072	C5J072_CHICK			chicken
C3U0I1	C3U0I1_ANAPL			duck
F1NHM3	F1NHM3_CHICK			chicken
F1NXM5	F1NXM5_CHICK			chicken
P13904	MYODA_XENLA			frog
Q8AVZ0	Q8AVZ0_XENLA			frog
Q7T109	Q7T109_XENTR			MyoD protein (<i>Xenopus laevis</i>)
				Myod1-a protein (<i>Xenopus laevis</i>)
				MyoD protein (<i>Xenopus tropicalis</i>)

...And 1000's more...

Accession	Entry name	Status	Protein names	Organism	Length
Q7T109	Q7T109_XENTR	★	MyoD protein	Xenopus tropicalis (Western clawed frog) / Osteopilus tropicalis	288

Alignment 1 against Q7T109

Some Details from #25

And #1000 is good too!

Score	964	E-value	1.0×10^{-102}
Identity	64.0%	Positives	74.0%
Query length	320	Match length	288
Position	Q7T109 matches from 1 to 288 (288AA), in the query sequence from 1 to 320 (320AA)		

Graphical			
1	MELLSPPPLRDVDLTAPDGSLCSFATTDDFYDDPCFDSPDLRFFEDLDPRLMHVGALLKPE	60	P15172
	MELL PPLRD+++T +GSLCSF T DDFYDDPCF++ D+ FFEDLDPRL+HV ALLKPE		
1	MELLPPPLRDMEVT--EGSLCSFPTPDDFYDDPCFNTSDMSFFEDLDPRLVHV-ALLKPE	57	Q7T109
61	EHSHFPAAVHPAPGAREDEHVRAPSGHQAGRCLLWACKACKRKTTNADRRKAATMRERR	120	P15172
	+ H EDEHVRAPSGHQAGRCLLWACKACKRKTTNADRRKAATMRERR		
58	DPHH-----NEDEHVRAPSGHQAGRCLLWACKACKRKTTNADRRKAATMRERR	106	Q7T109
121	RLSKVNEAFETLKRCTSSNPNQRLPKVEILRNAIRYIEGLQALLRDQDAAPPGAAAAYFA	180	P15172
	RLSKVNEAFETLKRCTS+NPNQRLPKVEILRNAIRYIE LQ+LLR Q+ +FY		
107	RLSKVNEAFETLKRCTSTNPNQRLPKVEILRNAIRYIESLQSLLRGQE-----ESFY-	158	Q7T109
181	PGPLPPGRGGEHYSGDSDASSPRSNCSDGMDYSGPPSGARRRN CYEGAYYNEAPSEPRP	240	P15172
	P+ EHYSGDSDASSPRSNCSDGM DYS PP G+RRRN Y+ ++Y+++P+ R		
159	--PVL-----EHYSGDSDASSPRSNCSDGMTDYS-PPCGSRRRN SYDSSFYSDSPNGLRL	210	Q7T109
241	GKSAAVSSLCLSSIVERISTESP AAPALL ADVPSSESPPRRQEAAAPSEGES---SGDP	297	P15172
	GKS+ +SSLCLSSIVERISTESP P + AD SE P +P +GE+ SG		
211	GKSSVISSLCLSSIVERISTESP VCPVIPAADSGSEGSP-----CSPLQGETLSESGII	265	Q7T109 ⁹

The foregoing search capability is a *huge* deal

the “google” of molecular biology

millions of searches daily

biologists (not just “computational”
biologists) use this routinely

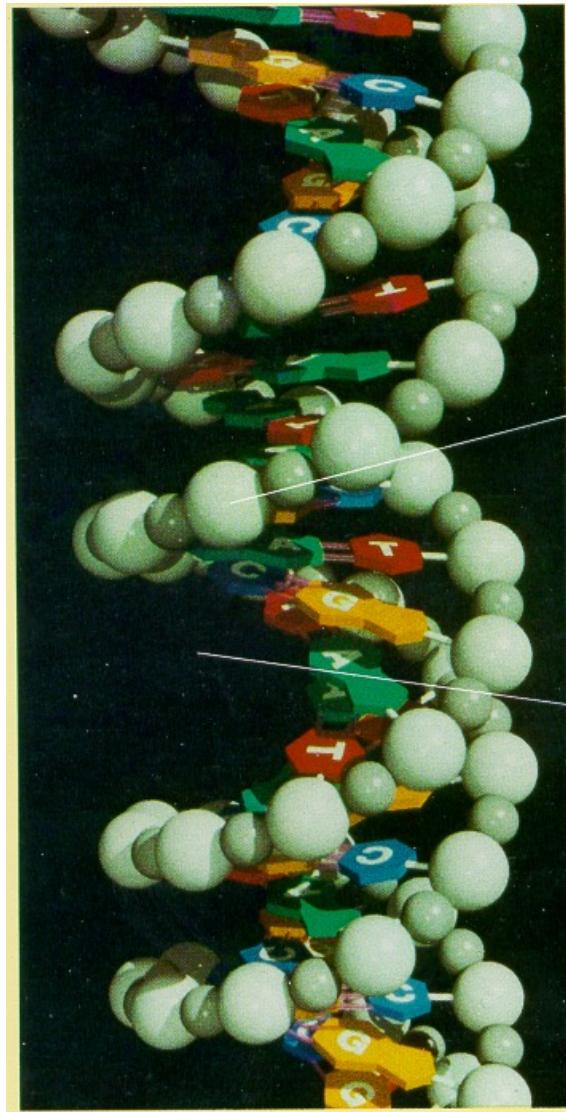
it connects information about *all* living things

(dynamic programming)

Time permitting, more on algorithm later ...

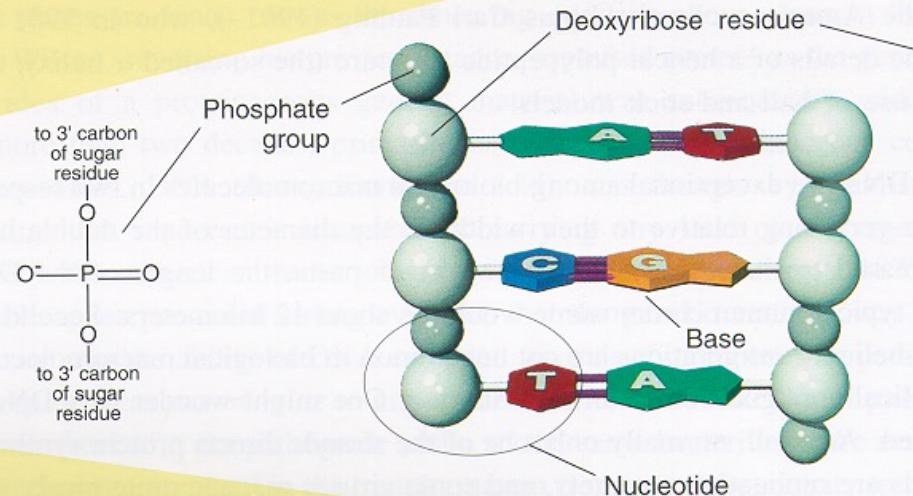
Application 2: RNA Structure

The Double Helix



(a) Computer-generated Image of DNA
(by Mel Prueitt)

(b) Uncoiled DNA Fragment



As shown, the two strands coil about each other in a fashion such that all the bases project inward toward the helix axis. The two strands are held together by hydrogen bonds (pink rods) linking each base projecting from one backbone to its so-called complementary base projecting from the other backbone. The base A always bonds to T (A and T are comple-

Shown in (b)
is an uncoiled fragment of (a
three complementary base pair.
In a chemist's viewpoint, each strand
is a polymer made up of four repeating
units called deoxyribonucleotides.

Central Dogma of Molecular Biology

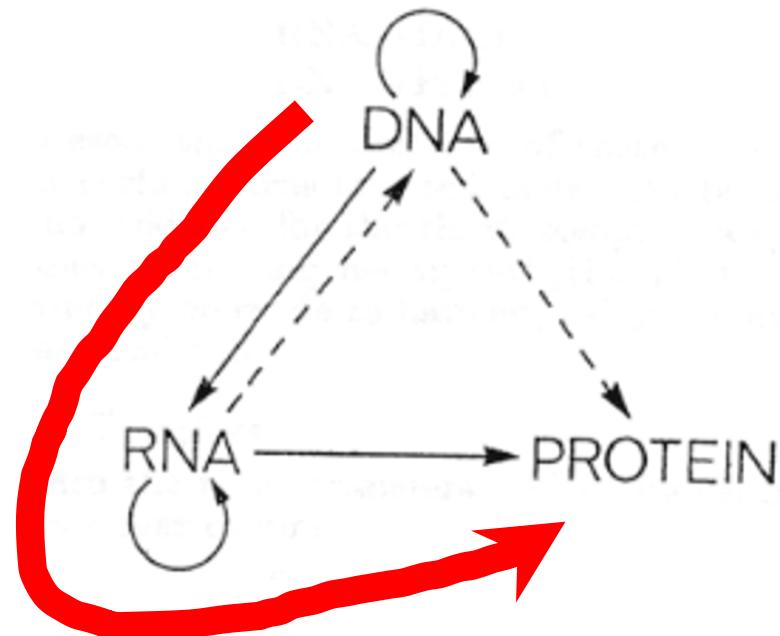
by

FRANCIS CRICK
MRC Laboratory
Hills Road,
Cambridge CB2 2QH

The central dogma of molecular biology deals with the detailed residue-by-residue transfer of sequential information. It states that such information cannot be transferred from protein to either protein or nucleic acid.

"The central dogma, enunciated by Crick in 1958 and the keystone of molecular biology ever since, is likely to prove a considerable over-simplification."

Fig. 2. The arrows show the situation as it seemed in 1958. Solid arrows represent probable transfers, dotted arrows possible transfers. The absent arrows (compare Fig. 1) represent the impossible transfers postulated by the central dogma. They are the three possible arrows starting from protein.



Non-coding RNA

Messenger RNA - codes for proteins

Non-coding RNA - all the rest

Before, say, mid 1990's, 1-2 dozen known
(critically important, but narrow roles: e.g., tRNA)

Since mid 90's dramatic discoveries

Regulation, transport, stability/degradation

E.g. "miRNA": >1000 in humans; regulate >50% of genes

E.g. "riboswitches": 10000's in bacteria

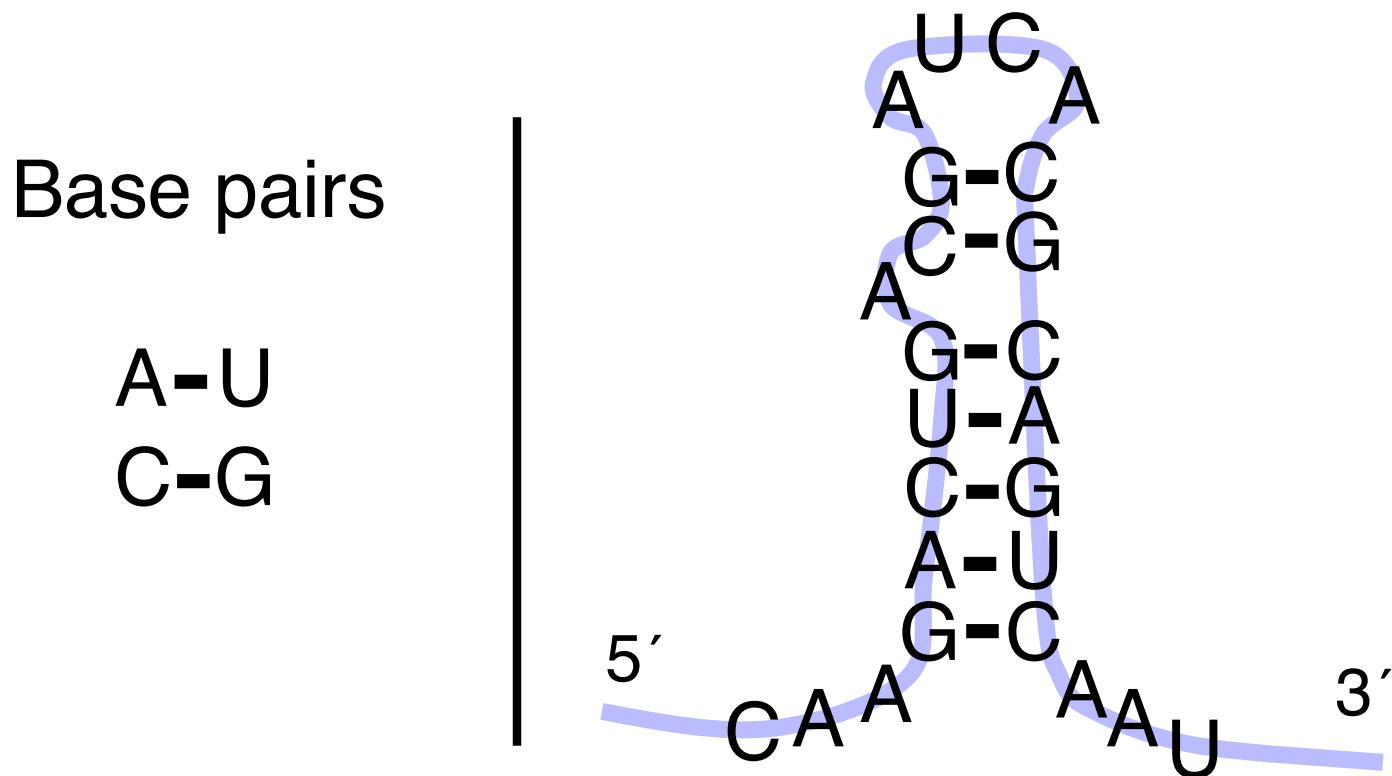
By some estimates, ncRNA >> mRNA

DNA structure: dull

5' ...ACCGCTAGATG...3'
| | | | | | | | | |
3' ...TGGCGATCTAC...5'



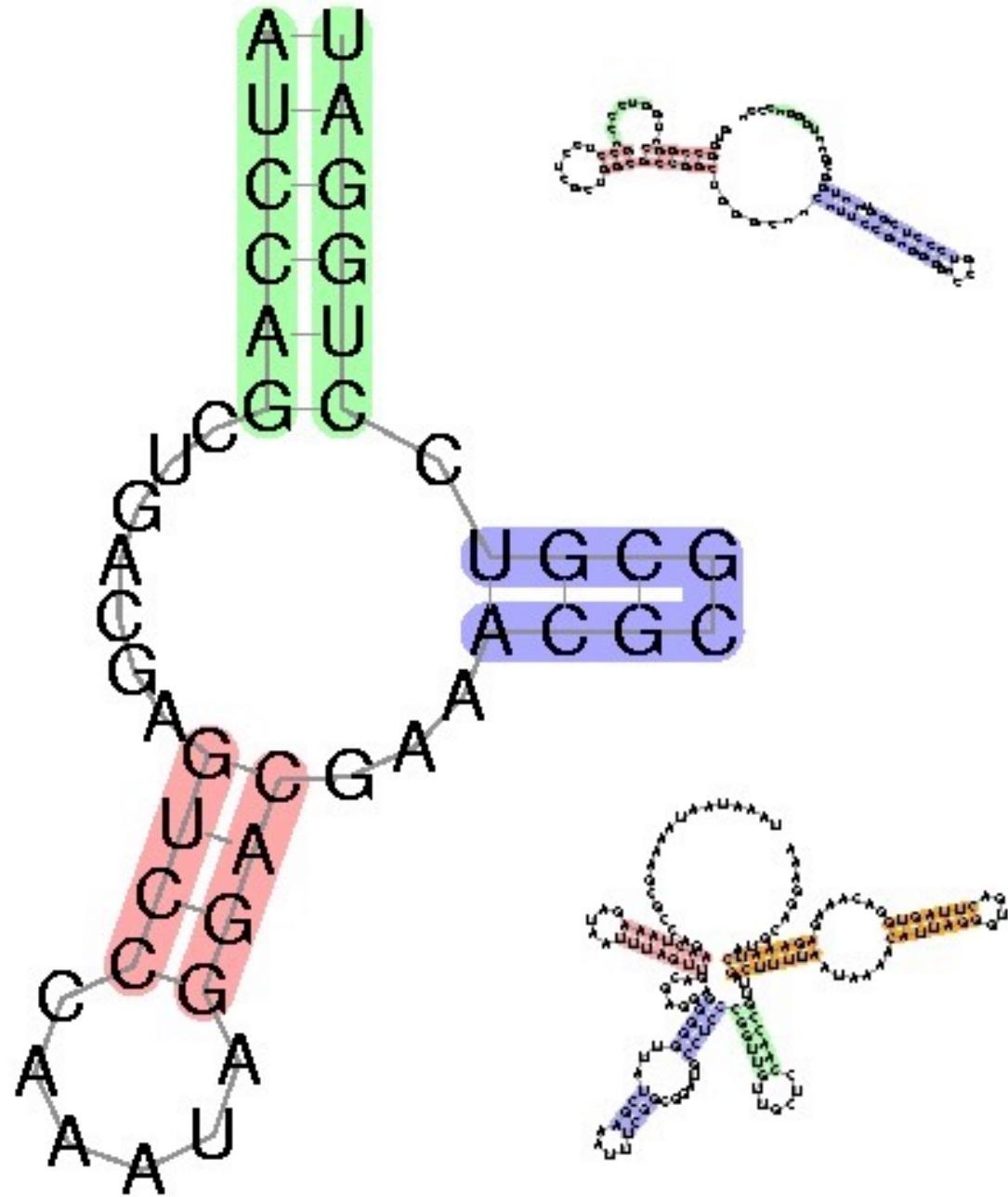
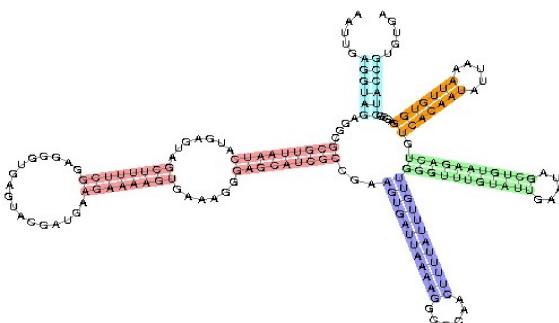
RNA Secondary Structure: RNA makes helices too



Usually *single* stranded

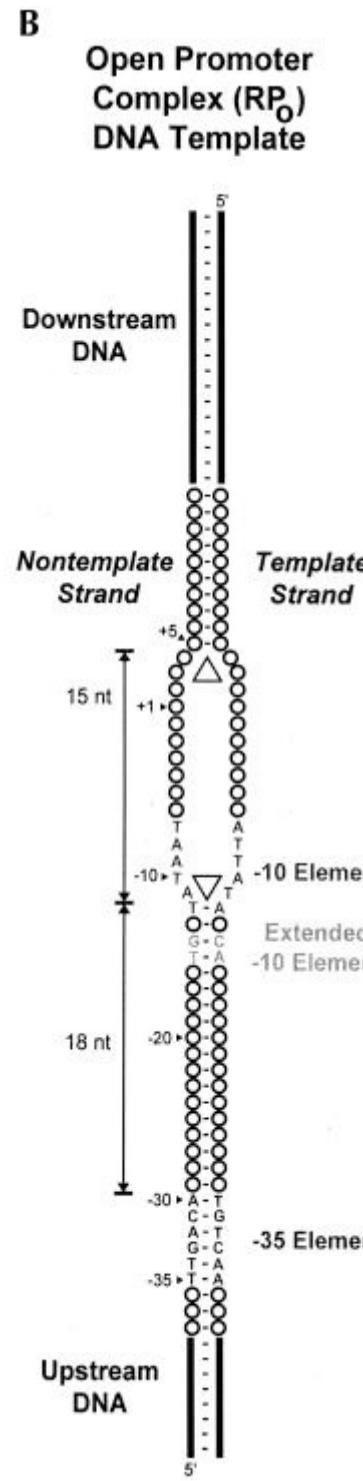
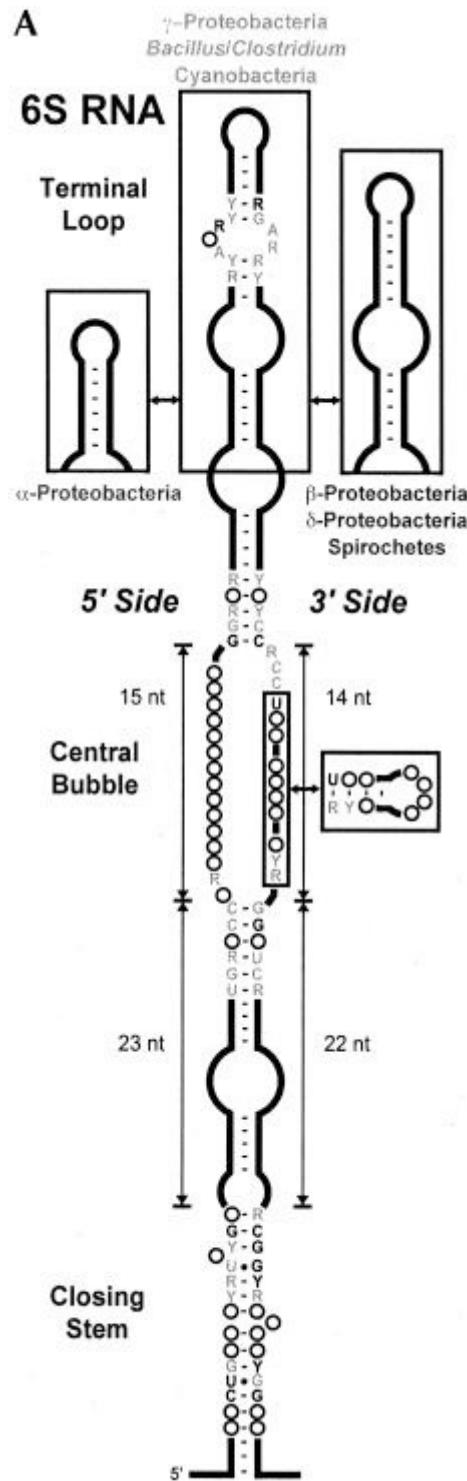
RNA Secondary Structure:

Not everything,
but important,
easier than 3d

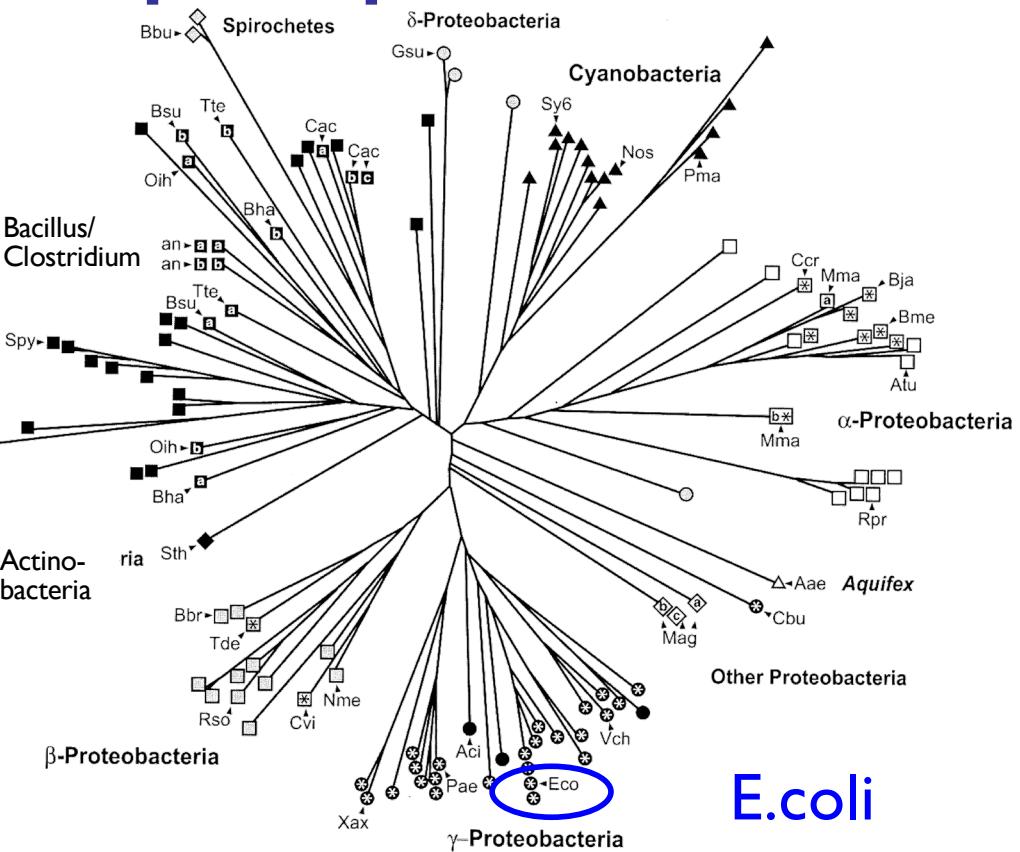


Why is structure important?

- For protein-coding, similarity in sequence is a powerful tool for finding related sequences
 - e.g. “hemoglobin,” “MyoD” and many others are easily recognized in all animals
- For many non-coding RNAs, *different sequences* can have the *same structure*, and structure is most important for function.
 - So, using structure plus sequence, can find related sequences at much greater evolutionary distances
 - 2 Examples below



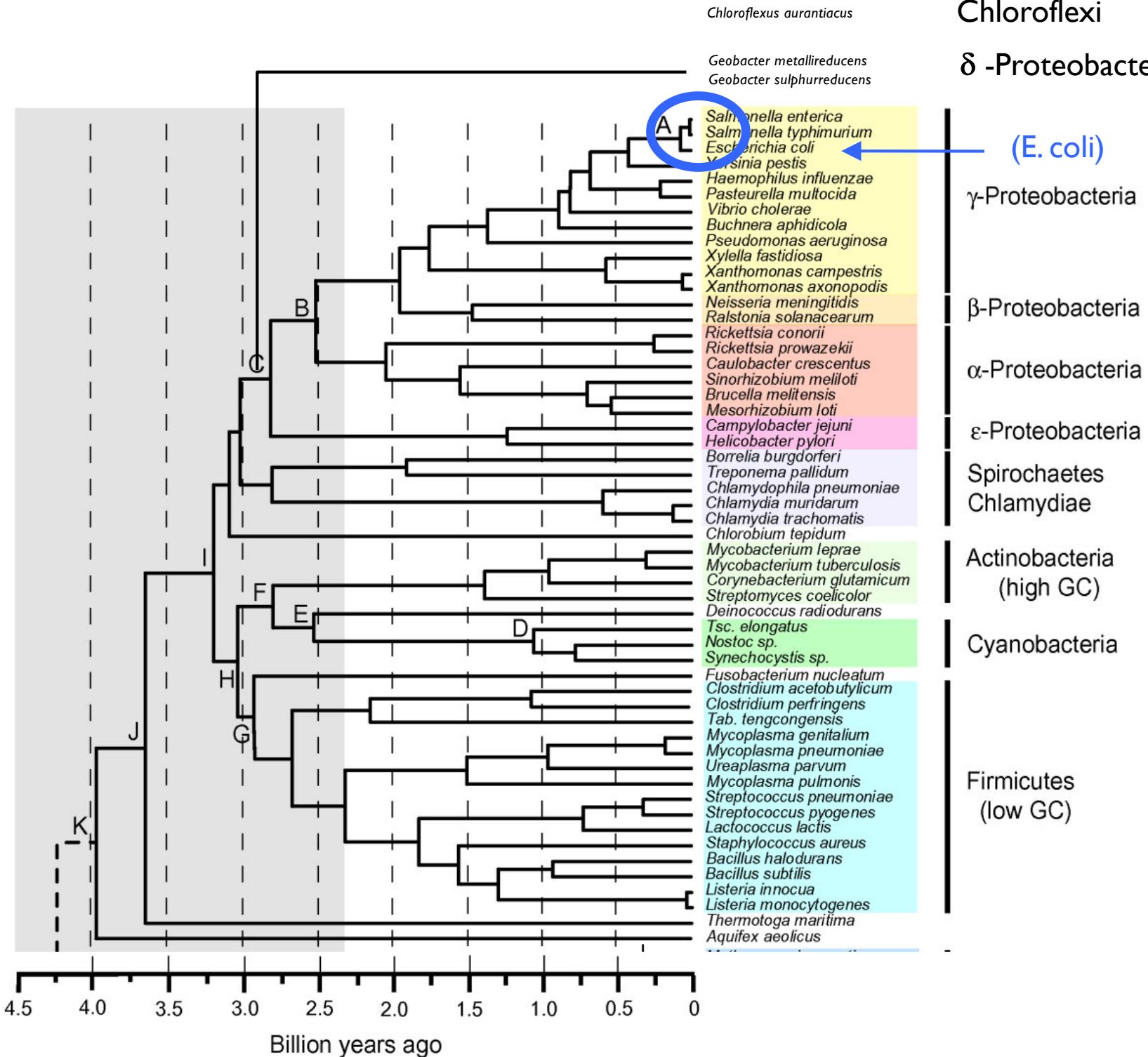
6S mimics an open promoter



Barrick et al. RNA 2005
Trotocaud et al. NSMB 2005
Willkomm et al. NAR 2005

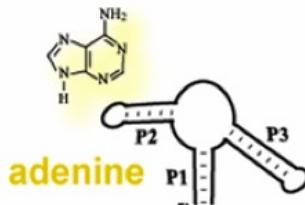
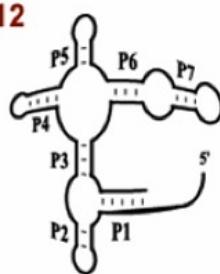
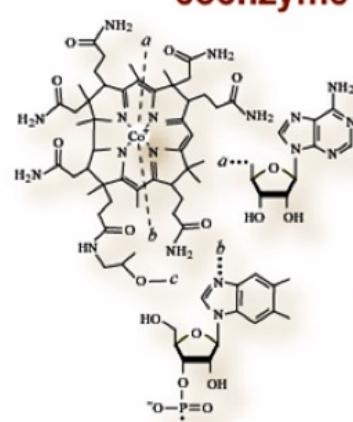
Chloroflexi

δ -Proteobacteria

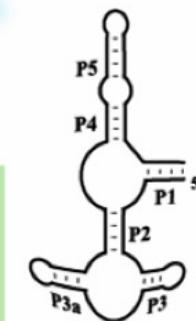
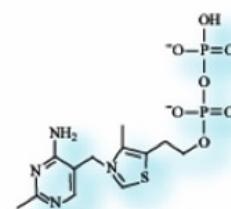


“Riboswitches”: small molecule sensors & gene on/off switches

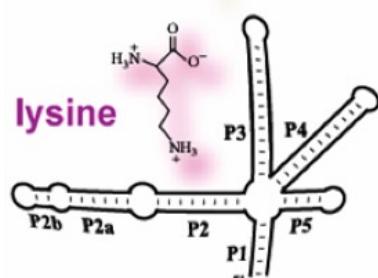
coenzyme B₁₂



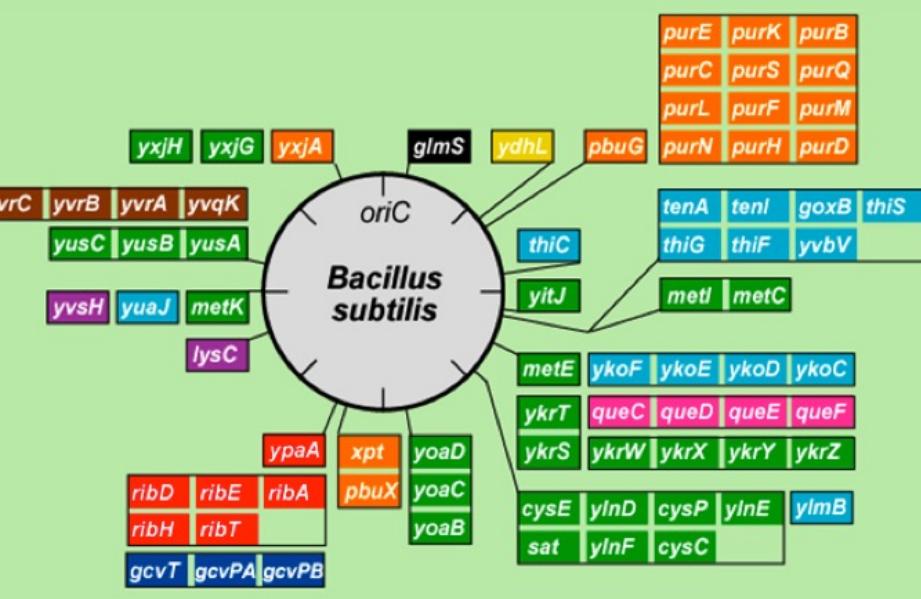
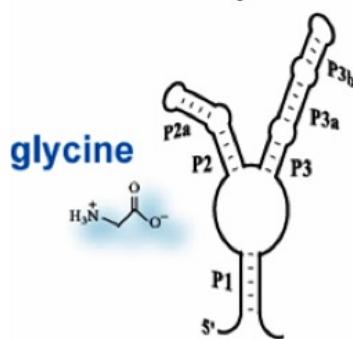
thiamine pyrophosphate



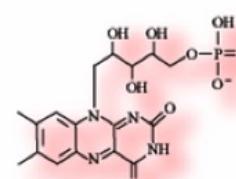
lysine



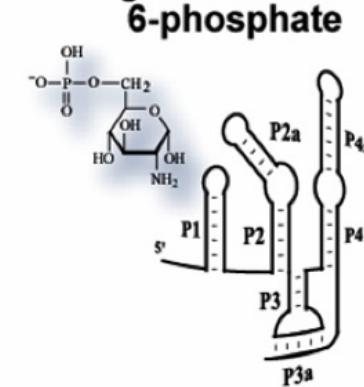
glycine



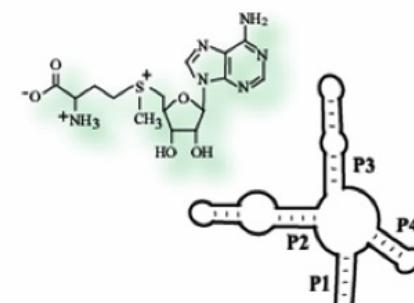
flavin mononucleotide



glucosamine-6-phosphate



S-adenosyl-methionine



Origin of Life?

Life needs

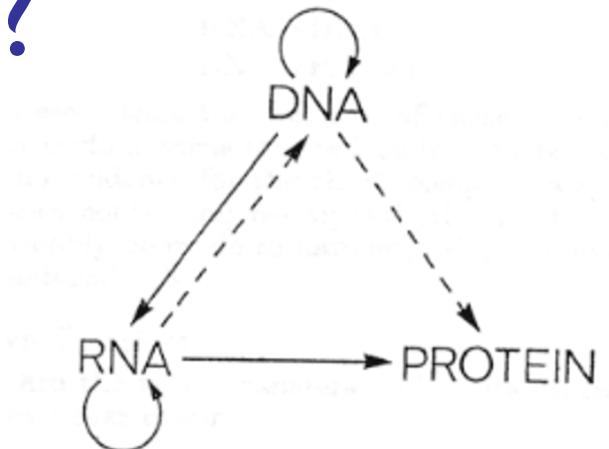
information carrier: DNA

molecular machines, like enzymes: Protein

making proteins needs DNA + RNA + proteins

making (duplicating) DNA needs proteins

Horrible circularities! How could it have arisen in an abiotic environment?



Origin of Life?

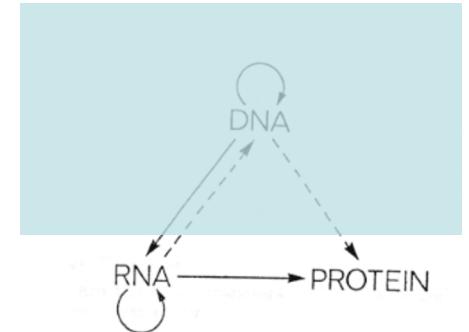
RNA can carry information, too

RNA double helix; RNA-directed RNA polymerase

RNA can form complex structures

RNA enzymes exist (ribozymes)

RNA can control, do logic (riboswitches)



The “RNA world” hypothesis:
1st life was RNA-based

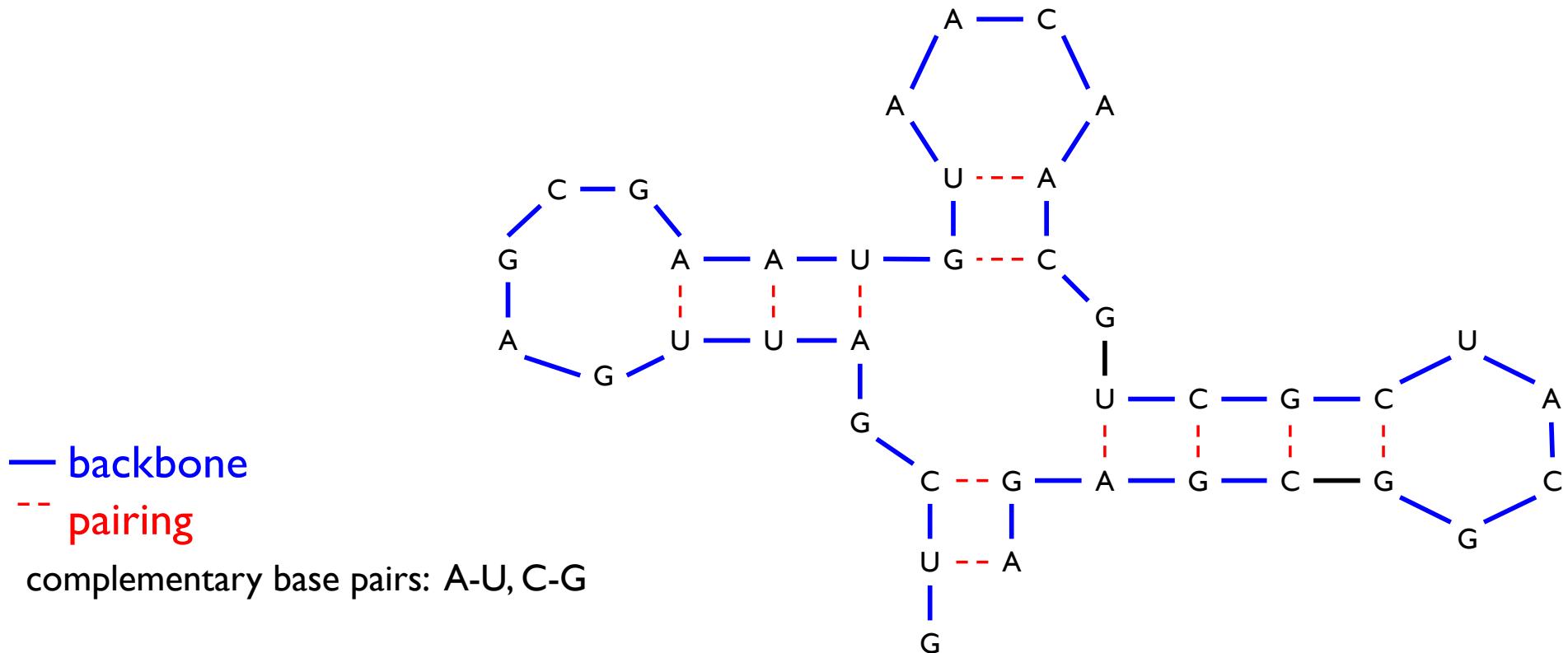
6.5 RNA Secondary Structure

Nussinov's Algorithm – core technology
for RNA structure prediction

RNA Secondary Structure

RNA. String $B = b_1b_2\dots b_n$ over alphabet { A, C, G, U }.

Secondary structure. RNA is usually single-stranded, and tends to loop back and form base pairs with itself. This structure is essential for understanding molecular behavior.



Ex: GUCGAUUGAGCGAAUGUAACAAACGUGGCUACGGCGAGA

RNA Secondary Structure (\approx oversimplified)

T in DNA →
U in RNA

RNA: String $B = b_1 b_2 \dots b_n$ over alphabet { A, C, G, U }.

Secondary structure: A set of pairs $S = \{ (b_i, b_j) \}$ satisfying:

- [Watson-Crick Pairing.]
 - S is a *matching*: each base pairs with \leq other, and
 - each pair in S is a Watson-Crick pair: A-U, U-A, C-G, or G-C.
- [No sharp turns.] Pairs are separated by ≥ 4 intervening bases.
 - If $(b_i, b_j) \in S$, then $i < j - 4$.
- [Non-crossing.] If (b_i, b_j) and (b_k, b_l) are two pairs in S , then we cannot have $i < k < j < l$. (Violation is called a *pseudoknot*.)

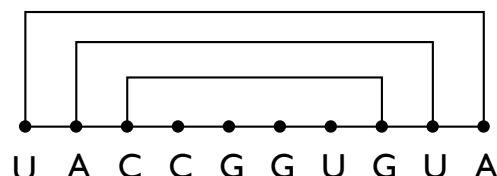
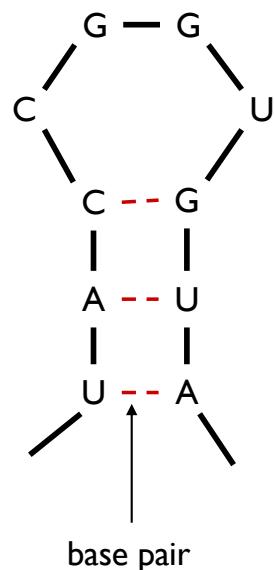
What's Best: RNA will form the structure that *minimizes* free energy.

approximated by maximizing
number of base pairs

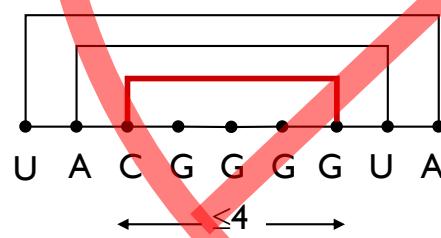
Goal: find a secondary structure S *maximizing* the number of base pairs.

RNA Secondary Structure: Examples

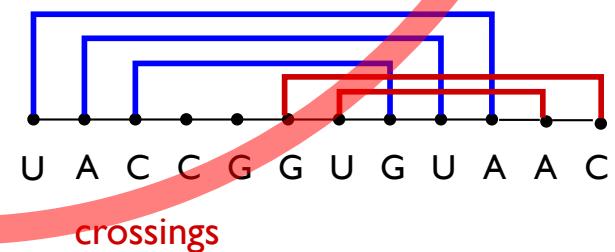
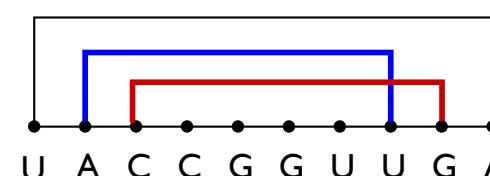
Examples.



ok

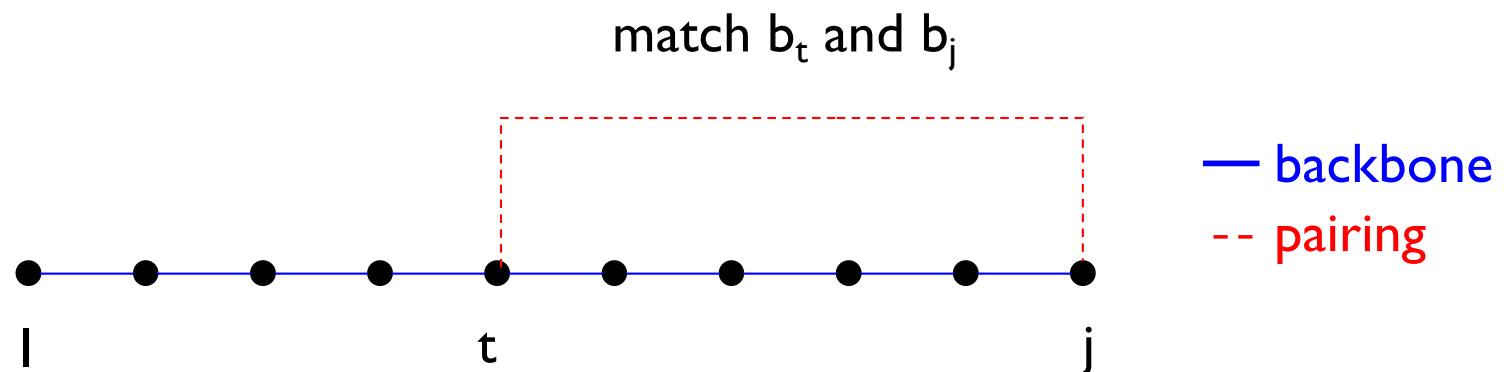


sharp turn



RNA Secondary Structure: Subproblems

First attempt. $\text{OPT}[j]$ = maximum number of base pairs in a secondary structure of the substring $b_1b_2\dots b_j$.



Results in two sub-problems.

- Find secondary structure in: $b_1b_2\dots b_{t-1}$. $\leftarrow \text{OPT}(t-1)$; good!
- Find secondary structure in: $b_{t+1}b_{t+2}\dots b_{j-1}$. \leftarrow DIFFICULTY: this isn't "OPT" of anything; need more flexible set of sub-problems

Dynamic Programming Over Intervals: (R. Nussinov's algorithm)

Notation. $\text{OPT}[i, j]$ = maximum number of base pairs in a secondary structure of the substring $b_i b_{i+1} \dots b_j$.

- Case Ia. If $i \geq j - 4$ (and base b_j is not paired):

$\text{OPT}[i, j] = 0$ by no-sharp turns condition.

- Case Ib. If $i < j - 4$, but base b_j is not paired:

$\text{OPT}[i, j] = \text{OPT}[i, j-1]$

- Case 2. Base b_j pairs with b_t for some $i \leq t < j - 4$.
non-crossing constraint decouples resulting sub-problems

$$\text{OPT}[i, j] = 1 + \max_t \{ \text{OPT}[i, t-1] + \text{OPT}[t+1, j-1] \}$$

take max over t such that $i \leq t < j-4$ and
 b_t and b_j are Watson-Crick complements

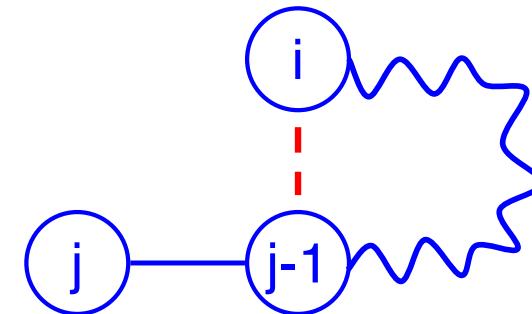
omit when $t == i$
(see next slide)

“Optimal pairing of $b_i \dots b_j$ ”

Two possibilities:

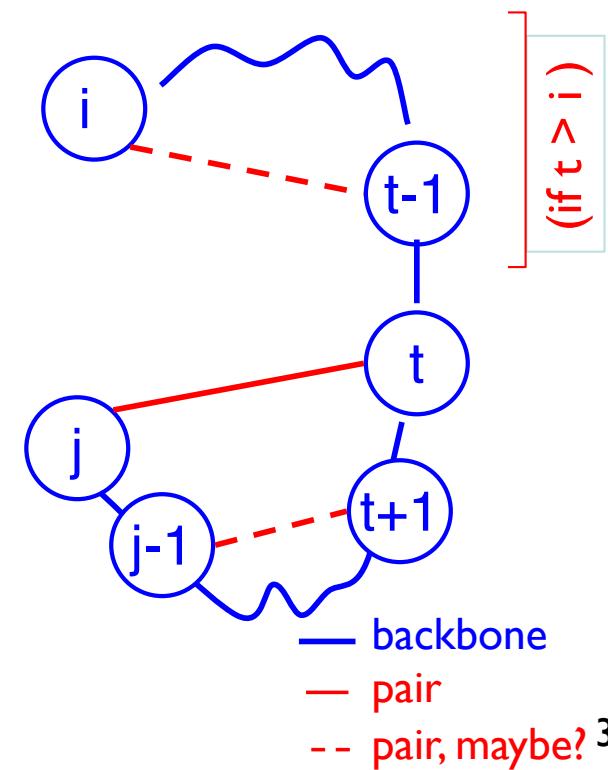
j Unpaired:

Find best pairing of $b_i \dots b_{j-1}$



j Paired (with some t):

Find best $b_i \dots b_{t-1} +$
best $b_{t+1} \dots b_{j-1}$ plus |



Why is it slow?

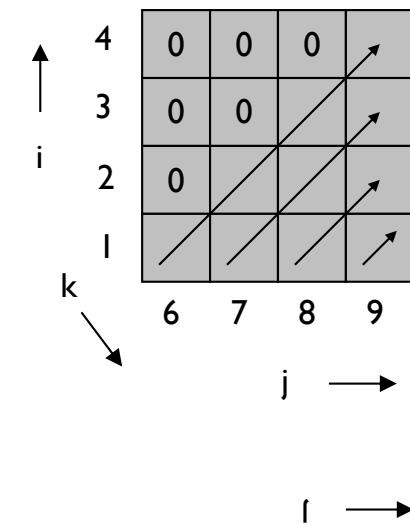
Why do pseudoknots matter?

Bottom Up Dynamic Programming Over Intervals

Q. What order to solve the sub-problems?

A1. Book way—do shortest intervals first, then earliest start:

```
RNA ( $b_1, \dots, b_n$ ) {  
Interval length →     for  $k = 5, 6, \dots, n-1$   
Start position →        for  $i = 1, 2, \dots, n-k$   
End position →             $j = i + k$   
                            Compute  $\text{OPT}[i, j]$   
                            ↑  
                            using recurrence  
                            return  $\text{OPT}[1, n]$   
}
```

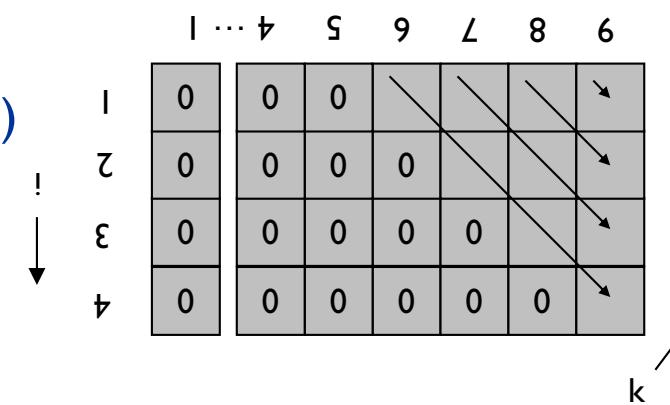


book

A2. Slides way: earliest start first,
then shortest intervals (next slides)

+ HW

Running time. $O(n^3)$ (either way)



slides

Nussinov: Max Pairing

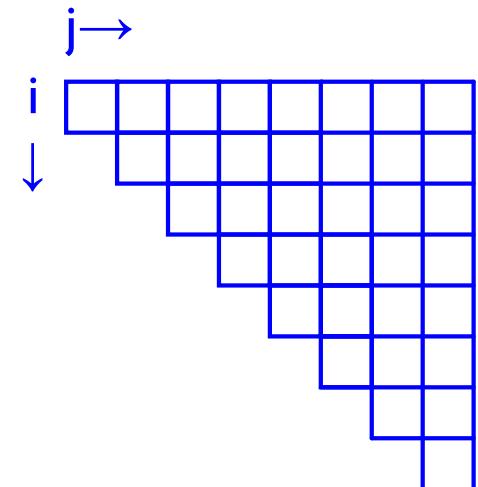
$\text{Opt}[i,j] = \# \text{ pairs in optimal pairing of } b_i \dots b_j$

$\text{Opt}[i,j] = 0$ for all i, j with $i \geq j-4$; otherwise

$\text{Opt}[i,j] = \max \text{ of:}$

$$\left\{ \begin{array}{l} \text{Opt}[i,j-1] \\ \max \{ \text{Opt}[i,t-1] + 1 + \text{Opt}[t+1,j-1] \mid \\ \quad i \leq t < j-4 \text{ and } b_t - b_j \text{ may pair} \} \end{array} \right.$$

if $t > i$



Another Computation Order

$\text{Opt}[i, j] = \text{optimal } \# \text{ pairs in } b_i \dots b_j$

```

→ for(j = 1 to n)
→ for(i = j downto 1)

 $\text{Opt}[i, j] = 0 \text{ if } i \geq j-4 \text{ else:}$ 

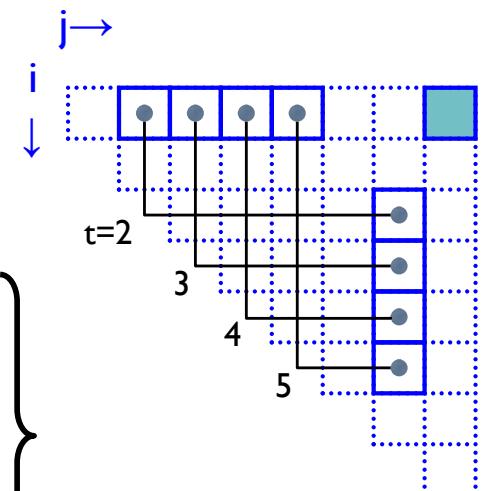
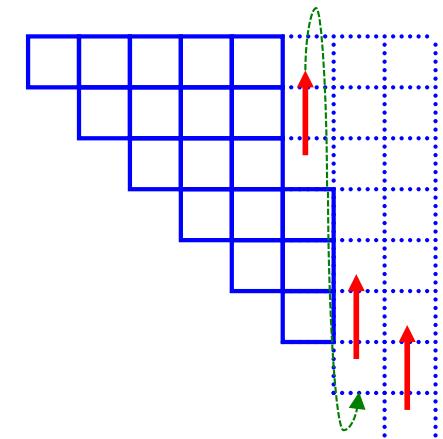
```

max of:

$$\left\{ \begin{array}{l} \text{Opt}[i, j-1] \\ \max \{ \text{Opt}[i, t-1] + 1 + \text{Opt}[t+1, j-1] \mid \\ \quad i \leq t < j-4 \text{ and } b_t - b_j \text{ may pair } \} \end{array} \right\}$$

if $t > i$

Time: $O(n^3)$



Which Pairs?

Usual dynamic programming “trace-back” tells you *which* base pairs are in the optimal solution, not just how many

Details? : homework

Computing one cell: $\text{OPT}[2, 18] = ?$

(Examples here and below
assume 1-based indexing)

G	G	G	A	A	A	A	C	C	C	A	A	A	G	G	G	U	U	U
(((.	.	.	.)))	(((.	.	.)))
0	0	0	0	0	0	0	1	2	3	3	3	3	3	3	3	4	5	6
0	0	0	0	0	0	0	1	2	2	2	2	2	3	3	3	4	5	6
0	0	0	0	0	0	0	1	1	1	1	1	1	2	2	3	3	4	5
0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1

n= 20

Case I:

$2 \geq 18-4$? no.

Case 2:

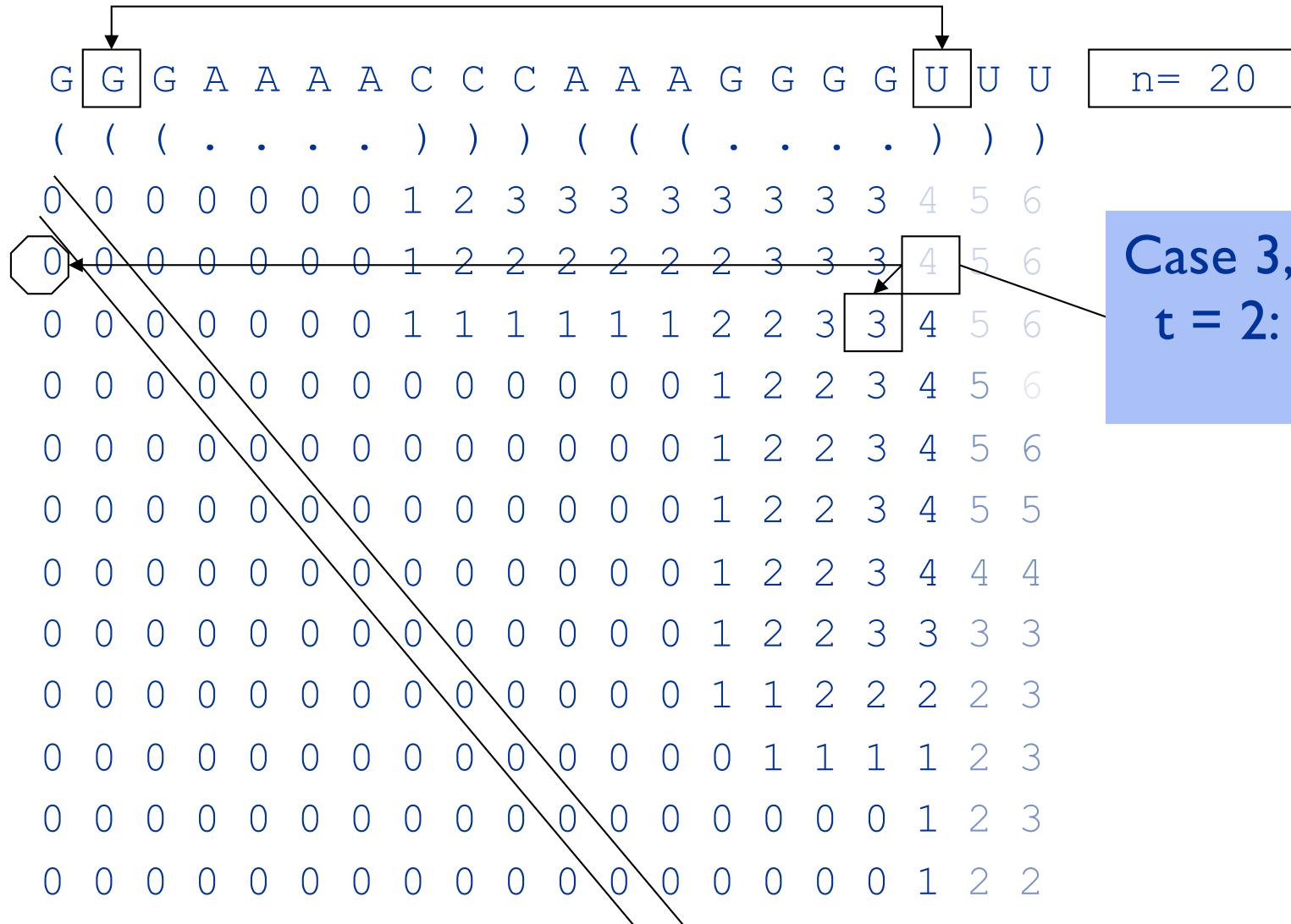
B_{18} unpaired?

Always a possibility;
then $\text{OPT}[2, 18] \geq 3$

GGAAAAACCCAAAGGGGU
((.....)) (.....)

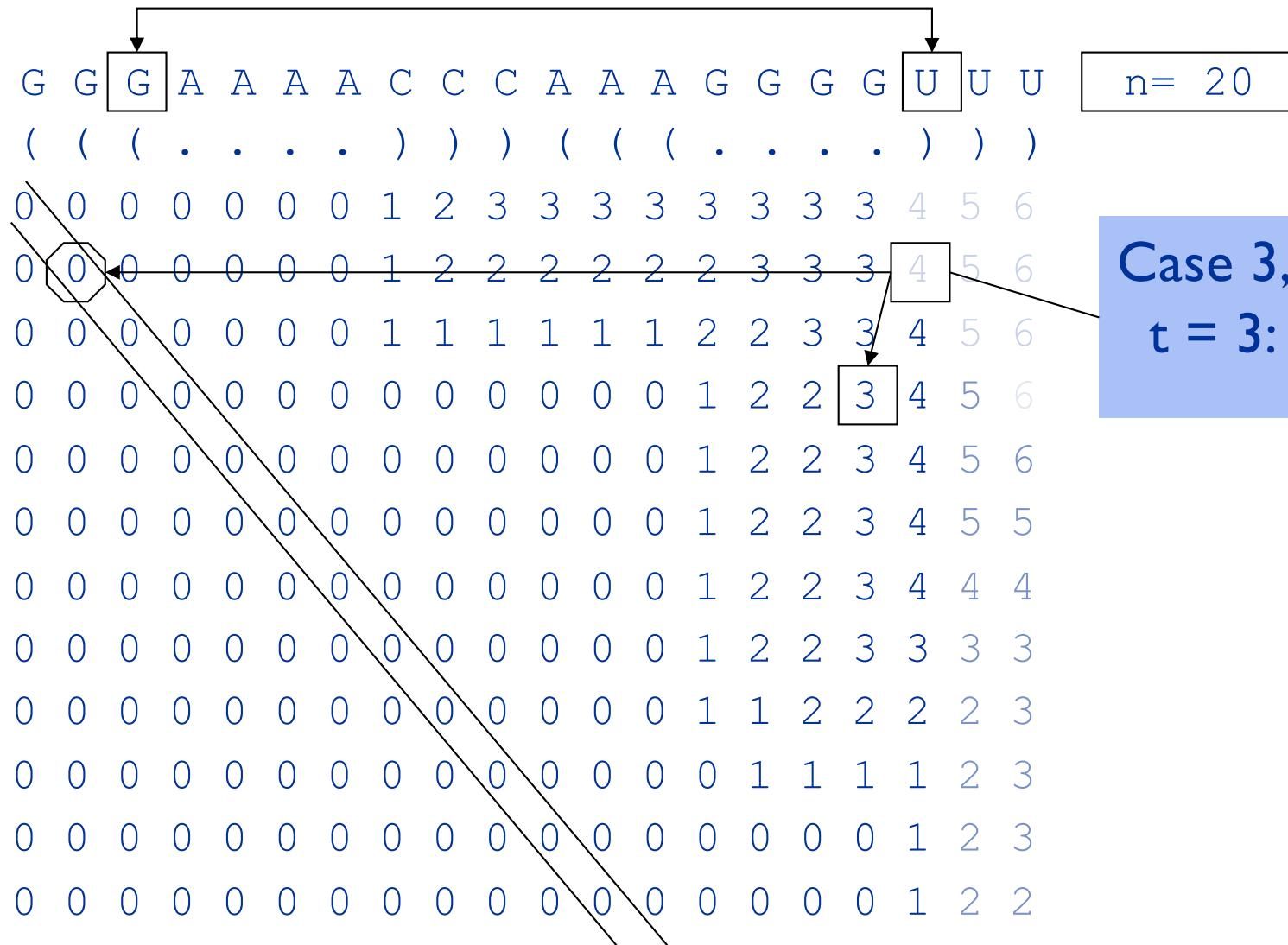
$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2, 18] = ?$



$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2, 18] = ?$



$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2, 18] = ?$



Case 3, $2 \leq t < 18-4$:
 $t = 4$: yes pair
 $\text{OPT}[2, 18] \geq 1 + 0 + 3$

GG**A**AAACCCAAAGGGG**U**
...(((((....))))

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$



Case 3, $2 \leq t < 18-4$:
 $t = 5$: yes pair
 $\text{OPT}[2,18] \geq 1 + 0 + 3$

GGA**A**AAACCCAAAGGGGG**U**
(....((....)))

$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2, 18] = ?$

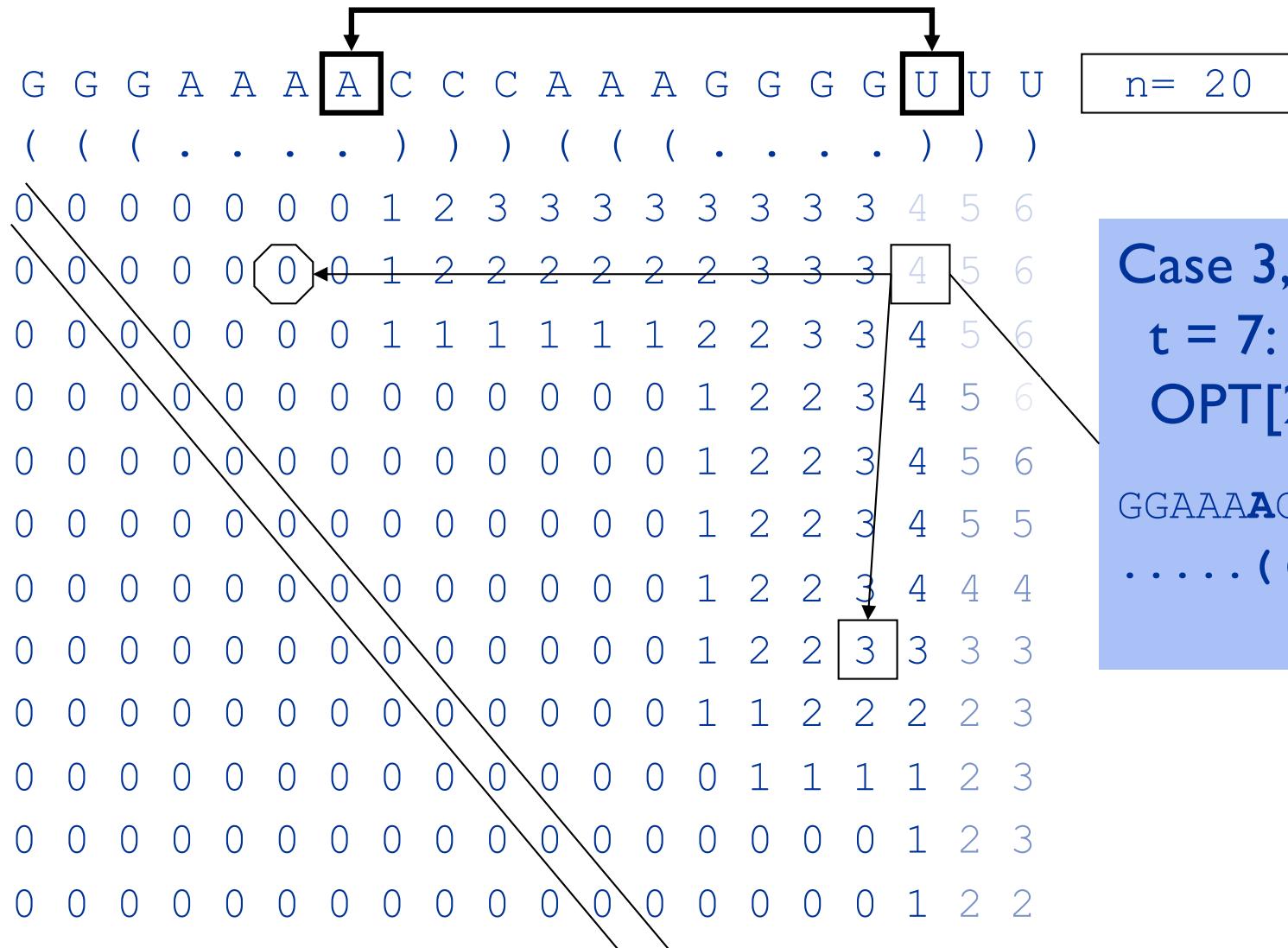


Case 3, $2 \leq t < 18 - 4$:
 $t = 6$: yes pair
 $\text{OPT}[2, 18] \geq 1 + 0 + 3$

GGAAAACCCAAAGGGGU
....(.(.(.(.)))

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$



Case 3, $2 \leq t < 18-4$:
 $t = 7$: yes pair
 $\text{OPT}[2,18] \geq 1 + 0 + 3$

GGAAA**A**CCC**A**AGGG**G**
 $\dots \dots (((((\dots))))$

$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

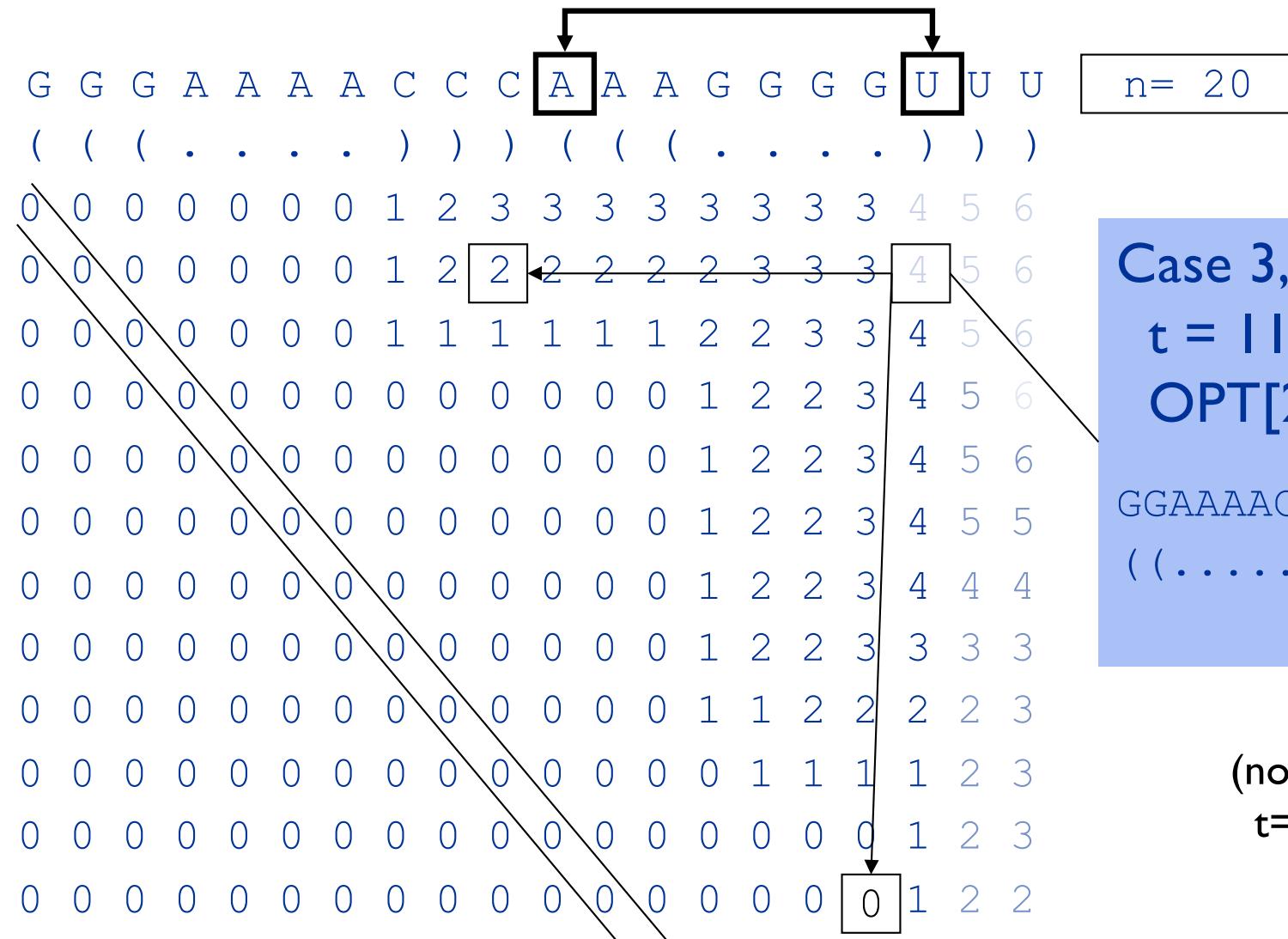
Computing one cell: $\text{OPT}[2, 18] = ?$



Case 3, $2 \leq t < 18 - 4$:
 $t = 8$: no pair

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell: $\text{OPT}[2,18] = ?$



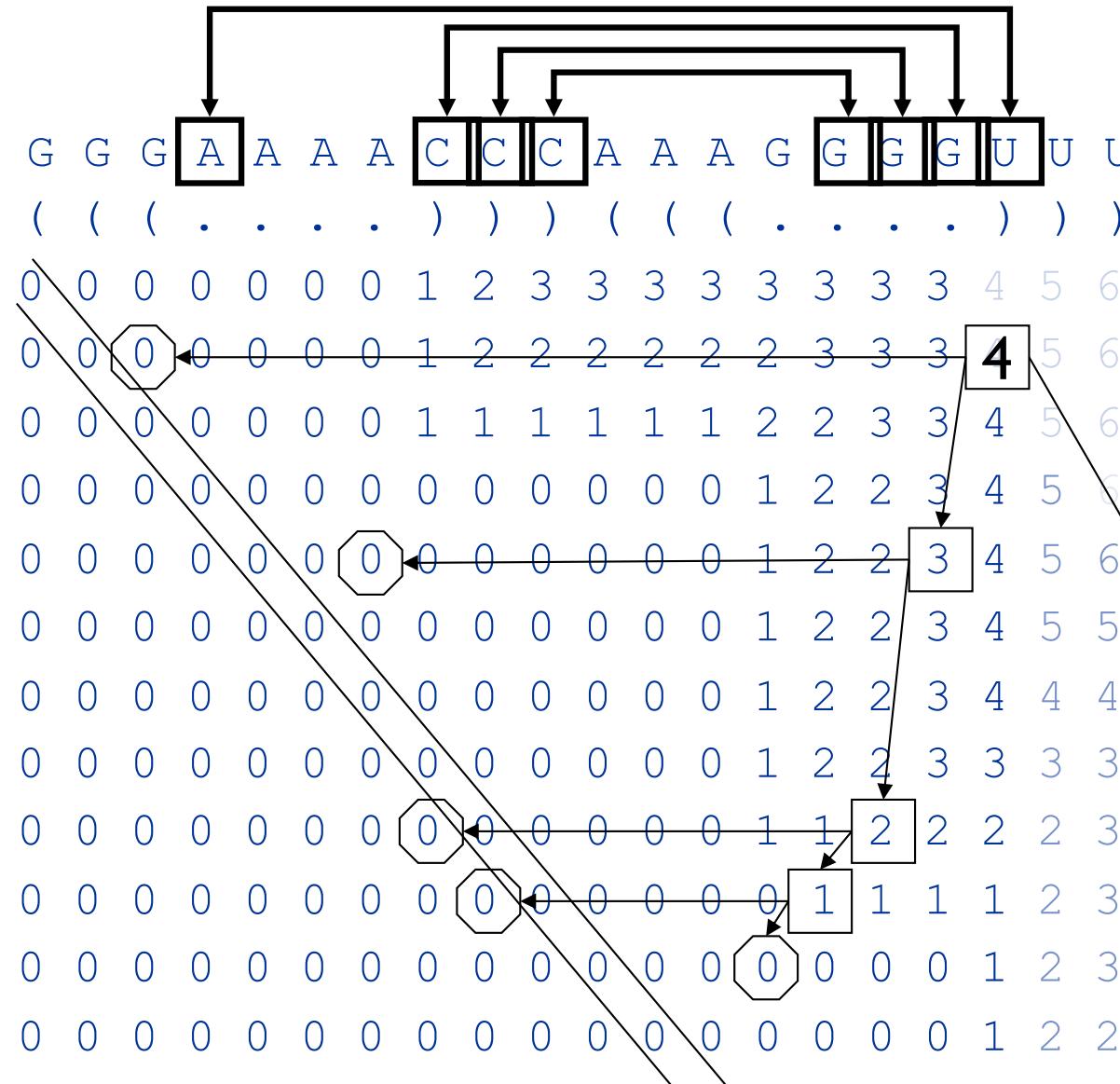
Case 3, $2 \leq t < 18-4$:
 $t = 11$: yes pair
 $\text{OPT}[2,18] \geq 1+2+0$

GGAAAACCC**A**AAGGGGG**U**
 ((.....)) (.....)

(not shown:
 $t=9, 10, 12, 13$)

$$\text{OPT}(i,j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Computing one cell:
 $\text{OPT}[2, 18] = 4$
 $n = 20$



Overall, Max = 4
several ways, e.g.:

GGAAAACCCAAAGGGGU
... (....((.....)))

tree shows trace back:

square = case 3
octagon = case 1

$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j-1] \\ 1 + \max_t (\text{OPT}[i, t-1] + \text{OPT}[t+1, j-1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

All 5 optimal structures on the above example

GGGAAAACCCAAAGGGGUUU

...((((.....))))
...((..(((.....))))
...(.((((.....))))
....((((.....))))
(((.....))((.....)))

0	0	0	0	0	0	0	1	2	3	3	3	3	3	3	3	4	5	6		
-7	0	0	0	0	0	0	1	2	2	2	2	2	2	3	3	3	4	5	6	
-7-7	0	0	0	0	0	0	1	1	1	1	1	1	2	2	3	3	4	5	6	
-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6	
-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	6	
-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5	
-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	5	5	
-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	4	4	4	
-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2	3	3	3	3	
-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	1	2	2	2	2	3	
-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	2	3	
-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	3	
-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	2	2
-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	
-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
-7-7-7-7-7-7-7-7-7-7-7-7-7-7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
n= 20 Pairs= 6 AltStructs= 5 0.000117 (sec. total)																				

Another Example

| n = 16

E.g.:
 $\text{OPT}[1,6] = 1:$

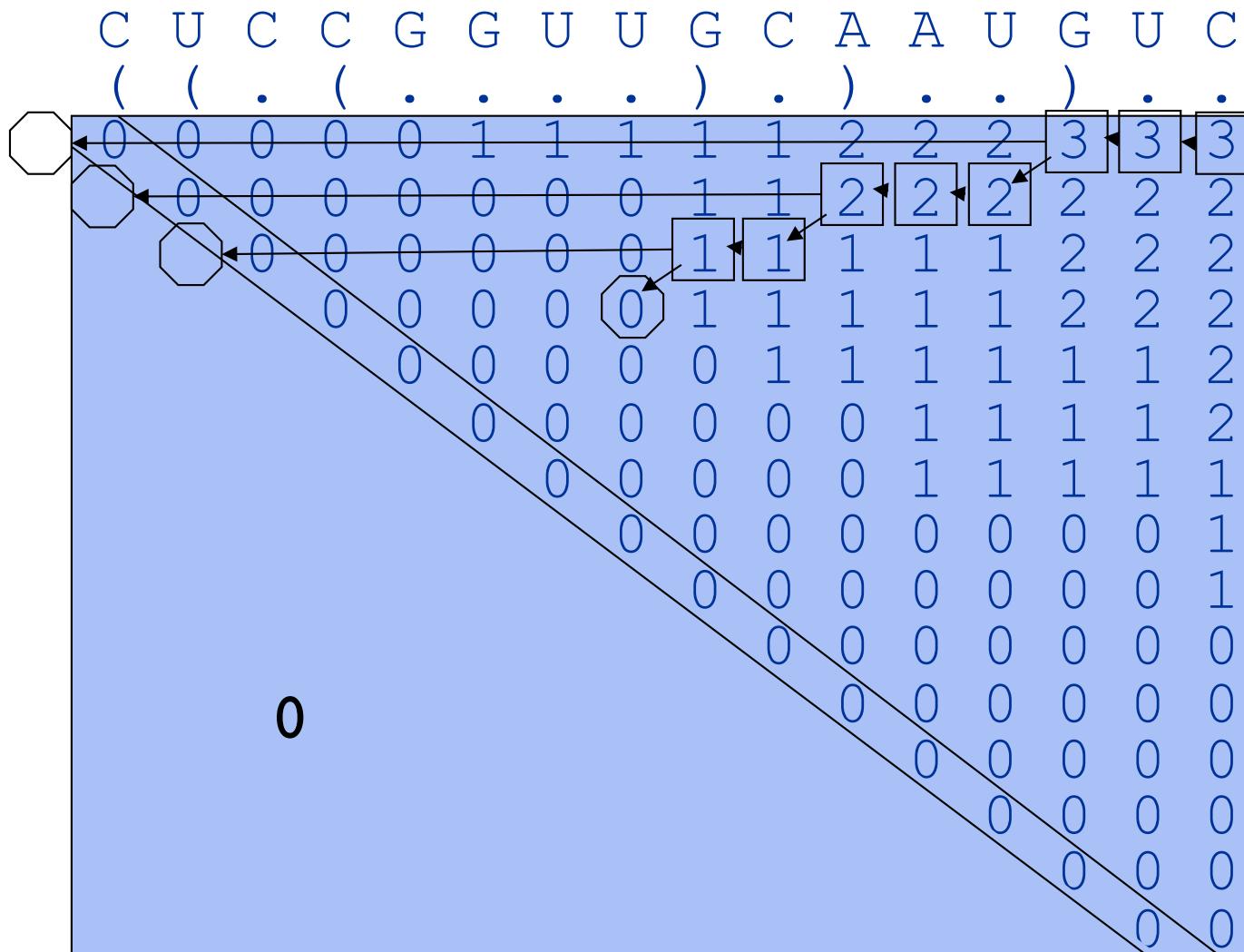
CUCCGG
(. . .)

E.g.:
 $\text{OPT}[6, | 6] = 2$:

GUUGCAAUGUC
((.....) . . .)

(Examples here and below assume 1-based indexing)

Another Trace Back Example



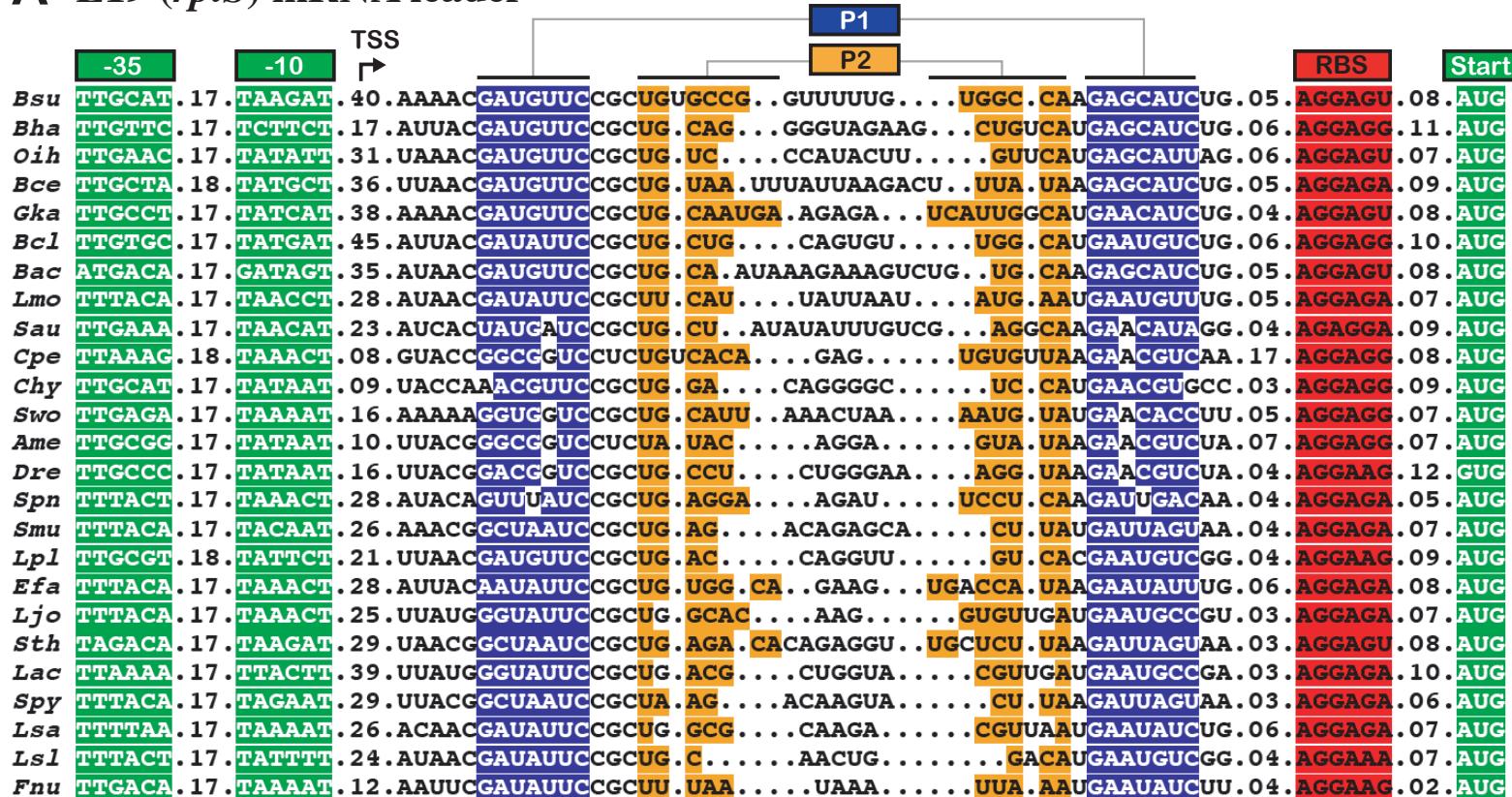
$n = 16$

E.g.:
 $\text{OPT}[1,16] = 3:$
 CUCCGGUUGCAAUGUC
 ((...(.) .) ...)

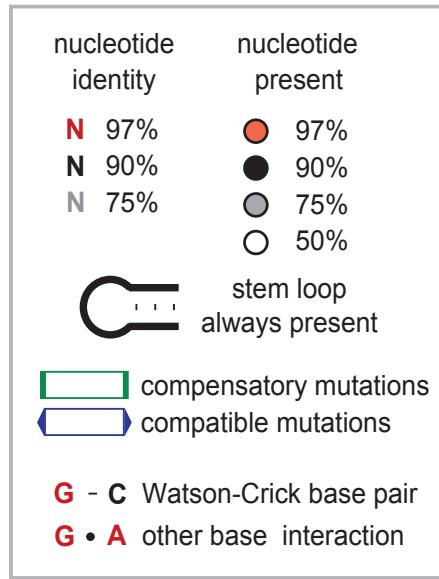
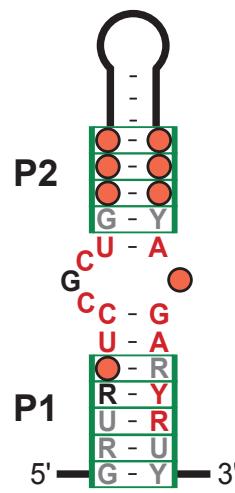
$$\text{OPT}(i, j) = \begin{cases} 0 & \text{if } i \geq j - 4 \\ \max \left\{ \begin{array}{l} \text{OPT}[i, j - 1] \\ 1 + \max_t (\text{OPT}[i, t - 1] + \text{OPT}[t + 1, j - 1]) \end{array} \right\} & \text{otherwise} \end{cases}$$

Example: Ribosomal Autoregulation: Excess L19 represses L19 (RF00556; 555-559 similar)

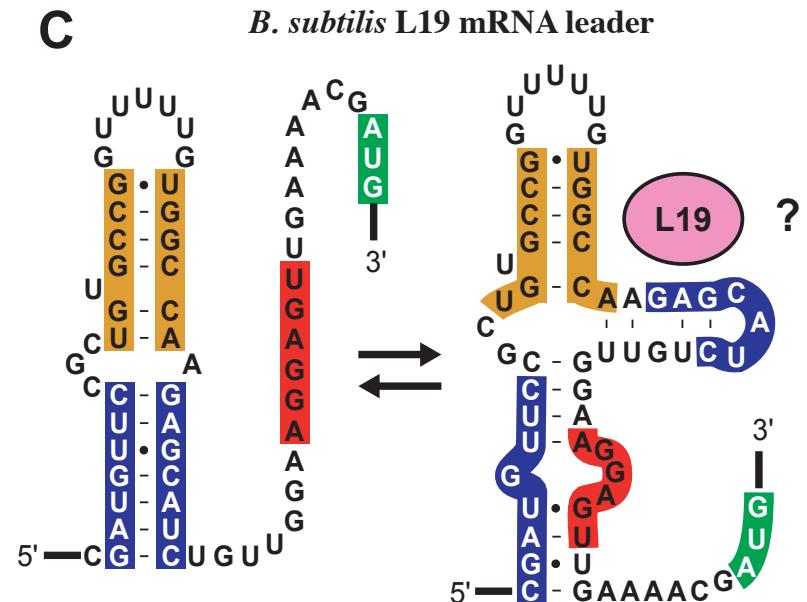
A L19 (*rplS*) mRNA leader

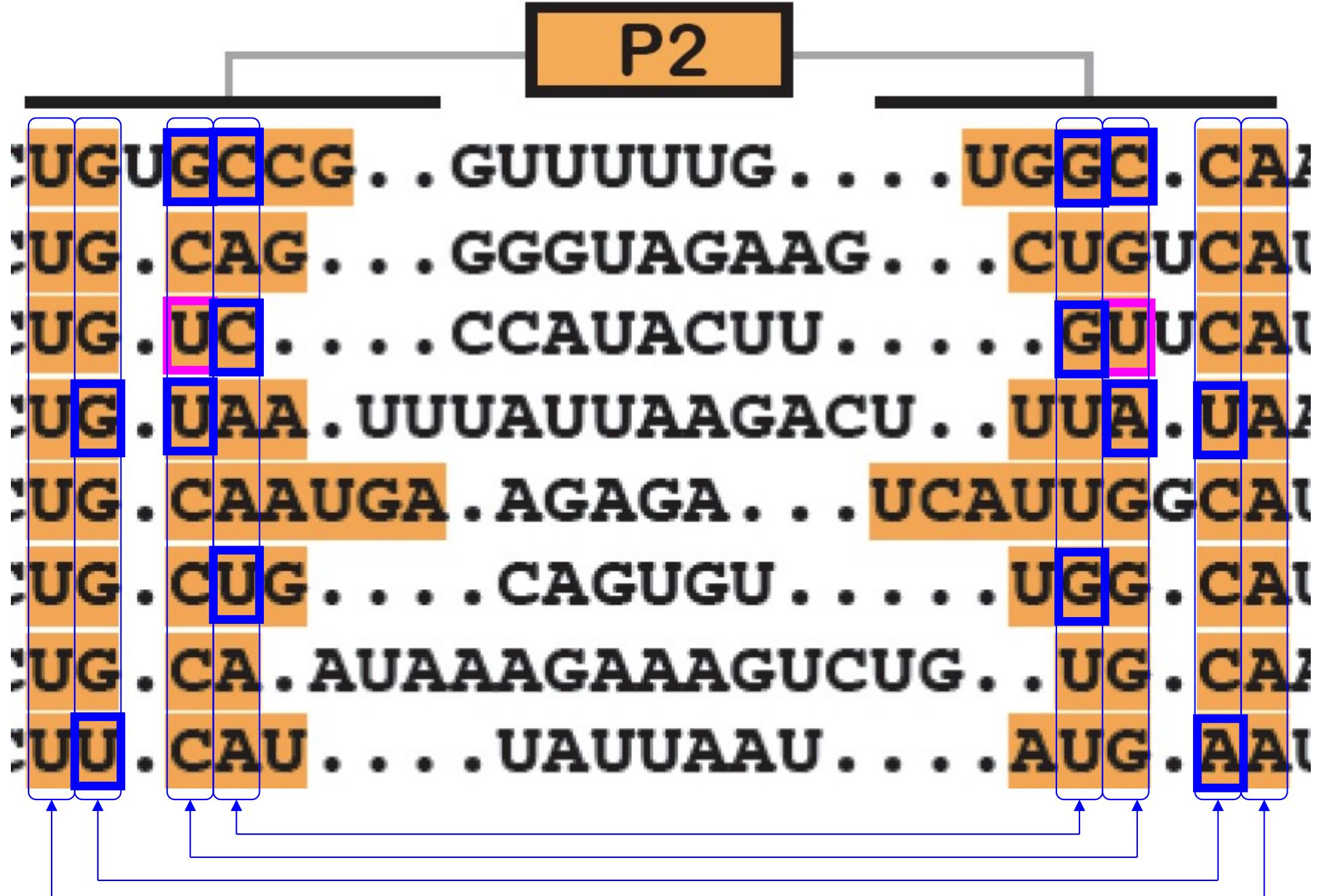


B



C





Covariation is strong evidence for base pairing

Summary

RNA has important roles

Beyond mRNA; many unexpected recent discoveries

Structure is critical to function

True of other molecules, too

RNA secondary structure prediction is a key tool

Dynamic programming—useful accuracy, $O(n^3)$ time:

Binary choice again: last base is paired or not

Optimal substructure again: given last pair, optimally fold inside & outside separately

Tabulate again: best folding of all substrings.