

# CSE 417: Algorithms and Computational Complexity

## Lecture I: Overview

Winter 2022

Larry Ruzzo



# University of Washington

## Computer Science & Engineering

### CSE 417, Wi '22: Algorithms & Computational Complexity

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#### Administrative

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#### Lecture Notes

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#### Lecture Recordings

[Here](#)

**Lecture:** (By Zoom, Week 1 at least), MWF 9:30-10:20

**Zoom:** <https://washington.zoom.us/j/94357067917> Requires login to Zoom with a UWNetID.

Instructor:	Office Hours	Location	Phone
<a href="#">Larry Ruzzo</a> , ruzzo@cs	TBD	<a href="#">Zoom</a>	
<b>TAs:</b> Nathan Akkaraphab, akkanath@cs	TBD		
Todor Dimitrov, todord@cs	TBD		
William Viet Nguyen, williamv@cs	TBD		
Lin Qiu, lq9@cs	TBD		
Luna Wang, zhengw28@cs	TBD		
Yifan Zhang, yifanz47@cs	TBD		
Yilin Zhang, yilinz24@cs	TBD		

**Course Email:** [cse417a\\_wi22@uw.edu](mailto:cse417a_wi22@uw.edu). Staff and TAs are subscribed to this list. [Emails are archived.](#)

#### Discussion

Discuss homework, etc.

structures. Efficient algorithms for manipulating graphs and strings. Fast Fourier Transform. Turing machines. Time and space complexity. NP-complete problems and undecidable problems.

**Late Policy:** TBD

**Required Text:** *Algorithm Design* by [Jon Kleinberg](#) and [Eva Tardos](#). Addison Wesley, 2006. (Available from [U Book Store](#), [Amazon](#), etc.)

**References:** See [Schedule & Reading](#).

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<http://courses.cs.washington.edu/417>

# What you'll have to do

Homework ~~(~55% of grade)~~ **100%**

Programming

Several small projects

Written homework assignments

English exposition and pseudo-code

Analysis and argument as well as design

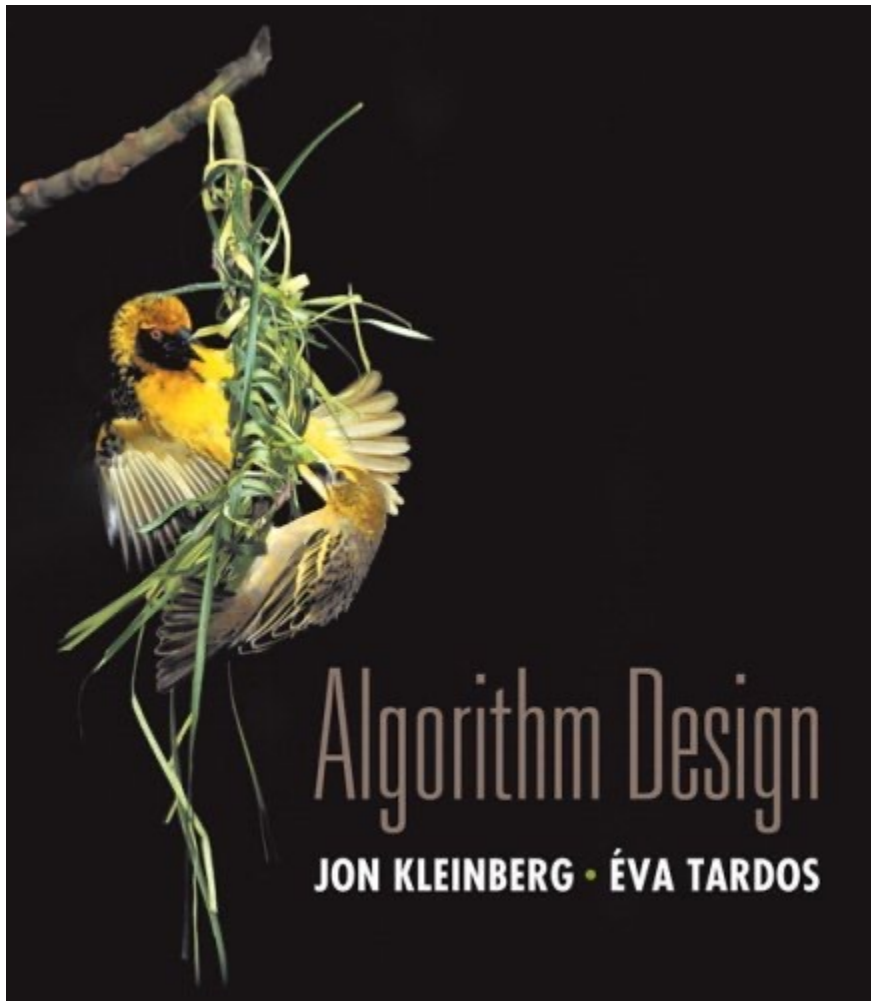
**“Mastery Grading”:**  
If you can solve it, you get credit for “mastering” the concept, with a chance to resubmit/regrade some work each week.  
Details TBD

~~Midterm / Final Exam (~15% / 30%)~~

Late Policy: **TBD**

~~Papers and/or electronic turnins are due at the *start* of class on the due date. 10% off for one day late; 15% per day thereafter.~~

# Textbook



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

# What the course is about

## Design of Algorithms

design methods

common or important types of problems

analysis of algorithms - efficiency

correctness proofs

# What the course is about

Complexity, NP-completeness and intractability

solving problems in principle is not enough

algorithms must be *efficient*

some problems have *no efficient solution*

NP-complete problems

important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

# Very Rough Division of Time

Algorithms (7-8 weeks)

Analysis of Algorithms

Basic Algorithmic Design Techniques

Applications

Complexity & NP-completeness (2-3 weeks)

Check online  
schedule page for  
(evolving) details



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CSE 417, Wi '06: *Approximate Schedule*

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	Due	Lecture Topic	Reading
<b>Week 1</b> 1/2-1/6	M	Holiday	
	W	Intro, Examples & Complexity	Ch. 1; Ch. 2
	F	Intro, Examples & Complexity	
<b>Week 2</b> 1/9-1/13	M	Intro, Examples & Complexity	Ch. 3 <b>7</b>
	W	Graph Algorithms	
	F	Graph Algorithms	

# Complexity Example

Cryptography (e.g., RSA, SSL in browsers)

Secret:  $p, q$  prime, say 512 bits each

Public:  $n$  which equals  $p \times q$ , 1024 bits

In principle

*there is an algorithm* that given  $n$  will find  $p$  and  $q$ :  
try all  $2^{512} > 1.3 \times 10^{154}$  possible  $p$ 's: kinda slow...

In practice

*no fast algorithm* known for this problem (on non-quantum computers)

security of RSA depends on this fact

(“quantum computing”: strongly driven by possibility of changing this)



# Algorithms versus Machines

You may know about Moore's Law and the exponential improvements in hardware...

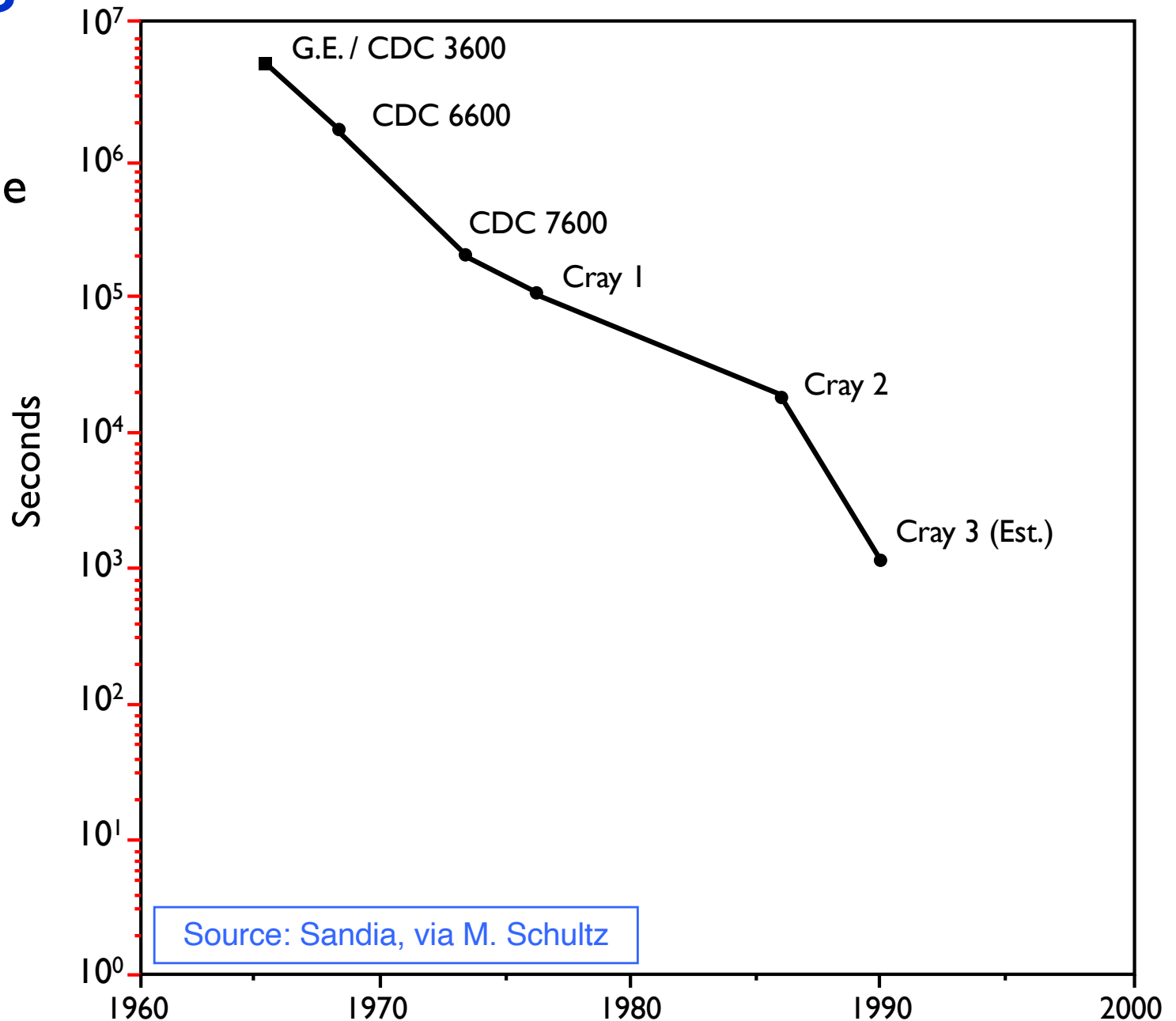
Ex: sparse linear equations over 25 years

10 orders of magnitude improvement!

# Algorithms or Hardware?

25 years  
progress  
solving sparse  
linear  
systems

hardware:  
4 orders of  
magnitude

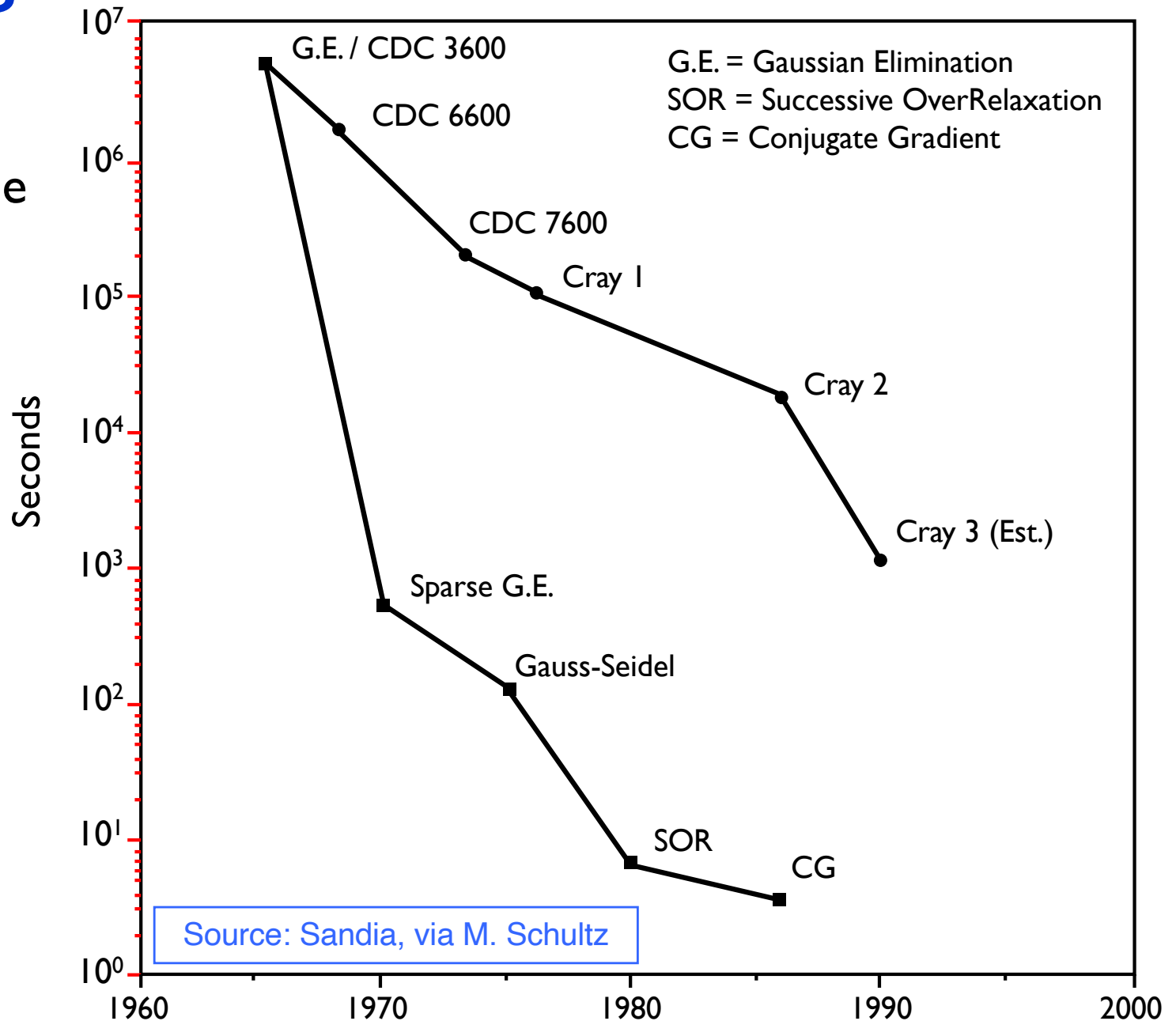


# Algorithms or Hardware?

25 years  
progress  
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hardware: 4  
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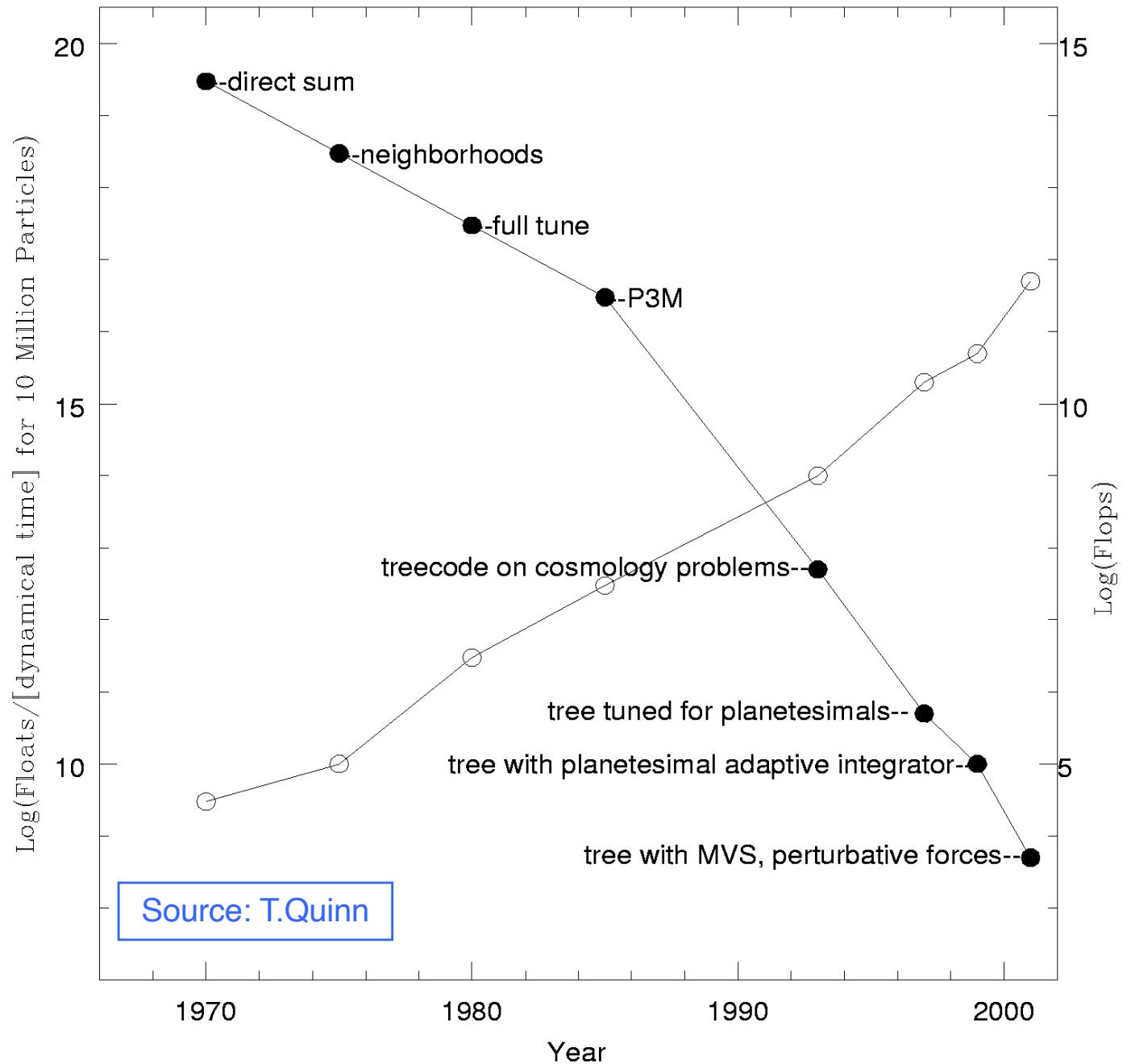
software: 6  
orders of  
magnitude



# Algorithms or Hardware?

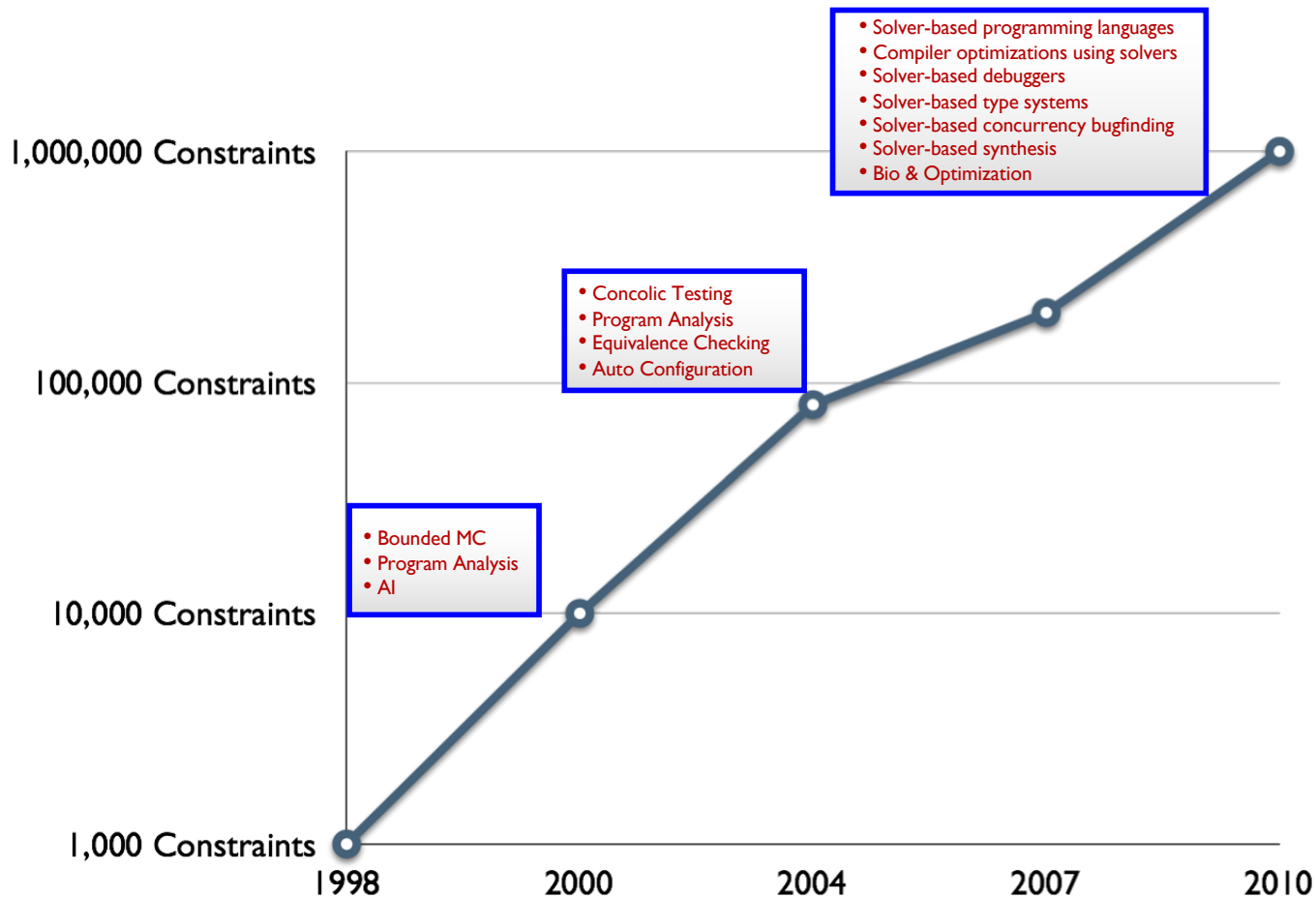
The  
N-Body  
Problem:

in 30 years  
 $10^7$  hardware  
 $10^{10}$  software



# Algorithms or Hardware?

SAT/SMT Solvers: 1000x improvement in a dozen years



Data courtesy of Dr. Vijay Ganesh, U. Waterloo

# Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them

“accomplish” via simple, well-defined steps

Ex: sorting names (via comparison)

Ex: checking for primality (via  $+$ ,  $-$ ,  $*$ ,  $/$ ,  $\leq$ )

# Goals

Correctness

often subtle

Analysis

often subtle

Generality, Simplicity, 'Elegance'

Efficiency

time, memory, network bandwidth, ...

# Algorithms: a sample problem

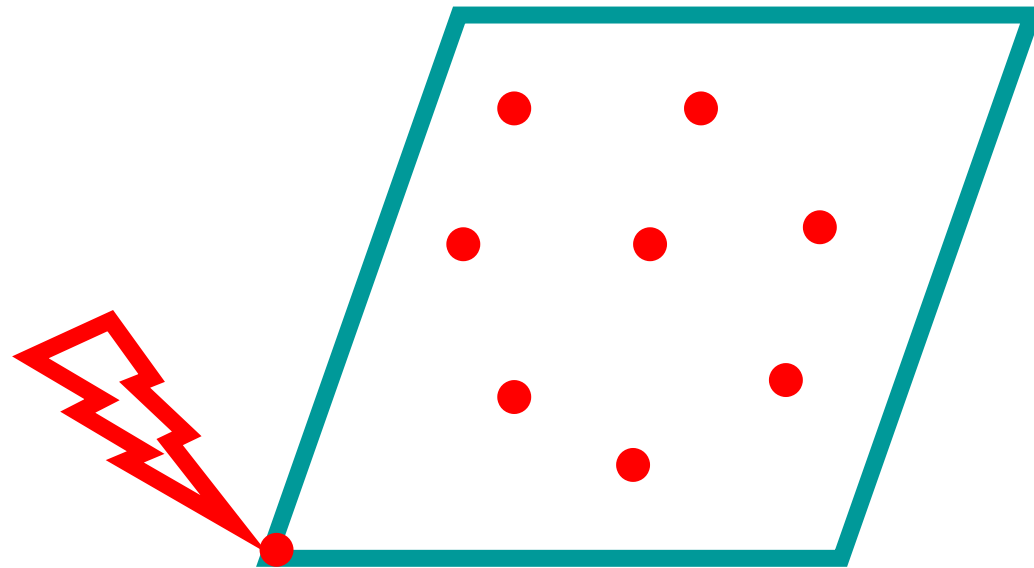
Printed circuit-board company has a robot arm that solders components to the board

Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

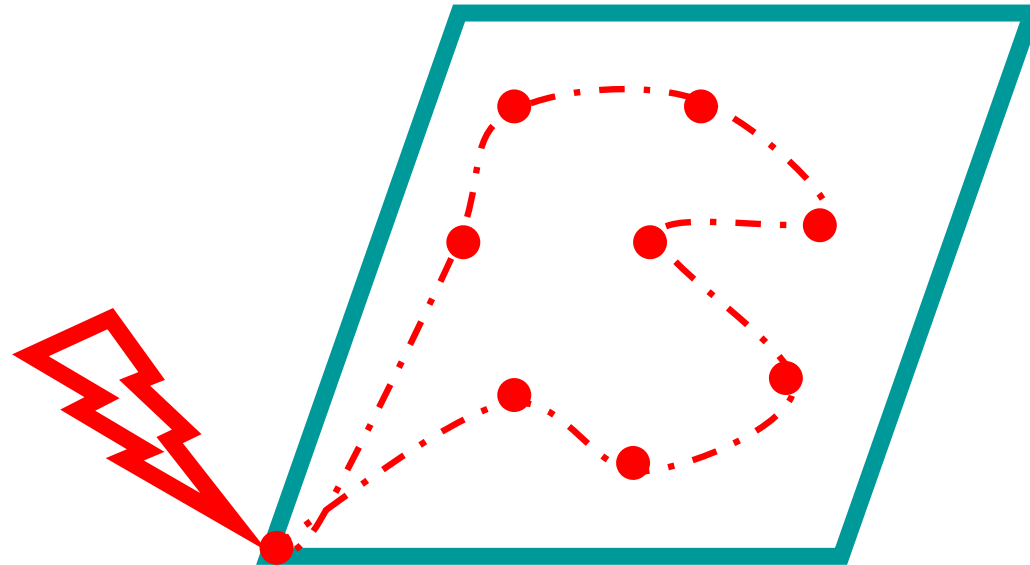
For each board design, find best order to do the soldering



# Printed Circuit Board



# Printed Circuit Board



# A Well-defined Problem

Input: Given a set  $S$  of  $n$  points in the plane

Output: The shortest cycle tour that visits each point in the set  $S$  once.

Better known as “TSP”

How might you solve it?

# Nearest Neighbor Heuristic

Start at some point  $p_0$

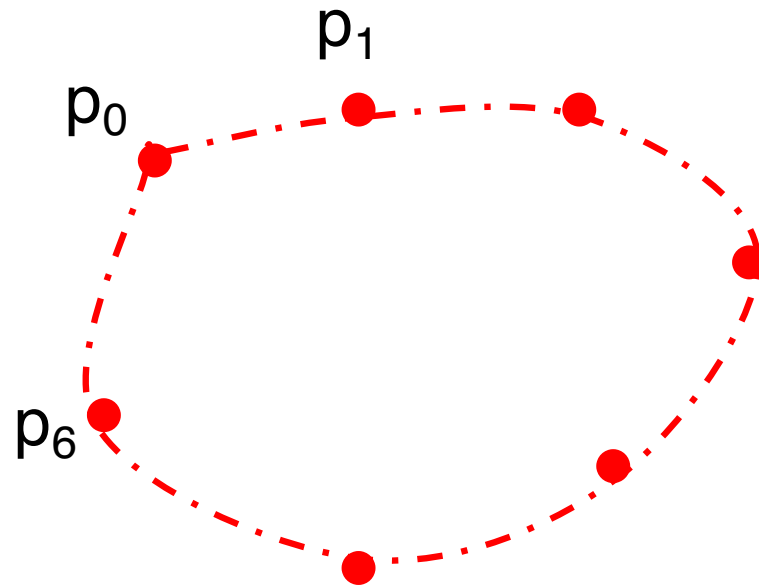
Walk first to its  
nearest neighbor  $p_1$

Repeatedly walk to the nearest unvisited neighbor  
 $p_2$ , then  $p_3, \dots$  until all points have been visited

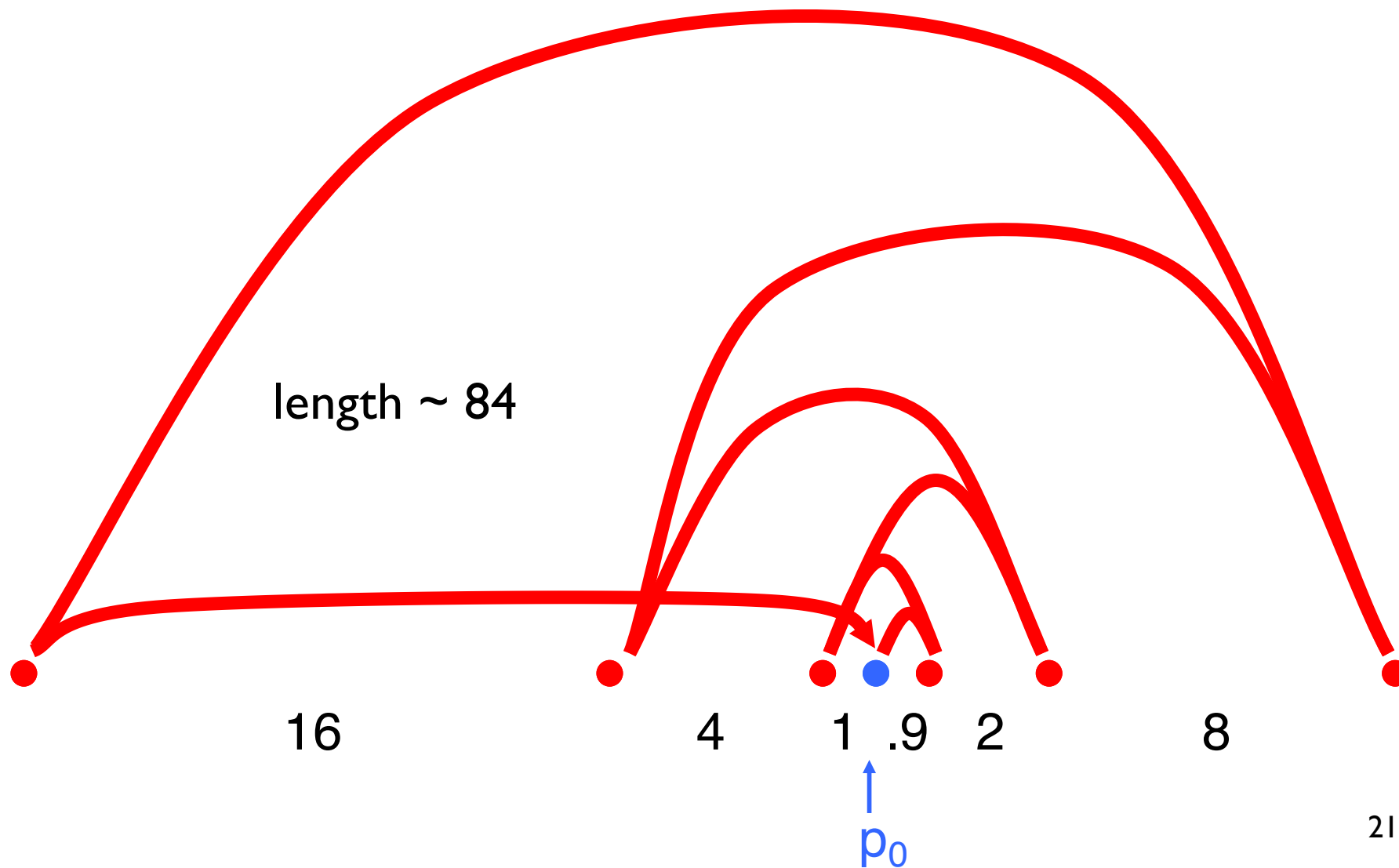
Then walk back to  $p_0$

**heuristic:** A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but *not* proven to give the best or fastest solution.

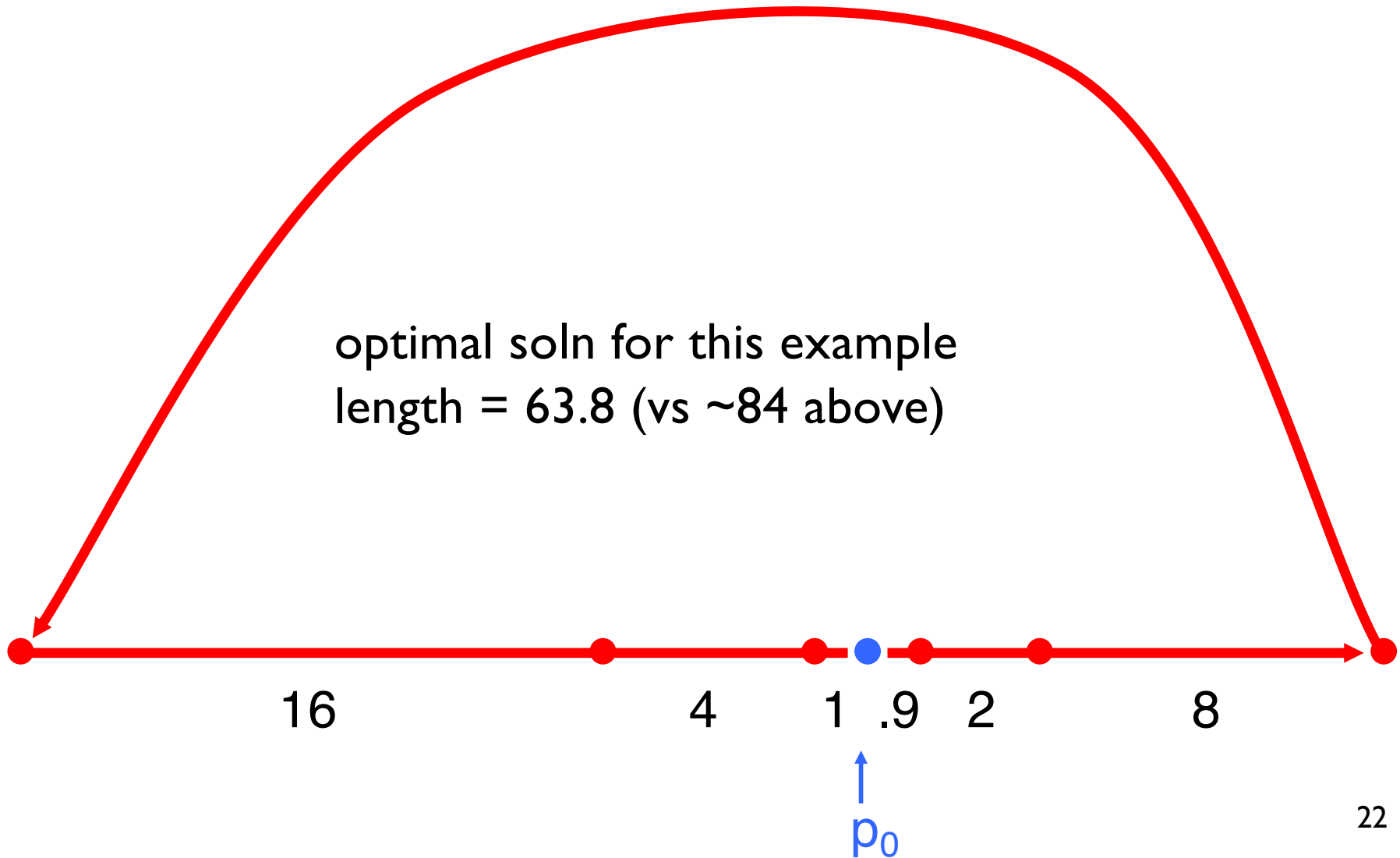
# Nearest Neighbor Heuristic



# An input where NN works badly



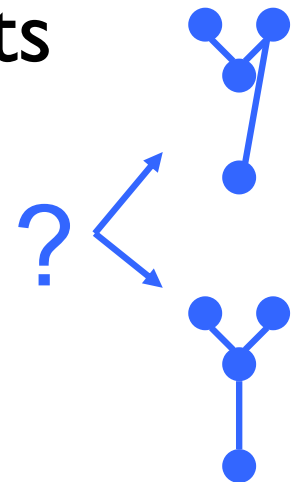
# An input where NN works badly



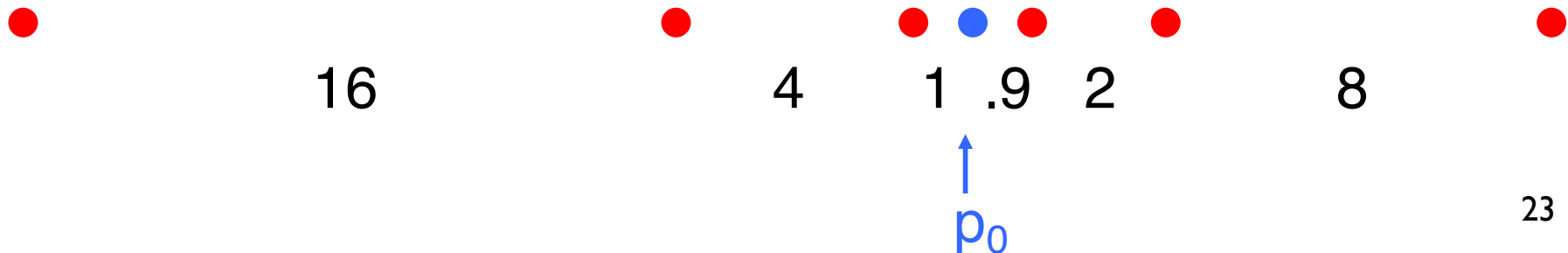
# Revised idea - Closest pairs first

Repeatedly join the closest pair of points

(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)

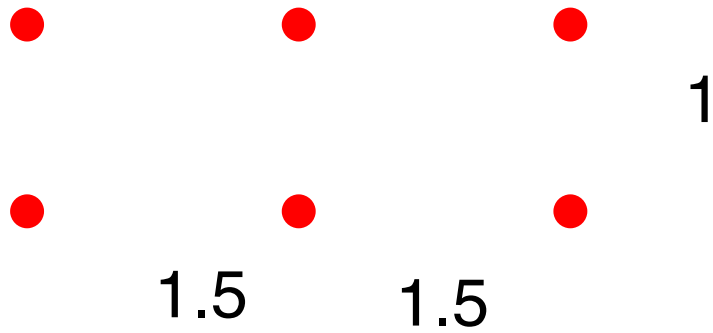


How does this work on our bad example?

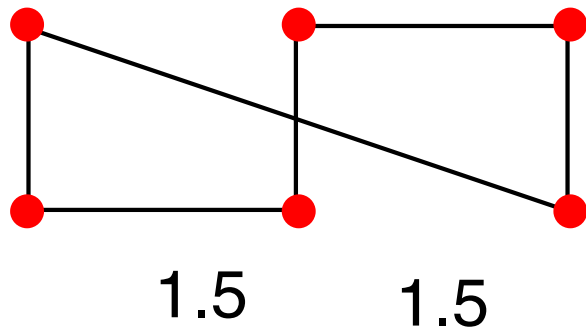




# A bad example for “close pairs”



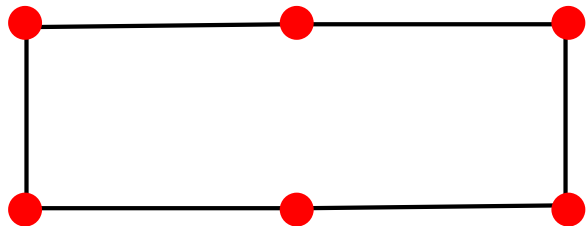
# A bad example for “close pairs”



1

$$6 + \sqrt{10} \approx 9.16$$

VS



8

# Something that works

“Brute Force Search”:

For each of the  $n! = n(n-1)(n-2)\dots 1$  orderings of the points, check the length of the cycle you get

Keep the best one

(Easy to see that it's correct, but slow!)

# Two Notes

The two *incorrect* algorithms were “greedy”

Often very natural & tempting ideas

They make choices that look great “locally” (and never reconsider them)

When greed works, the algorithms are typically efficient

BUT: often does not work - you get boxed in

Our correct alg avoids this, but is incredibly slow

20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)

And *growing*:  $n! \sim \sqrt{2\pi n} \cdot (n/e)^n \sim 2^{O(n \log n)}$

# The Morals of the Story

Algorithms are important

Many software gains outstrip hardware gains (Moore's law)

Simple problems can be hard

Factoring, TSP

Simple ideas don't always work

Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

Brute-force factoring, TSP

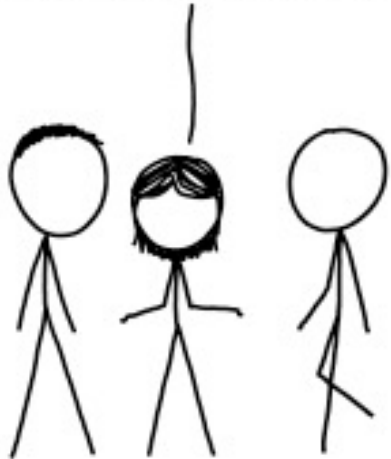
For some problems, even the *best* algorithms are slow

Course Goals:

formalize these ideas, and

develop more sophisticated approaches

OUR FIELD HAS BEEN STRUGGLING WITH THIS PROBLEM FOR YEARS.



STRUGGLE NO MORE! I'M HERE TO SOLVE IT WITH *ALGORITHMS!*



SIX MONTHS LATER:

WOW, THIS PROBLEM IS REALLY HARD.

YOU DON'T SAY.

