

## Algorithm

```

VertexCover(graph G, int k)
  if(G has no edges) //we've covered them all!
    return true
  if(k < 0)
    return false
  H1 = copy of G
  H2 = copy of G
  pick any edge (u,v)
  H1 = H1.remove(u) //removes u and all edges with u as
an endpoint
  H2 = H2.remove(v) //removes u and all edges with u as
an endpoint
  return VertexCover(H1,k-1) || VertexCover(H2, k-1)

```

## Running Time

$$\text{Recurrence: } T(k) = \begin{cases} 2T(k-1) + O(n+m) & \text{if } k \geq 1 \\ O(1) & \text{if } k < 0 \end{cases}$$

Running time? Unroll or use recursion tree

## Approximation Ratio

For a minimization problem (find the shortest/smallest/least/etc.)

If  $OPT(G)$  is the value of the best solution for  $G$ , and  $ALG(G)$  is the value that your algorithm finds, then  $ALG$  is an  $\alpha$  approximation algorithm if for every  $G$ ,

$$\alpha \cdot OPT(G) \geq ALG(G)$$

i.e. you're within an  $\alpha$  factor of the real best.

## Vertex Cover LP

Minimize  $\sum w(u) \cdot x_u$

Subject to:

$x_u + x_v \geq 1$  for all  $(u, v) \in E$

$0 \leq x_u \leq 1$  for all  $u$ .

Don't worry about the weights for today.

We got an exact solution for bipartite graphs....