# CSE 417: Algorithms and Computational Complexity

Lecture I: Overview

Spring 2014 Larry Ruzzo



## University of Washington Computer Science & Engineering

CSE 417, Sp '14: Algorithms & Computational Complexity

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FAQ

Schedule & Reading

#### Course Email/BBoard

**Subscription Options Class List Archive** E-mail Course Staff **GoPost BBoard** 

#### Homework

1: Assignment **Electronic Turnin** 

#### **Lecture Notes**

1: Intro

Lecture: MWF 12:30-1:20 EEB 037 (schematic)

> Office Hours Location Phone

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TAs: CSE xxx Yanling He, heyl@cs Jianghong Shi, jhshi@cs tba CSE xxx CSE xxx Tianhui Shi, tianhui@cs tba

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agorithms for manipulating graphs and strings. Fast and space complexity. NP-complete problems and

http://courses.cs.washington.edu/4/7 Midterm, Final. Homework will be a mix of paper & pencil exercises and programing. Overall weights 55%, 15%,

Late Policy: Papers and/or electronic turnins are due at the start of class on the due date. 10% off for up to one day late (business day, e.g., Monday for Friday due dates); additional 20% per day thereafter.

Textbooks: Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006. (Available from U Book Store, Amazon, etc.)

### What you'll have to do

Homework (~55% of grade)

**Programming** 

Several small projects

Written homework assignments

English exposition and pseudo-code

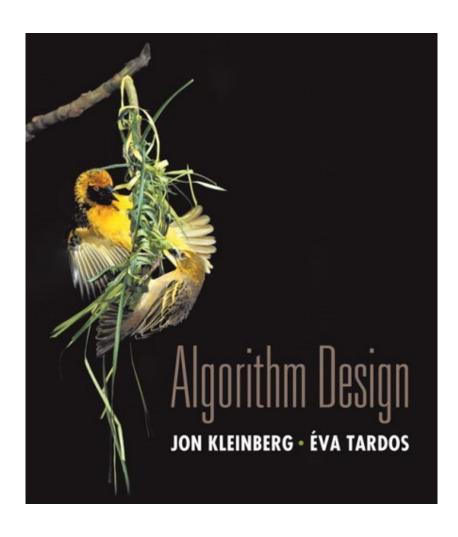
Analysis and argument as well as design

Midterm / Final Exam (~15% / 30%)

Late Policy:

Papers and/or electronic turnins are due at the *start* of class on the due date. 10% off for one day late (Monday, for Friday due dates); 20% per day thereafter.

### **Textbook**



Algorithm Design by Jon Kleinberg and Eva Tardos. Addison Wesley, 2006.

### What the course is about

#### Design of Algorithms

design methods
common or important types of problems
analysis of algorithms - efficiency
correctness proofs

### What the course is about

Complexity, NP-completeness and intractability solving problems in principle is not enough algorithms must be efficient

some problems have no efficient solution

NP-complete problems

important & useful class of problems whose solutions (seemingly) cannot be found efficiently, but *can* be checked easily

### Very Rough Division of Time

Algorithms (7 weeks)

Analysis of Algorithms

Basic Algorithmic Design Techniques

Graph Algorithms

Complexity & NP-completeness (3 weeks)

Check online schedule page for (evolving) details



### University of Washington Computer Science & Engineering

CSE 417, Wi '06: Approximate Schedule

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		Due	Lecture Topic	Reading
Week 1 1/2-1/6	М		Holiday	
	W		Intro, Examples & Complexity	Ch. 1; Ch. 2
	F		Intro, Examples & Complexity	
Week 2 1/9-1/13	М		Intro, Examples & Complexity	
	W		Graph Algorithms	Ch. 3
	F		Graph Algorithms	

### Complexity Example

```
Cryptography (e.g., RSA, SSL in browsers)
    Secret: p,q prime, say 512 bits each
    Public: n which equals p \times q, 1024 bits
In principle
    there is an algorithm that given n will find p and q:
   try all 2^{5/2} > 1.3 \times 10^{154} possible p's: kinda slow...
In practice
    no fast algorithm known for this problem (on non-quantum computers)
    security of RSA depends on this fact
    ("quantum computing": strongly driven by possibility of changing this)
```

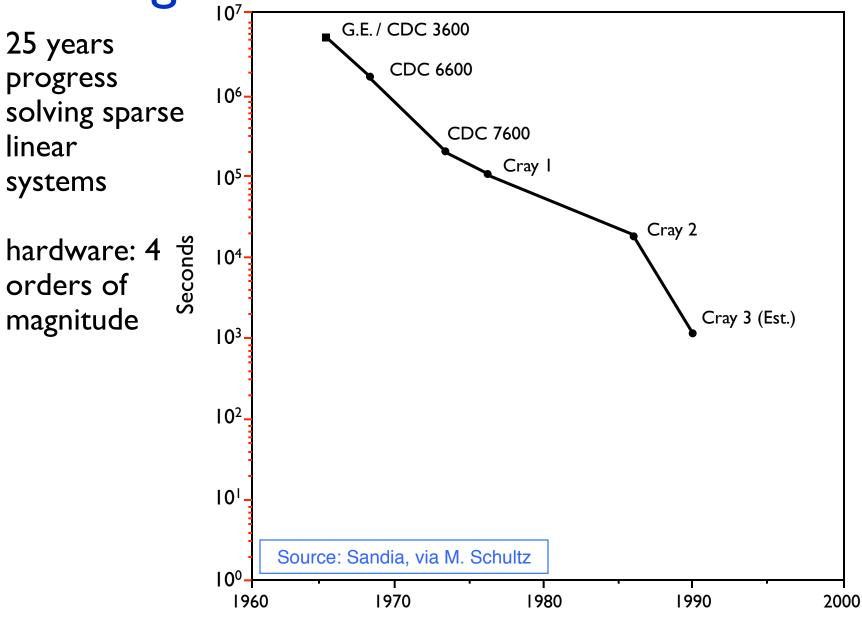
### Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

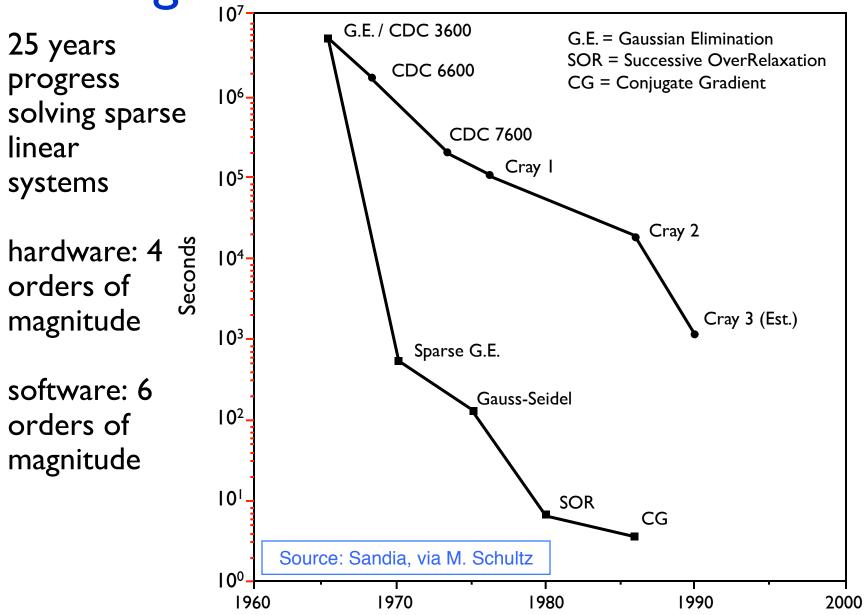
Ex: sparse linear equations over 25 years

10 orders of magnitude improvement!

Algorithms or Hardware?



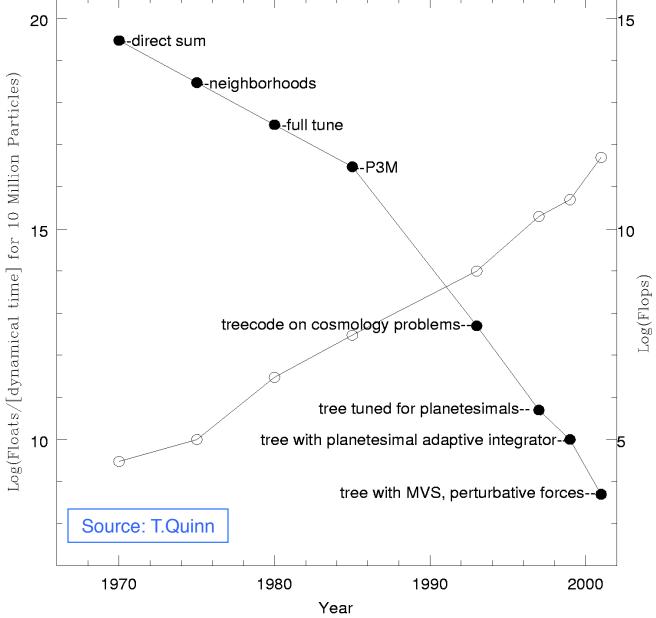
Algorithms or Hardware?



Algorithms or Hardware?

The N-Body Problem:

in 30 years 10<sup>7</sup> hardware 10<sup>10</sup> software



### Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them

"accomplish" via simple, well-defined steps

Ex: sorting names (via comparison)

Ex: checking for primality (via +, -, \*, /,  $\leq$ )

### Goals

```
Correctness
often subtle
Analysis
often subtle
Generality, Simplicity, 'Elegance'
Efficiency
time, memory, network bandwidth, ...
```

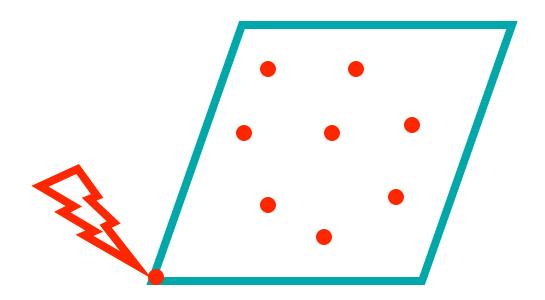
### Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board

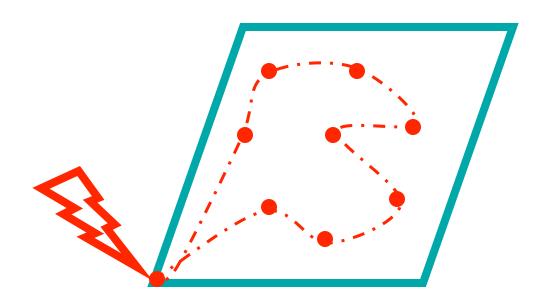
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position

For each board design, find best order to do the soldering

### Printed Circuit Board



### Printed Circuit Board



### A Well-defined Problem

Input: Given a set S of n points in the plane Output: The shortest cycle tour that visits each point in the set S once.

Better known as "TSP"

How might you solve it?

## Nearest Neighbor Heuristic

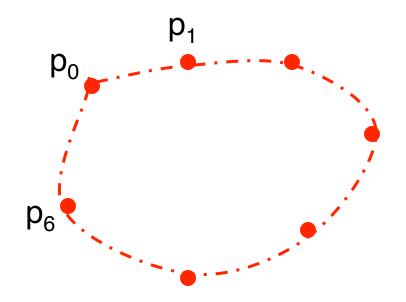
Start at some point  $p_0$ Walk first to its nearest neighbor  $p_1$ 

#### heuristic:

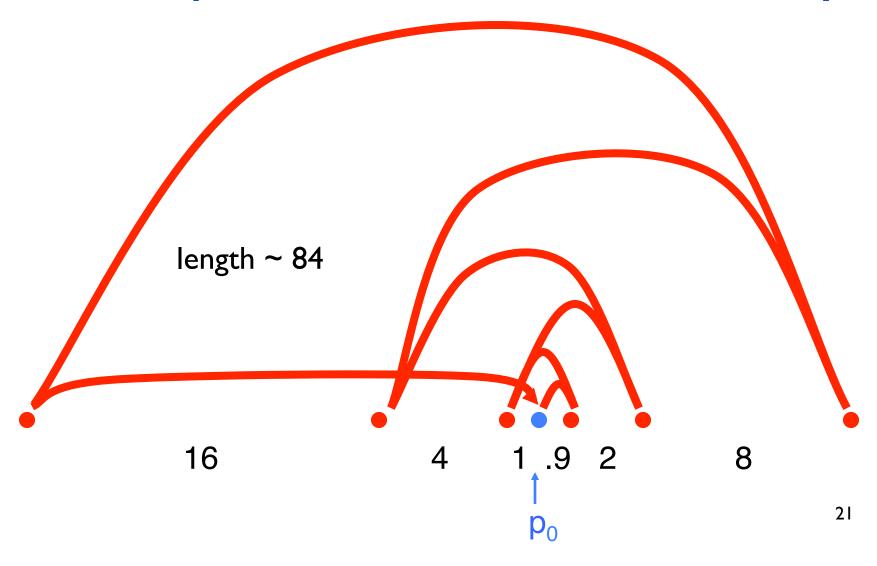
A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually *not* guaranteed to give the best or fastest solution.

Repeatedly walk to the nearest unvisited neighbor  $p_2$ , then  $p_3$ ,... until all points have been visited. Then walk back to  $p_0$ 

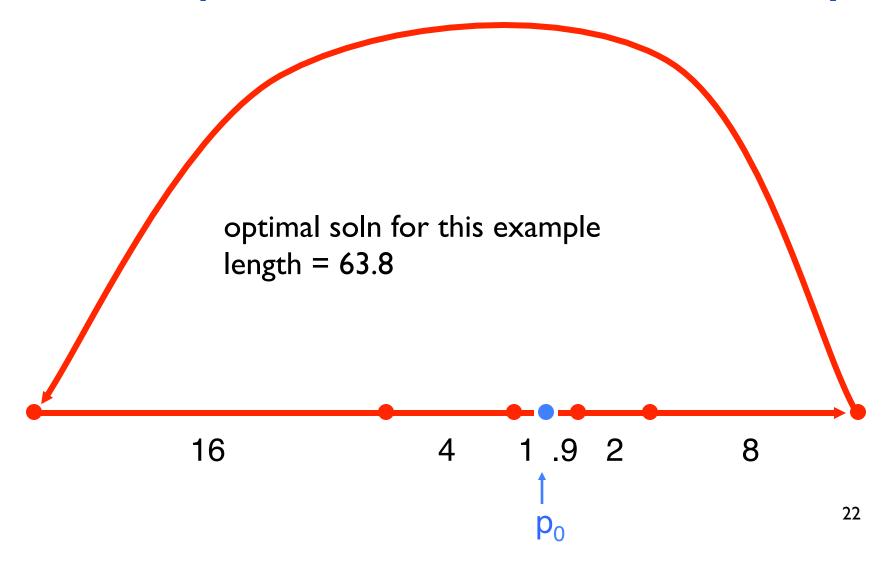
## Nearest Neighbor Heuristic



## An input where NN works badly



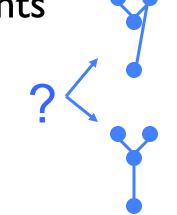
### An input where NN works badly



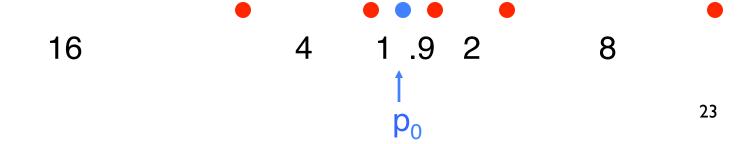
### Revised idea - Closest pairs first

Repeatedly join the closest pair of points

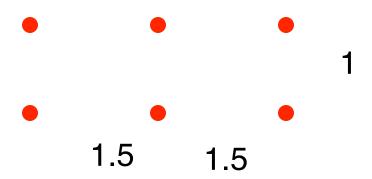
(s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)



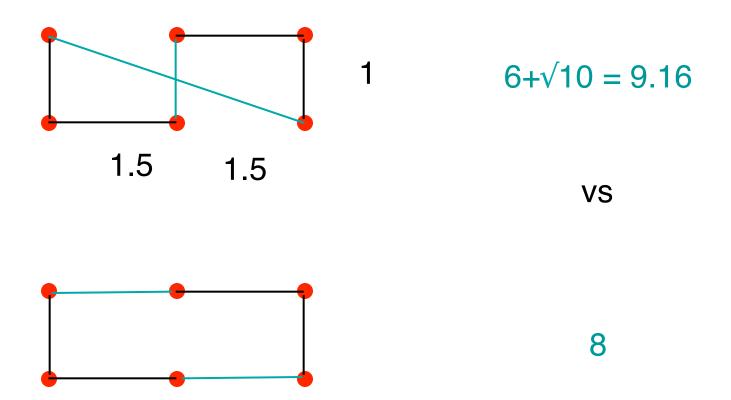
How does this work on our bad example?



## Another bad example



### Another bad example



### Something that works

"Brute Force Search":

For each of the n! = n(n-1)(n-2)...I orderings of the points, check the length of the cycle you get Keep the best one

### Two Notes

#### The two incorrect algorithms were greedy

Often very natural & tempting ideas

They make choices that look great "locally" (and never reconsider them)

When greed works, the algorithms are typically efficient

BUT: often does not work - you get boxed in

#### Our correct alg avoids this, but is incredibly slow

20! is so large that checking one billion orderings per second would take 2.4 billion seconds (around 70 years!)

And growing: n! ~  $\sqrt{2 \pi n} \cdot (n/e)^n \sim 2^{O(n \log n)}$ 

### The Morals of the Story

#### Algorithms are important

Many performance gains outstrip Moore's law

Simple problems can be hard

Factoring, TSP

Simple ideas don't always work

Nearest neighbor, closest pair heuristics

Simple algorithms can be very slow

Brute-force factoring, TSP

For some problems, even the best algorithms are slow

#### Course Goals:

formalize these ideas, and develop more sophisticated approaches