# CSE 4I7: Algorithms and Computational Complexity 

# Lecture I: Overview 

Winter 2009<br>Larry Ruzzo



## University of Washington <br> Computer Science \& Engineering

CSE 417, Wi '07: Algorithms \& Computational Complexity

Administrative
FAQ
Schedule \& Reading
Email
Class List Archive
E-mail Course Staff

## Assignments

## Solutions

## Lecture Notes

1: Overview \& Example ('06)
2-3: Analysis ('06)
4-6: Graphs, B/DFS ('06)
7-11: Greedy
Scheduling ('06)
Huffman ('06)
12-15: Dyn. Prog. Fibonaccl ('06)
Stamps ('06)
RNA Structure ('06)
Scheduling ('06)
16-19: Divide \& Conor-
20-26: $P$ \& Ar http:I

| Lecture: Low 101 (schematic) | MWF 2:30- 3:20 |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  |  |
|  |  | Office Hours | Phone |
| Instructor: Larry Ruzzo, zuzzo at cs | MF? 12:00-1:00? CSE 554 206-543-6298 |  |  |
| TA: | Zizhen Yao, yzizhen at cs | TBA |  |

Course Email: cse417a_wi07@u.washington.edu. Use this list to ask and/or answer questions about homework, lectures, etc. The instructor and TA are subscribed to this list. All messages are automatically archived. Questions not of general interest may be directed to the instructor and TA: cse417-staff or just to the instructor: ruzzo@cs. You can (and perhaps should) change voי- bscription options.
Catalog Description: Design and analysis of algorithms and do*manipulating graphs and strings. Fast Fourier Tranefmachines. Time and space complexity. $\mathrm{Nr}^{\mathrm{r}}$

## Prerequisite: $\operatorname{CSE} 373$

## Credits: 3

Iwww.cs.

## What you'll have to do

Homework (~55\% of grade)
Programming
Several small projects
Written homework assignments
English exposition and pseudo-code
Analysis and argument as well as design
Midterm / Final Exam (~15\% / 30\%)
Late Policy:
Papers and/or electronic turnins are due at the start of class on the due date. 10\% off for one day late (Monday, for Friday due dates); 20\% per day thereafter.

## Textbook



Algorithm Design by
Jon Kleinberg and
Eva Tardos. Addison Wesley, 2006.

## What the course is about

Design of Algorithms
design methods
common or important types of problems analysis of algorithms - efficiency correctness proofs

## What the course is about

Complexity, NP-completeness and intractability solving problems in principle is not enough
algorithms must be efficient
some problems have no efficient solution
NP-complete problems
important \& useful class of problems whose solutions (seemingly) cannot be found efficiently, but can be checked easily

## Very Rough Division of Time

Algorithms (7 weeks)
Analysis of Algorithms
Basic Algorithmic Design Techniques
Graph Algorithms
Complexity \& NP-completeness (3 weeks)

Check online schedule page for (evolving) details


CSE 417, Wi '06: Approximate Schedule
CSE Home
About Us $>$

|  |  | Due | Lecture Topic | Reading |
| :---: | :--- | :--- | :--- | :--- |
| Week 1 <br> $\mathbf{1 / 2 - 1 / 6}$ | M |  | Holiday |  |
|  | W |  | Intro, Examples \& Complexity | Ch. 1; Ch. 2 |
|  | F |  | Intro, Examples \& Complexity |  |
| Week 2 <br> $\mathbf{1 / 9 - 1 / 1 3}$ | M |  | Intro, Examples \& Complexity |  |
|  | W |  | Graph Algorithms | Ch. 3 |
|  | F |  | Graph Algorithms |  |

## Complexity Example

Cryptography (e.g. RSA, SSL in browsers)
Secret: p,q prime, say 512 bits each
Public: n which equals $\mathrm{p} \times \mathrm{q}, 1024$ bits
In principle
there is an algorithm that given $n$ will find $p$ and $q$ :
try all $2^{512}$ possible p's, an astronomical number
In practice
no efficient algorithm is known for this problem
security of RSA depends on this fact

## Algorithms versus Machines

We all know about Moore's Law and the exponential improvements in hardware...

Ex: sparse linear equations over 25 years

IO orders of magnitude improvement!

## Algorithms or Hardware?

25 years progress solving sparse linear systems
hardware: 4 orders of magnitude


## Algorithms or Hardware?

25 years progress solving sparse linear systems
hardware: $4 \stackrel{\text { ® }}{\stackrel{\circ}{6}}$
orders of
$\sim$ magnitude
software: 6 orders of magnitude


## Algorithms or Hardware?

## The N-Body Problem:

in 30 years
$10^{7}$ hardware $10^{10}$ software


## Algorithm: definition

Procedure to accomplish a task or solve a well-specified problem

Well-specified: know what all possible inputs look like and what output looks like given them
"accomplish" via simple, well-defined steps
Ex: sorting names (via comparison)
Ex: checking for primality (via,,+- , /, $\leq$ )

## Algorithms: a sample problem

Printed circuit-board company has a robot arm that solders components to the board
Time: proportional to total distance the arm must move from initial rest position around the board and back to the initial position
For each board design, find best order to do the soldering

## Printed Circuit Board



## Printed Circuit Board



## A Well-defined Problem

Input: Given a set $S$ of $n$ points in the plane Output: The shortest cycle tour that visits each point in the set $S$.

Better known as "TSP"

How might you solve it?

## Nearest <br> Neighbor Heuristic

Start at some point $\mathrm{P}_{0}$
Walk first to its
nearest neighbor $\mathrm{P}_{\mathrm{I}}$

## heuristic:

A rule of thumb, simplification, or educated guess that reduces or limits the search for solutions in domains that are difficult and poorly understood. May be good, but usually not guaranteed to give the best or fastest solution.

Repeatedly walk to the nearest unvisited neighbor $\mathrm{P}_{2}$, then $\mathrm{P}_{3}, \ldots$ until all points have been visited Then walk back to $\mathrm{P}_{0}$

## Nearest Neighbor Heuristic



An input where NN works badly


An input where NN works badly


## Revised idea - Closest pairs first

Repeatedly join the closest pair of points (s.t. result can still be part of a single loop in the end. I.e., join endpoints, but not points in middle, of path segments already created.)


How does this work on our bad example?

16
4


8

## Another bad example



## Another bad example



## Something that works

"Brute Force Search":
For each of the $n!=n(n-I)(n-2) \ldots I$ orderings of the points, check the length of the cycle you get Keep the best one

## Two Notes

The two incorrect algorithms were greedy
Often very natural \& tempting ideas
They make choices that look great "locally" (and never reconsider them)
When greed works, the algorithms are typically efficient BUT: often does not work - you get boxed in
Our correct alg avoids this, but is incredibly slow 20 ! is so large that checking one billion per second would take 2.4 billion seconds (around 70 years!)

## Something that "works" (differently)

I. Find Min Spanning Tree


## Something that "works" (differently)

2. Walk around it


## Something that "works" (differently)

3. Take shortcuts (instead of revisiting)


## Something that "works" (differently): Guaranteed Approximation

Does it seem wacky?
Maybe, but it's always within a factor of 2 of the best tour!
deleting one edge from best tour gives a spanning tree, so Min spanning tree < best tour best tour $\leq$ wacky tour $\leq 2 *$ MST $<2 *$ best

## The Morals of the Story

Simple problems can be hard
Factoring, TSP
Simple ideas don't always work
Nearest neighbor, closest pair heuristics
Simple algorithms can be very slow
Brute-force factoring, TSP
Changing your objective can be good
Guaranteed approximation for TSP

