

# Introduction to Database Systems

## CSE 414

### Lecture 12: NoSQL

# Class Overview

- Unit 1: Intro
- Unit 2: Relational Data Models and Query Languages
- Unit 3: Non-relational data
  - NoSQL
  - Json
  - SQL++
- Unit 4: RDMBS internals and query optimization
- Unit 5: Parallel query processing
- Unit 6: DBMS usability, conceptual design
- Unit 7: Transactions
- Unit 8: Advanced topics (time permitting)

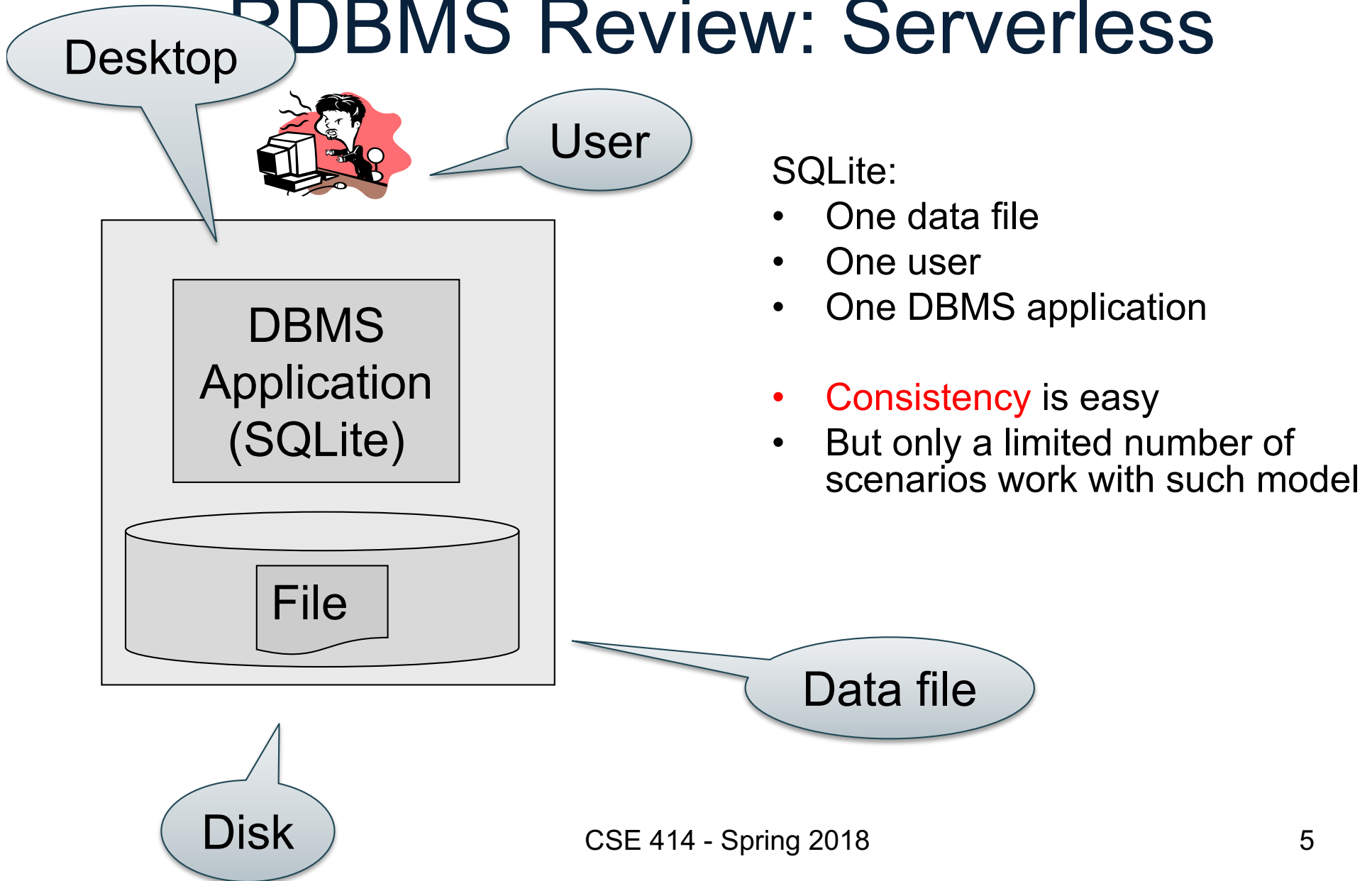
# Two Classes of Database Applications

- OLTP (Online Transaction Processing)
  - Queries are simple lookups: 0 or 1 join  
E.g., find customer by ID and their orders
  - Many updates. E.g., insert order, update payment
  - **Consistency** is critical: **transactions** (more later)
- OLAP (Online Analytical Processing)
  - aka “Decision Support”
  - Queries have many joins, and group-by’s  
E.g., sum revenues by store, product, clerk, date
  - No updates

# NoSQL Motivation

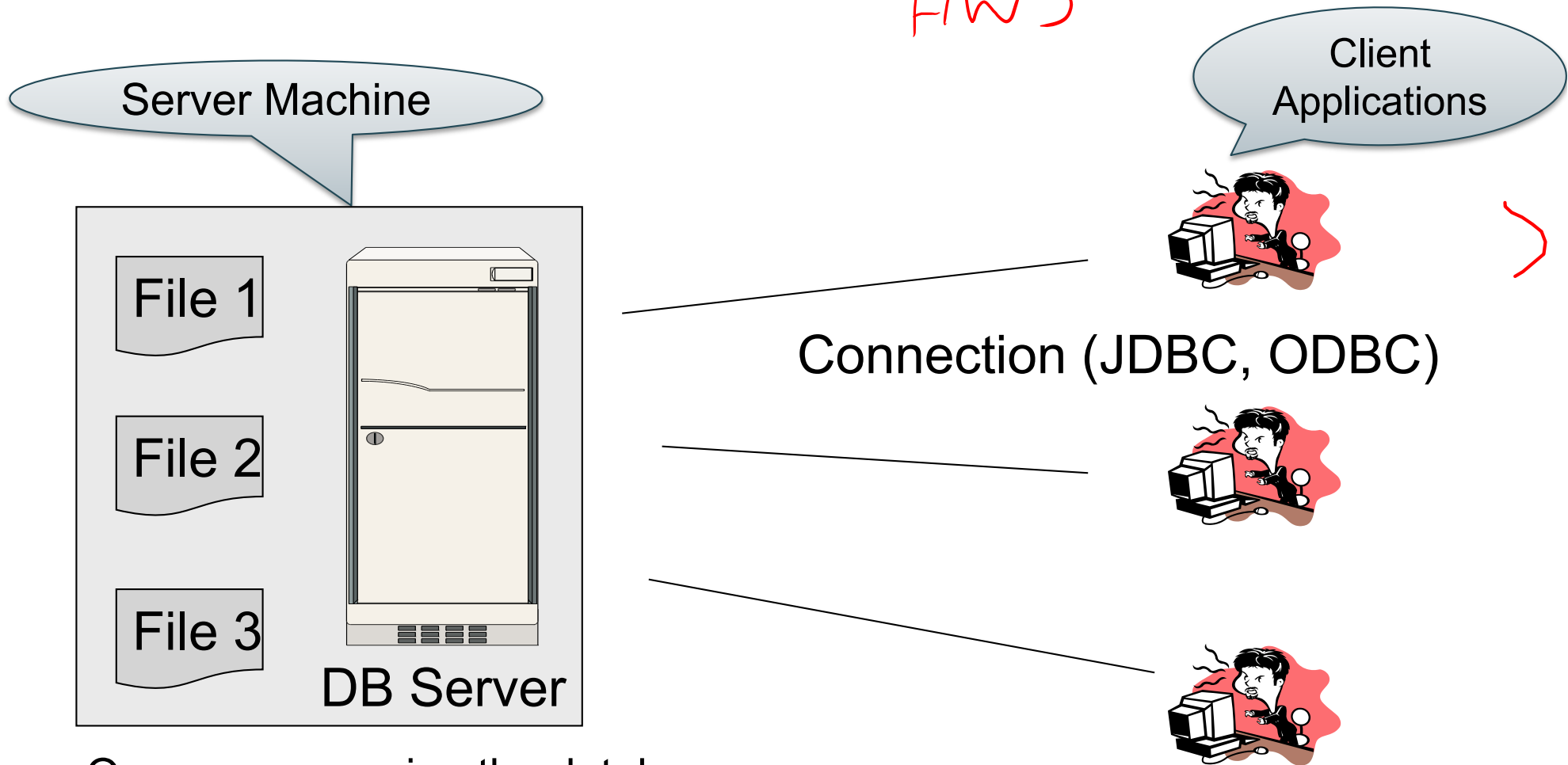
- Originally motivated by Web 2.0 applications
  - E.g. Facebook, Amazon, Instagram, etc
  - Web startups need to scaleup from 10 to 100000 users very quickly
- Needed: very large scale OLTP workloads
- Give up on consistency
- Give up OLAP

# DBMS Review: Serverless



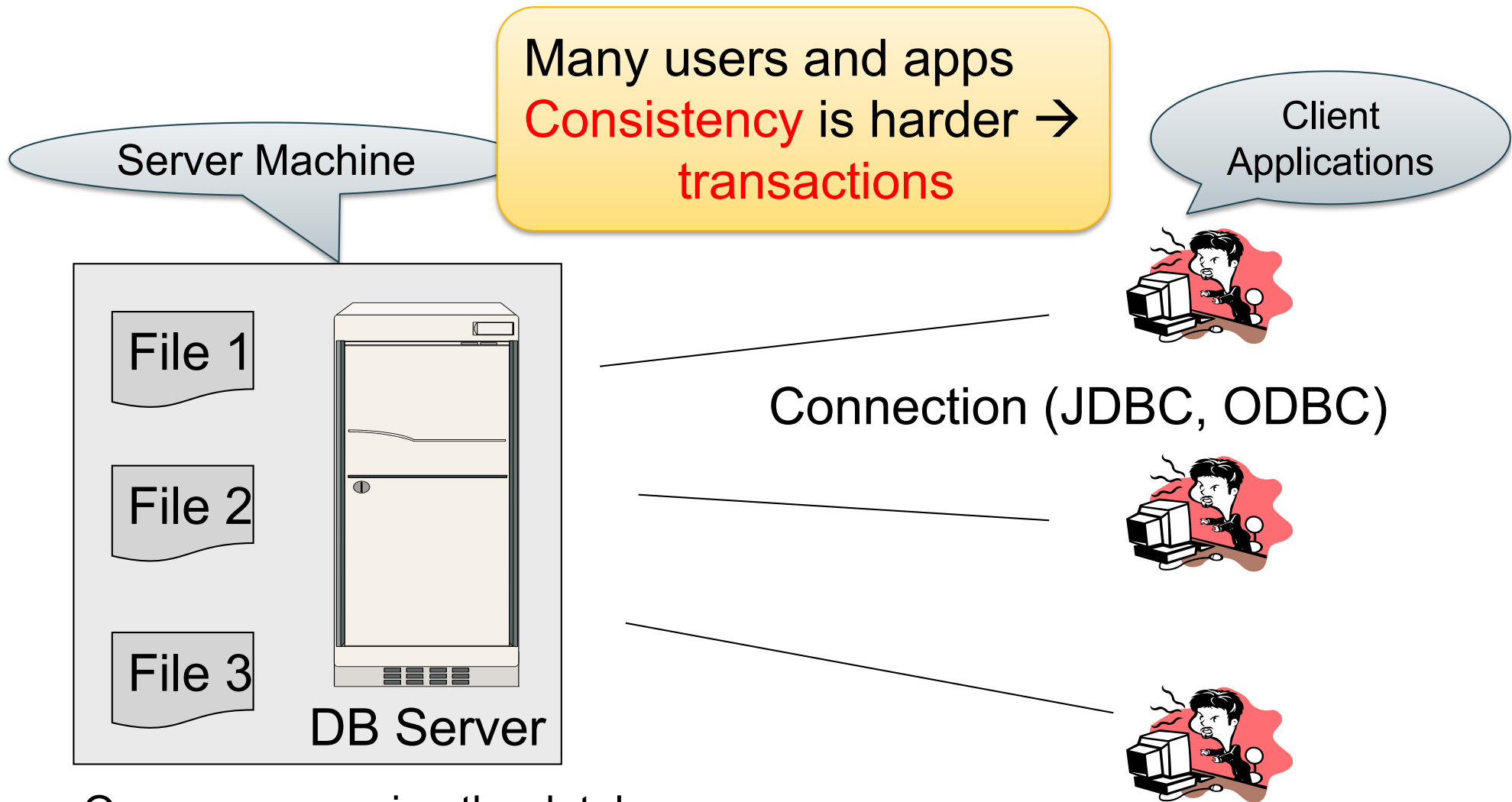
# RDBMS Review: Client-Server

HW 3



- One server running the database
- Many clients, connecting via the ODBC or JDBC (Java Database Connectivity) protocol

# RDBMS Review: Client-Server



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# Client-Server

- One *server* that runs the DBMS (or RDBMS):
  - Your own desktop, or
  - Some beefy system, or
  - A cloud service (SQL Azure)



# Client-Server

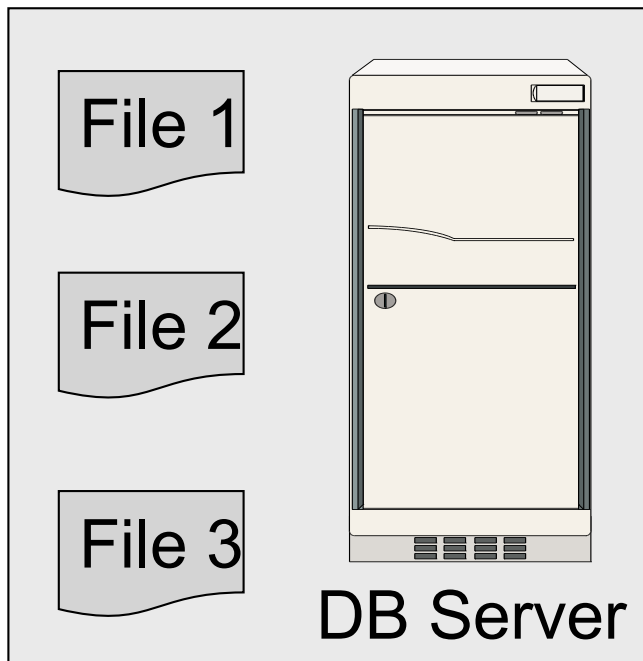
- One *server* that runs the DBMS (or RDBMS):
  - Your own desktop, or
  - Some beefy system, or
  - A cloud service (SQL Azure)
- Many *clients* run apps and connect to DBMS
  - Microsoft's Management Studio (for SQL Server), or
  - psql (for postgres)
  - Some Java program (HW8) or some C++ program

# Client-Server

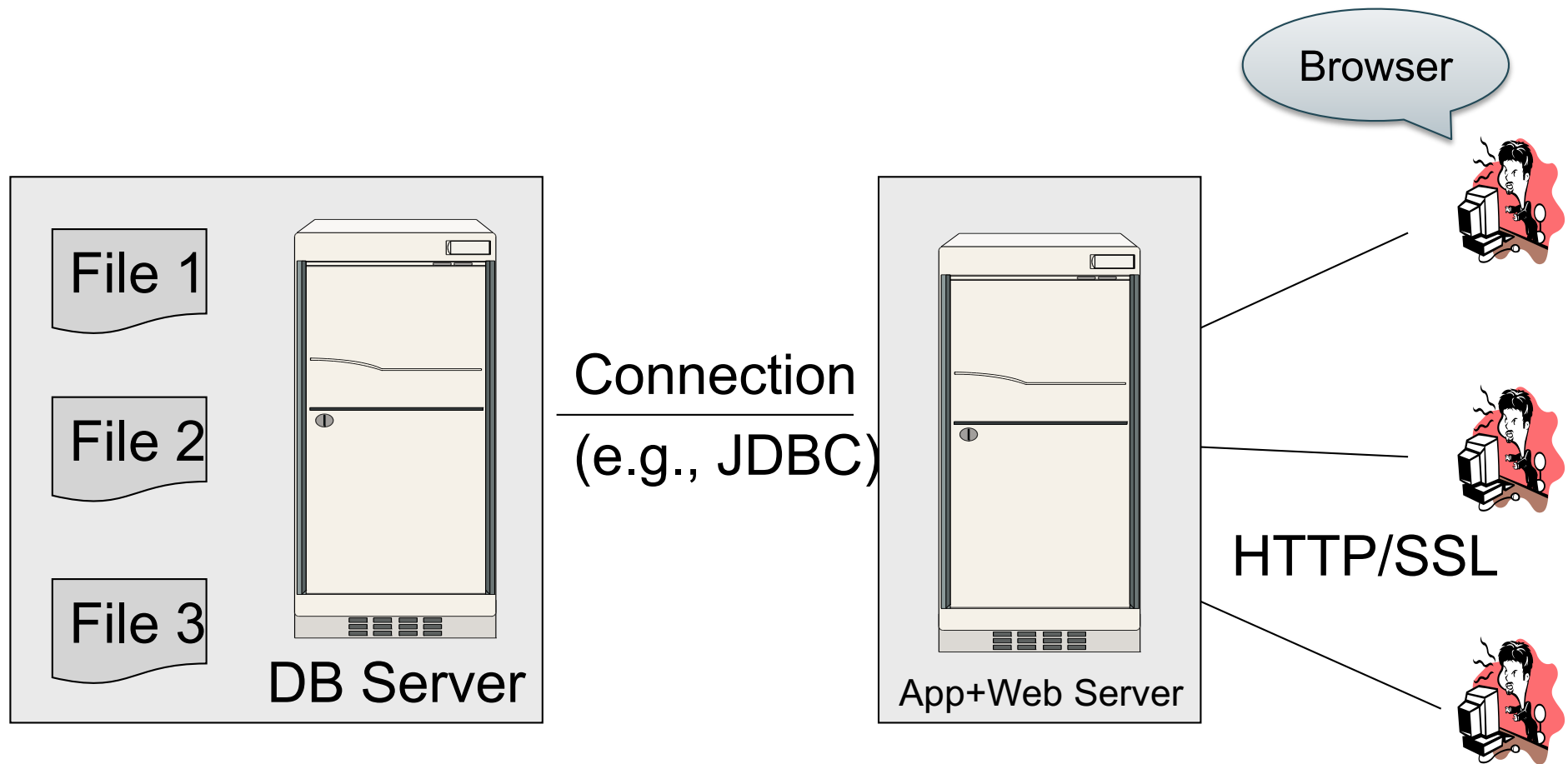
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- Clients “talk” to server using JDBC/ODBC protocol

# Web Apps: 3 Tier

Browser

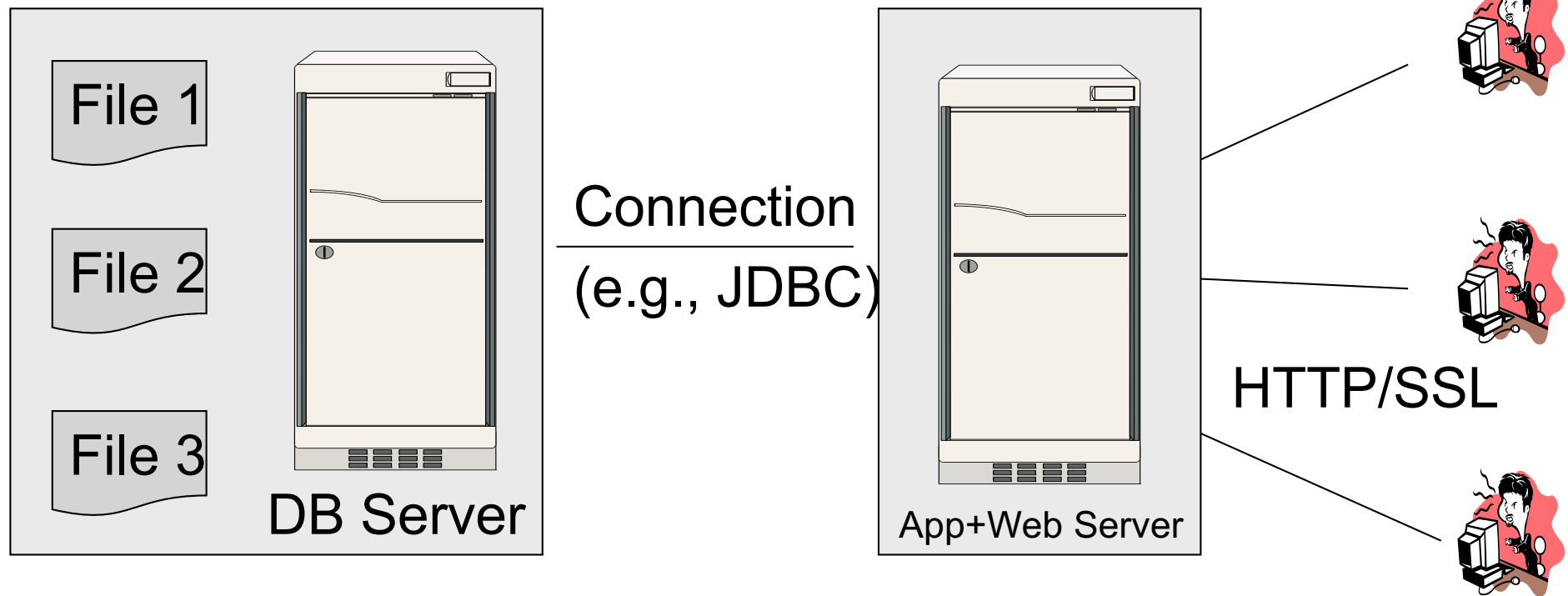


# Web Apps: 3 Tier



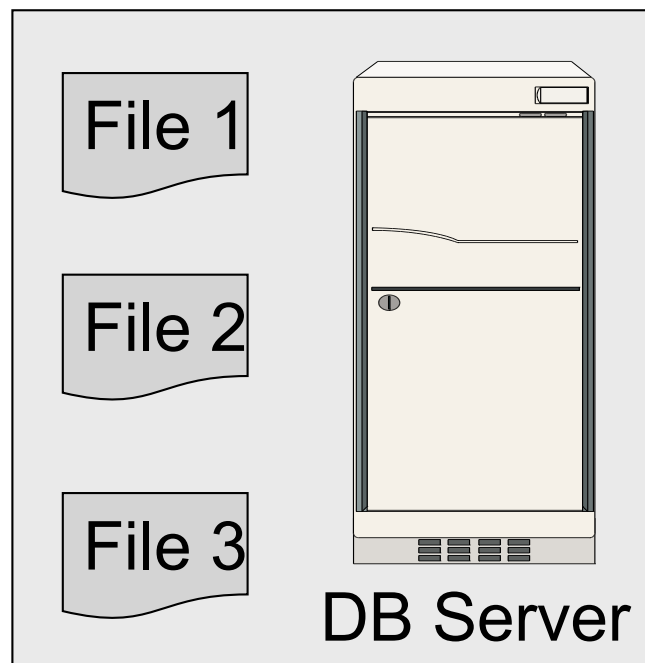
# Web Apps: 3 Tier

Web-based applications

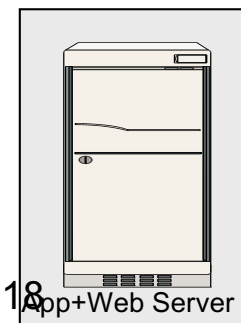
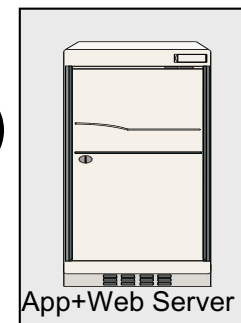
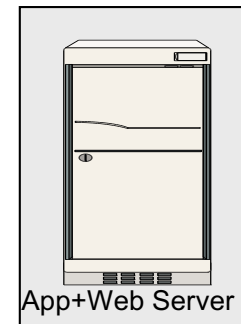


# Web Apps: 3 Tier

Web-based applications



Connection  
(e.g., JDBC)



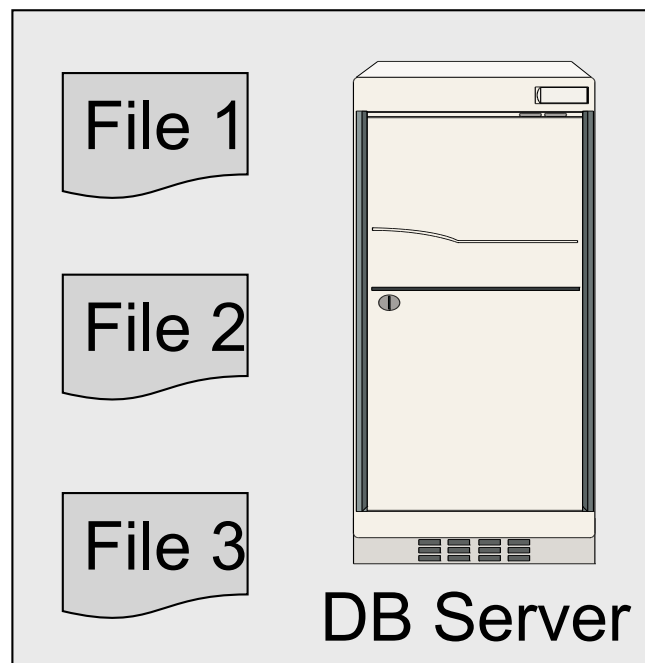
HTTP/SSL



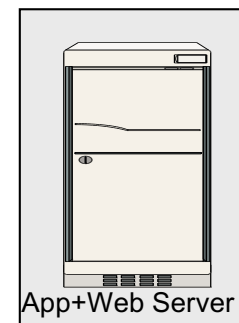
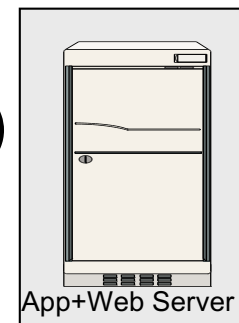
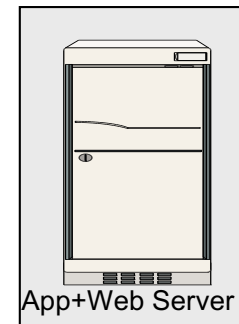
Replicate  
App server  
for scaleup

## os: 3 Tier

Web-based applications



Connection  
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HTTP/SSL

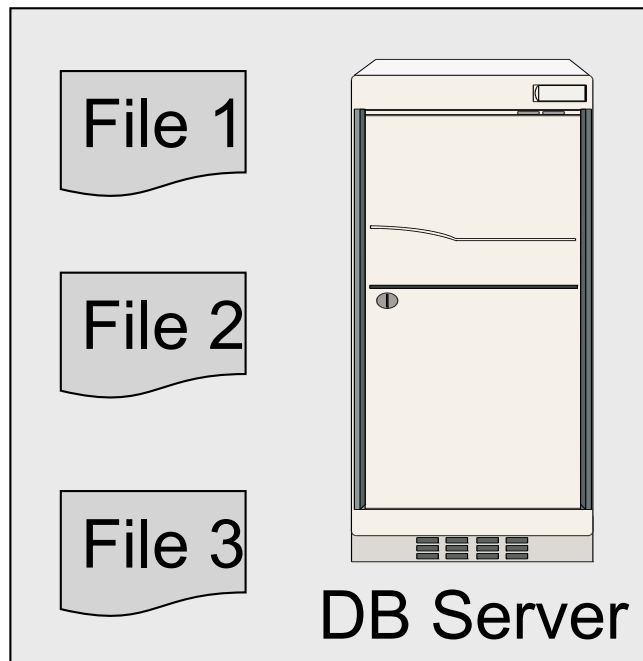
Why not replicate DB server?



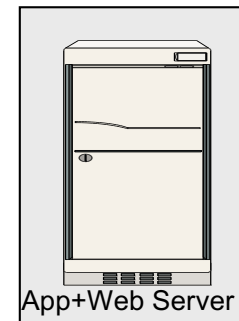
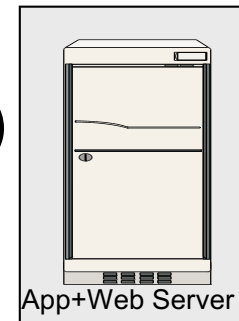
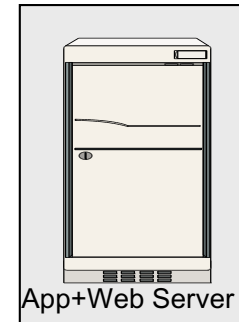
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HTTP/SSL

Why not replicate DB server?  
**Consistency!**



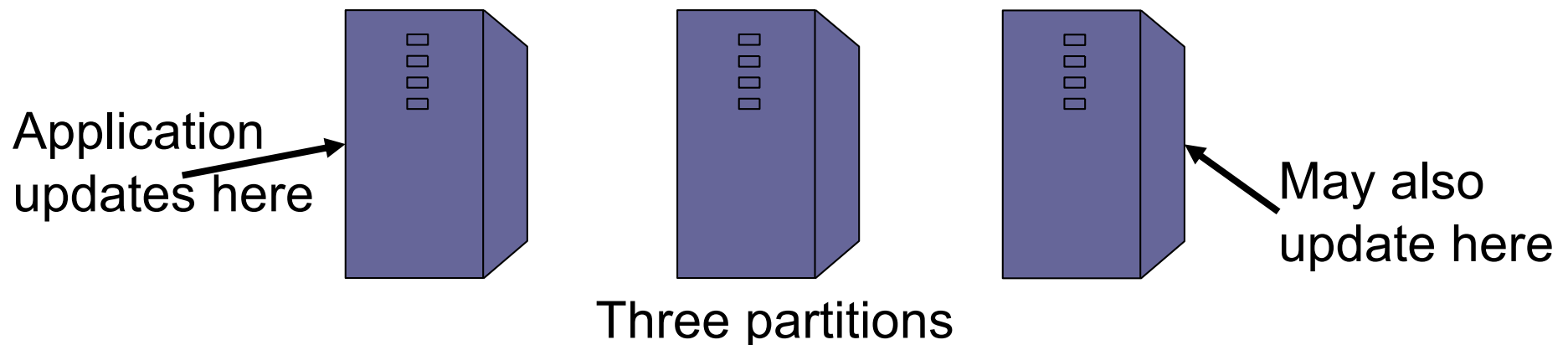


# Replicating the Database

- Two basic approaches:
  - Scale up through **partitioning**
  - Scale up through **replication**
- **Consistency** is much harder to enforce

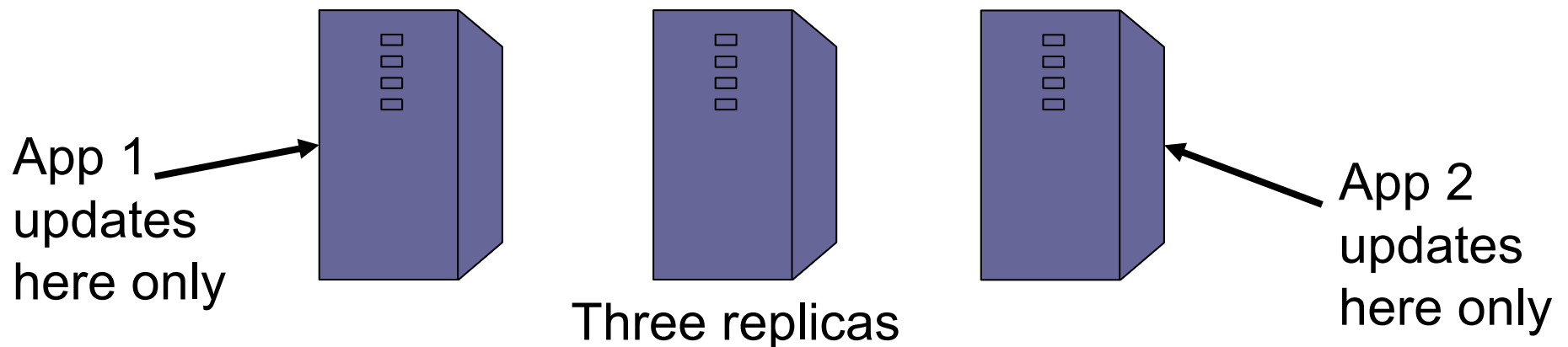
# Scale Through Partitioning

- Partition the database across many machines in a cluster
  - Database now fits in main memory
  - Queries spread across these machines
- Can increase throughput
- Easy for writes but reads become expensive!



# Scale Through Replication

- Create multiple copies of each database partition
- Spread queries across these replicas
- Can increase throughput and lower latency
- Can also improve fault-tolerance
- Easy for reads but writes become expensive!



# Relational Model → NoSQL

- Relational DB: difficult to replicate/partition
- Given `Supplier(sno,...)`, `Part(pno,...)`, `Supply(sno,pno)`
  - Partition: we may be forced to join across servers
  - Replication: local copy has inconsistent versions
  - Consistency is hard in both cases (why?)
- NoSQL: simplified data model
  - Given up on functionality
  - Application must now handle joins and consistency

# Data Models

Taxonomy based on data models:

- ☞ • **Key-value stores**
  - e.g., Project Voldemort, Memcached
- **Document stores**
  - e.g., SimpleDB, CouchDB, MongoDB

# Key-Value Stores Features

- **Data model:** (key,value) pairs
  - Key = string/integer, unique for the entire data
  - Value = can be anything (very complex object)

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  - `get(key)`, `put(key,value)`
  - Operations on value not supported

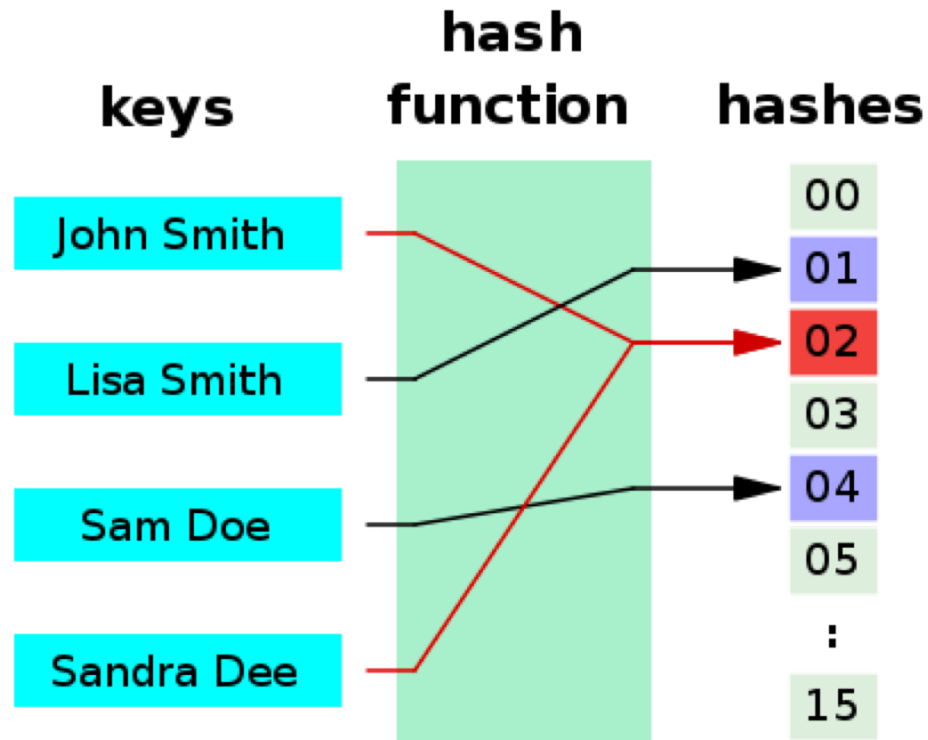
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  - Operations on value not supported
- **Distribution / Partitioning**



# Aside: Hash Functions

- A function that maps any data to a “hash value” (e.g., an integer)



# Aside: Hash Functions

- Example: data and hash value are integers
- Simple hash function:
  - $h(\text{key}) = \text{key} \% 42;$
  - $h(10) = 10$
  - $h(2) = 2$
  - $h(50) = 8$
- What does this have to do with data distribution?

# Key-Value Stores Features

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- **Operations**
  - `get(key)`, `put(key,value)`
  - Operations on value not supported
- **Distribution / Partitioning** – w/ hash function
  - No replication: key  $k$  is stored at server  $h(k)$
  - 3-way replication: key  $k$  stored at  $h1(k), h2(k), h3(k)$

How does `get(k)` work? How does `put(k,v)` work?

Flights(fid, date, carrier, flight\_num, origin, dest, ...)  
Carriers(cid, name)

## Example

- How would you represent the Flights data as key, value pairs?

How does query processing work?

Flights(fid, date, carrier, flight\_num, origin, dest, ...)  
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## Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record

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Flights(fid, date, carrier, flight\_num, origin, dest, ...)  
Carriers(cid, name)

## Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record
- Option 2: key=date, value=all flights that day

How does query processing work?

Flights(fid, date, carrier, flight\_num, origin, dest, ...)  
Carriers(cid, name)


## Example

- How would you represent the Flights data as key, value pairs?
- Option 1: key=fid, value=entire flight record
- Option 2: key=date, value=all flights that day
- Option 3: key=(origin,dest), value=all flights between

How does query processing work?

# Data Models

Taxonomy based on data models:

- Key-value stores
  - e.g., Project Voldemort, Memcached
-  • Document stores
  - e.g., SimpleDB, CouchDB, MongoDB



# Motivation

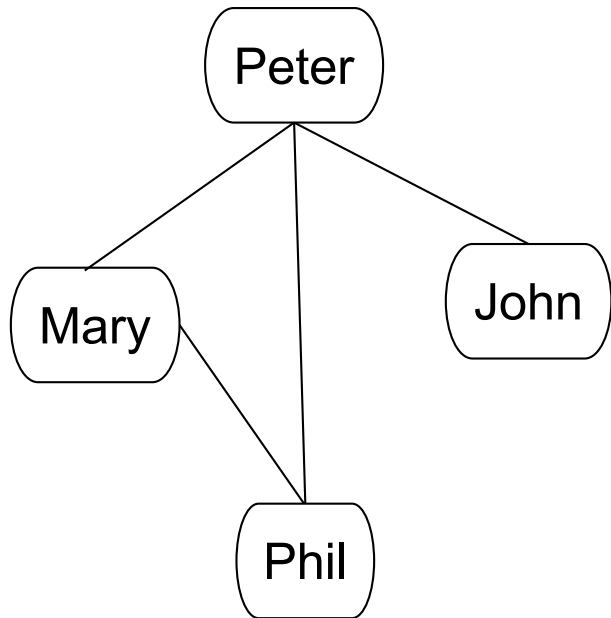
- In Key, Value stores, the Value is often a very complex object
  - Key = '2010/7/1', Value = [all flights that date]
- Better: allow DBMS to understand the *value*
  - Represent *value* as a JSON (or XML...) document
  - [all flights on that date] = a JSON file
  - May search for all flights on a given date

# Document Stores Features

- **Data model:** (key,document) pairs
  - Key = string/integer, unique for the entire data
  - Document = JSon, or XML
- **Operations**
  - Get/put document by key
  - Query language over JSon
- **Distribution / Partitioning**
  - Entire documents, as for key/value pairs

We will discuss JSon

# Example: storing FB friends



As a graph

Add new attributes?

Storing lists?

OR

Person1	Person2	is_friend
Peter	John	1
John	Mary	0
Mary	Phil	1
Phil	Peter	1
...	...	...

As a relation

We will learn the tradeoffs of different data models later this quarter

# JSON

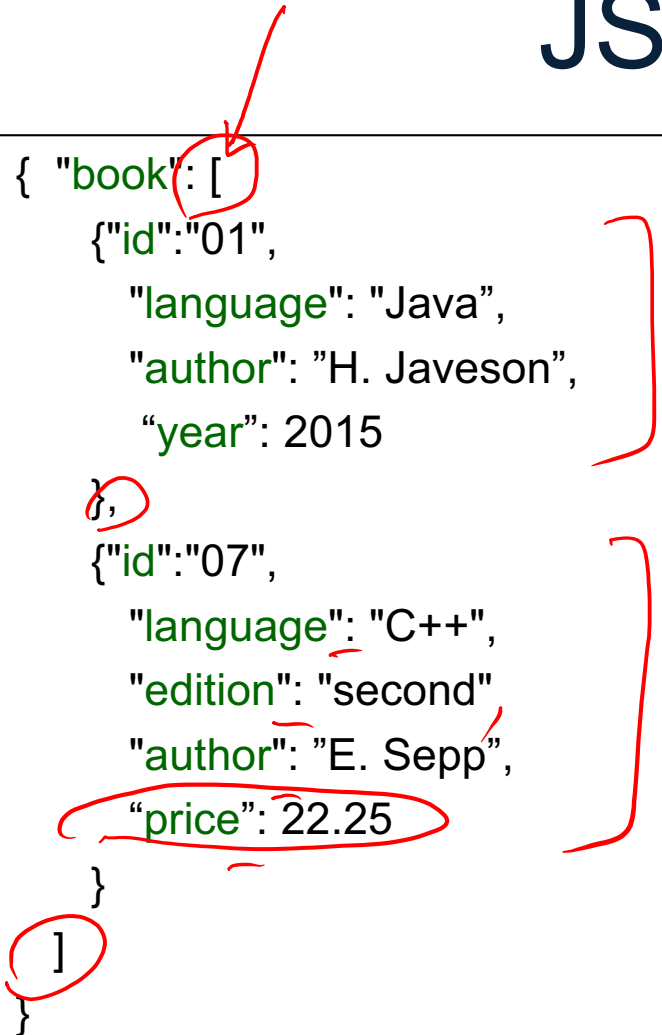
# JSON - Overview

- JavaScript Object Notation = lightweight text-based open standard designed for human-readable data interchange. Interfaces in C, C++, Java, Python, Perl, etc.
- The filename extension is .json.

We will emphasize JSon as semi-structured data

# JSon Syntax

```
{ "book": [  
  {"id": "01",  
   "language": "Java",  
   "author": "H. Javeson",  
   "year": 2015  
},  
  {"id": "07",  
   "language": "C++",  
   "edition": "second",  
   "author": "E. Sepp",  
   "price": 22.25  
}  
]  
}
```



# JSon vs Relational

- Relational data model
  - Rigid flat structure (tables)
  - Schema must be fixed in advanced
  - Binary representation: good for performance, bad for exchange
  - Query language based on Relational Calculus
- Semistructured data model / JSon
  - Flexible, nested structure (trees)
  - Does not require predefined schema ("self describing")
  - Text representation: good for exchange, bad for performance
  - Most common use: Language API; query languages emerging

# JSon Terminology

- Data is represented in name/value pairs.
- Curly braces hold objects
  - Each object is a list of name/value pairs separated by , (comma)
  - Each pair is a name is followed by ':'(colon) followed by the value
- Square brackets hold arrays and values are separated by ,(comma).



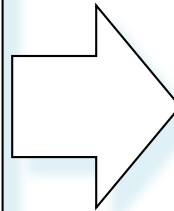
# JSON Data Structures

- Collections of name-value pairs:
  - {“name1”: value1, “name2”: value2, ...}
  - The “name” is also called a “key”
- Ordered lists of values:
  - [obj1, obj2, obj3, ...]

# Avoid Using Duplicate Keys

The standard allows them, but many implementations don't

```
{  
  "id": "07",  
  "title": "Databases",  
  "author": "Garcia-Molina",  
  "author": "Ullman",  
  "author": "Widom"  
}
```



```
{  
  "id": "07",  
  "title": "Databases",  
  "author": ["Garcia-Molina",  
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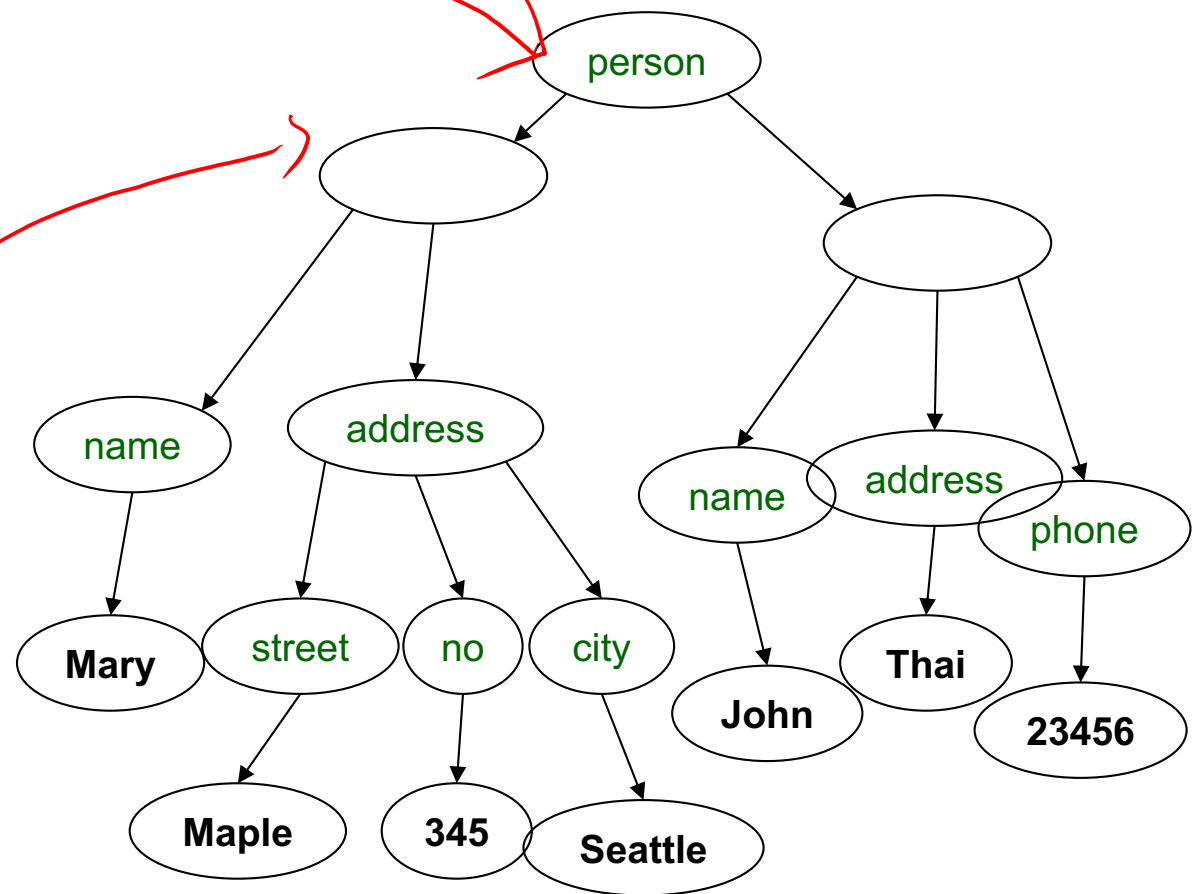
[1]

# JSon Datatypes

- Number
- String = double-quoted
- Boolean = true or false
- nullempty

# JSon Semantics: a Tree !

```
{  
  "person":  
    [  
      {  
        "name": "Mary",  
        "address": {  
          "street": "Maple",  
          "no": 345,  
          "city": "Seattle"  
        }  
      },  
      {  
        "name": "John",  
        "address": "Thailand",  
        "phone": 2345678  
      }  
    ]  
}
```



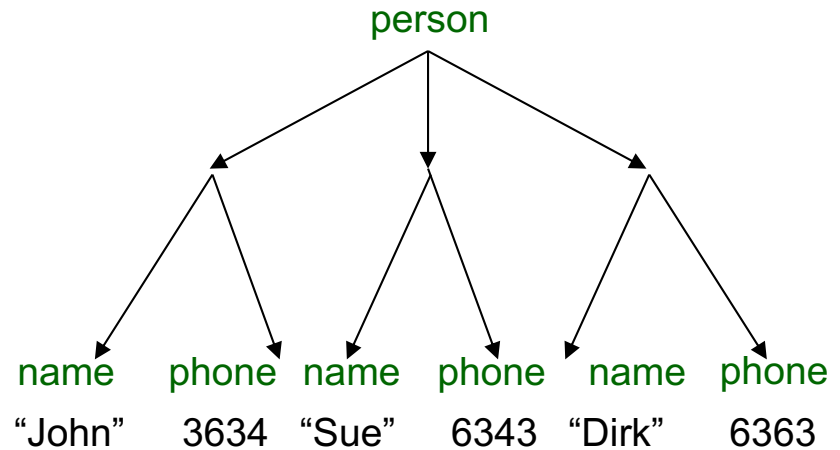
# JSon Data

- JSon is **self-describing**
- Schema elements become part of the data
  - Relational schema: `person(name,phone)`
  - In JSon “`person`”, “`name`”, “`phone`” are part of the data, and are repeated many times
- Consequence: JSon is much more flexible
- JSon = **semistructured** data

# Mapping Relational Data to JSON

## Person

name	phone
John	3634
Sue	6343
Dirk	6363



```
{ "person":
  [ { "name": "John", "phone": 3634 },
    { "name": "Sue", "phone": 6343 },
    { "name": "Dirk", "phone": 6383 }
  ]
}
```

# Mapping Relational Data to JSON

May inline foreign keys

## Person

name	phone
John	3634
Sue	6343

## Orders

personName	date	product
John	2002	Gizmo
John	2004	Gadget
Sue	2002	Gadget

```
{ "Person":  
  [{ "name": "John",  
    "phone": 3646,  
    "Orders": [{ "date": 2002,  
                  "product": "Gizmo" },  
                { "date": 2004,  
                  "product": "Gadget" }  
              ]  
    },  
    { "name": "Sue",  
      "phone": 6343,  
      "Orders": [{ "date": 2002,  
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                ]  
    }  
  ]  
}
```