# Introduction to Database Systems CSE 414

Lecture 6: SQL Subqueries

#### Announcements

Web Quiz 2 due Friday night

HW 2 due Tuesday at midnight

 Section this week important for HW 3, must attend

#### Announcements

 Many students did not turn in hw1 correctly – need to make sure your files are here:

https://gitlab.cs.washington.edu/cse414-2018au/cse414-[username]/tree/master/hw/hw[homework#]/submission

E.g. <a href="https://gitlab.cs.washington.edu/cse414-2018au/cse414-">https://gitlab.cs.washington.edu/cse414-2018au/cse414-</a> maas/tree/master/hw/hw1/submission

#### AND you have the hw1 tag here:

https://gitlab.cs.washington.edu/cse414-2018au/cse414-[username]/tags

- Commit, then use ./turnInHw.sh hw2 script.
- MUST have this correct for HW2

# Semantics of SQL With Group-By

**SELECT** S

FROM  $R_1, ..., R_n$ 

WHERE C1

**GROUP BY**  $a_1,...,a_k$ 

HAVING C2

FWGHOS

#### Evaluation steps:

- 1. Evaluate FROM-WHERE using Nested Loop Semantics
- 2. Group by the attributes a<sub>1</sub>,...,a<sub>k</sub>
- 3. Apply condition C2 to each group (may have aggregates)
- 4. Compute aggregates in S and return the result

# Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

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Problem: manufacturer is in Product, price is in Purchase...

# Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: manufacturer is in Product, price is in Purchase...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

# Aggregate + Join

For each manufacturer, compute how many products with price > \$100 they sold

Problem: manufacturer is in Product, price is in Purchase...

```
-- step 1: think about their join
SELECT ...
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
```

manu facturer	 price	
Hitachi	150	
Canon	300	
Hitachi	180	

```
-- step 2: do the group-by on the join SELECT x.manufacturer, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
and y.price > 100
GROUP BY x.manufacturer

CSE 414 - Autumn 2018
```

manu facturer	count(*)
Hitachi	2
Canon	1
	9

# Aggregate + Join

Variant:

For each manufacturer, compute how many products with price > \$100 they sold in each month

```
SELECT x.manufacturer, y.month, count(*)
FROM Product x, Purchase y
WHERE x.pid = y.product_id
  and y.price > 100
GROUP BY x.manufacturer, y.month
```

manu facturer	month	count(*)
Hitachi	Jan	2
Hitachi	Feb	1
Canon	Jan	3
		10



In the result of a group by query, there
is one row per group in the result

SELECT x.manufacturer, count(\*)
FROM Product x, Purchase y
WHERE x.pname = y.product
GROUP BY x.manufacturer

Count(\*) is not
0 because
there are no
tuples to count!

SELECT x.manufacturer, count(\*)
FROM Product x, Purchase y

WHERE x.pname = y.product

GROUP BY x.manufacturer

**FWGHOS** 

#### **Product**

#### **Purchase**

pname	manufacturer	
Gizmo	GizmoWorks	
Camera	Canon	
OneClick	Hitachi	

product	price	
Camera	150	
Camera	300	
OneClick	180	

#### Final results

manufacturer	Count(*)
Canon	2
Hitachi	1

#### Join(Product, Purchase)

pname	manu facturer	 manu facturer	price	
Camera	Canon	Canon	150	
Camera	Canon	Canon	300	
OneClick	Hitachi	Hitachi	180	

No GizmoWorks!



```
SELECT x.manufacturer, count(y.pid)
FROM Product x LEFT OUTER JOIN Purchase y
ON x.pname = y.product
GROUP BY x.manufacturer
```

Count(pid) is 0 when all pid's in the group are NULL

SELECT x.manufacturer, count(y.pid)

FROM Product x LEFT OUTER JOIN Purchase y

ON x.pname = y.product

GROUP BY x.manufacturer

#### **Product**

#### **Purchase**

pname	manufacturer	
Gizmo	GizmoWorks	
Camera	Canon	
OneClick	Hitachi	

product	price	
Camera	150	
Camera	300	
OneClick	180	

#### Left Outer Join(Product, Purchase)

pname	manufacturer	 product	price	
Camera	Canon	Camera	150	
Camera	Canon	Camera	300	
OneClick	Hitachi	OneClick	180	
Gizmo	GizmoWorks	 NULL	NULL	NULL

Why 0 for GizmoWorks?

#### Final results

manufact	urer	Count(y.pid)
Canoi	n	2
Hitach	ni	1
GizmoW	orks	0

GizmoWorks is paired with NULLs

```
SELECT x.manufacturer, count(*)
FROM Product x LEFT OUTER JOIN Purchase y
ON x.pname = y.product
GROUP BY x.manufacturer
```

#### **Product**

#### **Purchase**

price

150

300

180

pname	manufacturer			product	
Gizmo	GizmoWorks			Camera	
Camera	Canon			Camera	
OneClick	OneClick Hitachi			OneClick	

#### Left Outer Join(Product, Purchase)

pname	manufacturer	 product	price	
Camera	Canon	Camera	150	
Camera	Canon	Camera	300	
OneClick	Hitachi	OneClick	180	
Gizmo	GizmoWorks	 NULL	NULL	NULL

#### Final results

manufacturer	Count(*)			
Canon	2			
Hitachi	1			
GizmoWorks	1			

Probably not what we want!

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#### What we have in our SQL toolbox

- Projections (SELECT \* / SELECT c1, c2, ...)
- Selections (aka filtering) (WHERE cond, HAVING)
- Joins (inner and outer)
- Aggregates
- Group by
- Inserts, updates, and deletes

Make sure you read the textbook!

### Subqueries

 In the relational model, the output of a query is also a relation

 Can use output of one query as input to another

# Subqueries

- A subquery is a SQL query nested inside a larger query
- Such inner-outer queries are called nested queries
- A subquery may occur in:
  - A SELECT clause
  - A FROM clause
  - A WHERE clause
- Rule of thumb: avoid nested queries when possible
  - But sometimes it's impossible, as we will see



### Subqueries...

- Can return a single value to be included in a SELECT clause
- Can return a relation to be included in the FROM clause, aliased using a tuple variable
- Can return a single value to be compared with another value in a WHERE clause
- Can return a relation to be used in the WHERE or HAVING clause under an existential quantifier

# Subqueries...

#### Subqueries are often:

- Intuitive to write
- Slow

Be careful!

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

For each product return the city where it is manufactured

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
FROM Product X

"correlated subquery"
```

What happens if the subquery returns more than one city? We get a runtime error

```
(and SQLite simply ignores the extra values...)
```

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Whenever possible, don't use a nested queries:

```
SELECT X.pname, (SELECT Y.city
FROM Company Y
WHERE Y.cid=X.cid) as City
FROM Product X
```



```
SELECT X.pname, Y.city
FROM Product X, Company Y
WHERE X.cid=Y.cid
```

We have "unnested" the query

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Compute the number of products made by each company

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid (C.cid)
FROM Company C
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Compute the number of products made by each company

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

Better: we can unnest using a GROUP BY

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

#### But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

#### But are these really equivalent?

```
SELECT DISTINCT C.cname, (SELECT count(*)
FROM Product P
WHERE P.cid=C.cid)
FROM Company C
```

```
SELECT C.cname, count(*)
FROM Company C, Product P
WHERE C.cid=P.cid
GROUP BY C.cname
```

No! Different results if a company has no products

```
SELECT C.cname, count(pname)
FROM Company C LEFT OUTER JOIN Product P
ON C.cid=P.cid
GROUP BY C.cname
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

### 2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

```
SELECT X.pname
FROM (SELECT *
FROM Product AS Y
WHERE price > 20) as X
WHERE X.price < 500
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

#### 2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

```
SELECT X.pname
FROM (SELECT *
FROM Product AS Y
WHERE price > 20) as X
WHERE X.price < 500
```

Try to unnest this query!

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

#### 2. Subqueries in FROM

Find all products whose prices is > 20 and < 500

```
SELECT X.pname
FROM (SELECT *
        FROM Product AS Y
        WHERE price > 20) as X
WHERE X.price < 500</pre>
```

Side note: This is not a correlated subquery. (why?)

Try to unnest this query!

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

**Existential quantifiers** 

#### Using EXISTS:

```
SELECT DISTINCT C.cname

FROM Company C
WHERE EXISTS (SELECT *
FROM Product P
WHERE C.cid = P.cid and P.price < 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

#### Using IN

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price < 200)
```

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

#### Using ANY:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

**Existential quantifiers** 

#### Using ANY:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 > ANY (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

Not supported in sqlite

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

#### Now let's unnest it:

```
SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200
```

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Find all companies that make <u>some</u> products with price < 200

Existential quantifiers

#### Now let's unnest it:

```
SELECT DISTINCT C.cname
FROM Company C, Product P
WHERE C.cid = P.cid and P.price < 200</pre>
```

#### Existential quantifiers are easy!

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```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

Universal quantifiers

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

same as:

Find all companies that make only products with price < 200

Universal quantifiers

Universal quantifiers are hard! 😊

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies that make <u>some</u> product ≥ 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. <u>all</u> their products have price < 200

1. Find *the other* companies that make <u>some</u> product ≥ 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

2. Find all companies s.t. <u>all</u> their products have price < 200

```
SELECT DISTINCT C.cname
FROM Company C
WHERE C.cid NOT IN (SELECT P.cid
FROM Product P
WHERE P.price >= 200)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. all their products have price < 200

Universal quantifiers

### Using EXISTS:

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. all their products have price < 200

Universal quantifiers

### Using ALL:

```
SELECT DISTINCT C.cname
FROM Company C
WHERE 200 >= ALL (SELECT price
FROM Product P
WHERE P.cid = C.cid)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

Find all companies s.t. all their products have price < 200

Universal quantifiers

### Using ALL:

Not supported in sqlite

# Question for Database Theory Fans and their Friends

Can we unnest the universal quantifier query?

 We need to first discuss the concept of monotonicity

```
Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)
```

- Definition A query Q is monotone if:
  - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)

### Monotone Queries

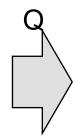
- Definition A query Q is monotone if:
  - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

#### **Product**

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

#### Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



pname	city
Gizmo	Lyon
Camera	Lodtz

Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)

### Monotone Queries

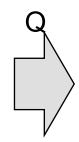
- Definition A query Q is monotone if:
  - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

#### **Product**

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

#### Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



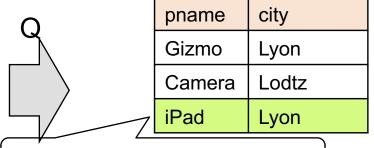
pname	city
Gizmo	Lyon
Camera	Lodtz

#### **Product**

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003
iPad	499.99	c001

#### Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



So far it looks monotone...

Product (<a href="mailto:pname">pname</a>, price, cid)
Company (<a href="mailto:cid">cid</a>, cname, city)

### Monotone Queries

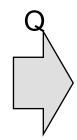
- Definition A query Q is monotone if:
  - Whenever we add tuples to one or more input tables, the answer to the query will not lose any of the tuples

#### **Product**

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003

#### Company

cid	cname	city
c002	Sunworks	Bonn
c001	DB Inc.	Lyon
c003	Builder	Lodtz



pname	city
Gizmo	Lyon
Camera	Lodtz

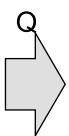
#### **Product**

pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c004
Camera	149.99	c003
iPad	499.99	c001

#### Company

cid	cname	city	
c002	Sunworks	Bonn	
c001	DB Inc.	Lyon	
c003	Builder <sub>414</sub> -	ALROHA 2	2018
c004	Crafter	Lodtz	

#### Q is not monotone!



$\overline{}$	
pname	city
Gizmo	Lodtz
Camera	Lodtz
iPad	Lyon

• Theorem: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.

- Theorem: If Q is a SELECT-FROM-WHERE query that does not have subqueries, and no aggregates, then it is monotone.
- Proof. We use the nested loop semantics: if we insert a tuple in a relation R<sub>i</sub>, this will not remove any tuples from the answer

```
SELECT a_1, a_2, ..., a_k
FROM R_1 AS x_1, R_2 AS x_2, ..., R_n AS x_n
WHERE Conditions
```

```
for X<sub>1</sub> in R<sub>1</sub> do
  for X<sub>2</sub> in R<sub>2</sub> do
  ...
  for X<sub>n</sub> in R<sub>n</sub> do
   if Conditions
    output (a<sub>1</sub>,...,a<sub>k</sub>)
```

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

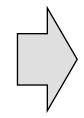
```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

pname	price	cid
Gizmo	19.99	c001

cid	cname	city
c001	Sunworks	Bonn



cname	
Sunworks	

Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)

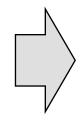
### Monotone Queries

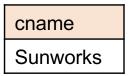
The query:

Find all companies s.t. <u>all</u> their products have price < 200 is not monotone

pname	price	cid
Gizmo	19.99	c001

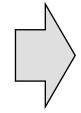
cid	cname	city
c001	Sunworks	Bonn

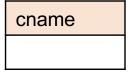




pname	price	cid
Gizmo	19.99	c001
Gadget	999.99	c001

cid	cname	city
c001	Sunworks	Bonn





 Consequence: If a query is not monotonic, then we cannot write it as a SELECT-FROM-WHERE query without nested subqueries

### Queries that must be nested

 Queries with universal quantifiers or with negation

### Queries that must be nested

- Queries with universal quantifiers or with negation
- Queries that use aggregates in certain ways
  - sum(..) and count(\*) are NOT monotone,
     because they do not satisfy set containment
  - select count(\*) from R is not monotone!

Author(<u>login</u>, name)
Wrote(login, url)

More Unnesting

Find authors who wrote ≥ 10 documents:

Author(login, name)
Wrote(login, url)

# More Unnesting

Find authors who wrote ≥ 10 documents:

Attempt 1: with nested queries

This is SQL by a novice

**SELECT DISTINCT Author.name** 

FROM Author

WHERE (SELECT count(Wrote.url)

**FROM** Wrote

WHERE Author.login=Wrote.login)

>= 10

Author(login, name)
Wrote(login, url)

# More Unnesting

Find authors who wrote ≥ 10 documents:

Attempt 1: with nested queries

Attempt 2: using GROUP BY and HAVING

SELECT Author.name
FROM Author, Wrote
WHERE Author.login=Wrote.login
GROUP BY Author.name
HAVING count(wrote.url) >= 10

```
Product (pname, price, cid)
Company (cid, cname, city)
Finding Witnesses
```

For each city, find the most expensive product made in that city

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

# Finding Witnesses

For each city, find the most expensive product made in that city Finding the maximum price is easy...

```
SELECT x.city, max(y.price)
FROM Company x, Product y
WHERE x.cid = y.cid
GROUP BY x.city;
```

But we need the witnesses, i.e., the products with max price

Product (<u>pname</u>, price, cid) Company (<u>cid</u>, cname, city)

# Finding Witnesses

To find the witnesses, compute the maximum price in a subquery (in FROM)

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

# Finding Witnesses

Or we can use a subquery in where clause

```
Product (<u>pname</u>, price, cid)
Company (<u>cid</u>, cname, city)
```

# Finding Witnesses

There is a more concise solution here: