# CSE 413: Programming Languages and their Implementation

Luke McDowell Autumn Quarter 2004

#### Today's Outline

- Administrative Info
- Survey
- Overview of the CourseIntroduction to Scheme

- Course Information
- Instructor: Luke McDowell, CSE 214 lucasm@cs.washington.edu
   Office hours: 1:30-2:20 Mon, 3:30-4:20 Wed
- Teaching Assistant: Lincoln Ritter lritter@cs.washington.edu
   Office hours: See web page
- Text: Concepts of Programming Languages, Robert W. Sebesta, Sixth Edition
   » Fifth edition is fine.
- Other references available from web page
  - » Revised<sup>5</sup> Report on the Algorithmic Language Scheme (R5RS)
  - » Link to Structure and Interpretation of Computer Programs

- **Course Policies**
- Homeworks
- » Turned in electronically before 11:59pm on due date
- » Late homework not accepted
- · Work in teams only on explicit team projects
- » Appropriate discussions encouraged see website
- » Must give credit for any such discussions on your homework
- Approximate Grading
- » Homework: 50%
- » Midterm:
  » Final:
- 25% Wed November 3, in class25% Tues December 14, 2:30-4:20

- Course Mechanics
- 413 Web page:
  - http://www.cs.washington.edu/413
- 413 mailing list
  - » cse413a\_au04@u.washington.edu
  - » You should automatically be included if enrolled
- Course labs : Math Science Computing Labs
  - » Basement of Communications building: B-022/027
  - » Or work from home all software available on course web

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#### Course Paper

- Slide handouts will be provided
  - » Also available on the web page
  - » Not...
- Homeworks not handed out, see the web page

#### That Survey Thing

- Why are you taking my picture?
- What if I forgot everything?
- What if I know this all already?
- What if I'm the famous one?

#### References

- Section 15.5, Concepts of Programming Languages
- Section 2, *Revised<sup>5</sup> Report on the Algorithmic Language Scheme (R5RS)*
- For more help:
  - » Sections 1-1.1.5, Structure and Interpretation of Computer Programs (Abelson, Sussman, & Sussman)

#### Elements of Programming

- Primitive expressions
  - » simplest entities of the language
- Means of combination
   » by which compound elements are built
- Means of abstraction
  - » by which compound elements can be named and manipulated as units

## There are many "languages"

- Computer programming
- Shell and scripting languages
- Applications
- Sciences

#### Training and Education

- Training
  - » learn the specifics of a known language
  - » build up a "tool chest" so that you can perform specific tasks in a particular field
- Education
  - » learn how to recognize valid abstractions and synthesize them in new and useful ways in many different knowledge domains
- We'll do some of both in this class

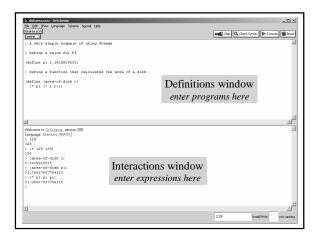
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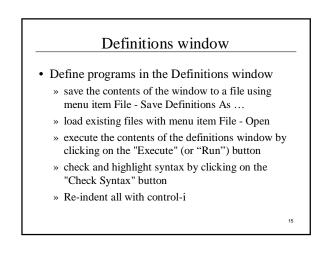
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#### Why Scheme?

- The simplicity of the language lets us work on problem solving, rather than just syntax issues
- Flexibility of the language lets us see that the structure of C/Java/Basic is not the only way to express problem solutions
- Variety is the spice of life
  - » study more than one language paradigm and study the relationship between design paradigms
  - » professional programmers switch languages every few years anyway, so start practicing now

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#### Interactions Window

- Evaluate simple expressions directly in the Interactions window
- Position the cursor after the ">", then type in your expression
  - » DrScheme responds by evaluating the expression and printing the result
  - » recall previous expression with escape-p
- Expressions can reference symbols defined when you executed the Definitions window

#### Think functionally

- Programming that makes extensive use of assignment is known as
- » The order of assignments changes the operation of the program because the state is changed by assignment
- Programming without the use of assignment statements is known as
  - » In such a language, all procedures implement welldefined mathematical functions of their arguments whose behavior does not change
- Scheme is heavily oriented towards *functional* style

#### Primitive Expressions

- constants
  - » integer :
  - » rational :
  - » real :
  - » boolean :
- variable names (symbols)
- » Names can contain almost any character except white space and parentheses
- » Stick with simple names like value, x, iter, ...

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# Compound Expressions

- Either a *combination* or a *special form*
- 1. Combination : (operator operand operand ...) » there are quite a few pre-defined operators
  - - 1
  - » We can define our own operators
- 2. Special form
  - » keywords in the language
  - » eg, define

#### Combinations

- (operator operand operand ...)
- this is *prefix* notation, the operator comes first
- a combination always denotes a procedure application
- the operator is a symbol or an expression, the applied procedure is the associated value
  - » +, -, abs, my-function, foop?
  - » characters like \* and + are not special; if they do not stand alone then they are part of some name

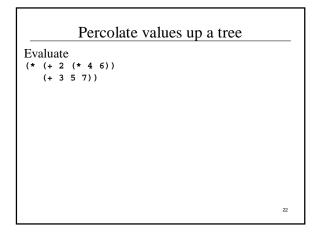
#### **Evaluating Combinations**

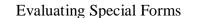
- To evaluate a combination
  - » Evaluate the subexpressions of the combination
  - » Apply the procedure that is the value of the leftmost subexpression (the operator) to the arguments that are the values of the other subexpressions (the operands)

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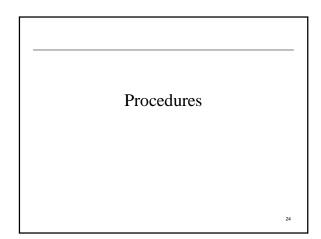
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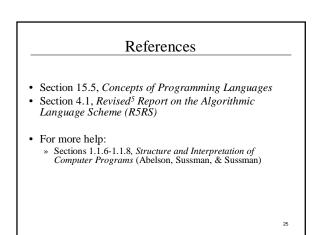
• For example





- Special forms have unique evaluation rules
- (define x 3) is an example of a special form; it is not a combination
  - » the evaluation rule for a simple define is "associate the given name with the given value"
- There are more special forms which we will encounter, but there are surprisingly few of them compared to other languages





#### Recall the *define* special form

- · Special forms have unique evaluation rules
- (define x 3) is an example of a special form; it is not a combination
  - » the evaluation rule for a simple define is "associate the given name with the given value"

#### Define and name a variable

- (define  $\langle name \rangle \langle expr \rangle$ )
  - » define special form
  - » name name that the value of expr is bound to
  - » *expr* expression that is evaluated to give the value for *name*
- define is valid only at the top level of a <program> and at the beginning of a <body>

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#### Define and name a procedure

- (define (*\(name\)\) \(formal params\)\)* > define special form
  - » *name* the name that the procedure is bound to
  - » *formal params* names used within the body of procedure
  - » body expression (or sequence of expressions) that will be evaluated when the procedure is called.
  - » The result of the last expression in the body will be returned as the result of the procedure call

#### Example definitions

(define pi 3.1415926535)

```
(define (area-of-disk r)
  (* pi (* r r)))
```

```
(define (area-of-ring outer inner)
 (- (area-of-disk outer)
      (area-of-disk inner)))
```

#### Defined procedures are "first class"

- Compound procedures that we define are used exactly the same way the primitive procedures provided in Scheme are used
  - » names of built-in procedures are not treated specially; they are simply names that have been pre-defined
  - » you can't tell whether a name stands for a primitive (built-in) procedure or a compound (defined) procedure by looking at the name or how it is used

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# Evaluation example

• (area-of-ring 4 1)

#### Booleans

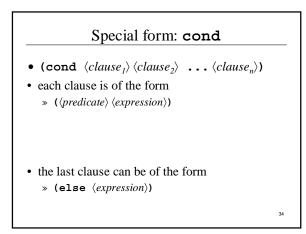
- Recall that one type of data object is boolean
   » #t (true) or #f (false)
- We can use these explicitly or by calculating them in expressions that yield boolean values
- An expression that yields a true or false value is called a predicate

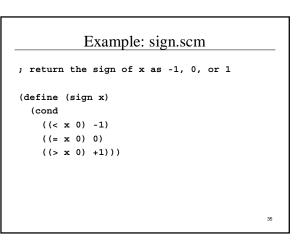
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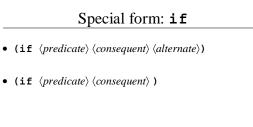
```
» #t =>
» (< 5 5) =>
```

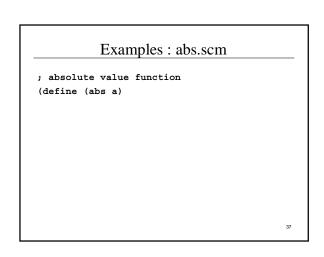
» (> pi 0) =>

# Conditional expressions • As in all languages, we need to be able to make decisions based on inputs and do something depending on the result Predicate Consequent









#### Examples : true-false.scm

; return 1 if arg is true, 0 if arg is false (define (true-false arg)

#### Logical composition

- (and  $\langle e_1 \rangle \langle e_2 \rangle ... \langle e_n \rangle$ )
- (or  $\langle e_1 \rangle \langle e_2 \rangle ... \langle e_n \rangle$ )
- (not  $\langle e \rangle$ )

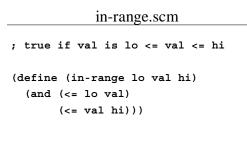
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• Scheme interprets the expressions *e<sub>i</sub>* one at a time in left-to-right order until it can tell the correct answer

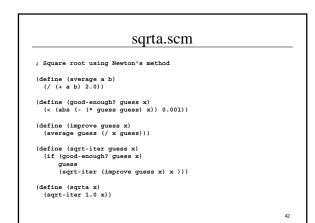
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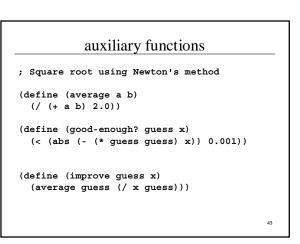
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#### Newton's method for square root

- Guess a value y for the square root of x
- Is it close enough to the desired value <sup>2</sup>√x ?
  » ie, is y<sup>2</sup> close to x?
- If yes, then done. Return recent guess.
- If no, then new guess is average of current guess and  $\frac{x}{guess}$
- Repeat with new guess





# iterator and main functions

```
(define (sqrt-iter guess x)
 (if (good-enough? guess x)
    guess
    (sqrt-iter (improve guess x) x )))
```

(define (sqrta x)
 (sqrt-iter 1.0 x))

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## sqrt-iter

- Our first example of recursion
- Note that this recursion is used to implement a loop (an iteration)
  - » We will see this over and over in Scheme
- Iteration is calling the same block of code with a changing set of parameters
- The syntax of the procedure is recursive but the resulting process is iterative

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» more on this later