
Testing and Branching

CSE 410, Spring 2006
Computer Systems

<http://www.cs.washington.edu/education/courses/410/06sp/>

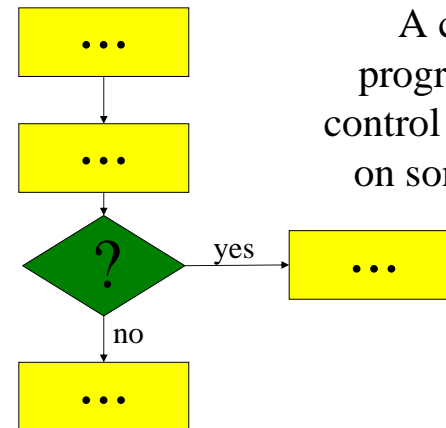
Reading and References

- *Computer Organization and Design*
 - » Section 2.5, Logical Operations
 - » Section 2.6, Instructions for Making Decisions
 - » (CD) Section A.9, SPIM
 - » (CD) Section A.10 through page A-50, MIPS R2000 Assembly Language

goto considered harmful

- “Oh what a tangled web we weave, When first we practice to deceive!”
 - » Sir Walter Scott
- Branching in assembly language can turn your program into a rat’s nest that cannot be debugged
- Keep control flow simple and logical
- Use comments describing the overall logic

Conditional Branch



A change in the program’s flow of control that depends on some condition

Branch instructions

- Branch instructions are I-format instructions
 - » op code field
 - » two register fields
 - » 16-bit offset field
- Simplest branches check for equality
 - » `beq $t0, $t1, address`
 - » `bne $t0, $t1, address`

`if (i==j) then a=b;`

- Assume all values are in registers
- Note that the test is inverted!

```
# $t0=i, $t1=j, $s0=a, $s1=b
```

```
    bne $t0, $t1, skip
```

```
    move $s0, $s1
```

```
skip:
```

`while (s[i]==k) i = i+j;`

```
# $s0=addr(s), $v1=i, $a0=k, $a1=j
```

```
loop:
```

```
    sll    $v0,$v1,2    # v0 = 4*i
    addu   $v0,$s0,$v0  # v0 = addr(s[i])
    lw     $v0,0($v0)   # v0 = s[i]
    addu   $v1,$v1,$a1  # i = i+j
    beq    $v0,$a0,loop # loop if equal
    subu   $v1,$v1,$a1  # i = i-j
```

`for (i=0; i<10; i++) s[i] = i;`

```
# $s0=addr(s), $t1=i
```

```
move    $t1,$zero    # i = 0
```

```
loop:
```

```
    sll    $t0,$t1,2    # t0 = i*4
    addu   $t0,$s0,$t0  # t0 = addr(s[i])
    sw     $t1,0($t0)   # s[i] = i
    addu   $t1,$t1,1    # i++
    slt    $t0,$t1,10   # if (i<10) $t0=1
    bnez   $t0,loop     # loop if (i<10)
```

How do we encode the destination?

- Calculating the destination address
 - » $4 \times$ (the 16-bit offset value)
 - » is added to the Program Counter (PC)
- The offset is a word offset in this case
- The base register is always the PC, so we don't need to specify it in the instruction
- Covers a range of 2^{16} words (64 KW)

Comparison instructions

- For comparisons other than equality
 - » `slt` : set less than
 - » `sltu` : set less than unsigned
 - » `slti` : set less than constant value
 - » `sltiu` : set less than unsigned constant
- set t0 to 1 if t1 < t2
`slt $t0, $t1, $t2`

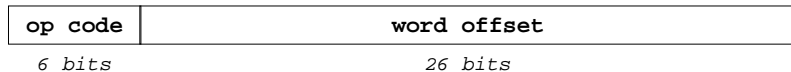
Pseudo-instructions

- The assembler is your friend and will build instruction sequences for you
- Original code:
`bge $a0,$t1,end # if a0>=t1 skip`
- Actual instructions:
`slt $at,$a0,$t1 # if a0<t1 at=true`
`beq $at,$0,end # skip if at==false`

Jump Instructions

- Jump instructions provide longer range than branch instructions
- 26-bit word offset in J-format instructions
 - » `j` : jump
 - » `jal` : jump and link (store return address)
- 32-bit address in register jumps
 - » `jr` : jump through register
 - » `jalr` : jump through register and link

J-format fields



- The word offset value is multiplied by 4 to create a byte offset
 - » the result is 28 bits wide
- Then concatenated with top 4 bits of PC to make a 32 bit destination address

Important Jumps

- Jump and link (**jal**)
 - » call procedure and store return address in \$ra
- Jump through register (**jr**)
 - » return to caller using the address in \$ra
- We will talk about procedure calls in excruciating detail next lecture