## Procedure Detail

CSE 410 - Computer Systems
October 10, 2001

## Readings and References

- Reading
- Other References
  - D. Sweetman, See MIPS Run, Morgan Kauffman, Publishers
    - Chapter 10, C Programming on MIPS

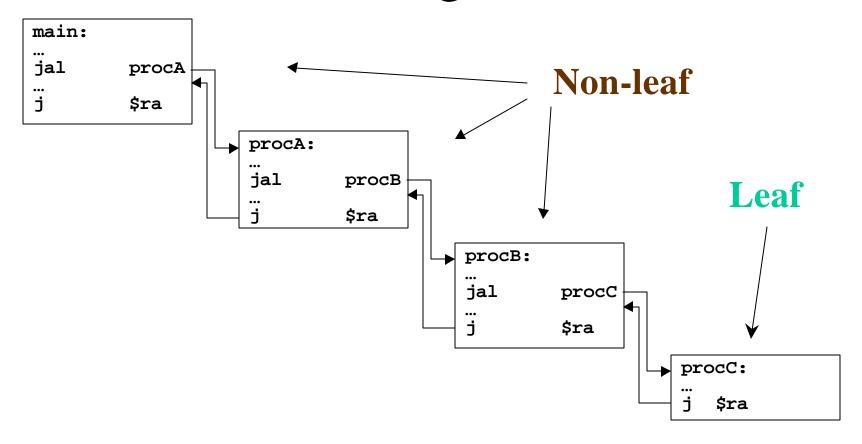
## Leaf procedures

- A leaf procedure is one that does not call another procedure
- Relatively simple register usage since the procedure doesn't call anyone else
- Little or no memory access requirements because you are not saving and restoring as many registers from the stack

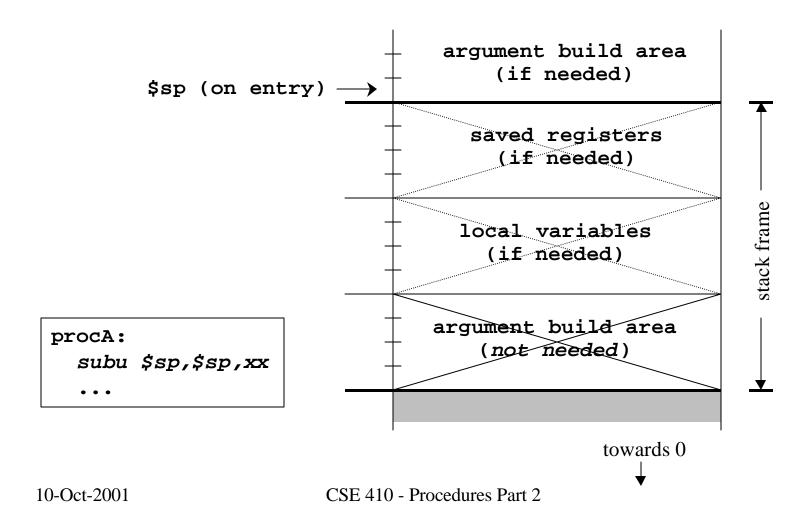
## Non-leaf procedure

- A non-leaf procedure is one that calls another procedure
- You must save at least register \$ra, since that register is overwritten by the jal when you call another procedure

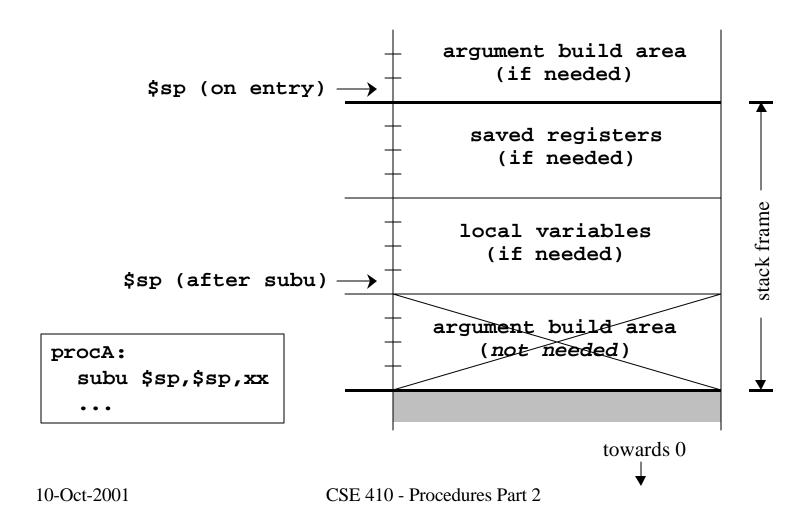
# Calling tree



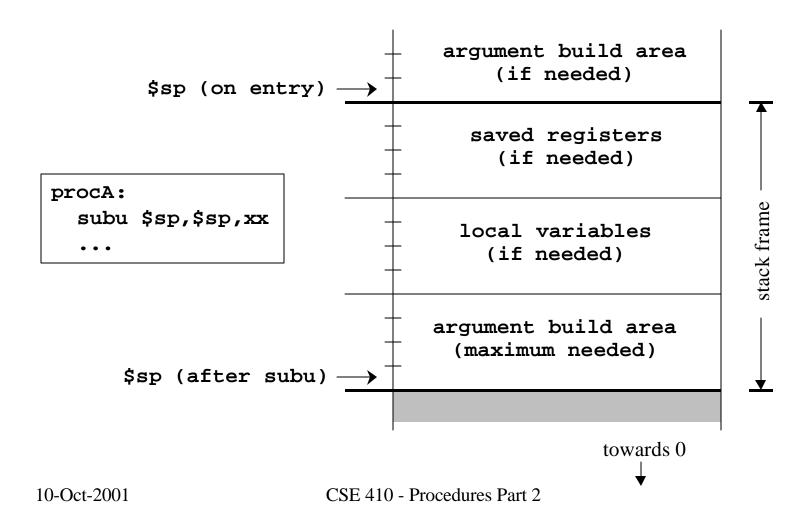
## Layout of stack frame (little leaf)



## Layout of stack frame (big leaf)



## Layout of stack frame (non-leaf)



## Little leaf example - swap.c

```
/* Swap two integer array elements */
void swap(int a[], int i, int j)
{
    int T;
    T = a[i];
    a[i] = a[j];
    a[j] = T;
}
```

## Little leaf example - swap.s

#### swap:

```
sll
                       # \$a1 = 4*i
       $a1,$a1,2
addu
       $a1,$a1,$a0
                     # $a1 = addr(a[i])
                    # $v1 = a[i]
lw
    $v1,0($a1)
                       \# a2 = 4*j
$11 $a2,$a2,2
                  # $a2 = addr(a[j])
addu $a2,$a2,$a0
                  # $v0 = a[j]
       $v0,0($a2)
lw
       v0,0(a1) # a[i] = old a[j]
SW
       $v1,0($a2)
                       # a[j] = old a[i]
SW
j
       $ra
                       # return
```

## Non-leaf example - QuickSort.c

```
void QuickSort(int a[], int lo0, int hi0)
{
    int lo = lo0;
    int hi = hi0;
    int mid;

    if ( hi0 > lo0)
    {
    ...
```

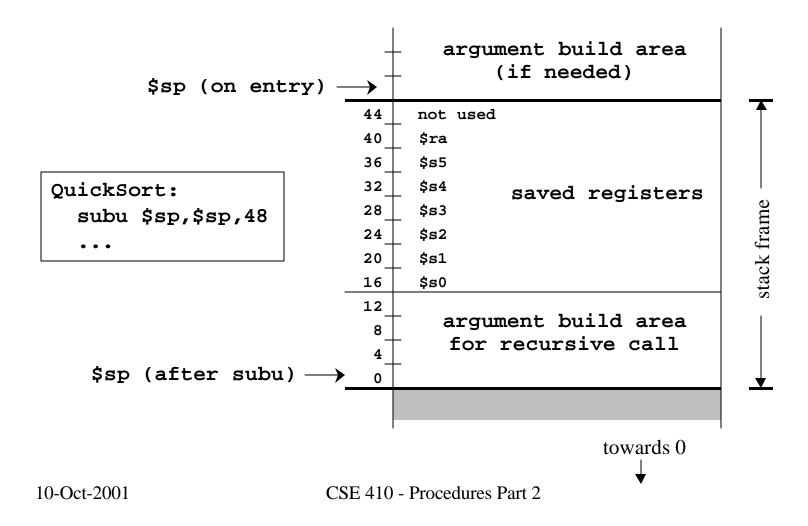
## Non-leaf example - QuickSort.s

#### QuickSort:

```
subu
        $sp,$sp,48
                             # create stack frame
                             #
        $ra,40($sp)
SW
                             #
        $s5,36($sp)
SW
        $s4,32($sp)
                             #
SW
                             #
        $s3,28($sp)
SW
                             #
        $s2,24($sp)
SW
        $s1,20($sp)
                             #
SW
        $s0,16($sp)
                             #
SW
        $s3,$a0
                             # $s3 = address(a)
move
        $s5,$a1
                             \# $s5 = 100
move
```

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## Layout of QuickSort stack frame



## \$ra - Return Address

- Return address register
  - written with jal, jalr instructions
  - must be saved if procedure calls another

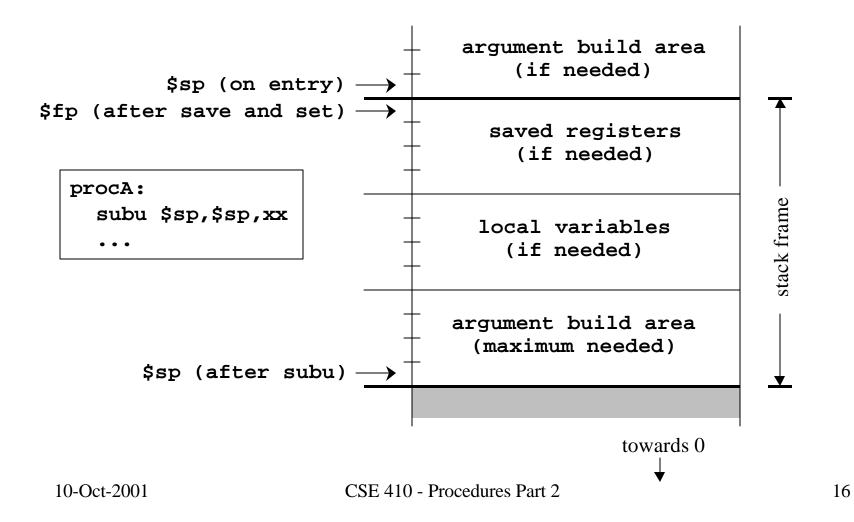
### QuickSort:

```
subu $sp,$sp,48  # create stack frame
sw $ra,40($sp) #
...
lw $ra,40($sp) # restore from stack ...
addu $sp,$sp,48 #
j $ra # return
```

## \$fp - Frame Pointer

- Frame pointer points to the largest address in the stack frame
- Stack pointer points to the smallest address in the stack frame
  - no advantage to \$fp if \$sp does not change during procedure's execution
- Consider \$fp to be \$s8
  - save and restore required if you use it

# Layout of stack frame (with \$fp)



## \$s0-\$s7 - Save and Restore

- These registers are available for unlimited use
- Must save immediately on procedure entry and restore just before procedure exit if you are going to use them
- As a result of this convention, the registers will have the same values after a procedure call as they had before

# \$t0-\$t9 - Temporary registers

- Use however you like
- No save and restore required or expected
- As a result of this convention, the registers have no guaranteed values when you get back from calling another procedure

# \$a0-\$a3, \$v0-\$v1 - Args/Return

- The argument registers can be changed in a procedure without restriction
- No guarantee that they will be the same upon return from a called procedure
- The result registers will contain whatever the function prototype says they will
  - undefined value in \$v1 if not used for return

## \$gp - Global Pointer

- Initialized so that it points to the middle of a 64KB section of the data segment
  - address 0x10008000
- Variables placed in this section can be accessed without loading a 32-bit address
  - lw \$t0,-32768(\$gp)
- Assembler directive
  - .extrn symbol bytecount

# Layout of program memory

7FFF	FF <u>FF</u>	reserved (4KB)
7FFF	EFFF	stack (grows down)
		+
		~1792 MB
		<b>^</b>
1001	0000	heap (grows up)
1000 1000		global data (64 KB)
OFFF	FFFF	
		program (252 MB)
0040	0000	
003F	FFFF	magamused (4 MD)
0000	0000	reserved (4 MB)

Not to Scale!

## Using the global pointer - gp.s

```
# global area symbol
  .extern common 4
  .data
local:
                        # non-global symbol
                        # data value
  .word
          0xAAAA
  .text
main:
            $t0,local # load word
  lw
            $t0,common # store word
  SW
  j
            $ra
                       # return
```

## A reference through \$gp

```
# lw $t0,local

0x3c011001 lui $1, 4097

0x8c280000 lw $8, 0($1)

# sw $t0,common

0xaf888000 sw $8, -32768($28)
```

### 0xAF888000 <=> sw \$t0,-32768(\$gp)

	<b>A</b>				I	<b></b> 		8					8			8				0				0			0				
1	0	1	0	1	1	1	1	1	0	0	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
												•			•		•							•							
	op			base					src					offset																	
6 bits				5 bits				5 bits					16 bits																		
1	0	1	0	1	1	1	1	1	0	0	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-																															
43=sw			28=\$gp				8=\$t0					0x8000=offset																			

## Pearls of wisdom from Sweetman

- These calling conventions can look very complex
  - but partly that's just appalling documentation
  - and the inclusion of debugging conventions
- Most functions that you may write in assembler for tuning reasons will be leaf functions
  - the declaration of such a function is very simple