## Not My Problem

Crowdsourcing Code Verification With Video Games

## Section 8.2 Program Correctness (for imperative programs)

A theory of program correctness needs wffs, axioms, and inference rules.

Wffs (called Hoare triples) are of the form

$$\{P\} S \{Q\}$$

Where S is a program statement and P (a *precondition*) and Q (a *postcondition*) are logical statements about the variables of S.

Semantics: The meaning of  $\{P\}$  S  $\{Q\}$  is the truth value of the statement:

If *P* is true before *S* executes, then *Q* is true after *S* halts.

Note that it is assumed that S halts. If  $\{P\}$  S  $\{Q\}$  is true, then S is said to be *correct* wrt to precondition P and postcondition Q.

Assignment Axiom (AA):  $\{P(x/t)\}\ x := t\ \{P\}.$ 

Example. 
$$\{x = 4\} \ x := x - 1 \ \{x = 3\}.$$
 Example.  $\{x < 4\} \ x := x - 1 \ \{x < 3\}.$ 

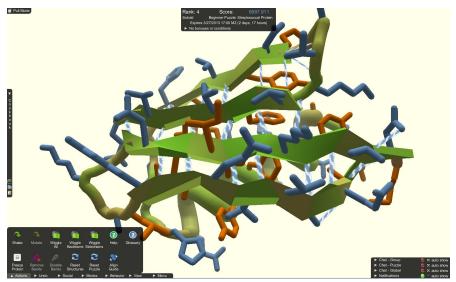
Inference Rules	$P \to R$ and $\{R\} S \{Q\}$	and	$\{P\} S \{T\} \text{ and } T \to Q$
Consequence Rule:	${P} S {O}$		$\{P\} S \{O\}$

*Example.* Prove the correctness of  $\{x < 3\}$  x := x - 1  $\{x < 3\}$ .

Proof: 1. 
$$\{x < 4\} \ x := x - 1 \ \{x < 3\}$$
 AA  
2.  $x < 3$   $P \ [for (x < 3) \rightarrow (x < 4)]$   
3.  $x < 4$  2,  $T$   
4.  $(x < 3) \rightarrow (x < 4)$  2-3, CP  
5.  $\{x < 3\} \ x := x - 1 \ \{x < 3\}$  1, 4, Consequence. QED.

https://www.google.com/amp/slideplayer.com/amp/9852342/?source=images





http://fold.it/portal/info/about



http://www.theesa.com/article/crowdsourced-video-games-prompt-breakthroughs-in-science-and-technology/