

Section 3: LR Parsing

CSE 401/M501

Adapted from Spring 2021

Announcements

- Scanner is due tonight
 - Be sure to test, push, and tag!
- Every person has 4 project late days and 4 assignment late days
 - Up to 2 can be used per assignment
 - Each late day gives an extra 24 hour chunk (including weekend days)
 - We recommend not using your late days this early if possible!
 - **Submitting project components a day late uses a project late day from each partner!**



11:45-12:45 OH (Rachel) CSE2 150	10	14:00-15:00 OH (Rachel) zoom	11	13:00-14:30 OH (John) CSE2 150	12	Section <i>LR parser construction</i>	13	12:30-13:30 OH (Randy) CSE2 151 + zoom	14
14:30-15:20 Lecture CSE2 G10 <i>LR parsing (concl.); LR table construction (start)</i> (3.5) slides				14:30-15:20 Lecture CSE2 G10 <i>LR table construction (cont.) (3.5)</i>		12:30-13:30 OH (Randy) CSE2 151 + zoom		14:30-15:20 Lecture CSE2 G10 <i>LR conflicts, first/follow</i>	
				16:30-17:30 OH (Robert) CSE2 152 + zoom		16:30-17:30 OH (Robert) CSE2 152 + zoom		15:30-16:30 OH (John) CSE2 151	
						23:00 Project: scanner due			

Agenda

- (Fast) LR terminology review
- Worksheet

Get Your LR Jargon On

- Frontier
 - The upper “layer” of the current parse tree (held in the stack)

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 - The upper “layer” of the current parse tree (held in the stack)
- Sentential Form
 - A string that can be generated at any point in a derivation (can be reached using any number of productions from the start symbol)

Get Your LR Jargon On

- Frontier
 - The upper “layer” of the current parse tree (held in the stack)
- Sentential Form
 - A string that can be generated at any point in a derivation (can be reached using any number of productions from the start symbol)
- Handle
 - An occurrence of the right side of a production in the frontier that is used in the rightmost derivation to arrive at the current string
 - Given the derivation ... $\Rightarrow \underline{aA}bcde \Rightarrow \underline{ab}bcde$, using the production $A ::= b:$
 - The production ‘ $A ::= b$ ’ at index 1 would be a handle of $\underline{ab}bcde$

Get Your LR Jargon On - Example

Shift-Reduce Example

$S ::= aABe$
 $A ::= Abc \mid b$
 $B ::= d$

Frontier

Stack	Input	Action
\$	abbcde\$	<i>shift</i>
\$a	bbcde\$	<i>shift</i>
\$ab	bcde\$	<i>reduce</i>
\$aA	bcde\$	<i>shift</i>
\$aAb	cde\$	<i>shift</i>
\$aAbc	de\$	<i>reduce</i>
\$aA	de\$	<i>shift</i>
\$aAd	e\$	<i>reduce</i>
\$aAB	e\$	<i>shift</i>
\$aABe	\$	<i>reduce</i>
\$S	\$	<i>accept</i>

Get Your LR Jargon On - Example

Shift-Reduce Example

$S ::= aABe$
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Stack	Input	Action
\$	abbcde\$	<i>shift</i>
\$a	bbcde\$	<i>shift</i>
\$ab	bcde\$	<i>reduce</i>
\$aA	bcde\$	<i>shift</i>
\$aAb	cde\$	<i>shift</i>
\$aAbc	de\$	<i>reduce</i>
\$aA	de\$	<i>shift</i>
\$aAd	e\$	<i>reduce</i>
\$aAB	e\$	<i>shift</i>
\$aABe	\$	<i>reduce</i>
\$S	\$	<i>accept</i>

Sentential
Forms

Get Your LR Jargon On - Example

Shift-Reduce Example

$S ::= aABe$
 $A ::= Abc \mid b$
 $B ::= d$

Handles	Stack	Input	Action
	\$	abbcde\$	<i>shift</i>
	\$a	bbcde\$	<i>shift</i>
A ::= b at index 2	\$ab	bcde\$	<i>reduce</i>
	\$aA	bcde\$	<i>shift</i>
	\$aAb	cde\$	<i>shift</i>
	\$aAbc	de\$	<i>reduce</i>
A ::= Abc at index 4	\$aA	de\$	<i>shift</i>
B ::= d at index 3	\$aAd	e\$	<i>reduce</i>
	\$aAB	e\$	<i>shift</i>
S ::= aABe at index 4	\$aABe	\$	<i>reduce</i>
	\$S	\$	<i>accept</i>

A Little Bit More Jargon

- Viable Prefix
 - The prefixes of a right sentential form that do not extend beyond the end of its handle
 - Perhaps less confusing -> the set of prefixes of strings that can appear on the stack of a shift-reduce parser

A Little Bit More Jargon

- Viable Prefix
 - The prefixes of a right sentential form that do not extend beyond the end of its handle
 - Perhaps less confusing -> the set of prefixes of strings that can appear on the stack of a shift-reduce parser
- Item
 - A marked production (a production with a '.' in it)
 - [A ::= .XY], [A ::= X.Y], [A ::= XY.]

Get Your LR Jargon On - Example

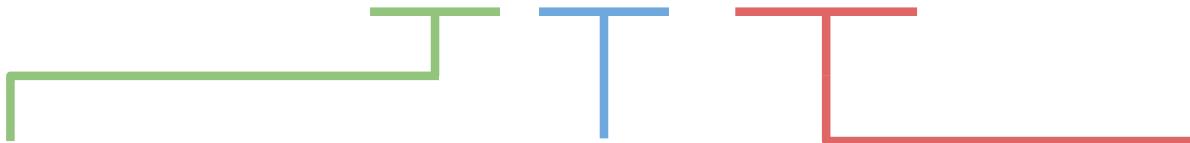
Shift-Reduce Example

$$\begin{aligned}S &::= aABe \\A &::= Abc \mid b \\B &::= d\end{aligned}$$

Viable
Prefix
(all prefixes of
these strings
are also viable
prefixes)

Stack	Input	Action
\$	abbcde\$	<i>shift</i>
\$a	bbcde\$	<i>shift</i>
\$ab	bcde\$	<i>reduce</i>
\$aA	bcde\$	<i>shift</i>
\$aAb	cde\$	<i>shift</i>
\$aAbc	de\$	<i>reduce</i>
\$aA	de\$	<i>shift</i>
\$aAd	e\$	<i>reduce</i>
\$aAB	e\$	<i>shift</i>
\$aABe	\$	<i>reduce</i>
\$S	\$	<i>accept</i>

L R (0)



Left-to-Right

Only takes one pass,
performed from the left

Rightmost

At each point, finds the
derivation for the rightmost
handle (bottom-up)

No Lookahead

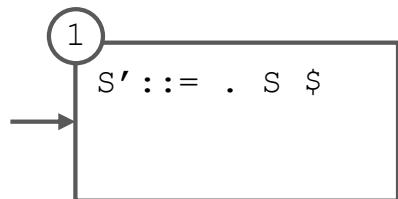
Decide what to do based on
current parser state and
stack, ignoring next input

Problem 1 (On Worksheet)

0. $S' ::= S \$$
1. $S ::= a z$
2. $S ::= b$
3. $z ::= a$
4. $z ::= b s$

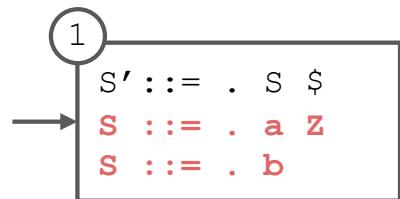
State Diagram Construction

0. $S' ::= S \$$
1. $S ::= a Z$
2. $S ::= b$
3. $Z ::= a$
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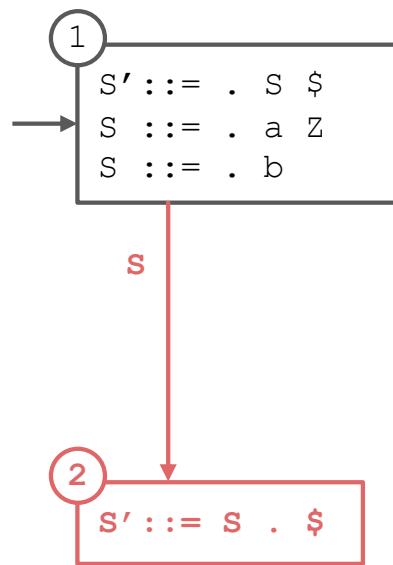
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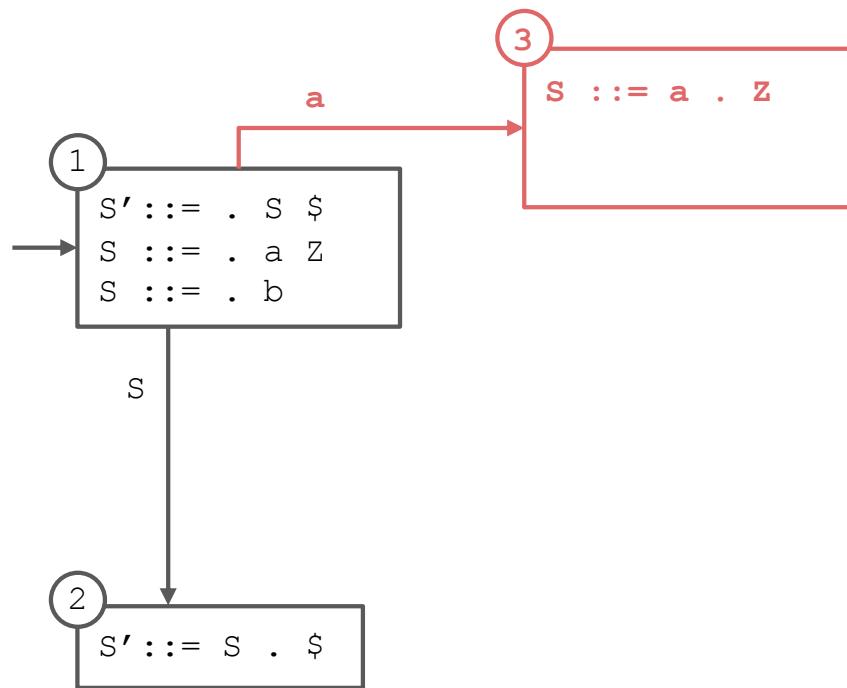
State Diagram Construction

0. $S' ::= S \$$
1. $S ::= . a Z$
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4. $Z ::= b S$



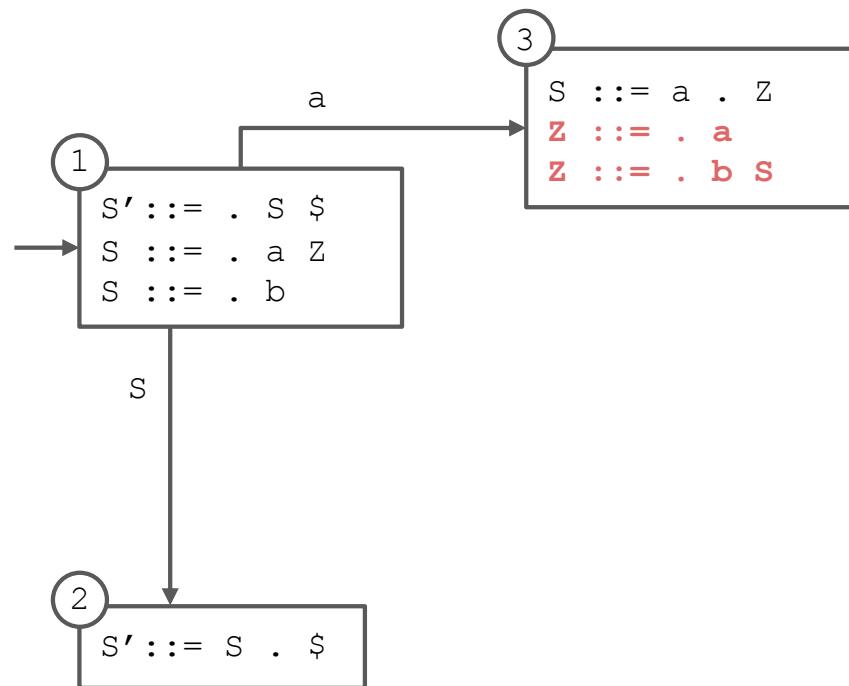
State Diagram Construction

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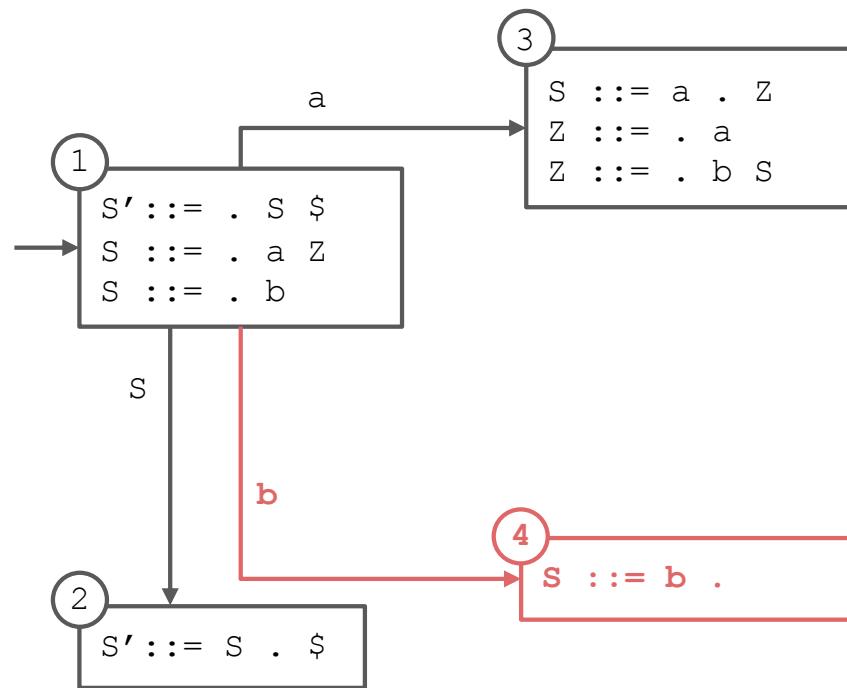
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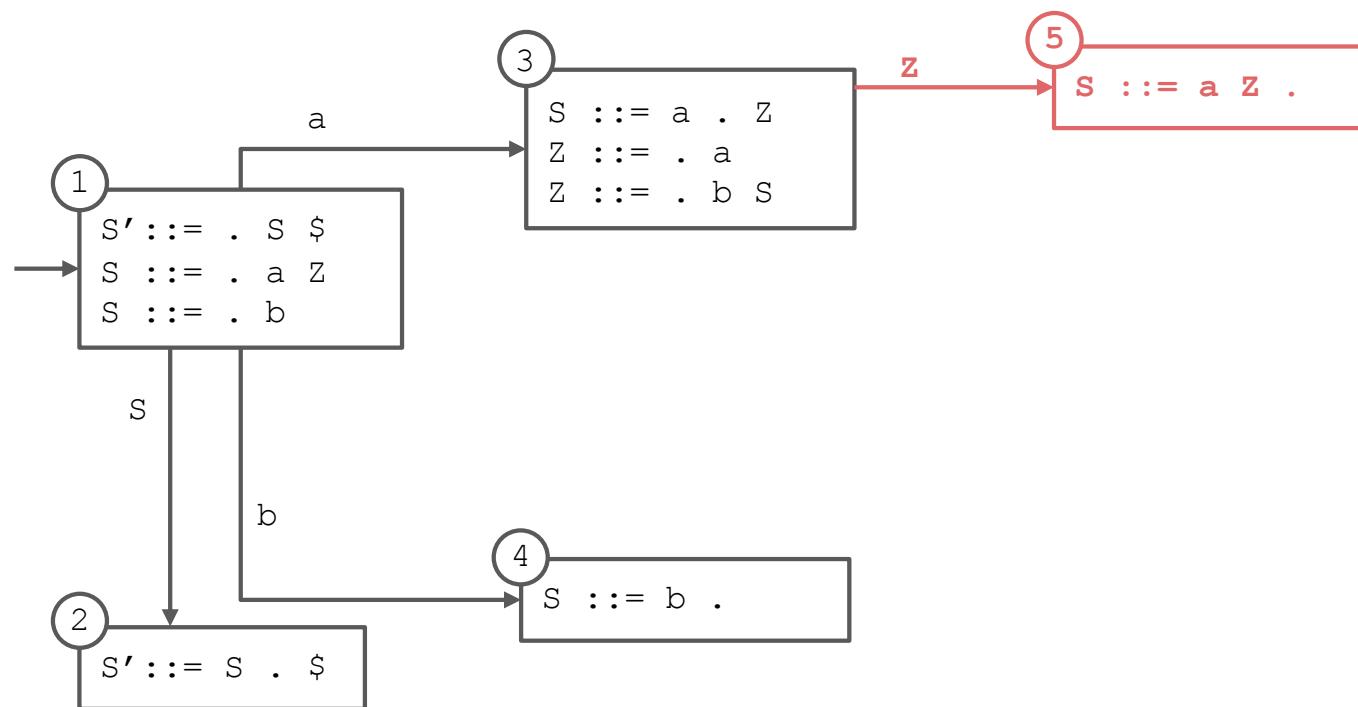
State Diagram Construction

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1. $S ::= . a Z$
2. $S ::= . b$
3. $Z ::= a . Z$
 $Z ::= . a$
 $Z ::= . b S$
4. $Z ::= b S$



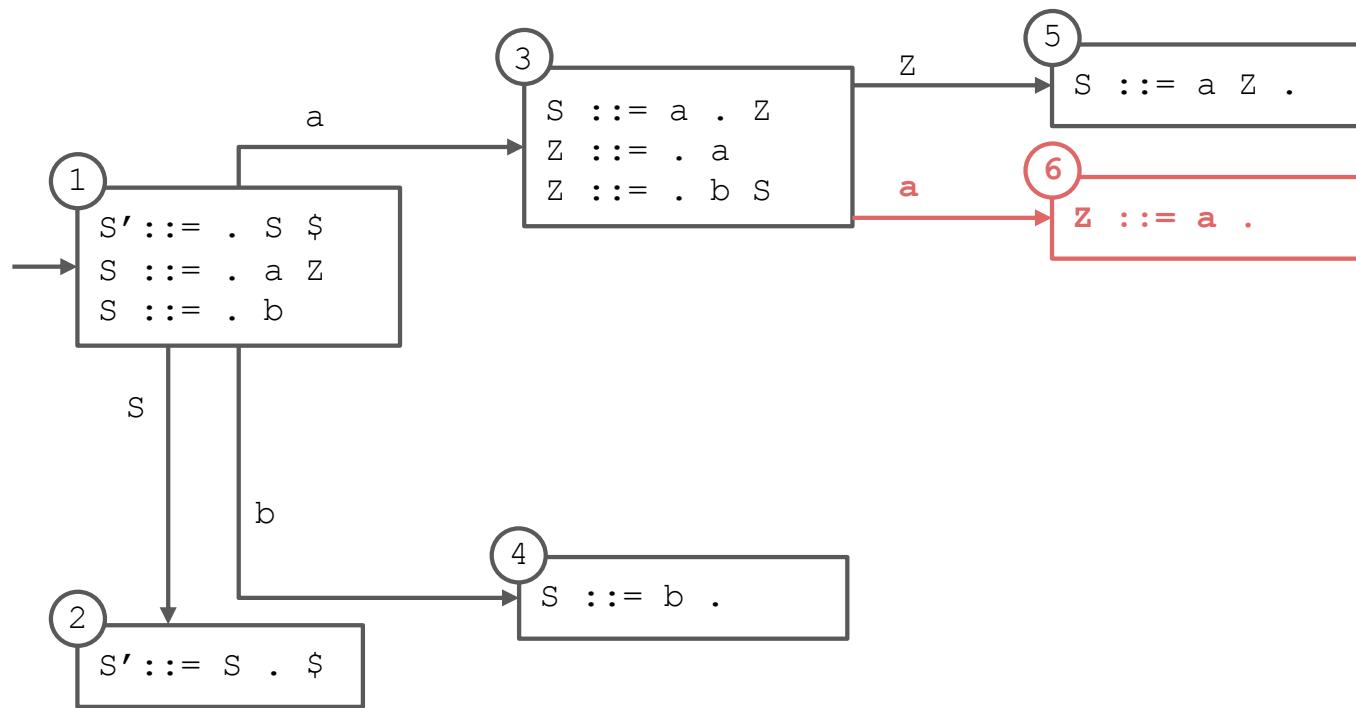
State Diagram Construction

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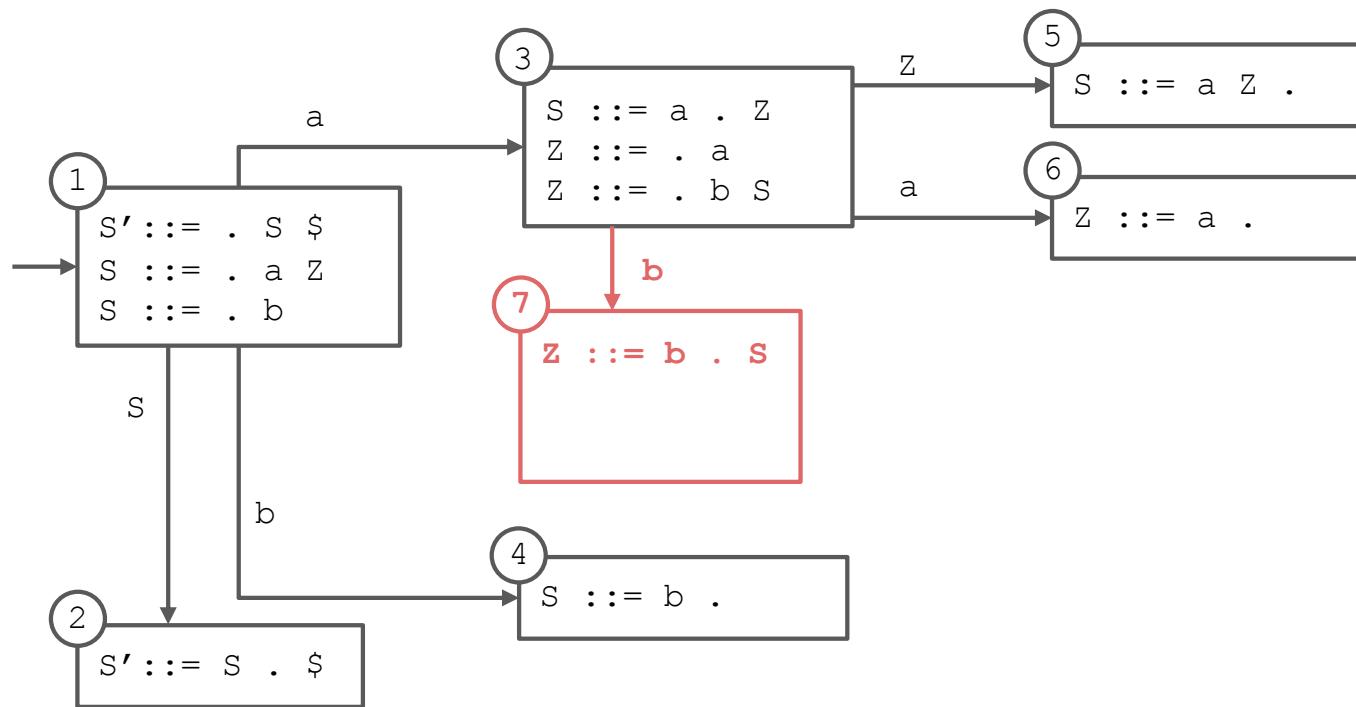
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- 4. $Z ::= b S$

State Diagram Construction



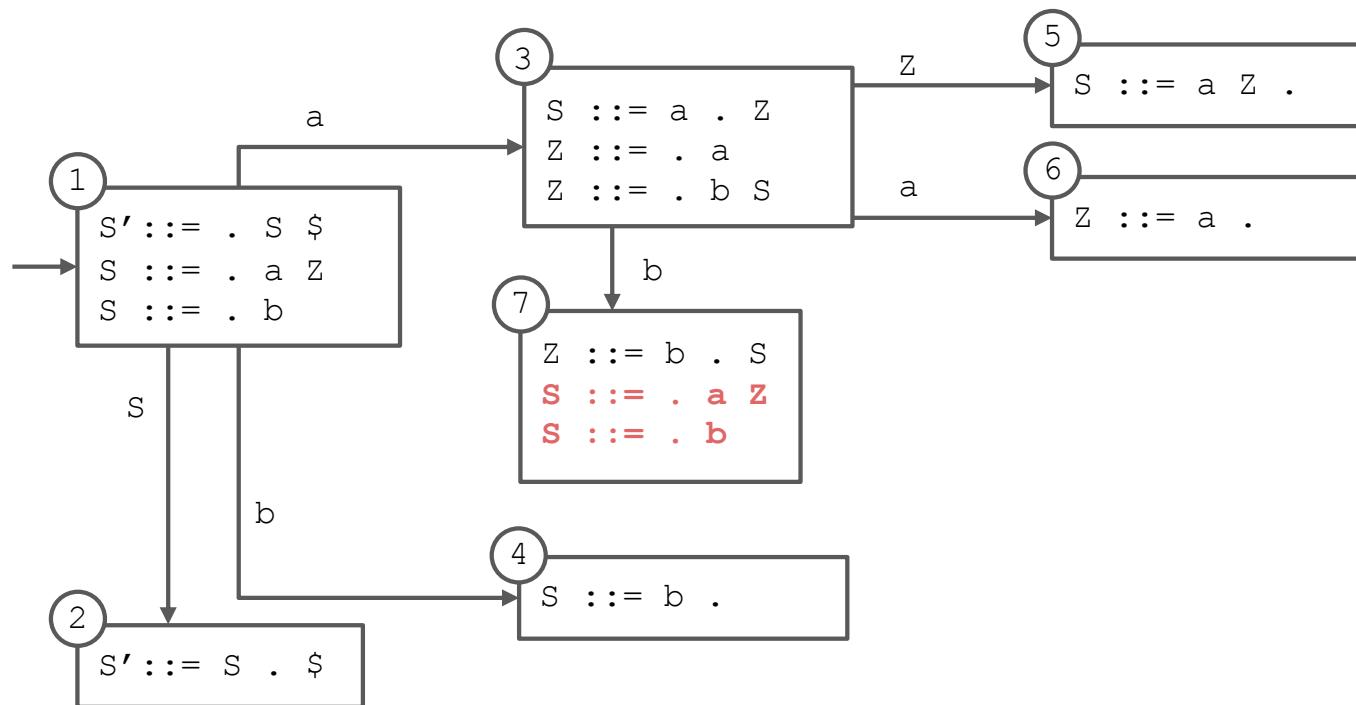
State Diagram Construction

0. $S' ::= S \$$
1. $S ::= . a Z$
2. $S ::= . b$
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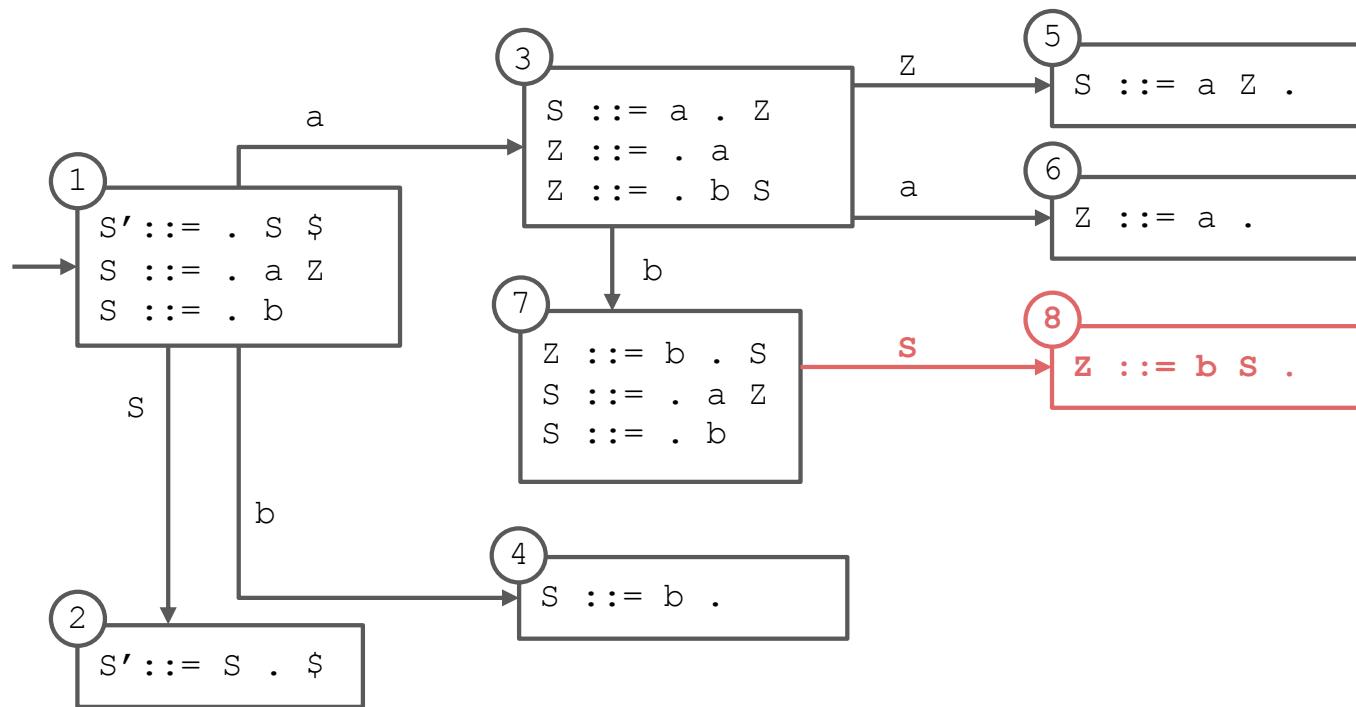
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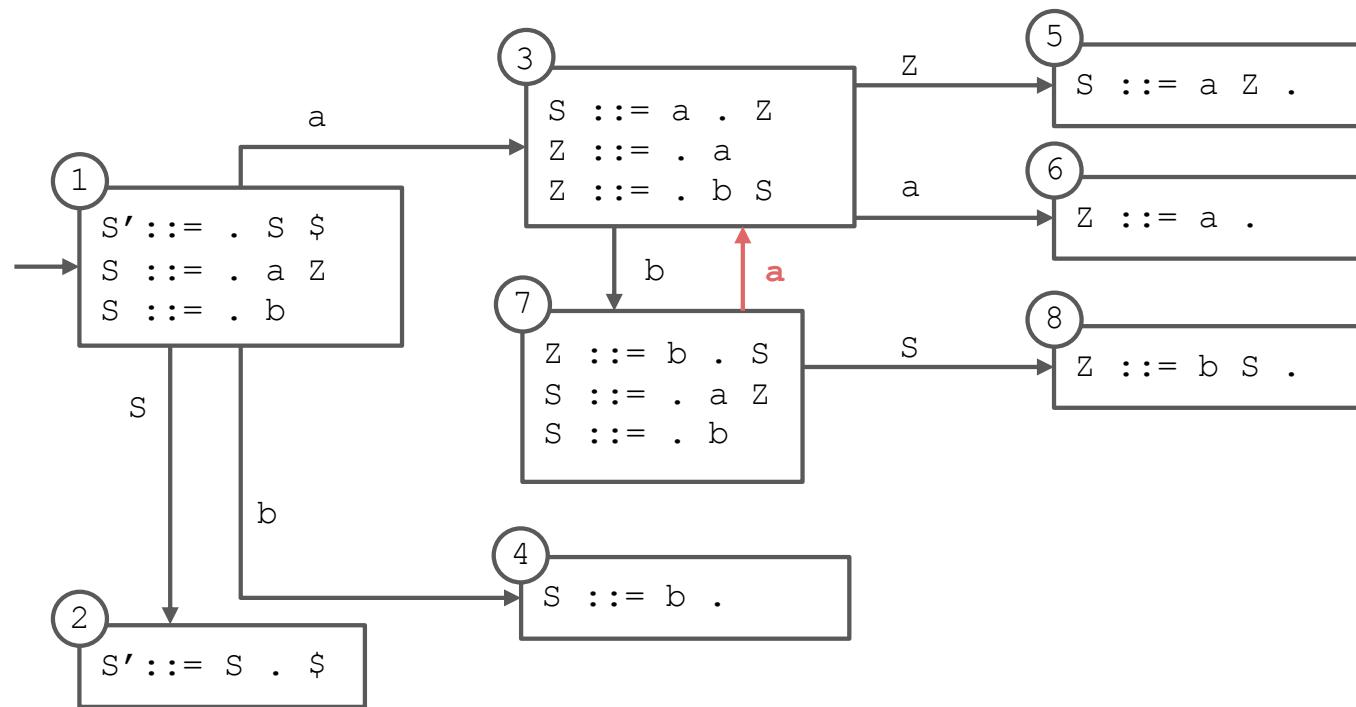
State Diagram Construction

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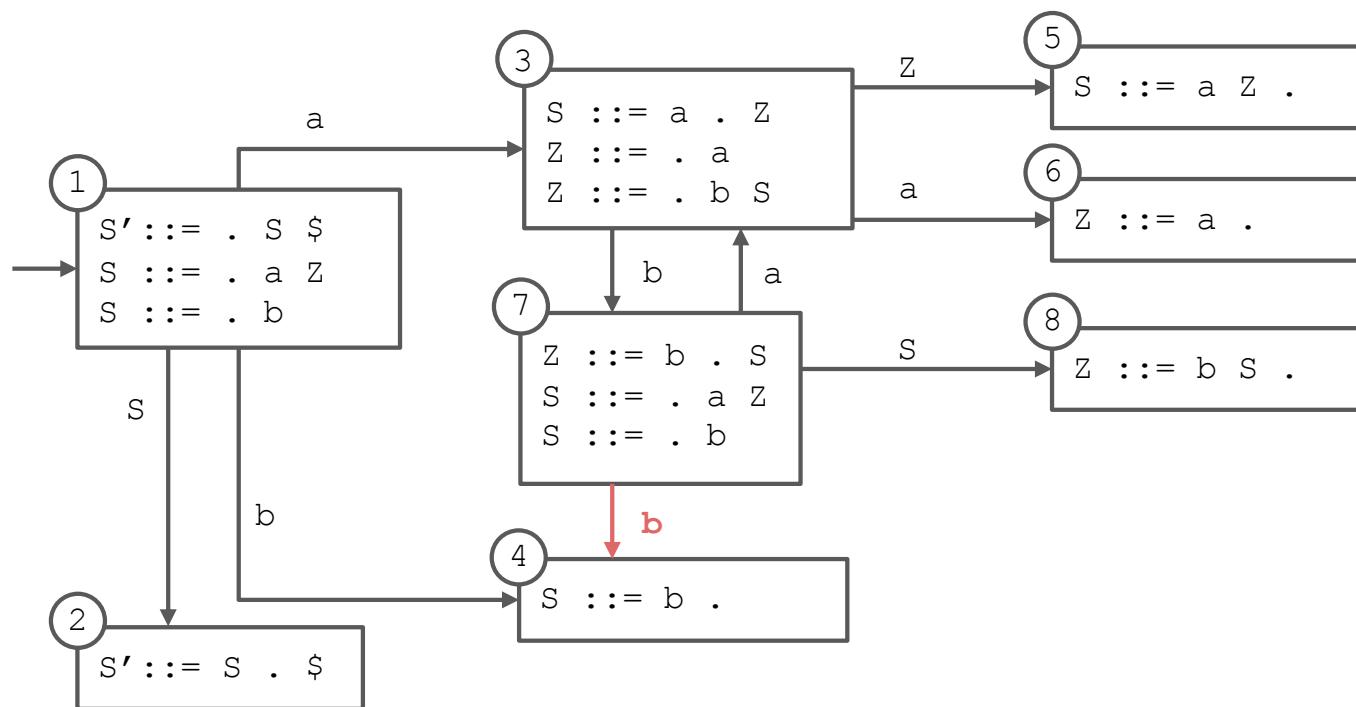
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1. $S ::= a Z$
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State Diagram Construction



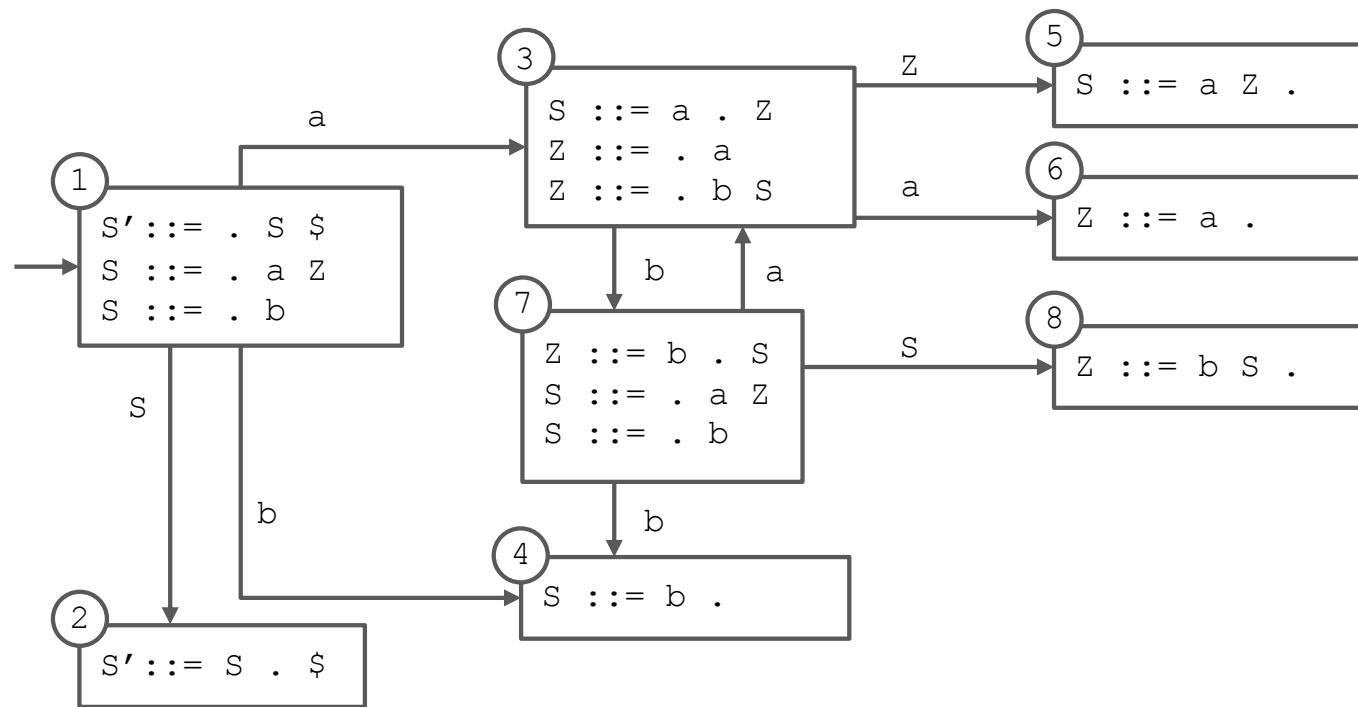
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State Diagram Construction



0. $S' ::= S \$$
1. $S ::= a Z$
2. $S ::= b$
3. $Z ::= a$
4. $Z ::= b S$

Completed State Diagram



Converted to Table

s# means “shift and enter state #”

- occurs when there is a transition on a terminal

r# means “reduce using production #”

- occurs when a state contains an item with the location at the end of the right-hand side

g# means “go to state #”

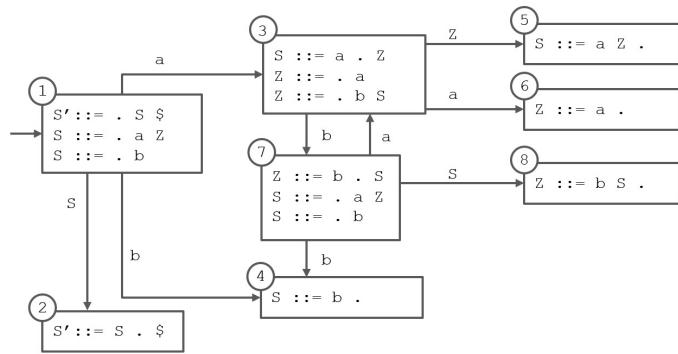
- occurs when there is a transition on a nonterminal

acc means “accept”

- occurs when the start symbol (S here) has been completed and there is no more input

STATE	ACTION			GOTO	
	a	b	\$	S	Z
1	s3	s4		g2	
2			acc		
3	s6	s7			g5
4	r2	r2	r2		
5	r1	r1	r1		
6	r3	r3	r3		
7	s3	s4		g8	
8	r4	r4	r4		

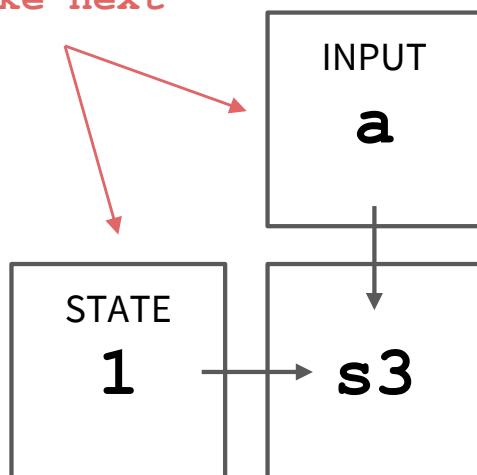
Parse Trace



STACK	INPUT	ACTION
\$ 1	a b a b b \$	

Parse Trace

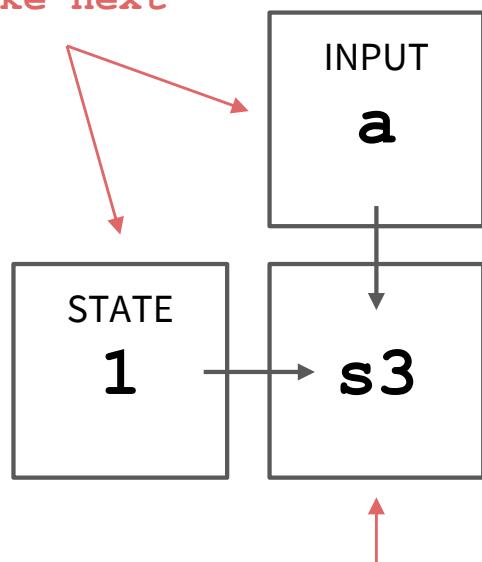
Row and column of table to look up: decides what action to take next



STACK	INPUT	ACTION
\$ 1 \$ 1 a	a b a b b \$ b a b b \$	SHIFT

Parse Trace

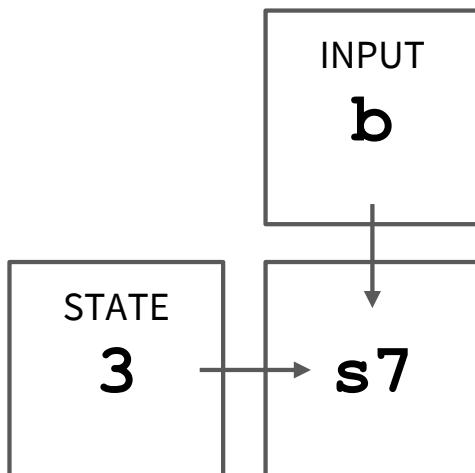
Row and column of table to look up: decides what action to take next



Shift, and enter state 3

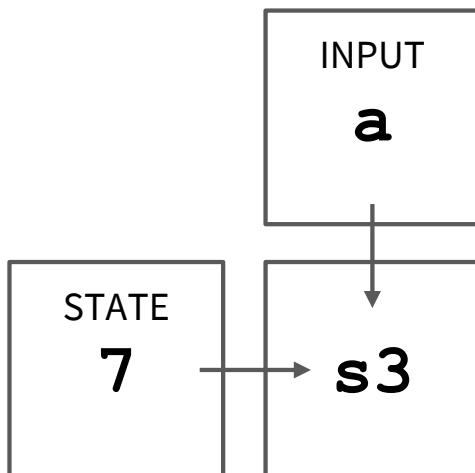
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3	a b a b b \$ b a b b \$	SHIFT

Parse Trace



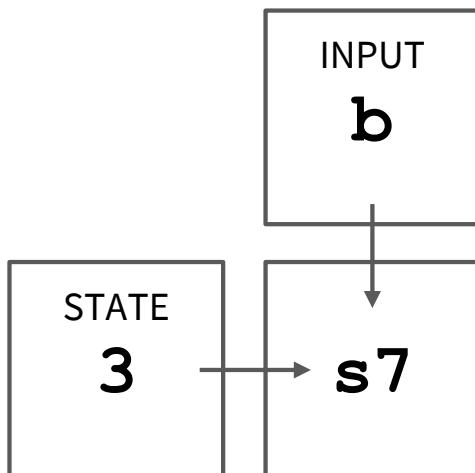
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3 \$ 1 a 3 b 7	a b a b b \$ b a b b \$ a b b \$	SHIFT SHIFT

Parse Trace



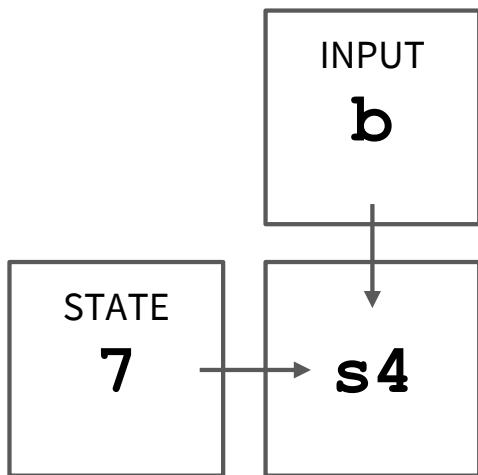
STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	
\$ 1 a 3 b 7 a 3	b b \$	SHIFT

Parse Trace



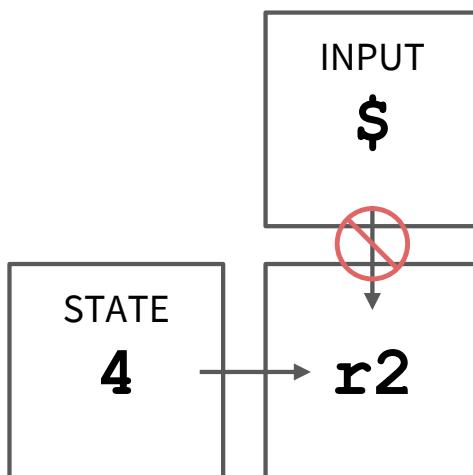
STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 7	a	SHIFT
\$ 1 a 3 b 7 a 3	b	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b	SHIFT

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 7	a	SHIFT
\$ 1 a 3 b 7 a 3	b	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	SHIFT

Parse Trace

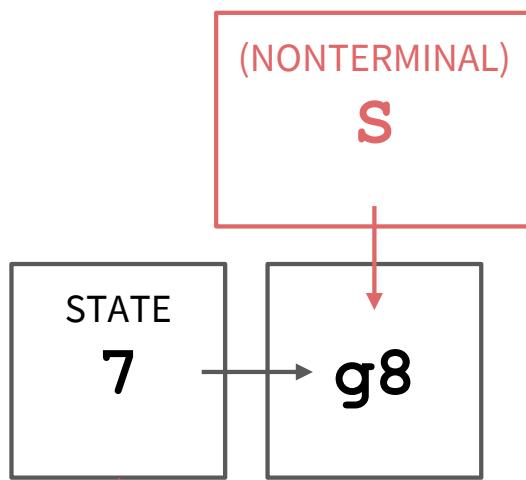


2. $S ::= b$

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 S	\$	REDUCE

For LR(0), the input doesn't technically matter here

Parse Trace

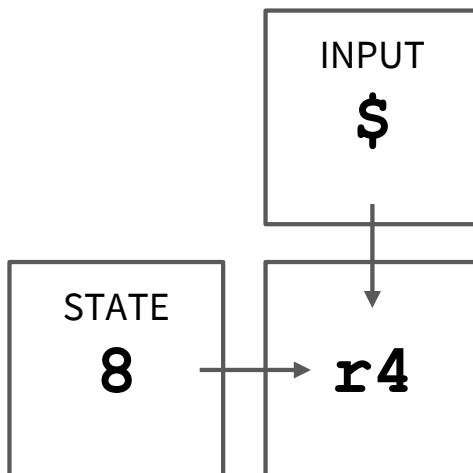


STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	

After a reduction, we go back to a previous state on the stack and use the reduced non-terminal to determine what state to GOTO.

This allows the parser to run in $O(n)$ time, since it doesn't have to re-evaluate the entire stack!

Parse Trace

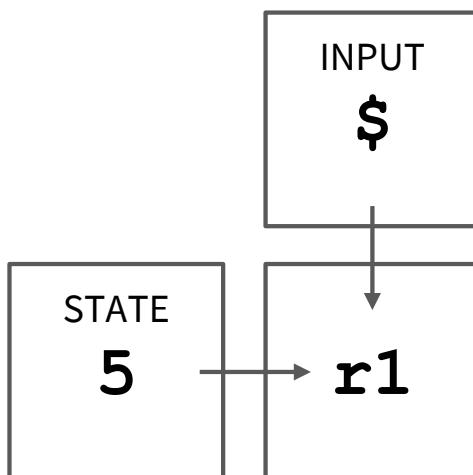


4. **Z ::= b S**

(and GOTO step:
s3 & Z -> g5)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	

Parse Trace

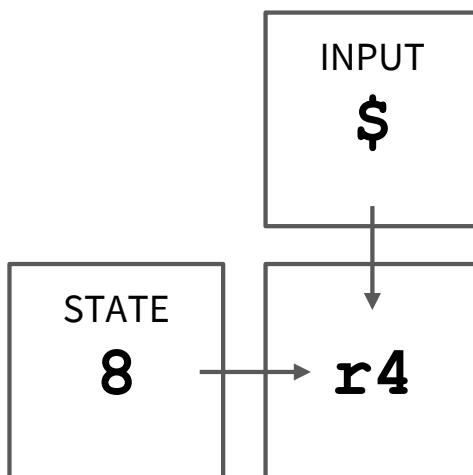


1. S ::= a Z

(and GOTO step:
s7 & S -> g8)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 s 8	\$	

Parse Trace

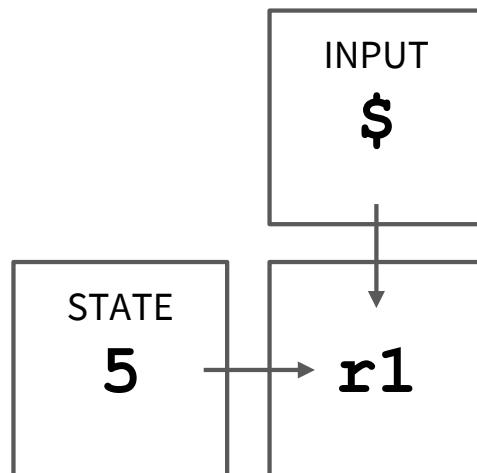


4. Z ::= b S

(and GOTO step:
s3 & Z -> g5)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 z 5	\$	REDUCE

Parse Trace

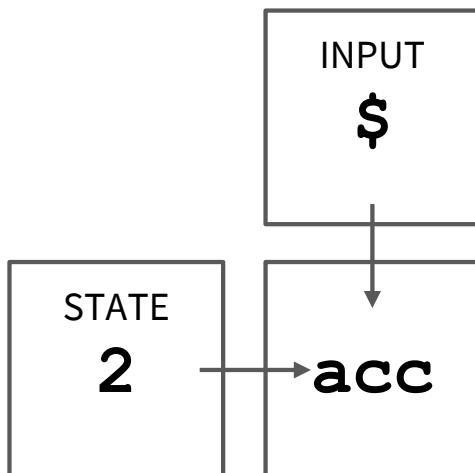


1. S ::= a Z

(and GOTO step:
s1 & S -> g2)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 z 5	\$	REDUCE
\$ 1 s 2	\$	REDUCE

Parse Trace



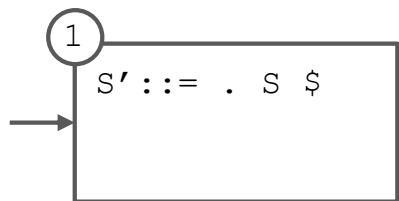
STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 s 8	\$	REDUCE
\$ 1 a 3 z 5	\$	REDUCE
\$ 1 s 2	\$	ACCEPT

Problem 2 (On Worksheet)

0. $S' ::= S \$$
1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$

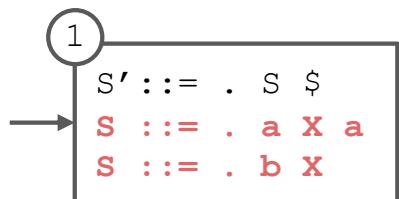
State Diagram Construction

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1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$



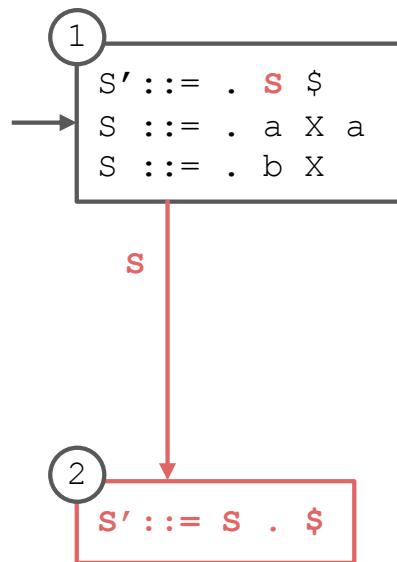
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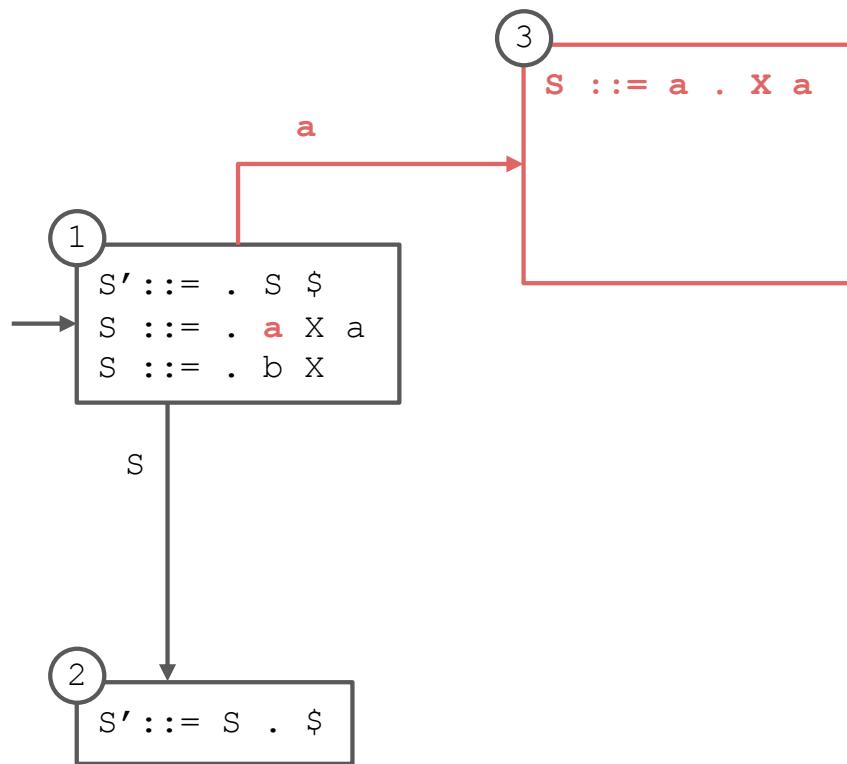


State Diagram Construction

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1. $S ::= . a X a$
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4. $X ::= S c$

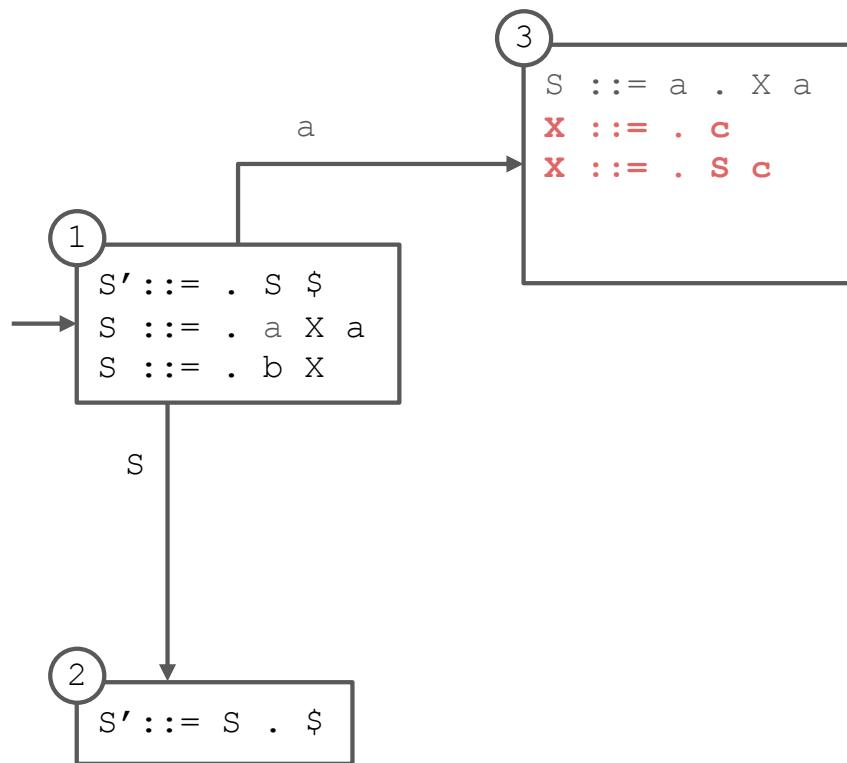


State Diagram Construction



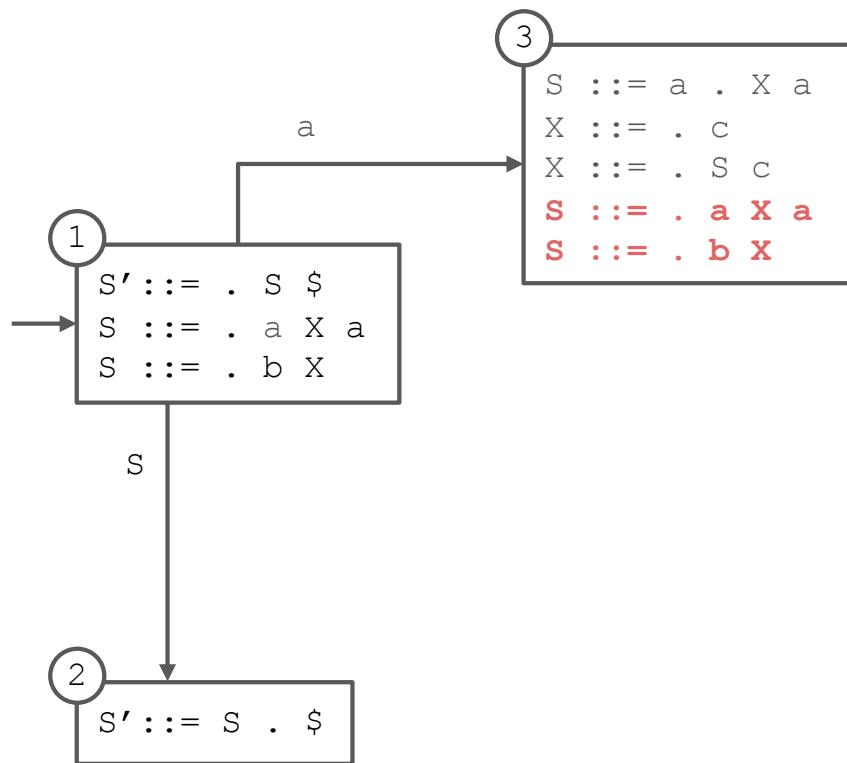
0. $S' ::= S \$$
1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$

State Diagram Construction



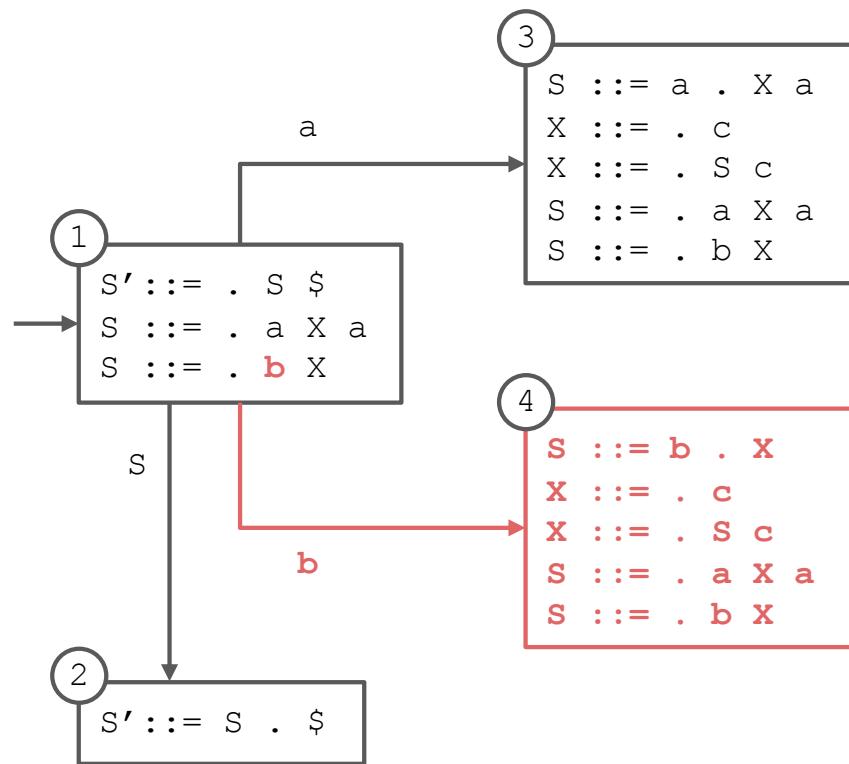
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State Diagram Construction



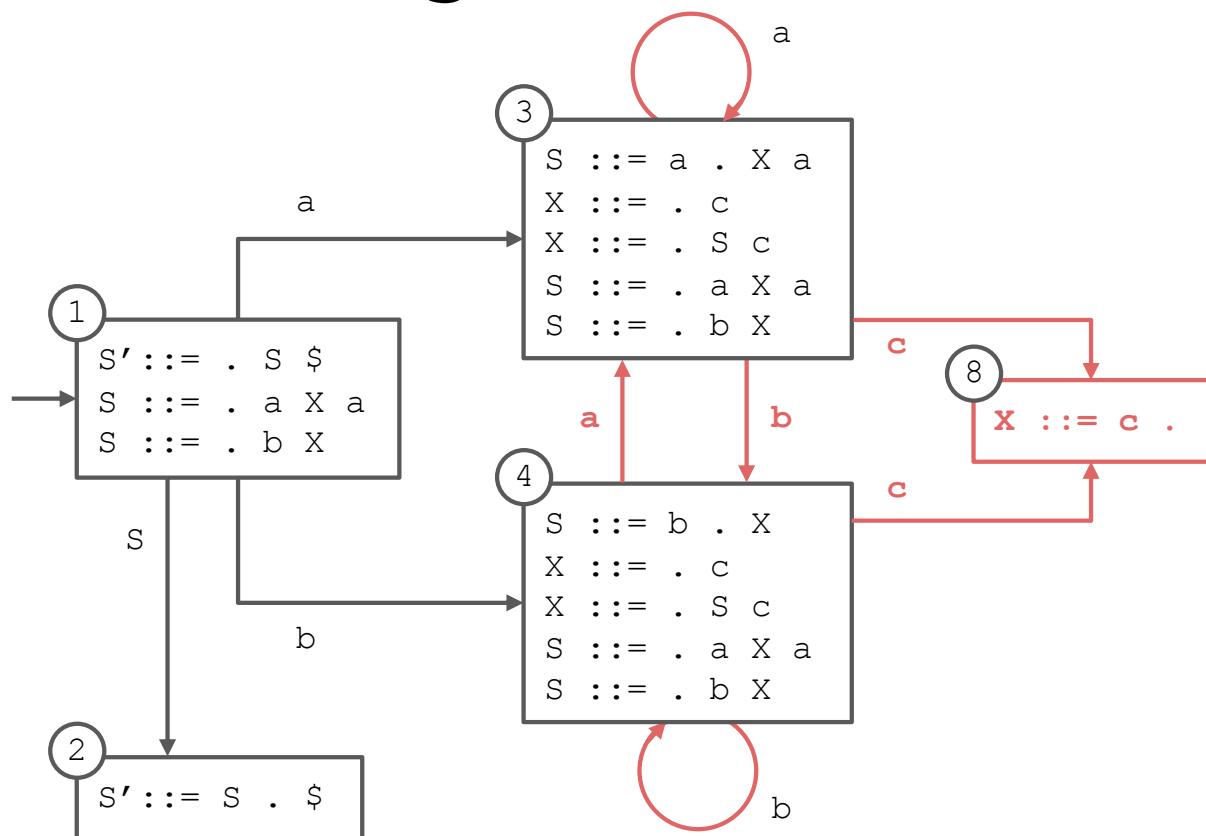
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State Diagram Construction



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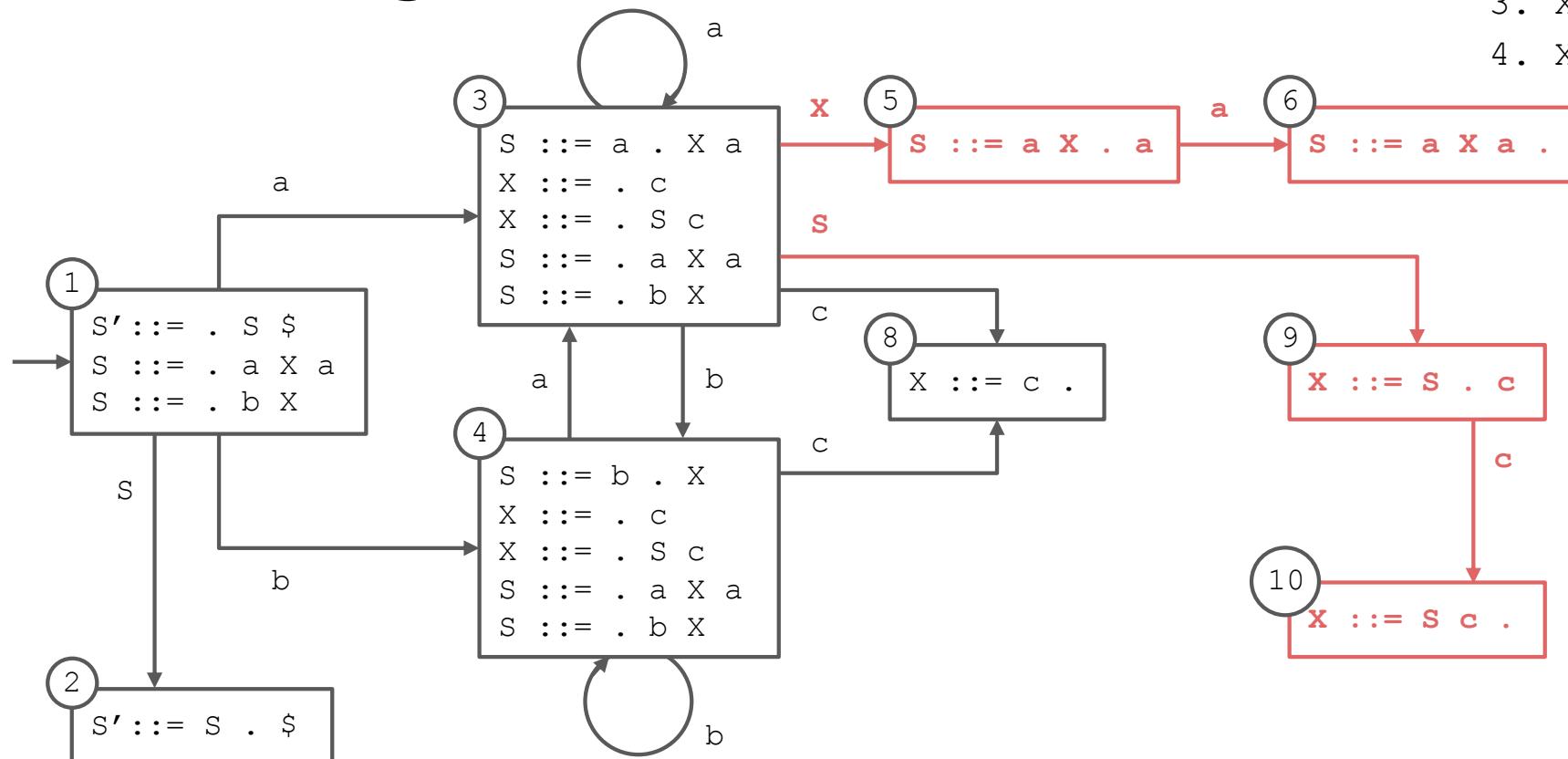
State Diagram Construction



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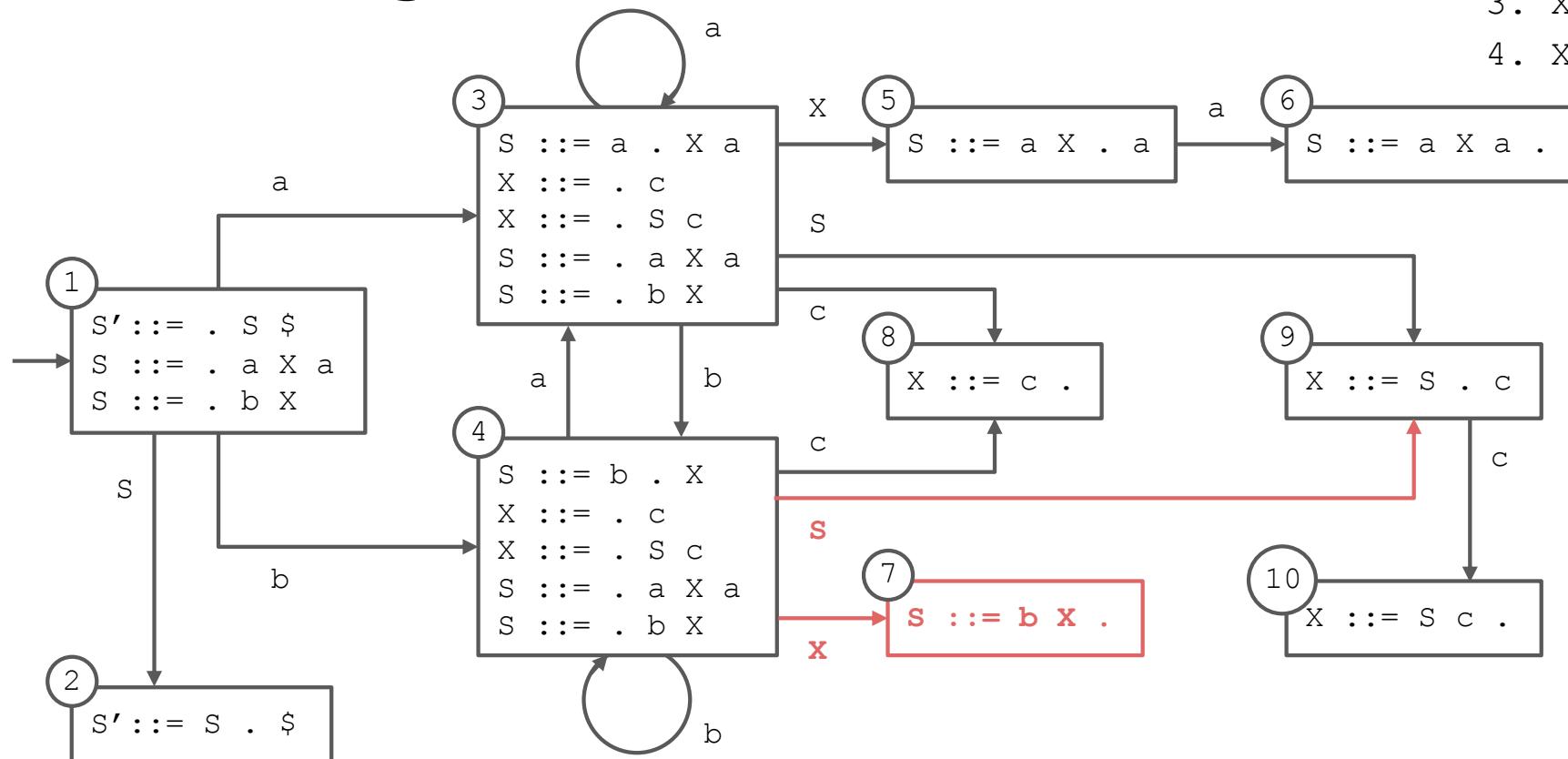
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State Diagram Construction



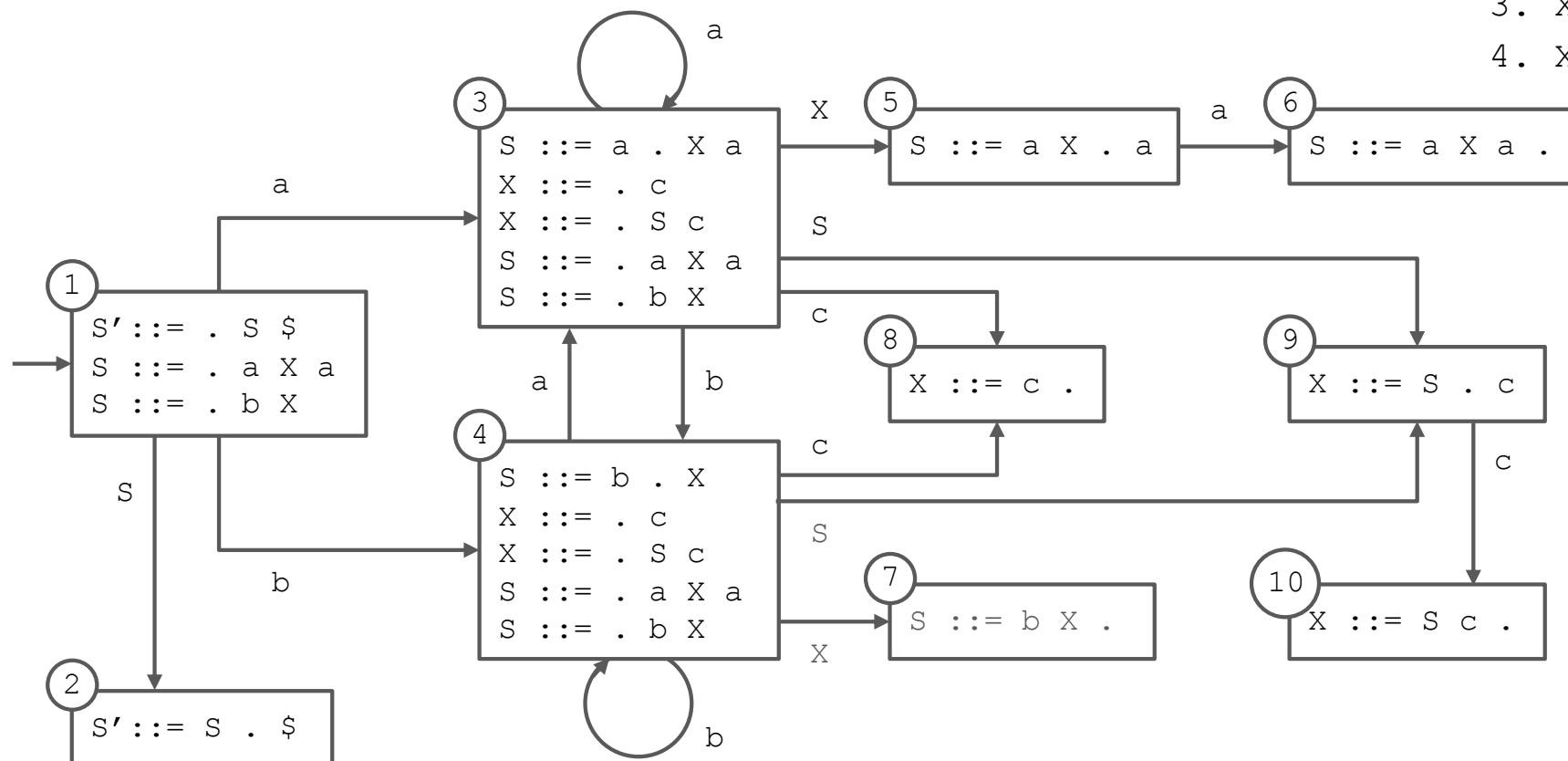
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State Diagram Construction



0. $S' ::= S \$$
1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$

Completed State Diagram



Converted to Table

s# means “shift and enter state #”

- occurs when there is a transition on a terminal

r# means “reduce using production #”

- occurs when a state contains an item with the location at the end of the right-hand side

g# means “go to state #”

- occurs when there is a transition on a nonterminal

acc means “accept”

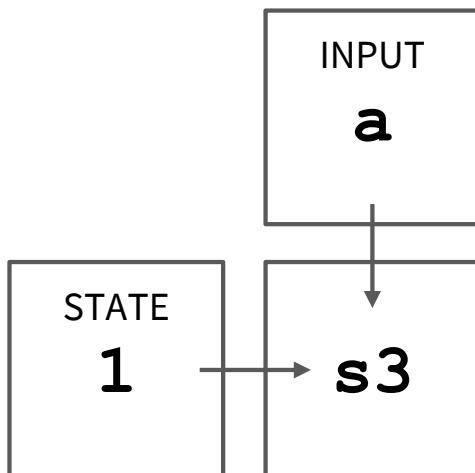
- occurs when the start symbol (S here) has been completed and there is no more input

STATE	ACTION				GOTO	
	a	b	c	\$	S	X
1	s3	s4			g2	
2				acc		
3	s3	s4	s8		g9	g5
4	s3	s4	s8		g9	g7
5	s6					
6	r1	r1	r1	r1		
7	r2	r2	r2	r2		
8	r3	r3	r3	r3		
9				s10		
10	r4	r4	r4	r4		

Parse Trace

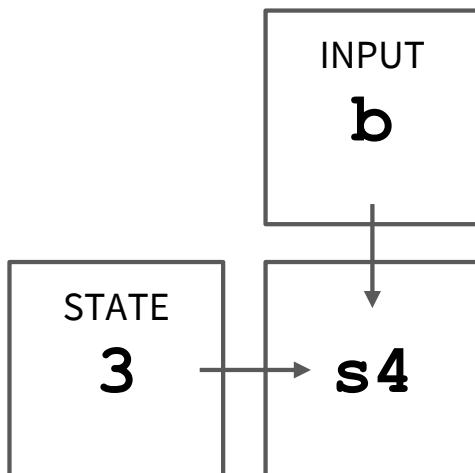
STACK	INPUT	ACTION
\$ 1	a b c c a \$	

Parse Trace



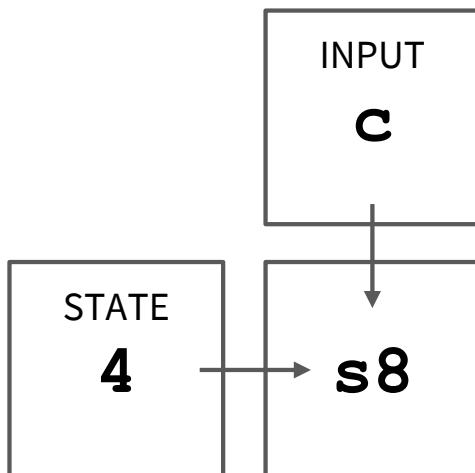
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3	a b c c a \$ b c c a \$	SHIFT

Parse Trace



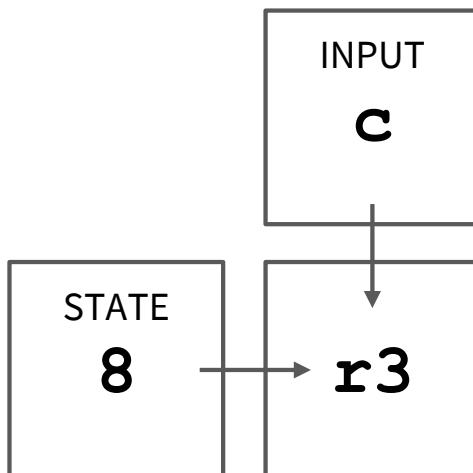
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3 \$ 1 a 3 b 4	a b c c a \$ b c c a \$ c c a \$	SHIFT SHIFT

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	a	

Parse Trace

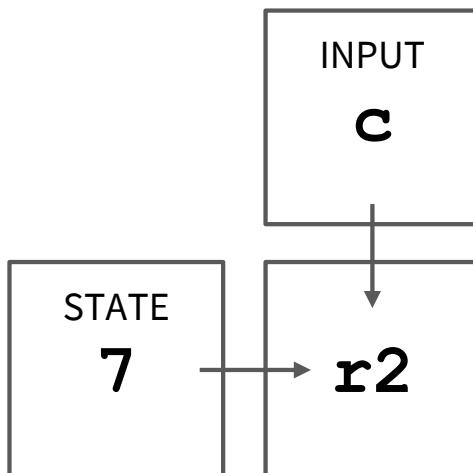


3. X ::= c

(and GOTO step:
s4 & X -> g7)

STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	

Parse Trace

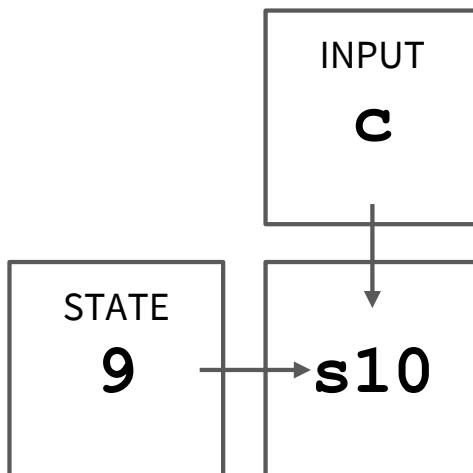


2. $S ::= b \ X$

(and GOTO step:
 $s3 \ \& \ S \rightarrow g9$)

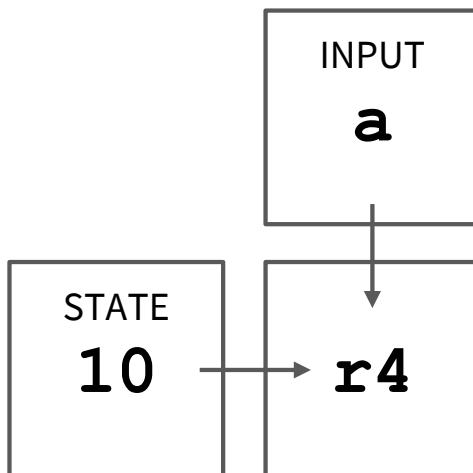
STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	

Parse Trace



STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	SHIFT
\$ 1 a 3 S 9 c 10	a \$	

Parse Trace

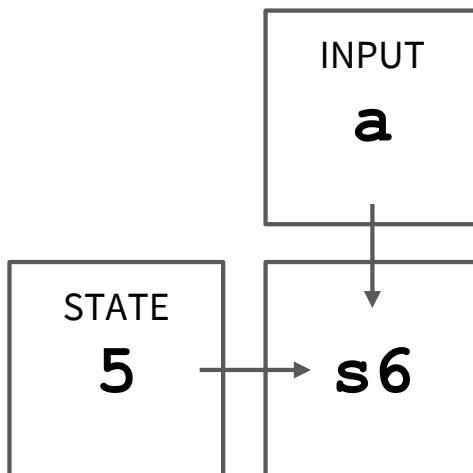


4. $X ::= S \ C$

(and GOTO step:
 $s3 \ \& \ X \rightarrow g5$)

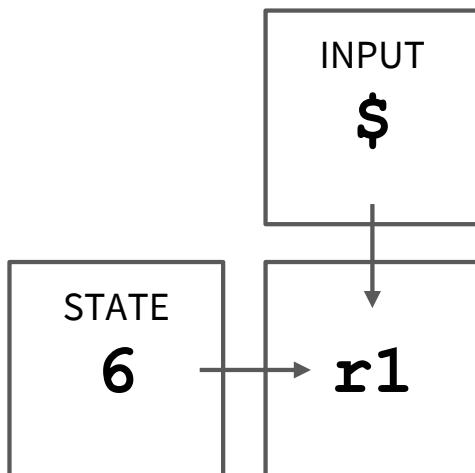
STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	SHIFT
\$ 1 a 3 S 9 c 10	a \$	
\$ 1 a 3 X 5	a \$	REDUCE

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	c	REDUCE
\$ 1 a 3 b 4 X 7	a	REDUCE
\$ 1 a 3 S 9	c	SHIFT
\$ 1 a 3 S 9 c 10	a	REDUCE
\$ 1 a 3 X 5	a	SHIFT
\$ 1 a 3 X 5 a 6	\$	

Parse Trace



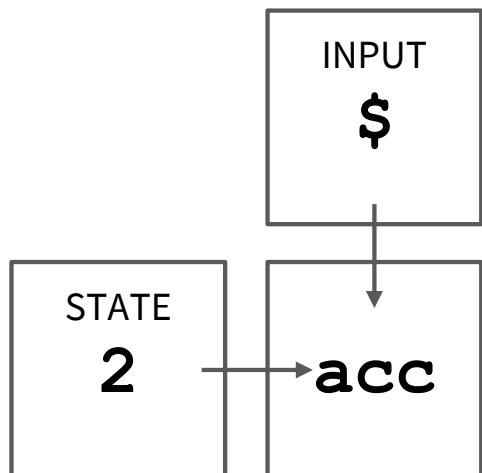
1. S ::= a X a

(and GOTO step:

s1 & S -> g2)

STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	SHIFT
\$ 1 a 3 S 9 c 10	a \$	REDUCE
\$ 1 a 3 X 5	a \$	SHIFT
\$ 1 a 3 X 5 a 6	\$	REDUCE
\$ 1 S 2	\$	

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	c	REDUCE
\$ 1 a 3 b 4 X 7	a	REDUCE
\$ 1 a 3 S 9	c	SHIFT
\$ 1 a 3 S 9 c 10	a	REDUCE
\$ 1 a 3 X 5	a	SHIFT
\$ 1 a 3 X 5 a 6	\$	REDUCE
\$ 1 S 2	\$	ACCEPT