

Section 3: LR Parsing

CSE 401/M501

Adapted from Spring 2021

Announcements

- Scanner is due tonight
 - Be sure to test, push, and tag!
- Every person has 4 project late days and 4 assignment late days
 - Up to 2 can be used per assignment
 - Each late day gives an extra 24 hour chunk (including weekend days)
 - We recommend not using your late days this early if possible!
 - **Submitting project components a day late uses a project late day from each partner!**

Agenda

- (Fast) LR terminology review
- Worksheet

Get Your LR Jargon On

- Frontier
 - The upper “layer” of the current parse tree (held in the stack)

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- Sentential Form
 - A string that can be generated at any point in a derivation (can be reached using any number of productions from the start symbol)

Get Your LR Jargon On

- Frontier
 - The upper “layer” of the current parse tree (held in the stack)
- Sentential Form
 - A string that can be generated at any point in a derivation (can be reached using any number of productions from the start symbol)
- Handle
 - An occurrence of the right side of a production in the frontier that is used in the rightmost derivation to arrive at the current string
 - Given the derivation ... $\Rightarrow \underline{a} \underline{A} \underline{b} \underline{c} \underline{d} \underline{e} \Rightarrow \underline{a} \underline{b} \underline{c} \underline{d} \underline{e}$, using the production $A ::= b$:
 - The production ‘ $A ::= b$ ’ at index 1 would be a handle of $\underline{a} \underline{b} \underline{c} \underline{d} \underline{e}$

Get Your LR Jargon On - Example

Shift-Reduce Example

$$\begin{aligned}S &::= aABe \\A &::= Abc \mid b \\B &::= d\end{aligned}$$

	Stack	Input	Action
Frontier	\$	abbcde\$	<i>shift</i>
	\$a	bbcde\$	<i>shift</i>
	\$ab	bcde\$	<i>reduce</i>
	\$aA	bcde\$	<i>shift</i>
	\$aAb	cde\$	<i>shift</i>
	\$aAbc	de\$	<i>reduce</i>
	\$aA	de\$	<i>shift</i>
	\$aAd	e\$	<i>reduce</i>
	\$aAB	e\$	<i>shift</i>
	\$aABe	\$	<i>reduce</i>
	\$S	\$	<i>accept</i>

Get Your LR Jargon On - Example

Shift-Reduce Example

$$\begin{aligned}S &::= aABe \\A &::= Abc \mid b \\B &::= d\end{aligned}$$

Sentential
Forms

Stack	Input	Action
\$	abbcde\$	<i>shift</i>
\$a	bbcde\$	<i>shift</i>
\$ab	bcde\$	<i>reduce</i>
\$aA	bcde\$	<i>shift</i>
\$aAb	cde\$	<i>shift</i>
\$aAbc	de\$	<i>reduce</i>
\$aA	de\$	<i>shift</i>
\$aAd	e\$	<i>reduce</i>
\$aAB	e\$	<i>shift</i>
\$aABe	\$	<i>reduce</i>
\$S	\$	<i>accept</i>

Get Your LR Jargon On - Example

Shift-Reduce Example

$$\begin{aligned}S &::= aABe \\A &::= Abc \mid b \\B &::= d\end{aligned}$$

	Stack	Input	Action
Handles A ::= b at index 2	\$	abbcde\$	shift
	\$a	bbcde\$	shift
	\$ab	bcde\$	reduce
	\$aA	bcde\$	shift
A ::= Abc at index 4	\$aAb	cde\$	shift
	\$aAbc	de\$	reduce
	\$aA	de\$	shift
B ::= d at index 3	\$aAd	e\$	reduce
S ::= aABe at index 4	\$aAB	e\$	shift
	\$aABe	\$	reduce
	\$S	\$	accept

A Little Bit More Jargon

- Viable Prefix
 - The prefixes of a right sentential form that do not extend beyond the end of its handle
 - Perhaps less confusing -> the set of prefixes of strings that can appear on the stack of a shift-reduce parser

A Little Bit More Jargon

- Viable Prefix
 - The prefixes of a right sentential form that do not extend beyond the end of its handle
 - Perhaps less confusing -> the set of prefixes of strings that can appear on the stack of a shift-reduce parser
- Item
 - A marked production (a production with a '.' in it)
 - [A ::= .XY], [A ::= X.Y], [A ::= XY.]

Get Your LR Jargon On - Example

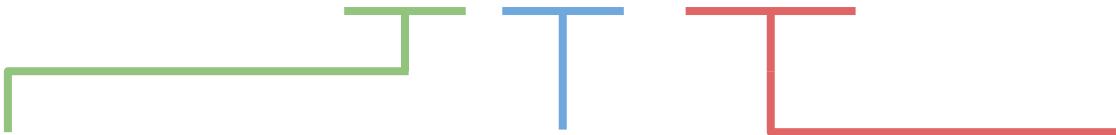
Shift-Reduce Example

$$\begin{aligned}S &::= aABe \\A &::= Abc \mid b \\B &::= d\end{aligned}$$

Stack	Input	Action
\$	abbcde\$	<i>shift</i>
\$a	bbcde\$	<i>shift</i>
\$ab	bcde\$	<i>reduce</i>
\$aA	bcde\$	<i>shift</i>
\$aAb	cde\$	<i>shift</i>
\$aAbc	de\$	<i>reduce</i>
\$aA	de\$	<i>shift</i>
\$aAd	e\$	<i>reduce</i>
\$aAB	e\$	<i>shift</i>
\$aABe	\$	<i>reduce</i>
\$S	\$	<i>accept</i>

Viable
Prefix
(all prefixes of
these strings
are also viable
prefixes)

L R (0)



Left-to-Right

Only takes one pass,
performed from the left

Rightmost

At each point, finds the
derivation for the rightmost
handle (bottom-up)

No Lookahead

Decide what to do based on
current parser state and
stack, ignoring next input

Problem 1 (On Worksheet)

0. $S' ::= S \$$

1. $S ::= a z$

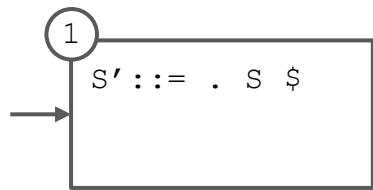
2. $S ::= b$

3. $z ::= a$

4. $z ::= b s$

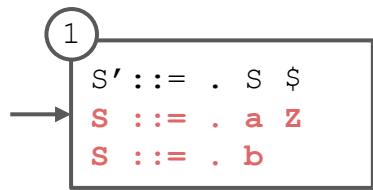
State Diagram Construction

0. $S' ::= S \$$
1. $S ::= a Z$
2. $S ::= b$
3. $Z ::= a$
4. $Z ::= b S$



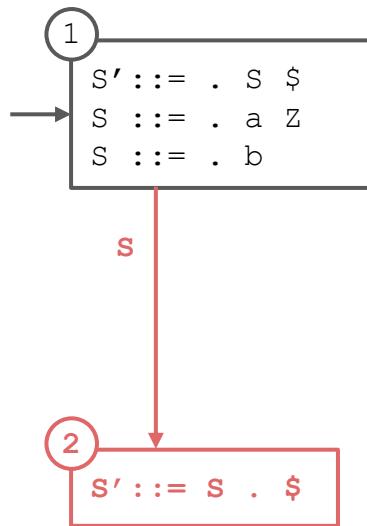
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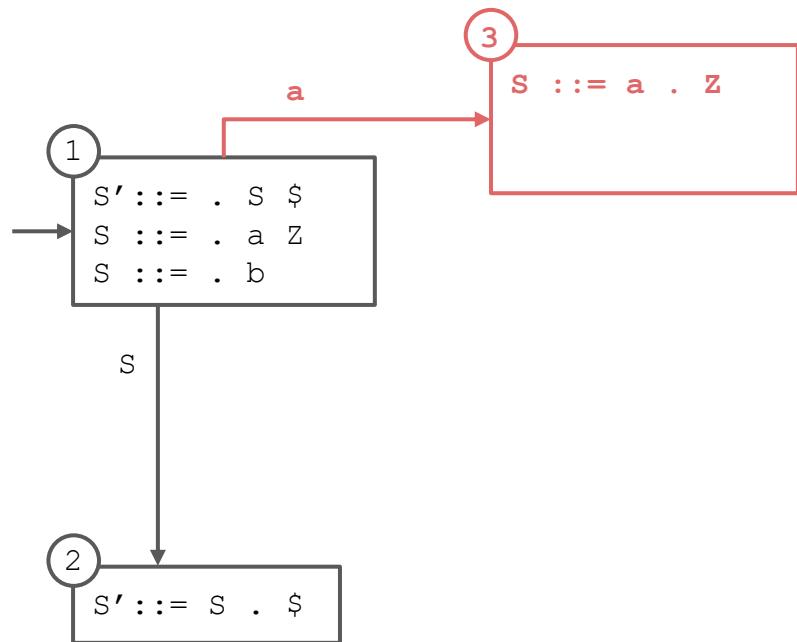
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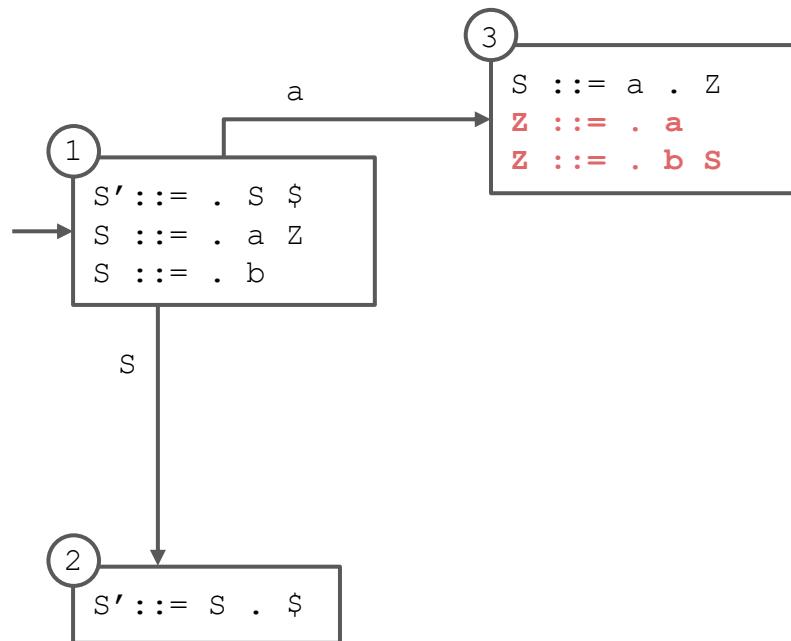
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1. $S ::= a Z$
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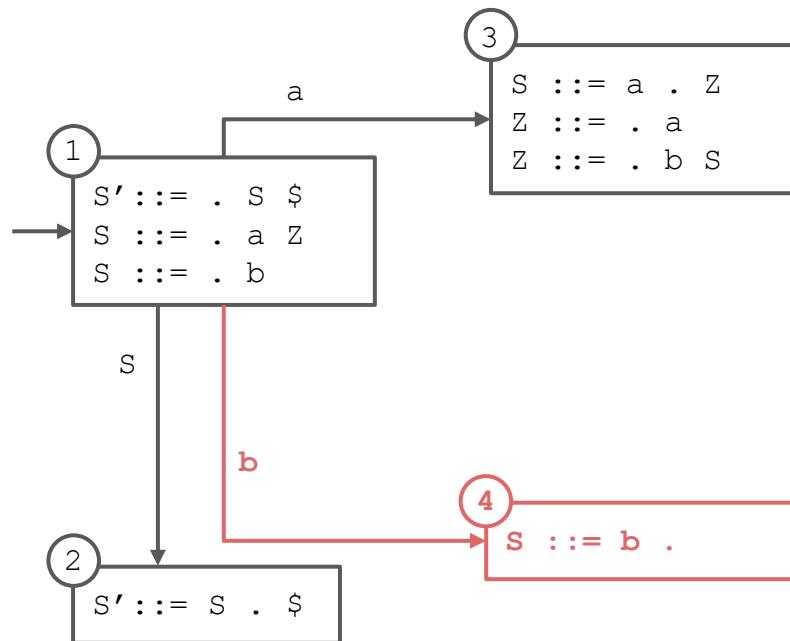
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State Diagram Construction



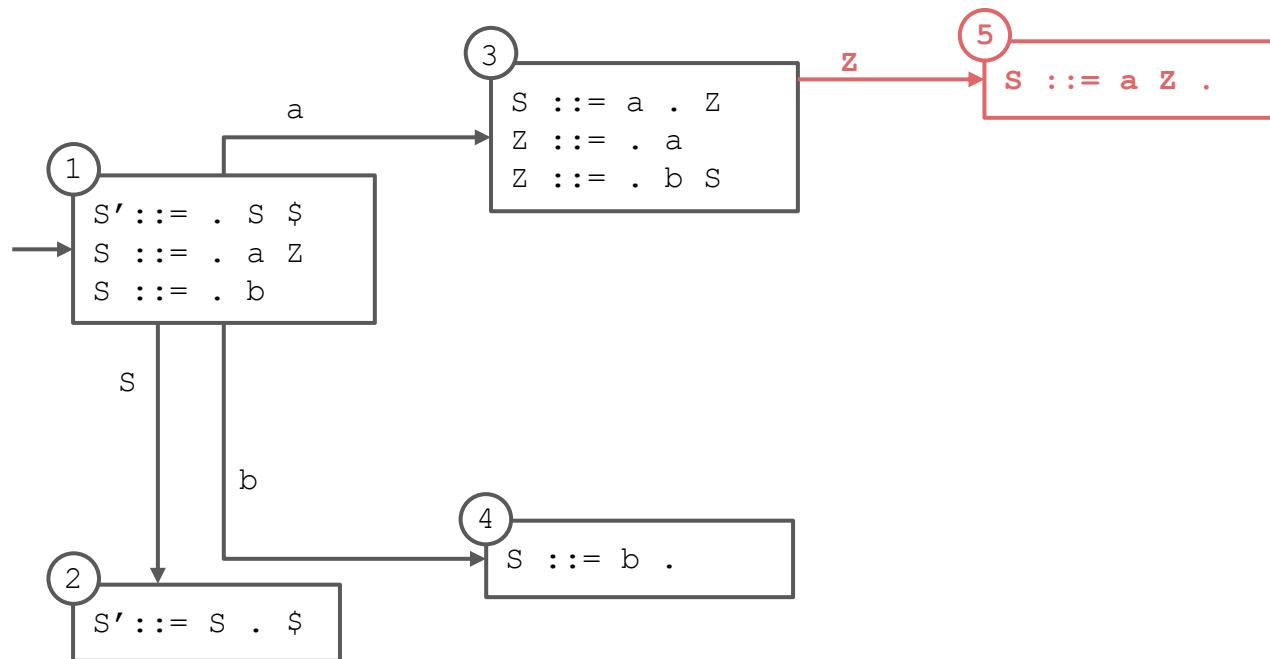
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State Diagram Construction



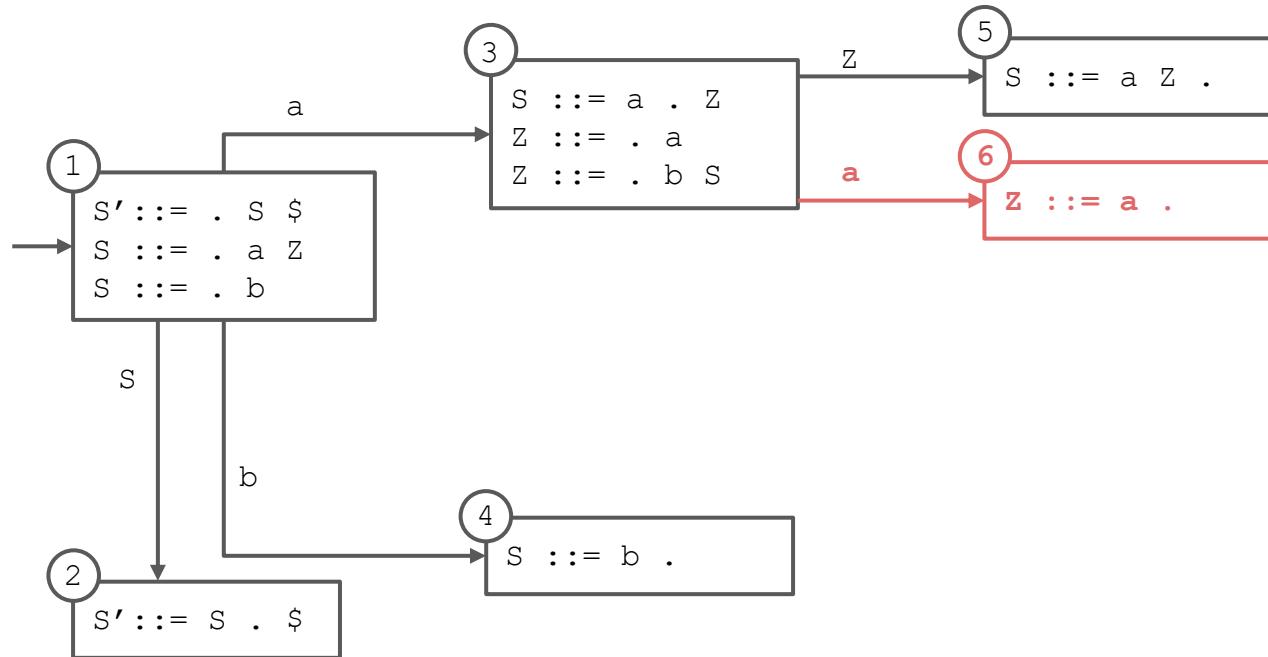
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State Diagram Construction



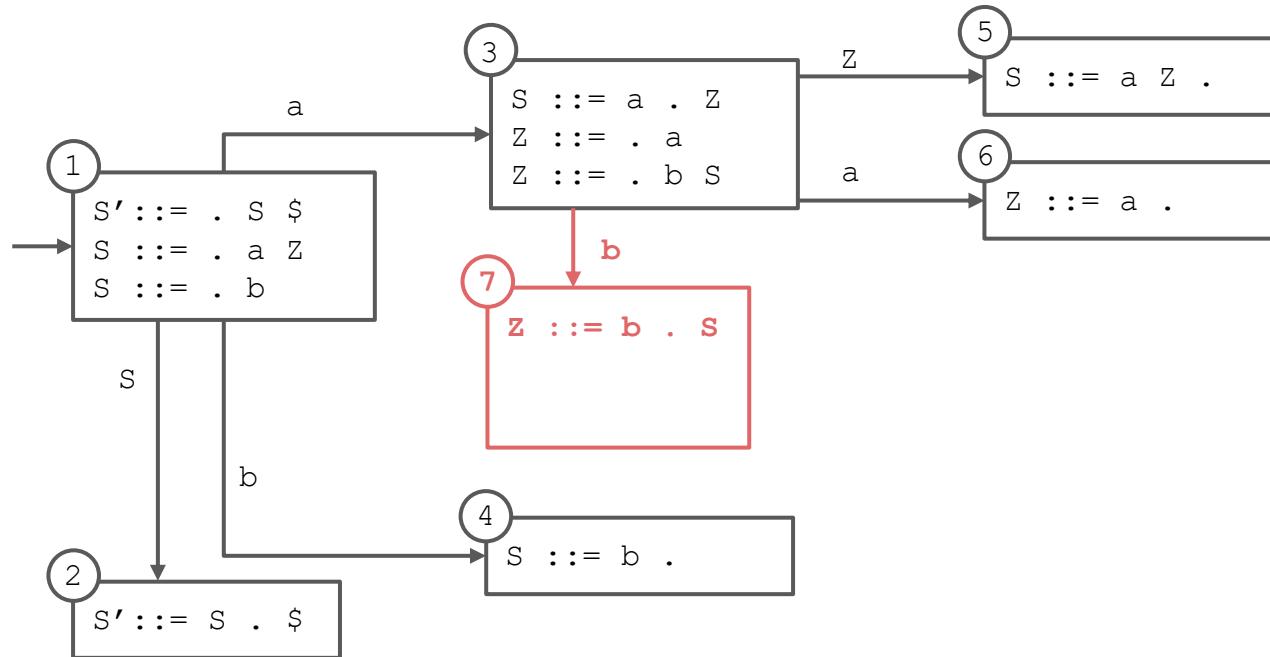
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State Diagram Construction



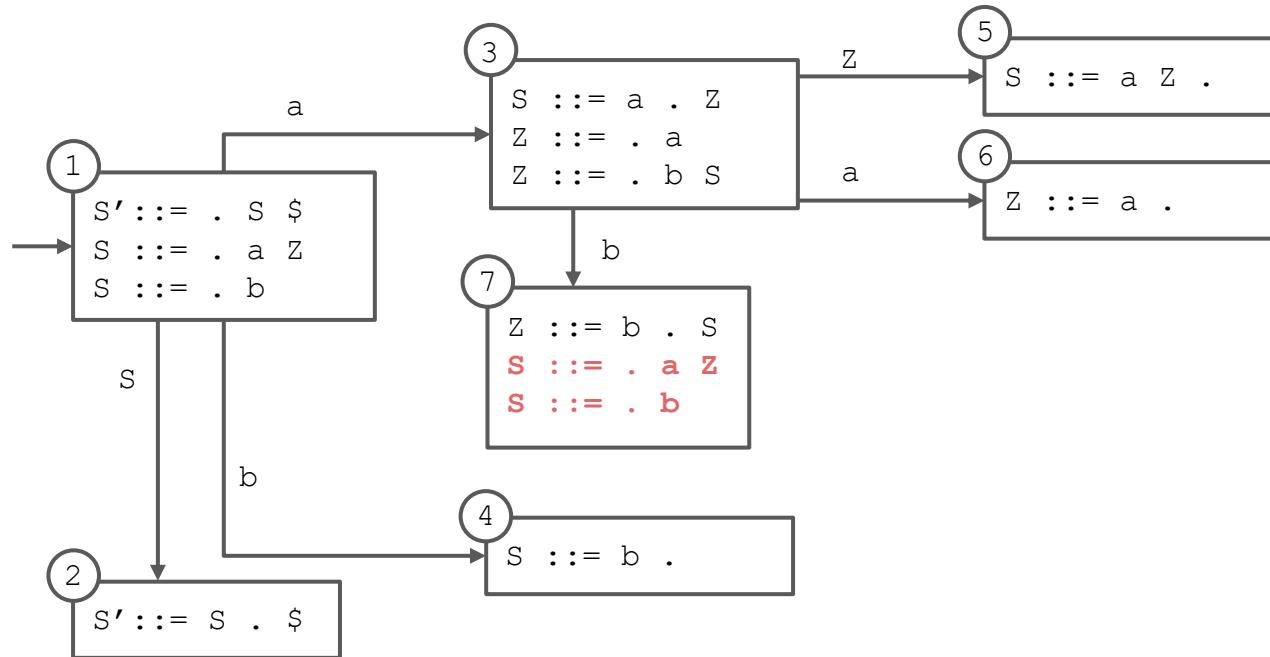
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State Diagram Construction



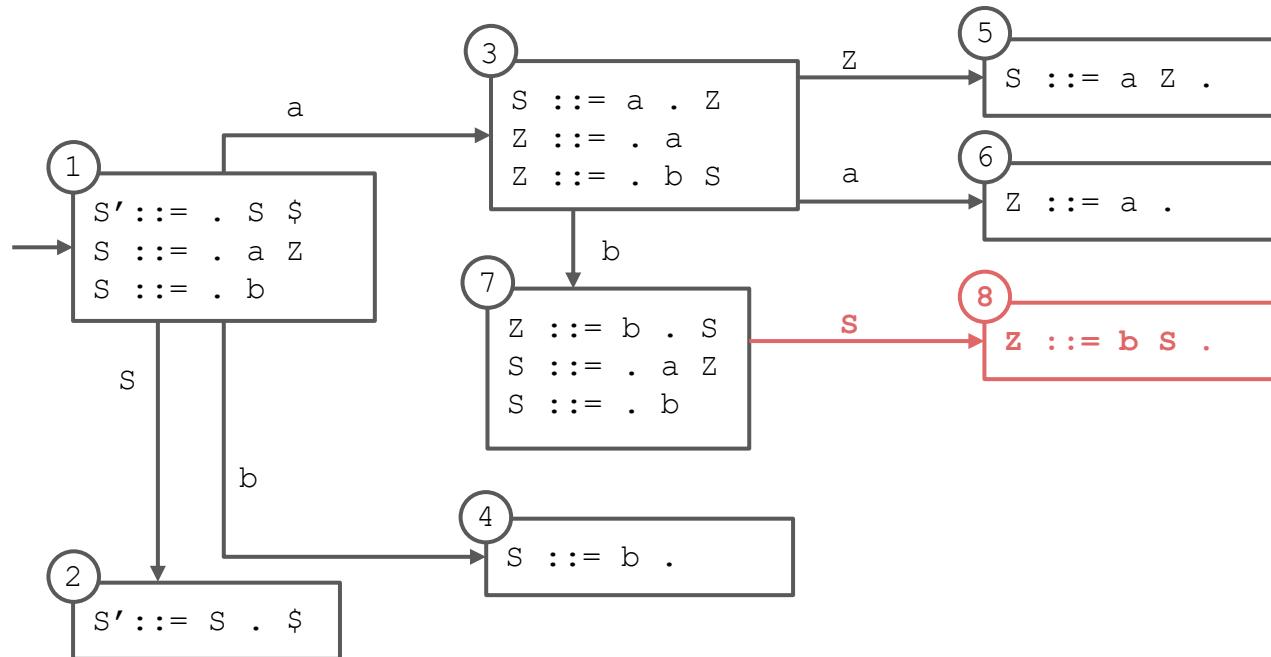
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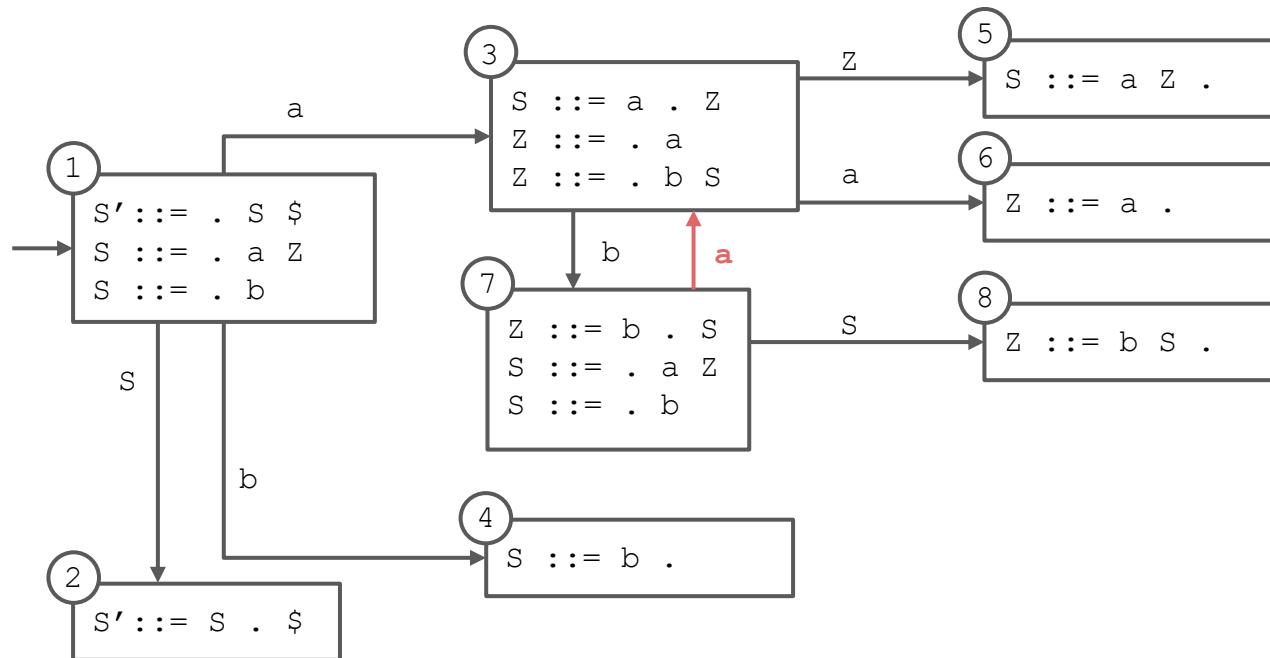
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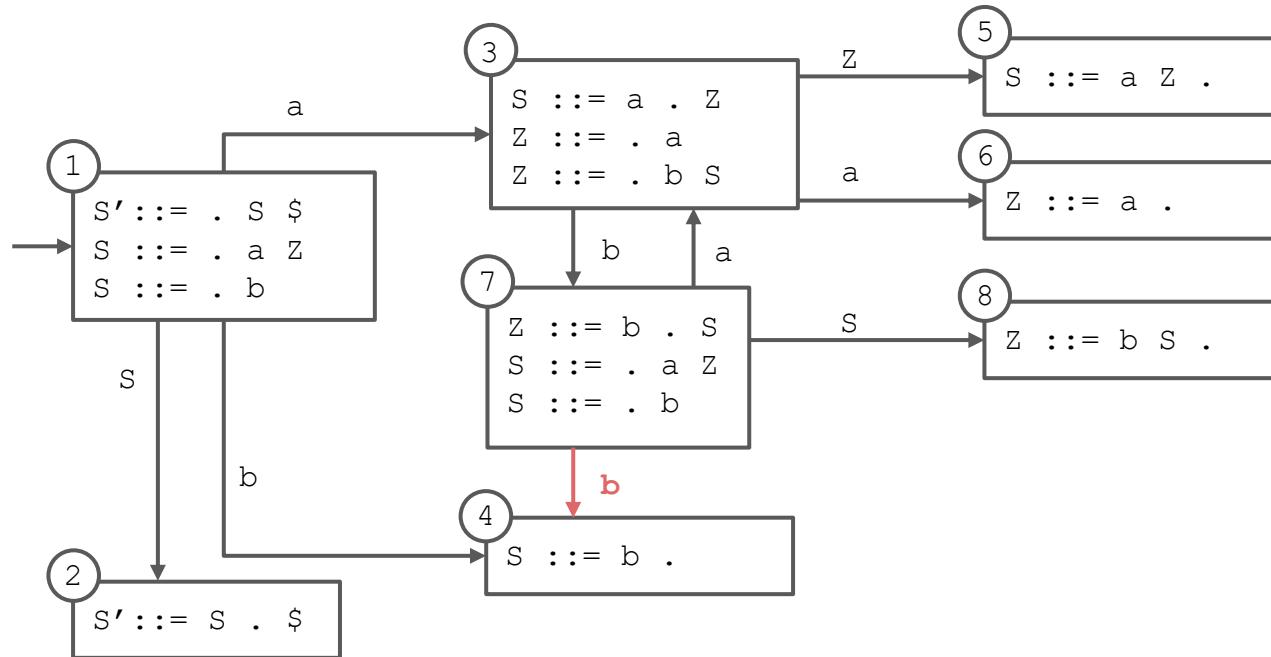
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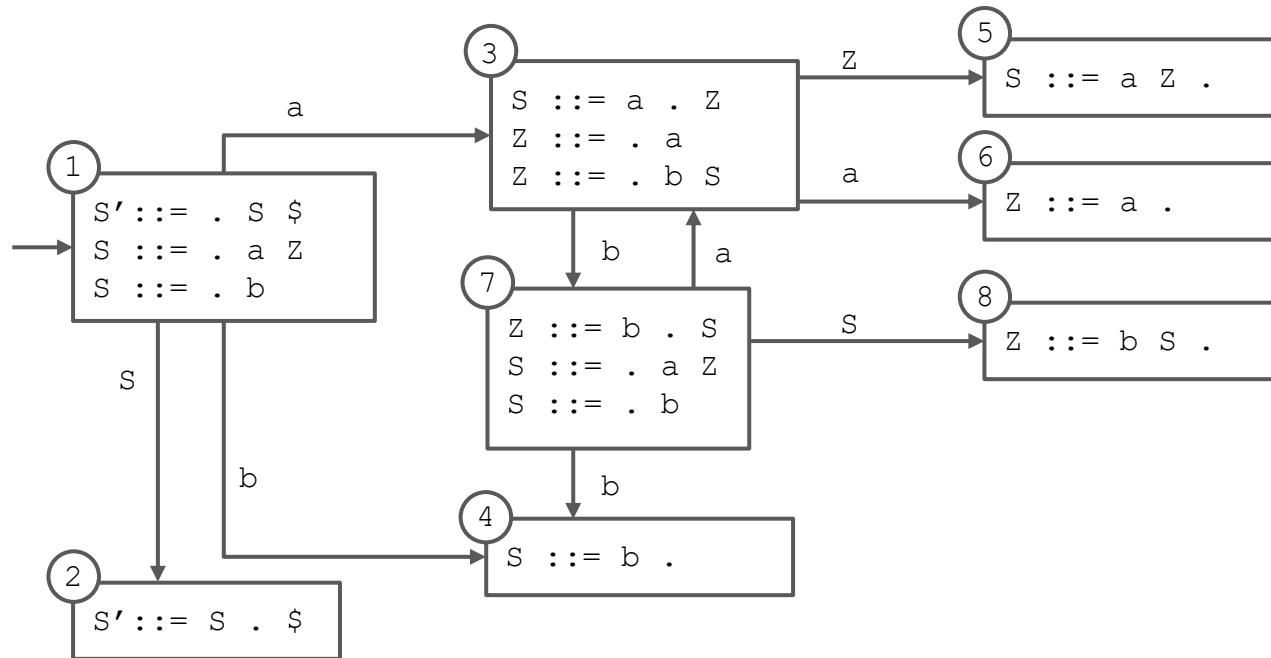
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State Diagram Construction



0. $S' ::= S \$$
1. $S ::= a Z$
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4. $Z ::= b S$

Completed State Diagram



Converted to Table

s# means “shift and enter state #”

- occurs when there is a transition on a terminal

r# means “reduce using production #”

- occurs when a state contains an item with the location at the end of the right-hand side

g# means “go to state #”

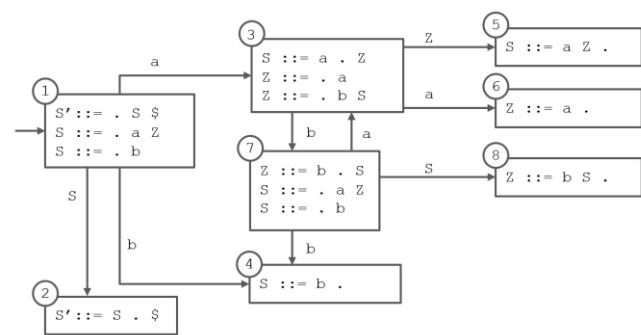
- occurs when there is a transition on a nonterminal

acc means “accept”

- occurs when the start symbol (S here) has been completed and there is no more input

STATE	ACTION			GOTO	
	a	b	\$	S	Z
1	s3	s4		g2	
2			acc		
3	s6	s7			g5
4	r2	r2	r2		
5	r1	r1	r1		
6	r3	r3	r3		
7	s3	s4		g8	
8	r4	r4	r4		

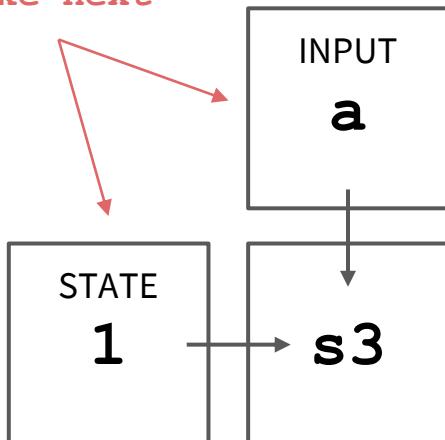
Parse Trace



STACK	INPUT	ACTION
\$ 1	a b a b b \$	

Parse Trace

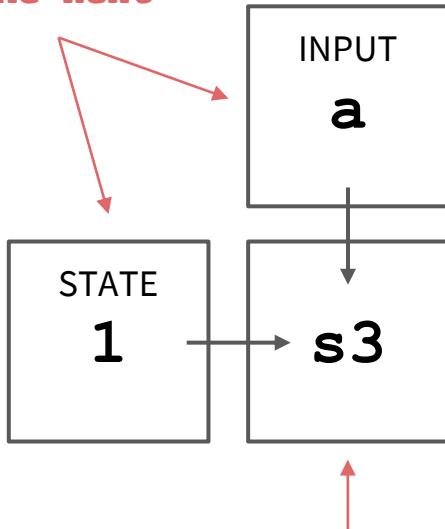
Row and column of table to look up: decides what action to take next



STACK	INPUT	ACTION
\$ 1 \$ 1 a	a b a b b \$ b a b b \$	SHIFT

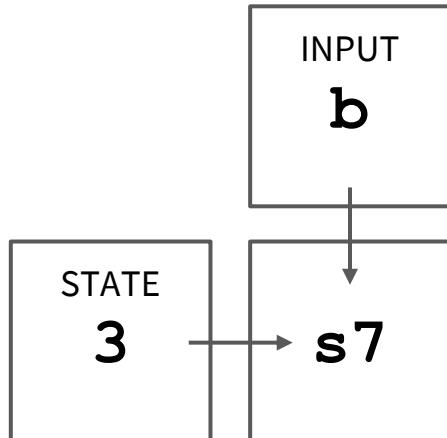
Parse Trace

Row and column of table to look up: decides what action to take next



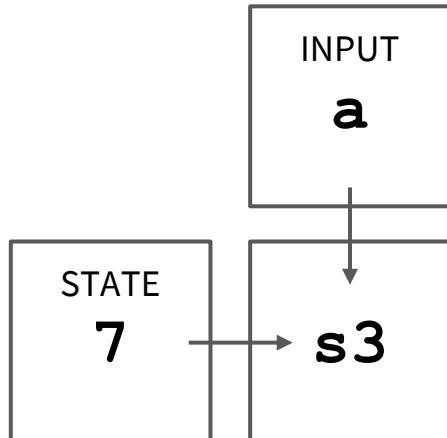
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3	a b a b b \$ b a b b \$	SHIFT

Parse Trace



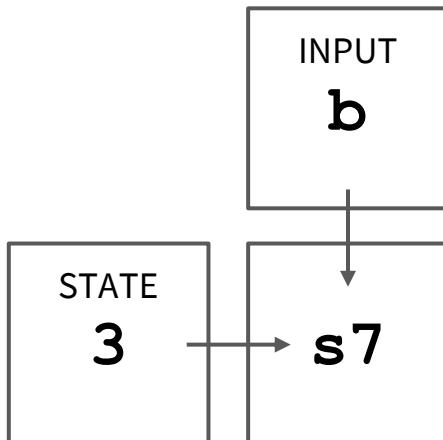
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3 \$ 1 a 3 b 7	a b a b b \$ b a b b \$ a b b \$	SHIFT SHIFT

Parse Trace



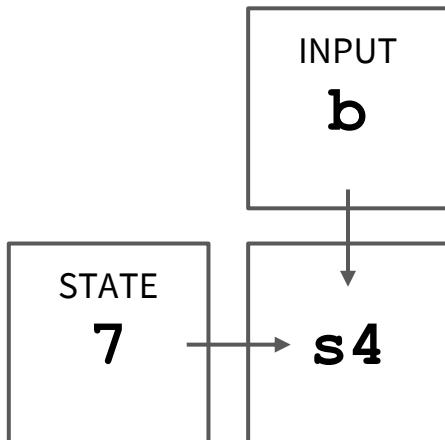
STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 7	a	SHIFT
\$ 1 a 3 b 7 a 3	b	SHIFT

Parse Trace



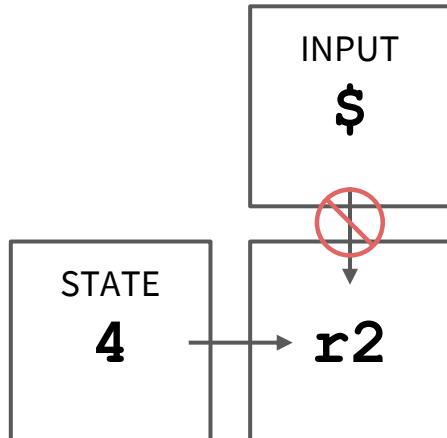
STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT

Parse Trace



STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	SHIFT

Parse Trace

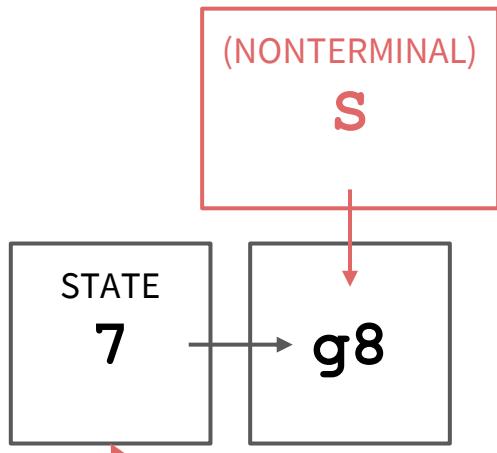


2. S ::= b

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 S	\$	REDUCE

For LR(0), the input doesn't technically matter here

Parse Trace

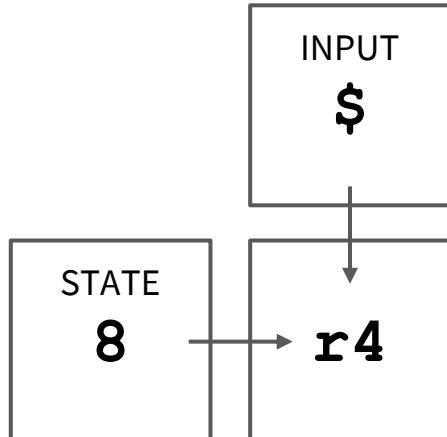


STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE

After a reduction, we go back to a previous state on the stack and use the reduced non-terminal to determine what state to GOTO.

This allows the parser to run in $O(n)$ time, since it doesn't have to re-evaluate the entire stack!

Parse Trace

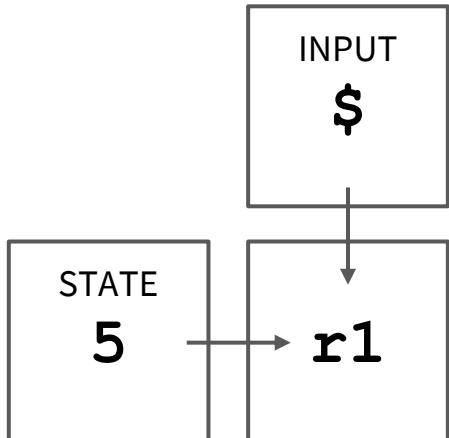


4. Z ::= b S

(and GOTO step:
s3 & Z -> g5)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	

Parse Trace



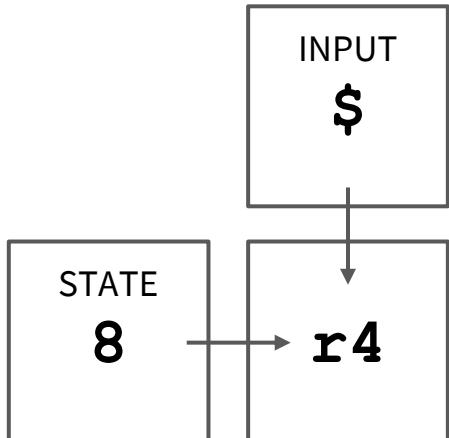
1. S ::= a z

(and GOTO step:

s7 & S -> g8)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 Z 5	\$	REDUCE
\$ 1 a 3 b 7 S 8	\$	

Parse Trace



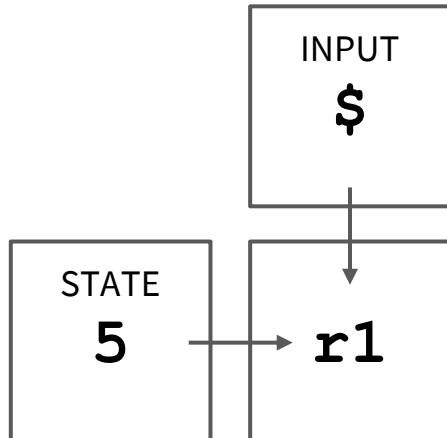
4. **Z ::= b S**

(and GOTO step:

s3 & Z -> g5)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 z 5	\$	REDUCE

Parse Trace



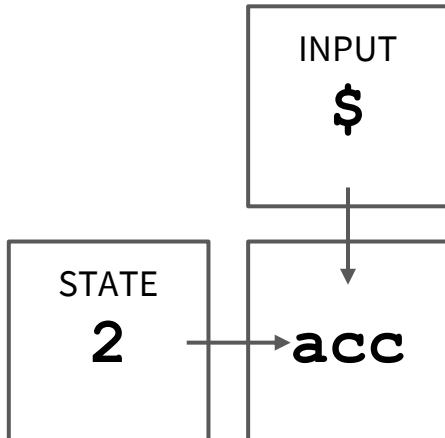
1. $S ::= a \cdot z$

(and GOTO step:

$s1 \& S \rightarrow g2$)

STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 z 5	\$	REDUCE
\$ 1 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 z 5	\$	REDUCE
\$ 1 S 2	\$	REDUCE

Parse Trace



STACK	INPUT	ACTION
\$ 1	a b a b b \$	SHIFT
\$ 1 a 3	b a b b \$	SHIFT
\$ 1 a 3 b 7	a b b \$	SHIFT
\$ 1 a 3 b 7 a 3	b b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7	b \$	SHIFT
\$ 1 a 3 b 7 a 3 b 7 b 4	\$	REDUCE
\$ 1 a 3 b 7 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 b 7 a 3 Z 5	\$	REDUCE
\$ 1 a 3 b 7 S 8	\$	REDUCE
\$ 1 a 3 Z 5	\$	REDUCE
\$ 1 S 2	\$	ACCEPT

Problem 2 (On Worksheet)

0. $S' ::= S \$$

1. $S ::= a \times a$

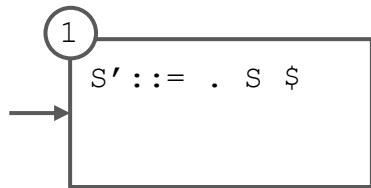
2. $S ::= b \times$

3. $X ::= c$

4. $X ::= S \ c$

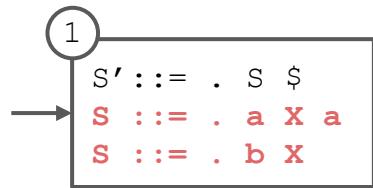
State Diagram Construction

0. $S' ::= S \$$
1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$



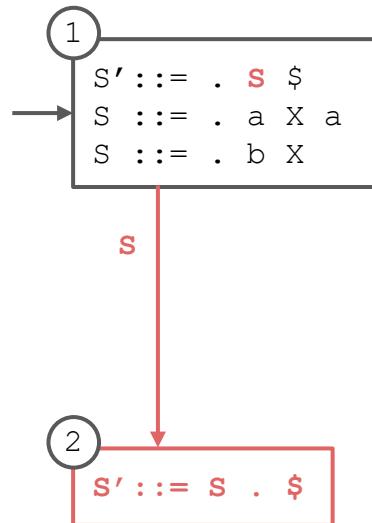
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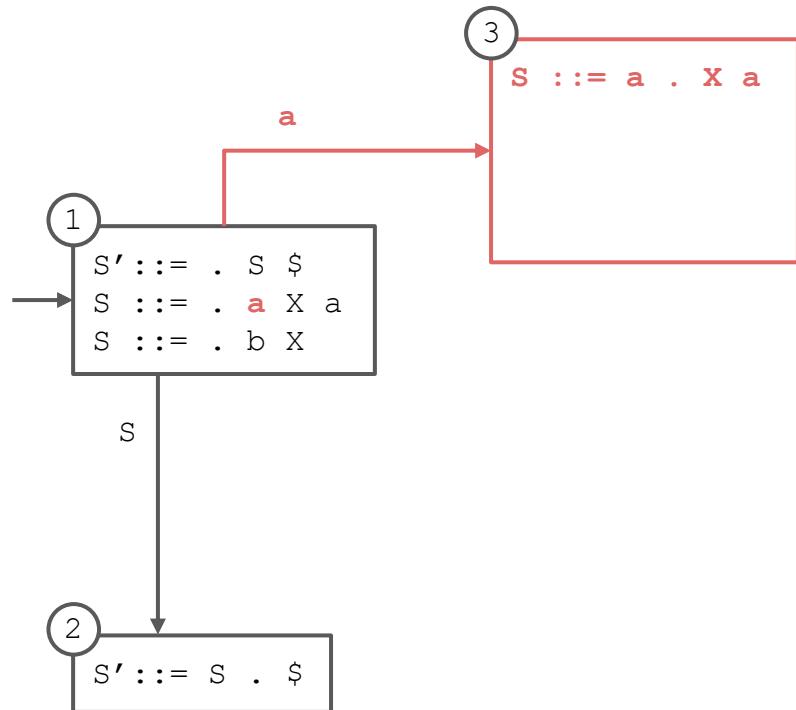
State Diagram Construction

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4. $X ::= S c$



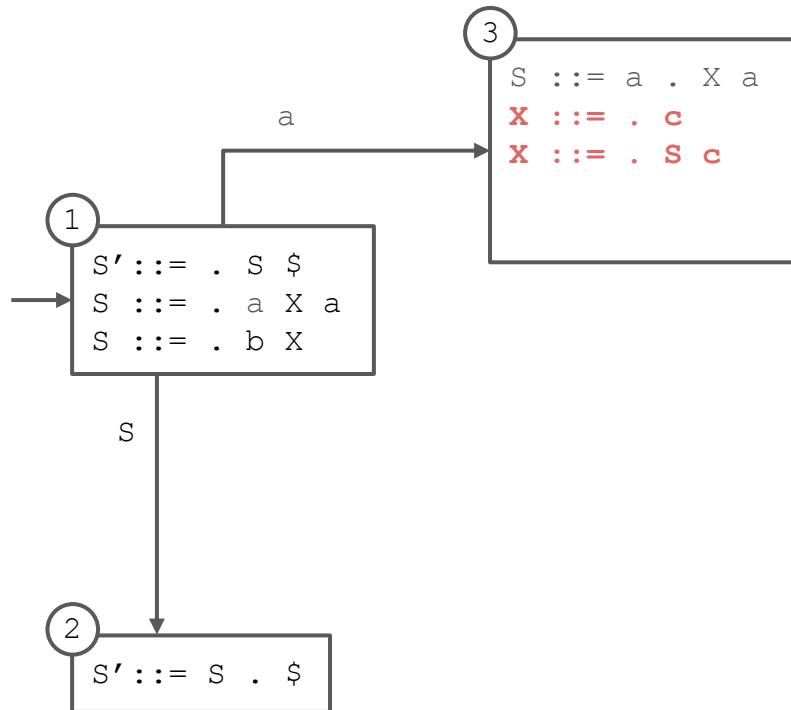
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- 1. $S ::= a X a$
- 2. $S ::= b X$
- 3. $X ::= c$
- 4. $X ::= S c$

State Diagram Construction



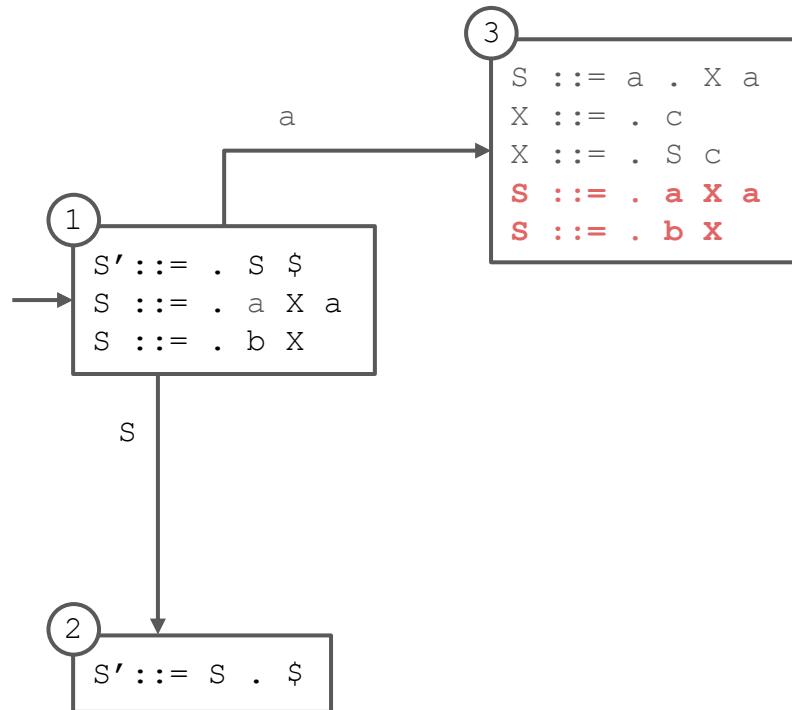
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State Diagram Construction



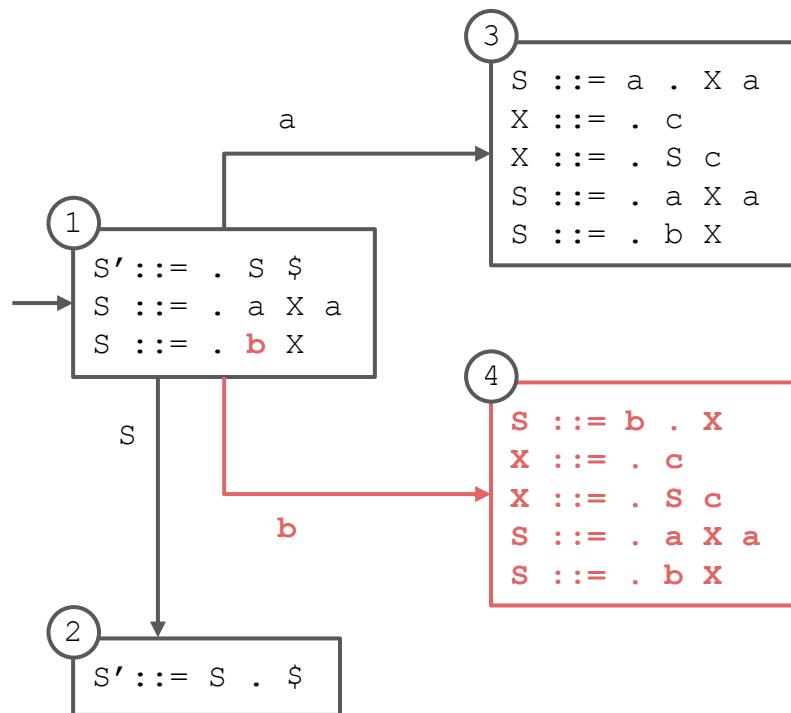
0. $S' ::= S \$$
1. $S ::= a X a$
2. $S ::= b X$
3. $X ::= c$
4. $X ::= S c$

State Diagram Construction



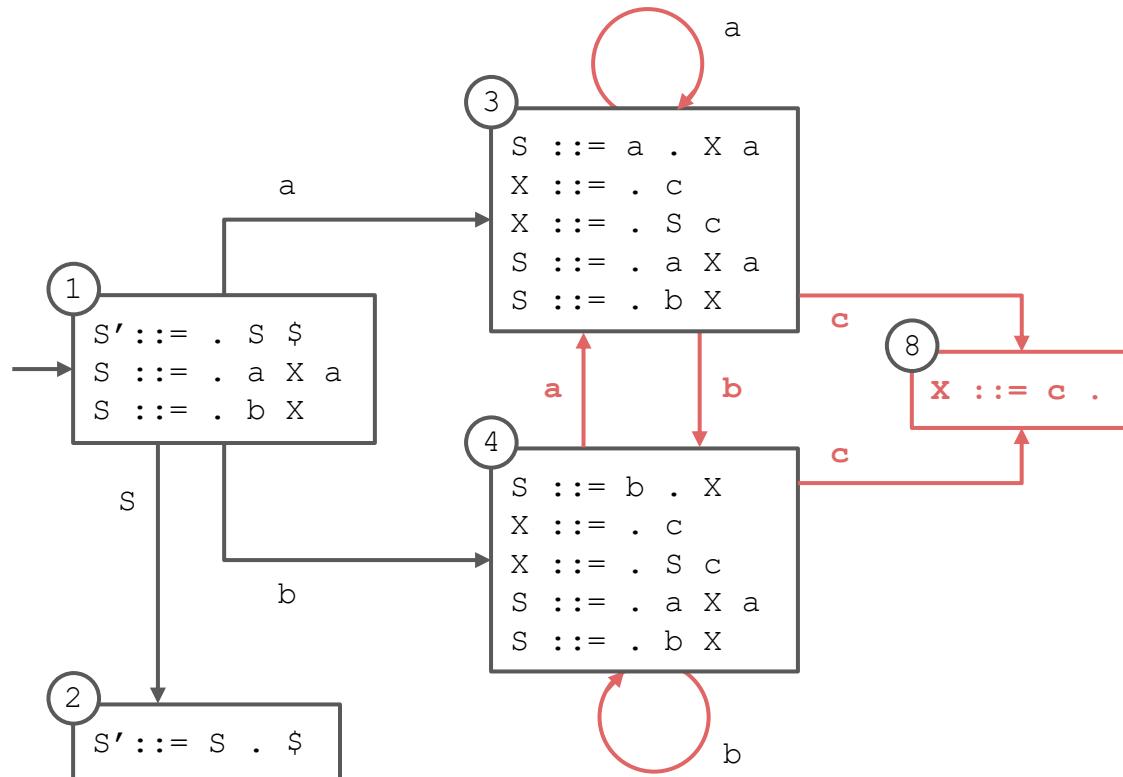
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State Diagram Construction



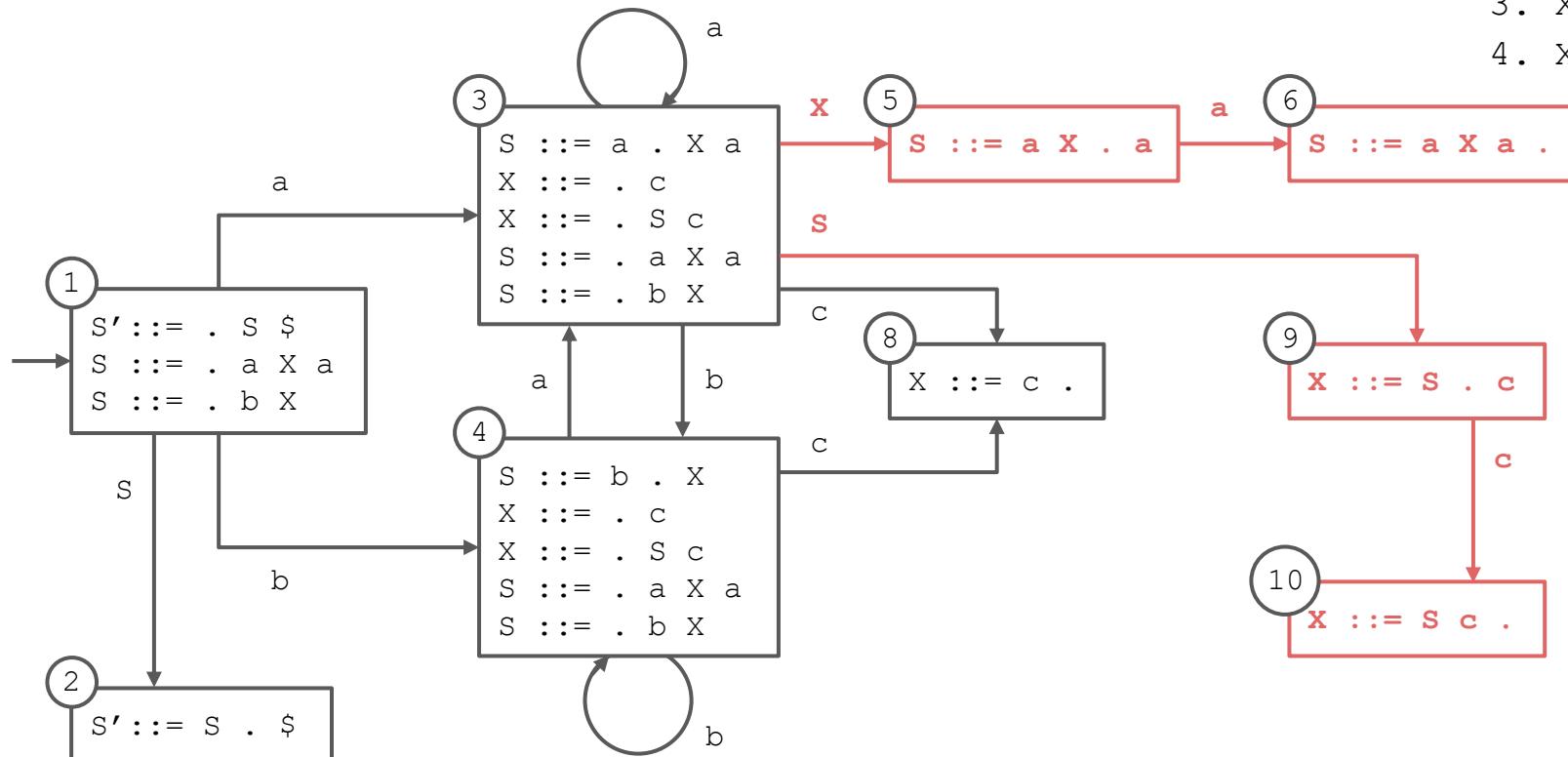
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State Diagram Construction



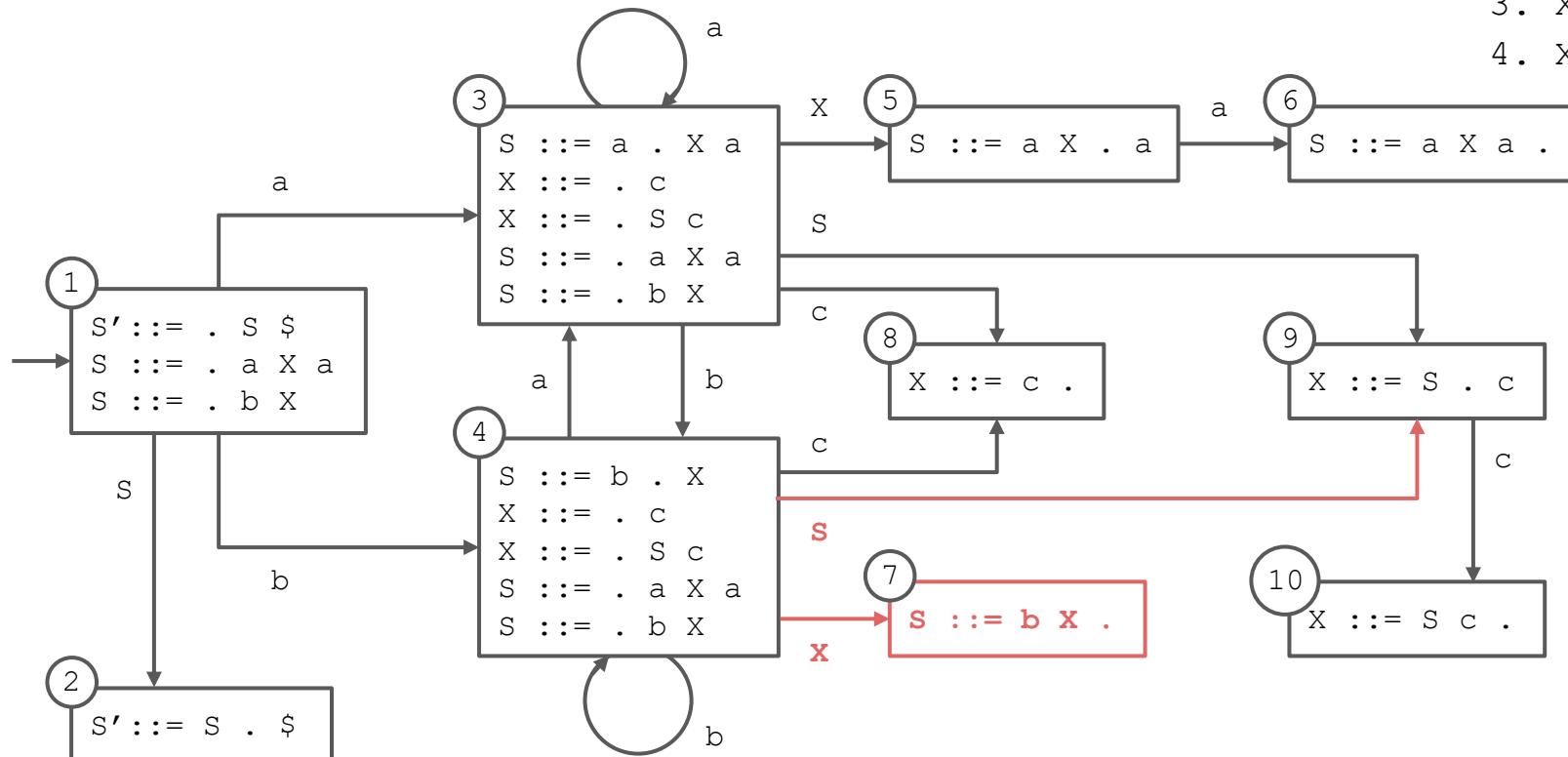
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State Diagram Construction



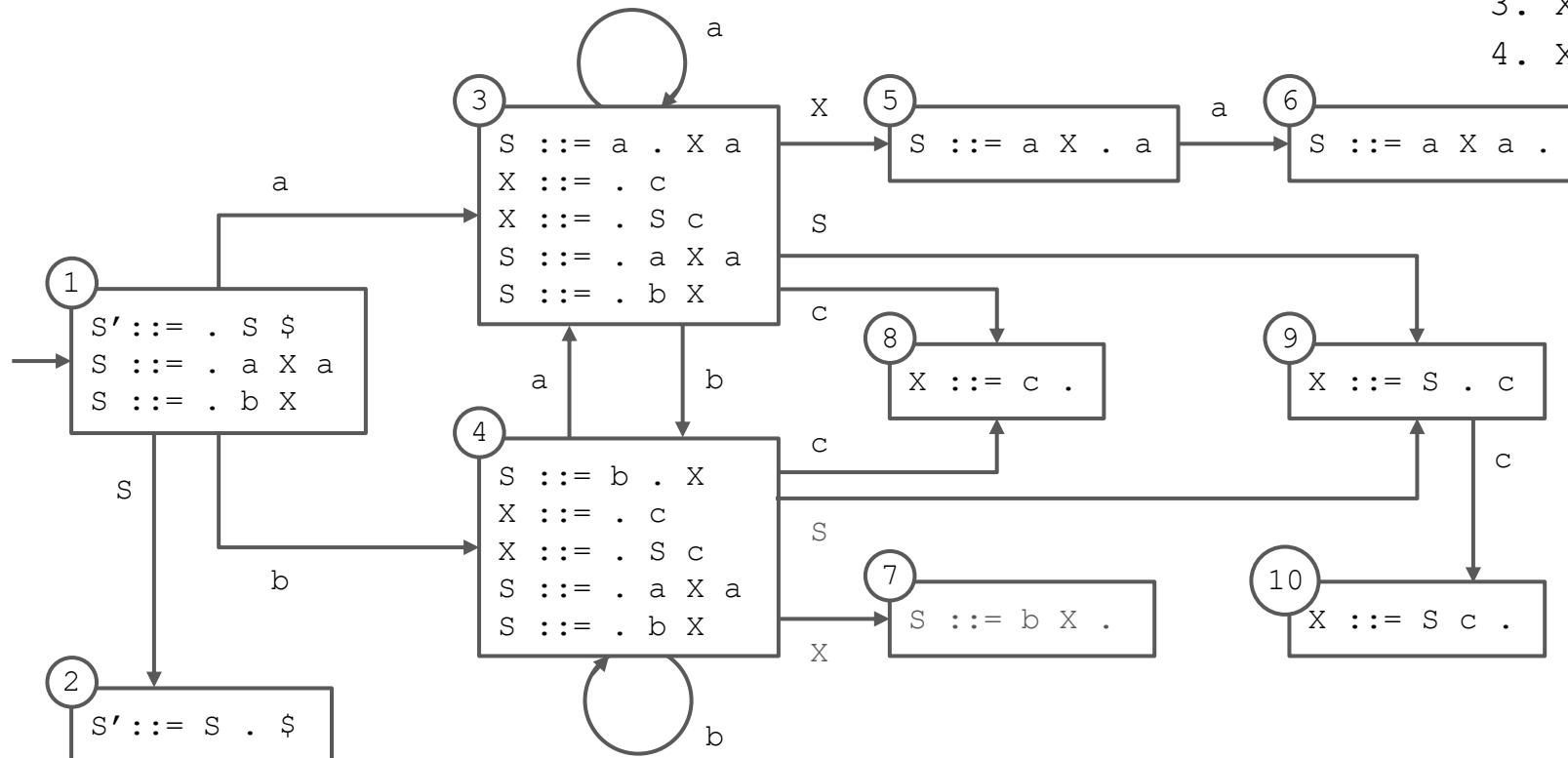
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State Diagram Construction



0. $S' ::= S \$$
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Completed State Diagram



Converted to Table

s# means “shift and enter state #”

- occurs when there is a transition on a terminal

r# means “reduce using production #”

- occurs when a state contains an item with the location at the end of the right-hand side

g# means “go to state #”

- occurs when there is a transition on a nonterminal

acc means “accept”

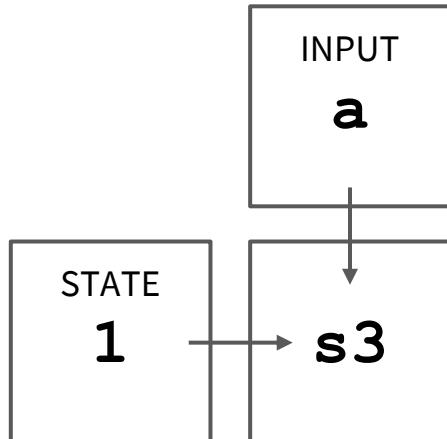
- occurs when the start symbol (S here) has been completed and there is no more input

STATE	ACTION				GOTO	
	a	b	c	\$	S	X
1	s3	s4			g2	
2				acc		
3	s3	s4	s8		g9	g5
4	s3	s4	s8		g9	g7
5	s6					
6	r1	r1	r1	r1		
7	r2	r2	r2	r2		
8	r3	r3	r3	r3		
9				s10		
10	r4	r4	r4	r4		

Parse Trace

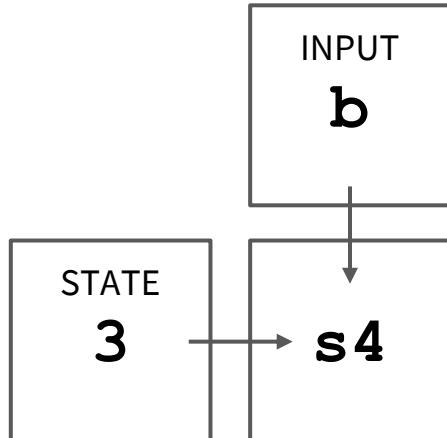
STACK	INPUT	ACTION
\$ 1	a b c c a \$	

Parse Trace



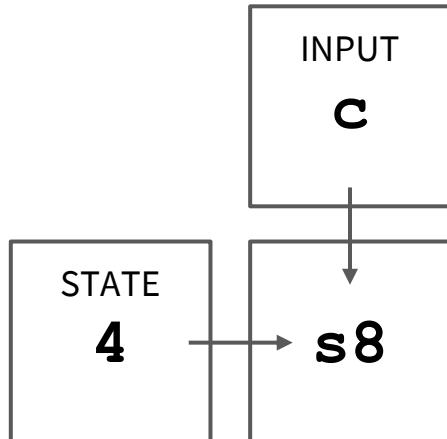
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3	a b c c a \$ b c c a \$	SHIFT

Parse Trace



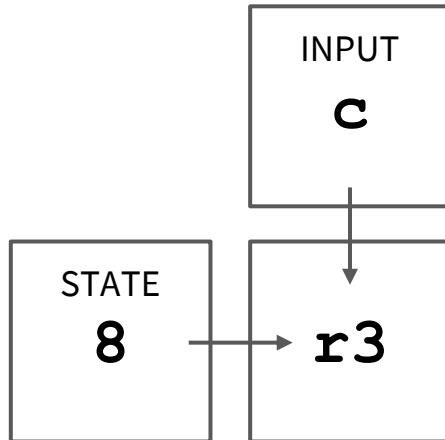
STACK	INPUT	ACTION
\$ 1 \$ 1 a 3 \$ 1 a 3 b 4	a b c c a \$ b c c a \$ c c a \$	SHIFT SHIFT SHIFT

Parse Trace



STACK	INPUT	ACTION
\$ 1 \$ 1 a 3 \$ 1 a 3 b 4 \$ 1 a 3 b 4 c 8	a b c c a \$ b c c a \$ c c a \$ c a \$	SHIFT SHIFT SHIFT

Parse Trace

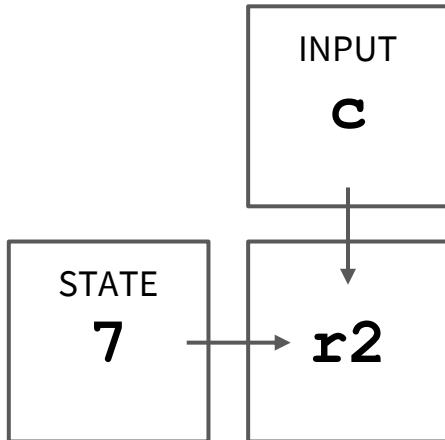


3. $x ::= c$

(and GOTO step:
s4 & x -> g7)

STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	
\$ 1 a 3 b 4 x 7	c a \$	REDUCE

Parse Trace



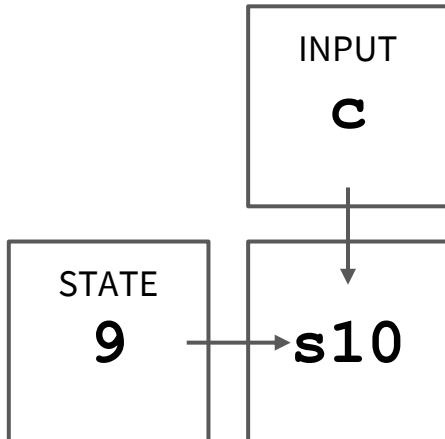
2. $S ::= b \ X$

(and GOTO step:

$s3 \ \& \ S \rightarrow g9$)

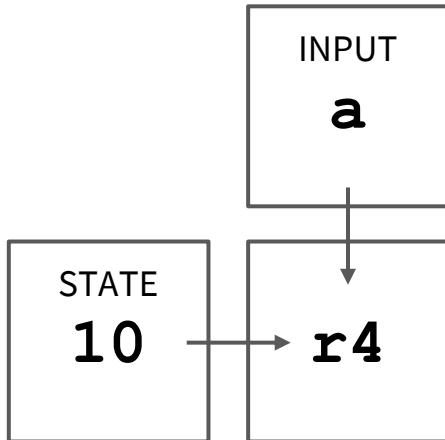
STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	

Parse Trace



STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	SHIFT
\$ 1 a 3 S 9 c 10	a \$	

Parse Trace

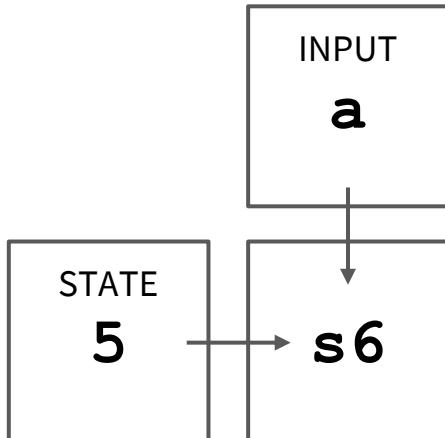


4 . X ::= S C

(and GOTO step:
s3 & X -> g5)

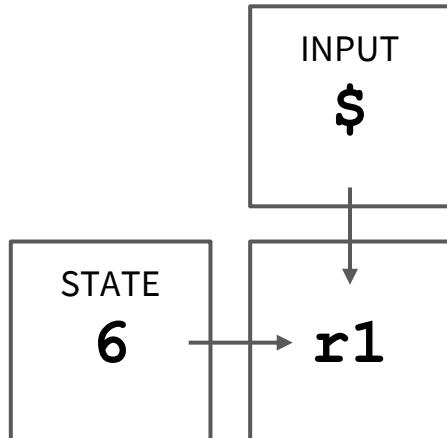
STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	c	REDUCE
\$ 1 a 3 b 4 X 7	a	REDUCE
\$ 1 a 3 S 9	c	SHIFT
\$ 1 a 3 S 9 c 10	a	SHIFT
\$ 1 a 3 X 5	a	REDUCE

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	c	REDUCE
\$ 1 a 3 b 4 X 7	a	REDUCE
\$ 1 a 3 S 9	c	SHIFT
\$ 1 a 3 S 9 c 10	a	REDUCE
\$ 1 a 3 X 5	a	SHIFT
\$ 1 a 3 X 5 a 6	\$	

Parse Trace

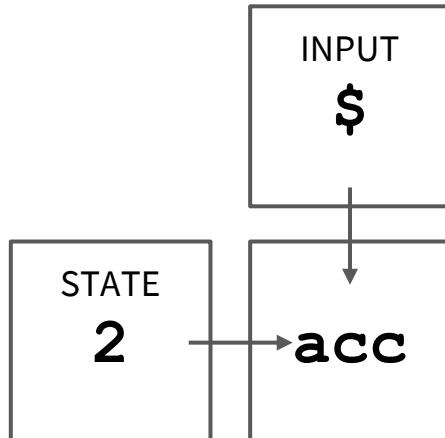


1. $S ::= a X a$

(and GOTO step:
 $s1 \& S \rightarrow g2$)

STACK	INPUT	ACTION
\$ 1	a b c c a \$	SHIFT
\$ 1 a 3	b c c a \$	SHIFT
\$ 1 a 3 b 4	c c a \$	SHIFT
\$ 1 a 3 b 4 c 8	c a \$	REDUCE
\$ 1 a 3 b 4 X 7	c a \$	REDUCE
\$ 1 a 3 S 9	c a \$	SHIFT
\$ 1 a 3 S 9 c 10	a \$	REDUCE
\$ 1 a 3 X 5	a \$	SHIFT
\$ 1 a 3 X 5 a 6	\$	REDUCE
\$ 1 S 2	\$	

Parse Trace



STACK	INPUT	ACTION
\$ 1	a	SHIFT
\$ 1 a 3	b	SHIFT
\$ 1 a 3 b 4	c	SHIFT
\$ 1 a 3 b 4 c 8	c	REDUCE
\$ 1 a 3 b 4 X 7	a	REDUCE
\$ 1 a 3 S 9	c	SHIFT
\$ 1 a 3 S 9 c 10	a	REDUCE
\$ 1 a 3 X 5	a	SHIFT
\$ 1 a 3 X 5 a 6	\$	REDUCE
\$ 1 S 2	\$	ACCEPT