

CSE 401/M501 – Compilers

Parsing & Context-Free Grammars

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Administrivia

- Reminders:
 - **Project partner signup**. Please fill out the form with your and your partner's names and cse netids
 - Only one form per group, please; +1 point each if it's right
 - Please finish by tomorrow, 11 pm
 - Who's still looking for a partner? 401? M501?
 - How do we match people up? Speed dating? Discussion group postings? Random lottery? ???
 - **hw1** due Thur. night (regexps, etc.) via gradescope
 - Use `cse401-staff@cs...` if you need to send email – helps us keep organized. Thanks.

Agenda for Today

- Parsing overview
- Context free grammars
- Ambiguous grammars
- Reading: Cooper & Torczon 3.1-3.2
 - Dragon book is also particularly strong on grammars and languages

Syntactic Analysis / Parsing

- Goal: Convert token stream to an **abstract syntax tree**
- Abstract syntax tree (AST):
 - Captures the structural features of the program
 - Primary data structure for next phases of compilation
- Plan
 - Study how context-free grammars specify syntax
 - Study algorithms for parsing and building ASTs

Context-free Grammars

- The syntax of most programming languages can be specified by a context-free grammar (CGF)
- Compromise between
 - REs: can't nest or specify recursive structure
 - General grammars: too powerful, undecidable
- Context-free grammars are a sweet spot
 - Powerful enough to describe nesting, recursion
 - Easy to parse; but also allow restrictions for speed
- Not perfect
 - Cannot capture semantics, like “must declare every variable” or “must be **int**” – requires later semantic pass
 - Can be ambiguous

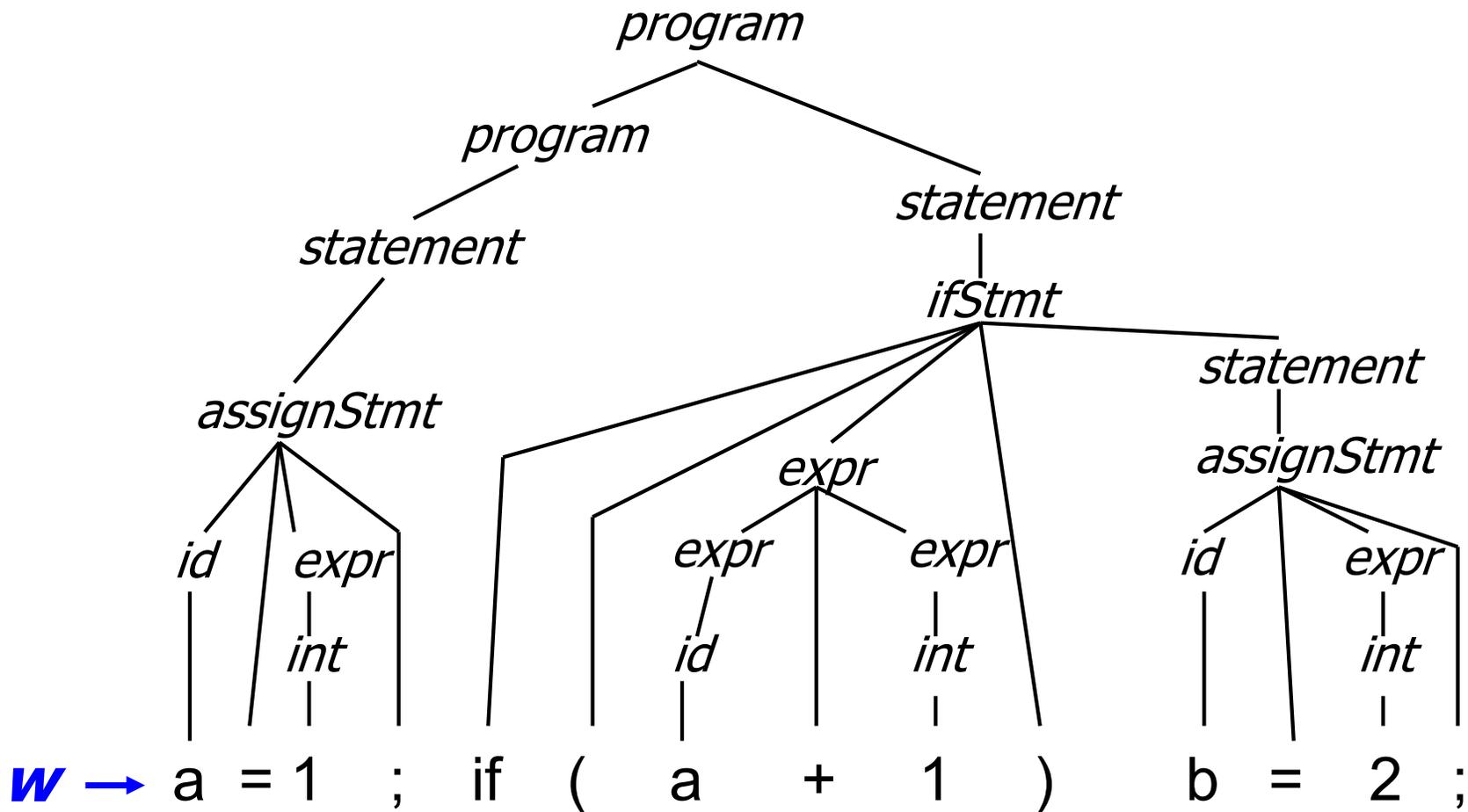
Derivations and Parse Trees

- Derivation: a sequence of expansion steps, beginning with a start symbol and leading to a sequence of terminals
- Parsing: inverse of derivation
 - Given a sequence of terminals (aka tokens) want to recover (discover) the nonterminals and structure, i.e., the parse (concrete syntax) tree

Old Example

G

$program ::= statement \mid program \ statement$
 $statement ::= assignStmt \mid ifStmt$
 $assignStmt ::= id = expr ;$
 $ifStmt ::= if (expr) statement$
 $expr ::= id \mid int \mid expr + expr$
 $id ::= a \mid b \mid c \mid i \mid j \mid k \mid n \mid x \mid y \mid z$
 $int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$



Parsing

- Parsing: Given a grammar G and a sentence w in $L(G)$, traverse the derivation (parse tree) for w in some *standard order* and do *something useful* at each node
 - The tree might not be produced explicitly, but the control flow of the parser will correspond to a traversal

“Standard Order”

- For practical reasons we want the parser to be *deterministic* (no backtracking), and we want to examine the source program from *left to right*.
 - (i.e., parse the program in linear time in the order it appears in the source file)

Common Orderings

- Top-down
 - Start with the root
 - Traverse the parse tree depth-first, left-to-right (leftmost derivation)
 - LL(k), recursive-descent
- Bottom-up
 - Start at leaves and build up to the root
 - Effectively a rightmost derivation in reverse(!)
 - LR(k) and subsets (LALR(k), SLR(k), etc.)

“Something Useful”

- At each point (node) in the traversal, perform some semantic action
 - Construct nodes of full parse tree (rare)
 - Construct abstract syntax tree (AST) (common)
 - Construct linear, lower-level representation (often produced by traversing initial AST in later phases of production compilers)
 - Generate target code on the fly (used in 1-pass compiler; not common in production compilers)
 - Can't generate great code in one pass, but useful if you need a quick 'n dirty working compiler

Context-Free Grammars

- Formally, a *grammar* G is a tuple $\langle N, \Sigma, P, S \rangle$ where
 - N is a finite set of *non-terminal* symbols
 - Σ is a finite set of *terminal* symbols (alphabet)
 - P is a finite set of *productions*
 - A subset of $N \times (N \cup \Sigma)^*$
 - S is the *start symbol*, a distinguished element of N
 - If not specified otherwise, this is usually assumed to be the non-terminal on the left of the first production

Standard Notations

a, b, c elements of Σ

w, x, y, z elements of Σ^*

A, B, C elements of N

X, Y, Z elements of $N \cup \Sigma$

α, β, γ elements of $(N \cup \Sigma)^*$

$A \rightarrow \alpha$ or $A ::= \alpha$ if $\langle A, \alpha \rangle$ in P

Derivation Relations (1)

- $\alpha A \gamma \Rightarrow \alpha \beta \gamma$ iff $A ::= \beta$ in P
 - derives
- $A \Rightarrow^* \alpha$ if there is a chain of productions starting with A that generates α
 - transitive closure

Derivation Relations (2)

- $w A \gamma \Rightarrow_{lm} w \beta \gamma$ iff $A ::= \beta$ in P
 - derives **leftmost**
- $\alpha A w \Rightarrow_{rm} \alpha \beta w$ iff $A ::= \beta$ in P
 - derives **rightmost**
- We will only be interested in leftmost and rightmost derivations – not random orderings

Languages

- For A in N , define $L(A) = \{ w \mid A \Rightarrow^* w \}$
- If S is the start symbol of grammar G , define $L(G) = L(S)$
 - Nonterminal on left of first rule is taken to be the start symbol if one is not specified explicitly

Reduced Grammars

- Grammar G is *reduced* iff for every production $A ::= \alpha$ in G there is a derivation
$$S \Rightarrow^* x A z \Rightarrow x \alpha z \Rightarrow^* xyz$$
 - i.e., no production is useless
- Convention: we will use only reduced grammars
 - There are algorithms for pruning useless productions from grammars – see a formal language or compiler book for details

Ambiguity

- Grammar G is *unambiguous* iff every w in $L(G)$ has a unique leftmost (or rightmost) derivation
 - Fact: unique leftmost or unique rightmost implies the other
- A grammar without this property is *ambiguous*
 - Note that other grammars that generate the same language may be unambiguous
- We need unambiguous grammars for parsing

Example: Ambiguous Grammar for Arithmetic Expressions

$$\text{expr} ::= \text{expr} + \text{expr} \mid \text{expr} - \text{expr} \\ \mid \text{expr} * \text{expr} \mid \text{expr} / \text{expr} \mid \text{int}$$
$$\text{int} ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

- Exercise: show that this is ambiguous
 - How? Show two different leftmost or rightmost derivations for the same string
 - Equivalently: show two different parse trees for the same string

Example (cont)

$$\begin{aligned} \text{expr} &::= \text{expr} + \text{expr} \mid \text{expr} - \text{expr} \\ &\quad \mid \text{expr} * \text{expr} \mid \text{expr} / \text{expr} \mid \text{int} \\ \text{int} &::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \end{aligned}$$

- Give a leftmost derivation of $2+3*4$ and show the parse tree

Example (cont)

$$\begin{aligned} \text{expr} &::= \text{expr} + \text{expr} \mid \text{expr} - \text{expr} \\ &\quad \mid \text{expr} * \text{expr} \mid \text{expr} / \text{expr} \mid \text{int} \\ \text{int} &::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \end{aligned}$$

- Give a different leftmost derivation of $2+3*4$ and show the parse tree

Another example

$expr ::= expr + expr \mid expr - expr$
 $\mid expr * expr \mid expr / expr \mid int$
 $int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$

- Give two different derivations of $5+6+7$

What's going on here?

- The grammar has no notion of precedence or associativity
- Traditional solution
 - Create a non-terminal for each level of precedence
 - Isolate the corresponding part of the grammar
 - Force the parser to recognize higher precedence subexpressions first
 - Use left- or right-recursion for left- or right-associative operators (non-associative operators are not recursive)

Classic Expression Grammar

(first used in ALGOL 60)

$expr ::= expr + term \mid expr - term \mid term$

$term ::= term * factor \mid term / factor \mid factor$

$factor ::= int \mid (expr)$

$int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7$

Check:

Derive $2 + 3 * 4$

$expr ::= expr + term \mid expr - term \mid term$
 $term ::= term * factor \mid term / factor \mid factor$
 $factor ::= int \mid (expr)$
 $int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7$

Check: Derive $5 + 6 + 7$

$expr ::= expr + term \mid expr - term \mid term$
 $term ::= term * factor \mid term / factor \mid factor$
 $factor ::= int \mid (expr)$
 $int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7$

- Note interaction between left- vs right-recursive rules and resulting associativity

Check:

Derive $5 + (6 + 7)$

$expr ::= expr + term \mid expr - term \mid term$
 $term ::= term * factor \mid term / factor \mid factor$
 $factor ::= int \mid (expr)$
 $int ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7$

Another Classic Example

- Grammar for conditional statements

$stmt ::= \text{if } (cond) stmt$

$\quad | \text{if } (cond) stmt \text{ else } stmt$

- Exercise: show that this is ambiguous
 - How?

$stmt ::= \text{if } (cond) stmt$
 $\quad | \text{if } (cond) stmt \text{ else } stmt$

One Derivation

$\text{if } (cond) \text{ if } (cond) stmt \text{ else } stmt$

$stmt ::= \text{if } (cond) stmt$
 $\quad | \text{if } (cond) stmt \text{ else } stmt$

Another Derivation

$\text{if } (cond) \text{ if } (cond) stmt \text{ else } stmt$

Solving “if” Ambiguity

- Fix the grammar to separate if statements with else clause and if statements with no else
 - Done in Java reference grammar
 - Adds lots of non-terminals
- or, Change the language
 - But it’d better be ok to do this
- or, Use some ad-hoc rule in the parser
 - “else matches closest unpaired if”

Resolving Ambiguity with Grammar (1)

Stmt ::= MatchedStmt | UnmatchedStmt

MatchedStmt ::= ... |

if (Expr) MatchedStmt **else** MatchedStmt

UnmatchedStmt ::= ... |

if (Expr) Stmt |

if (Expr) MatchedStmt **else** UnmatchedStmt

- formal, no additional rules beyond syntax
- can be more obscure than original grammar

Check

```
Stmt ::= MatchedStmt | UnmatchedStmt
MatchedStmt ::= ... |
    if ( Expr ) MatchedStmt else MatchedStmt
UnmatchedStmt ::= if ( Expr ) Stmt |
    if ( Expr ) MatchedStmt else UnmatchedStmt
```

if (cond) if (cond) stmt else stmt

Resolving Ambiguity with Grammar (2)

- If you can (re-)design the language, just avoid the problem entirely

```
Stmt ::= ... |  
       if Expr then Stmt end |  
       if Expr then Stmt else Stmt end
```

- formal, clear, elegant
- allows sequence of Stmts in then and else branches, no { , } needed
- extra end required for every if
(But maybe this is a good idea anyway?)

Parser Tools and Operators

- Most parser tools can cope with ambiguous grammars
 - Makes life simpler if used with discipline
- Usually can specify precedence & associativity
 - Allows simpler, ambiguous grammar with fewer nonterminals as basis for parser – let the tool handle the details (but only when it makes sense)
 - (i.e., $\text{expr} ::= \text{expr} + \text{expr} \mid \text{expr} * \text{expr} \mid \dots$ with assoc. & precedence declarations can be the best solution)
- Take advantage of this to simplify the grammar when using parser-generator tools

Parser Tools and Ambiguous Grammars

- Possible rules for resolving other problems
 - Earlier productions in the grammar preferred to later ones (danger here if parser input changes)
 - Longest match used if there is a choice (good solution for dangling if and similar things)
- Parser tools normally allow for this
 - But be sure that what the tool does is really what you want
 - And that it's part of the tool spec, so that v2 won't do something different (that you *don't* want!)

Coming Attractions

- Next topic: LR parsing
 - Continue reading ch. 3