

# CSE 390Z: Mathematics for Computation Workshop

## Week 6 Workshop Solutions

### 0. Induction: Warm-Up

Prove by induction that  $5 \mid (6^n - 1)$  for all  $n \in \mathbb{N}$ .

#### Solution:

Let  $P(n)$  be " $5 \mid 6^n - 1$ " for all  $n \in \mathbb{N}$ . We will show  $P(n)$  for all  $n$  by induction.

**Base Case ( $n = 0$ ).**  $6^0 - 1 = 1 - 1 = 0 = 0 \cdot 5$ , so  $5 \mid 6^0 - 1$ .

**Induction Hypothesis.** Assume  $P(k)$  holds for some arbitrary integer  $k \geq 0$ .

**Induction Step.** Goal: We aim to show  $P(k + 1)$ , i.e. that  $5 \mid (6^{k+1} - 1)$ .

By the Inductive Hypothesis, we have that  $5 \mid (6^k - 1)$ . Then by definition of divides,  $6^k - 1 = 5j$  for some  $j \in \mathbb{Z}$ . Observe:

$$\begin{aligned}6^k - 1 &= 5j \\6^{k+1} - 6 &= 30j \\6^{k+1} - 1 &= 30j + 5 \\6^{k+1} - 1 &= 5(6j + 1)\end{aligned}$$

Since  $j$  is an integer,  $j + 1$  is an integer. They by definition of divides, we have that  $5 \mid (6^{k+1} - 1)$ , as desired. So  $P(k + 1)$  holds.

**Conclusion.**  $P(n)$  is true for all  $n \in \mathbb{N}$  by induction.

### 1. Induction: Equality

Prove by induction that for every  $n \in \mathbb{N}$ , the following equality is true:

$$0 \cdot 2^0 + 1 \cdot 2^1 + 2 \cdot 2^2 + \dots + n \cdot 2^n = (n - 1)2^{n+1} + 2.$$

#### Solution:

Let  $P(n)$  be " $0 \cdot 2^0 + 1 \cdot 2^1 + 2 \cdot 2^2 + \dots + n \cdot 2^n = (n - 1)2^{n+1} + 2$ ". We will prove  $P(n)$  for all  $n \in \mathbb{N}$  by induction.

**Base Case.**  $0 \cdot 2^0 = 0$  and  $(0 - 1)2^{0+1} + 2 = -2 + 2 = 0$ , so  $P(0)$  is true.

**Induction Hypothesis.** Assume that  $P(k)$  holds true for some arbitrary integer  $k \geq 0$ .

**Induction Step** We show  $P(k + 1)$ :

$$\begin{aligned}0 \cdot 2^0 + 1 \cdot 2^1 + 2 \cdot 2^2 + \dots + (k + 1) \cdot 2^{k+1} & \\= 0 \cdot 2^0 + 1 \cdot 2^1 + 2 \cdot 2^2 + \dots + k \cdot 2^k + (k + 1) \cdot 2^{k+1} & \quad \text{[Show another term inside "..."]} \\= (k - 1)2^{k+1} + 2 + (k + 1)2^{k+1} & \quad \text{[Induction Hypothesis]} \\= ((k - 1) + (k + 1))2^{k+1} + 2 & \quad \text{[Group multiples of } 2^{k+1}\text{]} \\= (2k)2^{k+1} + 2 & \quad \text{[Algebra]} \\= k2^{k+2} + 2 & \quad \text{[Algebra].}\end{aligned}$$

Therefore  $P(k + 1)$  holds.

**Conclusion.**  $P(n)$  holds for all  $n \in \mathbb{N}$  by induction.

## 2. Induction: Inequality

Prove by induction on  $n$  that for all integers  $n \geq 0$  the inequality  $(3 + \pi)^n \geq 3^n + n\pi 3^{n-1}$  is true.

### Solution:

Let  $P(n)$  be " $(3 + \pi)^n \geq 3^n + n\pi 3^{n-1}$ ". We will prove  $P(n)$  is true for all  $n \in \mathbb{N}$ , by induction.

**Base Case:** ( $n = 0$ ):  $(3 + \pi)^0 = 1$  and  $3^0 + 0 \cdot \pi \cdot 3^{-1} = 1$ , since  $1 \geq 1$ ,  $P(0)$  is true.

**Inductive Hypothesis:** Suppose that  $P(k)$  is true for some arbitrary integer  $k \in \mathbb{N}$ .

**Inductive Step:**

Goal: Show  $P(k+1)$ , i.e. show  $(3 + \pi)^{k+1} \geq 3^{k+1} + (k+1)\pi 3^{(k+1)-1} = 3^{k+1} + (k+1)\pi 3^k$

$$\begin{aligned} (3 + \pi)^{k+1} &= (3 + \pi)^k \cdot (3 + \pi) && \text{(Factor out } (3 + \pi)) \\ &\geq (3^k + k3^{k-1}\pi) \cdot (3 + \pi) && \text{(By I.H., } (3 + \pi) \geq 0) \\ &= 3 \cdot 3^k + 3^k\pi + 3k3^{k-1}\pi + k3^{k-1}\pi^2 && \text{(Distributive property)} \\ &= 3^{k+1} + 3^k\pi + k3^k\pi + k3^{k-1}\pi^2 && \text{(Simplify)} \\ &= 3^{k+1} + (k+1)3^k\pi + k3^{k-1}\pi^2 && \text{(Factor out } (k+1)) \\ &\geq 3^{k+1} + (k+1)\pi 3^k && (k3^{k-1}\pi^2 \geq 0) \end{aligned}$$

**Conclusion:** So by induction,  $P(n)$  is true for all  $n \in \mathbb{N}$ .

### 3. Inductively Odd

An 123 student learning recursion wrote a recursive Java method to determine if a number is odd or not, and needs your help proving that it is correct.

```
public static boolean oddr(int n) {
    if (n == 0)
        return False;
    else
        return !oddr(n-1);
}
```

Help the student by writing an inductive proof to prove that for all integers  $n \geq 0$ , the method `oddr` returns True if  $n$  is an odd number, and False if  $n$  is not an odd number (i.e.  $n$  is even). You may recall the definitions  $\text{Odd}(n) := \exists x \in \mathbb{Z}(n = 2x + 1)$  and  $\text{Even}(n) := \exists x \in \mathbb{Z}(n = 2x)$ ;  $!True = False$  and  $!False = True$ .

#### Solution:

Let  $P(n)$  be "oddr( $n$ ) returns True if  $n$  is odd, or False if  $n$  is even". We will show that  $P(n)$  is true for all integers  $n \geq 0$  by induction on  $n$ .

#### Base Case: ( $n = 0$ )

0 is even, so  $P(0)$  is true if `oddr(0)` returns False, which is exactly the base case of `oddr`, so  $P(0)$  is true.

**Inductive Hypothesis:** Suppose  $P(k)$  is true for an arbitrary integer  $k \geq 0$ .

#### Inductive Step:

- **Case 1:**  $k + 1$  is even.

If  $k + 1$  is even, then there is an integer  $x$  s.t.  $k + 1 = 2x$ , so then  $k = 2x - 1 = 2(x - 1) + 1$ , so therefore  $k$  is odd. We know that since  $k + 1 > 0$ , `oddr(k+1)` should return `!oddr(k)`. By the Inductive Hypothesis, we know that since  $k$  is odd, `oddr(k)` returns True, so `oddr(k+1)` returns `!oddr(k) = False`, and  $k + 1$  is even, therefore  $P(k+1)$  is true.

- **Case 2:**  $k + 1$  is odd.

If  $k + 1$  is odd, then there is an integer  $x$  s.t.  $k + 1 = 2x + 1$ , so then  $k = 2x$  and therefore  $k$  is even. We know that since  $k + 1 > 0$ , `oddr(k+1)` should return `!oddr(k)`. By the Inductive Hypothesis, we know that since  $k$  is even, `oddr(k)` returns False, so `oddr(k+1)` returns `!oddr(k) = True`, and  $k + 1$  is odd, therefore  $P(k+1)$  is true.

Then  $P(k + 1)$  is true for all cases. Thus, we have shown  $P(n)$  is true for all integers  $n \geq 0$  by induction.

## 4. Strong Induction: Stamp Collection

A store sells 3 cent and 5 cent stamps. Use strong induction to prove that you can make exactly  $n$  cents worth of stamps for all  $n \geq 10$ .

**Hint:** you'll need multiple base cases for this - think about how many steps back you need to go for your inductive step.

### Solution:

Let  $P(n)$  be defined as "You can buy exactly  $n$  cents of stamps". We will prove  $P(n)$  is true for all integers  $n \geq 10$  by strong induction.

**Base Cases:** ( $n = 10, 11, 12$ ):

- $n = 10$ : 10 cents of stamps can be made from two 5 cent stamps.
- $n = 11$ : 11 cents of stamps can be made from one 5 cent and two 3 cent stamps.
- $n = 12$ : 12 cents of stamps can be made from four 3 cent stamps.

**Inductive Hypothesis:** Suppose for some arbitrary integer  $k \geq 12$ ,  $P(10) \wedge P(11) \wedge \dots \wedge P(k)$  holds.

**Inductive Step:**

**Goal:** Show  $P(k+1)$ , i.e. show that we can make  $k+1$  cents in stamps.

We want to buy  $k+1$  cents in stamps. By the I.H., we can buy exactly  $(k+1) - 3 = k - 2$  cents in stamps. Then, we can add another 3 cent stamp in order to buy  $k+1$  cents in stamps, so  $P(k+1)$  is true.

**Note:** How did we decide how many base cases to have? Well, we wanted to be able to assume  $P(k-2)$ , and add 3 to achieve  $P(k+1)$ . Therefore we needed to be able to assume that  $k-2 \geq 10$ . Adding 2 to both sides, we needed to be able to assume that  $k \geq 12$ . So, we have to prove the base cases up to 12, that is: 10, 11, 12.

Another way to think about this is that we had to use a fact from 3 steps back from  $k+1$  to  $k-2$  in the IS, so we needed 3 base cases.

**Conclusion:** So by strong induction,  $P(n)$  is true for all integers  $n \geq 10$ .

## 5. Strong Induction: Recursively Defined Functions

Consider the function  $f(n)$  defined for integers  $n \geq 1$  as follows:

$$f(1) = 1 \text{ for } n = 1$$

$$f(2) = 4 \text{ for } n = 2$$

$$f(3) = 9 \text{ for } n = 3$$

$$f(n) = f(n-1) - f(n-2) + f(n-3) + 2(2n-3) \text{ for } n \geq 4$$

Prove by strong induction that for all  $n \geq 1$ ,  $f(n) = n^2$ .

**Complete the induction proof below.**

### Solution:

1 Let  $P(n)$  be defined as " $f(n) = n^2$ ". We will prove  $P(n)$  is true for all integers  $n \geq 1$  by strong induction.

2 **Base Cases** ( $n = 1, 2, 3$ ):

- $n = 1$ :  $f(1) = 1 = 1^2$ .

- $n = 2: f(2) = 4 = 2^2.$
- $n = 3: f(3) = 9 = 3^2$

So the base cases hold.

3 **Inductive Hypothesis:** Suppose for some arbitrary integer  $k \geq 3$ ,  $P(j)$  is true for  $1 \leq j \leq k$ .

4 **Inductive Step:**

**Goal:** Show  $P(k + 1)$ , i.e. show that  $f(k + 1) = (k + 1)^2$ .

$$\begin{aligned}
 f(k + 1) &= f(k + 1 - 1) - f(k + 1 - 2) + f(k + 1 - 3) + 2(2(k + 1) - 3) && \text{Definition of } f \\
 &= f(k) - f(k - 1) + f(k - 2) + 2(2k - 1) \\
 &= k^2 - (k - 1)^2 + (k - 2)^2 + 2(2k - 1) && \text{By IH} \\
 &= k^2 - (k^2 - 2k + 1) + (k^2 - 4k + 4) + 4k - 2 \\
 &= (k^2 - k^2 + k^2) + (2k - 4k + 4k) + (-1 + 4 - 2) \\
 &= k^2 + 2k + 1 \\
 &= (k + 1)^2
 \end{aligned}$$

So  $P(k + 1)$  holds.

5 **Conclusion:** So by strong induction,  $P(n)$  is true for all integers  $n \geq 1$ .

## 6. Strong Induction: A Variation of the Stamp Problem

A store sells candy in packs of 4 and packs of 7. Let  $P(n)$  be defined as "You are able to buy  $n$  packs of candy". For example,  $P(3)$  is not true, because you cannot buy exactly 3 packs of candy from the store. However, it turns out that  $P(n)$  is true for any  $n \geq 18$ . Use strong induction on  $n$  to prove this.

**Hint:** you'll need multiple base cases for this - think about how many steps back you need to go for your inductive step.

### Solution:

Let  $P(n)$  be defined as "You are able to buy  $n$  packs of candy". We will prove  $P(n)$  is true for all integers  $n \geq 18$  by strong induction.

**Base Cases:** ( $n = 18, 19, 20, 21$ ):

- $n = 18$ : 18 packs of candy can be made up of 2 packs of 7 and 1 pack of 4 ( $18 = 2 * 7 + 1 * 4$ ).
- $n = 19$ : 19 packs of candy can be made up of 1 pack of 7 and 3 packs of 4 ( $19 = 1 * 7 + 3 * 4$ ).
- $n = 20$ : 20 packs of candy can be made up of 5 packs of 4 ( $20 = 5 * 4$ ).
- $n = 21$ : 21 packs of candy can be made up of 3 packs of 7 ( $21 = 3 * 7$ ).

**Inductive Hypothesis:** Suppose for some arbitrary integer  $k \geq 21$ ,  $P(18) \wedge \dots \wedge P(k)$  hold.

**Inductive Step:**

**Goal:** Show  $P(k + 1)$ , i.e. show that we can buy  $k + 1$  packs of candy.

We want to buy  $k + 1$  packs of candy. By the I.H., we can buy exactly  $k - 3$  packs, so we can add another pack of 4 packs in order to buy  $k + 1$  packs of candy, so  $P(k + 1)$  is true.

**Note:** How did we decide how many base cases to have? Well, we wanted to be able to assume  $P(k - 3)$ , and add 4 to achieve  $P(k + 1)$ . Therefore we needed to be able to assume that  $k - 3 \geq 18$ . Adding 3 to both sides, we needed to be able to assume that  $k \geq 21$ . So, we have to prove the base cases up to 21, that is: 18, 19, 20, 21.

Another way to think about this is that we had to use a fact from 4 steps back from  $k + 1$  to  $k - 3$  in the IS, so we needed 4 base cases.

**Conclusion:** So by strong induction,  $P(n)$  is true for all integers  $n \geq 18$ .