CSE 390Z: Mathematics for Computation Workshop

Practice 311 Midterm Solutions

Name: _____

UW ID: _____

Instructions:

- This is a **simulated practice midterm**. You will **not** be graded on your performance on this exam.
- Nevertheless, please treat this as if it is a real exam. That means that you may not discuss with your neighbors, reference outside material, or use your devices during the next 60 minute period.
- If you get stuck on a problem, consider moving on and coming back later. In the actual exam, there will likely be opportunity for partial credit.
- There are 5 problems on this exam, totaling 80 points.

1. Predicate Translation [15 points]

Let the domain of discourse be novels, comic books, movies, and TV shows. Translate the following statements to predicate logic, using the following predicates:

Novel(x) := x is a novel Comic(x) := x is a comic book Movie(x) := x is a movie Show(x) := x is a TV show Adaptation(x, y) := x is an adaptation of y

(a) (5 points) A novel cannot be adapted into both a movie and a TV show.

Solution:

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\forall x (\mathsf{Novel}(x) \rightarrow \forall m \forall s ((\mathsf{Movie}(m) \land \mathsf{Show}(s)) \rightarrow \neg (\mathsf{Adaptation}(m, x) \land \mathsf{Adaptation}(s, x)))
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(b) (5 points) Every movie is an adaptation of a novel or a comic book.

Solution:

 $\forall m(\mathsf{Movie}(m) \to \exists x(\mathsf{Adaptation}(m, x) \land (\mathsf{Novel}(x) \lor \mathsf{Comic}(x))))$

(c) (5 points) Every novel has been adapted into exactly one movie.

Solution:

$$\begin{split} &\forall x (\mathsf{Novel}(x) \to \exists m (\mathsf{Movie}(m) \land \mathsf{Adaptation}(m, x) \land \forall n ((\mathsf{Movie}(n) \land (n \neq m)) \to \neg \mathsf{Adaptation}(n, x)))) \\ &\mathsf{OR} \\ &\forall x (\mathsf{Novel}(x) \to \exists m (\mathsf{Movie}(m) \land \mathsf{Adaptation}(m, x) \land \forall n (\mathsf{Adaptation}(n, x) \to (\neg \mathsf{Movie}(n) \lor n = m)))) \\ &\mathsf{OR} \\ &\forall x (\mathsf{Novel}(x) \to \exists m (\mathsf{Movie}(m) \land \mathsf{Adaptation}(m, x) \land \forall n ((\mathsf{Adaptation}(n, x) \land \mathsf{Movie}(n)) \to (n = m)))) \end{split}$$

*Note that a great exercise is to show that the above 3 solutions are all logically equivalent :)

2. Boolean Algebra [15 points]

Let f be the boolean function defined as $f(x,y,z)=(x+y)^\prime+(zy)$

(a) (5 points) Fill in the following table with the values of f(x, y, z) in the last column. Feel free to use the blank columns while doing your work.

x	y	z		f(x,y,z)
0	0	0		
0	0	1		
0	1	0		
0	1	1		
1	0	0		
1	0	1		
1	1	0		
1	1	1		

Solution:

x	y	z	(x+y)	(x+y)'	(zy)	f(x,y,z)
0	0	0	0	1	0	1
0	0	1	0	1	0	1
0	1	0	1	0	0	0
0	1	1	1	0	1	1
1	0	0	1	0	0	0
1	0	1	1	0	0	0
1	1	0	1	0	0	0
1	1	1	1	0	1	1

(b) (5 points) Write f(x, y, z) as a Boolean algebra expression in the canonical sum-of-products form in terms of the three input bits x, y, and z.

Solution:

f(x,y,z) = x'y'z' + x'y'z + x'yz + xyz

(c) (5 points) Write f(x, y, z) as a Boolean algebra expression in the canonical product-of-sums form in terms of the three input bits x, y, and z.

Solution:

$$f(x, y, z) = (x + y' + z)(x' + y + z)(x' + y + z')(x' + y' + z)$$

3. Formal Proof with Divides [10 points]

Prove the following statement using a formal proof:

For all integers a, b, c, if $a^2 \mid b$ and $b^3 \mid c$, then $a^6 \mid c$.

Solution:

1. Let x, y, z be arbitrary integers.

$2.1 \hspace{.1in} x^2 \mid y \wedge y^3 \mid z$	Assumption
2.2 $x^2 \mid y$	Elim ∧: 2.1
2.3 $y^3 \mid z$	Elim ∧: 2.2
2.4 $\exists k(y = x^2k)$	Definition of divides: 2.2
2.5 $\exists k(z=y^3k)$	Definition of divides: 2.3
2.6 $y = x^2 s$	Elim ∃: 2.4
2.7 $z = y^3 t$	Elim ∃: 2.5
2.8 $z = (x^2 s)^3 t = x^6 (s^3 t)$	Algebra
2.9 $\exists k(z=x^6k)$	Intro ∃: 2.8
2.10 $x^6 \mid z$	Definition of divides: 2.9
2. $x^2 \mid y \land y^3 \mid z \to x^6 \mid z$	Direct Proof: 2.1 - 2.10
3. $\forall a \forall b \forall c((a^2 \mid b \land b^3 \mid c) \rightarrow a^6 \mid c$	Intro \forall

4. English Proof with Mod [20 points]

Write an English proof to show that for all integers n, $n^2 \equiv_4 0$ or $n^2 \equiv_4 1$.

Hint: You'll need to break this proof into cases based on the following definitions: $Even(x):= \exists k(x = 2k)$ $Odd(x):= \exists k(x = 2k + 1)$

Solution:

Let n be an arbitrary integer.

Case 1: n is even. Then n = 2k for some integer k. Then $n^2 = (2k)^2 = 4k^2$. Since k is an integer, k^2 is an integer. So n^2 is 4 times an integer. Then by definition of divides, $4 \mid n^2 - 0$. Then by definition of congruence, $n^2 \equiv_4 0$.

Case 2: n is odd. Then n = 2k + 1 for some integer k. Then $n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$. So $n^2 - 1 = 4(k^2 + k)$. Since k is an integer, $k^2 + k$ is an integer. So $n^2 - 1$ is 4 times an integer. Then by definition of divides, $4 \mid n^2 - 1$. Then by definition of congruence, $n^2 \equiv_4 1$.

Thus in all cases, $n^2 \equiv_4 0$ or $n^2 \equiv_4 1$. Since n was arbitrary, the claim holds.

5. Induction [20 points] Prove by induction that $(1 + \pi)^n > 1 + n\pi$ for all integers $n \ge 2$.

Solution:

1. Let P(n) be the statement " $(1 + \pi)^n > 1 + n\pi$ ". We prove P(n) for all integers $n \ge 2$ by induction.

2. Base Case: When n = 2, the LHS is $(1 + \pi)^2 = 1 + 2\pi + \pi^2$. The RHS is $1 + 2\pi$. Since $\pi^2 > 0$, $1 + 2\pi + \pi^2 > 1 + 2\pi$, so the Base Case holds.

3. Inductive Hypothesis: Suppose that P(k) holds for some arbitrary integer $k \ge 2$. Then $(1 + \pi)^k > 1 + k\pi$.

4. Inductive Step:

Goal: Show P(k+1), i.e. show $(1+\pi)^{k+1} > 1 + (k+1)\pi$

 $\begin{array}{ll} (1+\pi)^{k+1} = (1+\pi)(1+\pi)^k & \mbox{Definition of Exponent} \\ > (1+\pi)(1+k\pi) & \mbox{By IH} \\ = 1+\pi+k\pi+k\pi^2 & \mbox{Algebra} \\ = 1+(k+1)\pi+k\pi^2 & \mbox{Algebra} \\ > 1+(k+1)\pi & \mbox{Since } k\pi^2 > 0 \end{array}$

Thus $(1 + \pi)^{k+1} > 1 + (k+1)\pi$. So P(k+1) holds.

5. Thus we have proven P(n) for all integers $n \ge 2$ by induction.