CSE 390Z: Mathematics for Computation Workshop

Week 6 Workshop Solutions

0. Induction: Warm-Up

Prove by induction that $5 \mid (6^n - 1)$ for all $n \in \mathbb{N}$. Solution:

Let P(n) be "5 | $6^n - 1$ " for all $n \in \mathbb{N}$. We will show P(n) for all n by induction.

Base Case (n = 0**).** $6^0 - 1 = 1 - 1 = 0 = 0 \cdot 5$, so $5 \mid 6^0 - 1$.

Induction Hypothesis. Assume P(k) holds for some arbitrary integer $k \ge 0$.

Induction Step. Goal: We aim to show P(k+1), i.e. that $5 \mid (6^{k+1}-1)$.

By the Inductive Hypothesis, we have that $5 \mid (6^k - 1)$. Then by definition of divides, $6^k - 1 = 5j$ for some $j \in \mathbb{Z}$. Observe:

$$6^{k} - 1 = 5j$$

$$6^{k+1} - 6 = 30j$$

$$6^{k+1} - 1 = 30j + 5$$

$$6^{k+1} - 1 = 5(6j + 1)$$

Since j is an integer, j + 1 is an integer. They by definition of divides, we have that $5 | (6^{k+1} - 1)$, as desired. So P(k + 1) holds.

Conclusion. P(n) is true for all $n \in \mathbb{N}$ by induction.

1. Induction: Equality

Prove by induction that for every $n \in \mathbb{N}$, the following equality is true:

$$0 \cdot 2^{0} + 1 \cdot 2^{1} + 2 \cdot 2^{2} + \dots + n \cdot 2^{n} = (n-1)2^{n+1} + 2.$$

Solution:

Let P(n) be " $0 \cdot 2^0 + 1 \cdot 2^1 + 2 \cdot 2^2 + \cdots + n \cdot 2^n = (n-1)2^{n+1} + 2$ ". We will prove P(n) for all $n \in \mathbb{N}$ by induction.

Base Case. $0 \cdot 2^0 = 0$ and $(0 - 1)2^{0+1} + 2 = -2 + 2 = 0$, so P(0) is true.

Induction Hypothesis. Assume that P(k) holds true for some arbitrary integer $k \ge 0$.

Induction Step We show P(k+1):

$$\begin{array}{ll} 0 \cdot 2^{0} + 1 \cdot 2^{1} + 2 \cdot 2^{2} + \dots + (k+1) \cdot 2^{k+1} \\ &= 0 \cdot 2^{0} + 1 \cdot 2^{1} + 2 \cdot 2^{2} + \dots + k \cdot 2^{k} + (k+1) \cdot 2^{k+1} \\ &= (k-1)2^{k+1} + 2 + (k+1)2^{k+1} \\ &= ((k-1) + (k+1))2^{k+1} + 2 \\ &= (2k)2^{k+1} + 2 \\ &= k2^{k+2} + 2 \end{array}$$
[Show another term inside "..."]
[Induction Hypothesis]
[Group multiples of 2^{k+1}]
[Algebra]
[Algebra].

Therefore P(k+1) holds.

Conclusion. P(n) holds for all $n \in \mathbb{N}$ by induction.

2. Induction: Inequality

Prove by induction on n that for all integers $n \ge 0$ the inequality $(3 + \pi)^n \ge 3^n + n\pi 3^{n-1}$ is true. Solution:

Let P(n) be " $(3 + \pi)^n \ge 3^n + n\pi 3^{n-1}$ ". We will prove P(n) is true for all $n \in \mathbb{N}$, by induction.

Base Case: (n = 0): $(3 + \pi)^0 = 1$ and $3^0 + 0 \cdot \pi \cdot 3^{-1} = 1$, since $1 \ge 1$, P(0) is true.

Inductive Hypothesis: Suppose that P(k) is true for some arbitrary integer $k \in \mathbb{N}$.

Inductive Step:

Goal: Show P(k+1), i.e. show $(3+\pi)^{k+1} \ge 3^{k+1} + (k+1)\pi 3^{(k+1)-1} = 3^{k+1} + (k+1)\pi 3^k$

$$\begin{aligned} (3+\pi)^{k+1} &= (3+\pi)^k \cdot (3+\pi) & (Factor out (3+\pi)) \\ &\geq (3^k + k3^{k-1}\pi) \cdot (3+\pi) & (By I.H., (3+\pi) \ge 0) \\ &= 3 \cdot 3^k + 3^k \pi + 3k3^{k-1}\pi + k3^{k-1}\pi^2 & (Distributive property) \\ &= 3^{k+1} + 3^k \pi + k3^k \pi + k3^{k-1}\pi^2 & (Simplify) \\ &= 3^{k+1} + (k+1)3^k \pi + k3^{k-1}\pi^2 & (Factor out (k+1)) \\ &\geq 3^{k+1} + (k+1)\pi 3^k & (k3^{k-1}\pi^2 \ge 0) \end{aligned}$$

Conclusion: So by induction, P(n) is true for all $n \in \mathbb{N}$.

3. Inductively Odd

An 123 student learning recursion wrote a recursive Java method to determine if a number is odd or not, and needs your help proving that it is correct.

```
public static boolean oddr(int n) {
    if (n == 0)
        return False;
    else
        return !oddr(n-1);
}
```

Help the student by writing an inductive proof to prove that for all integers $n \ge 0$, the method oddr returns True if n is an odd number, and False if n is not an odd number (i.e. n is even). You may recall the definitions $Odd(n) := \exists x \in \mathbb{Z}(n = 2x + 1)$ and $Even(n) := \exists x \in \mathbb{Z}(n = 2x)$; !True = False and !False = True.

Solution:

Let P(n) be "oddr(n) returns True if n is odd, or False if n is even". We will show that P(n) is true for all integers $n \ge 0$ by induction on n.

Base Case: $(n = \underline{0})$

0 is even, so P(0) is true if oddr(0) returns False, which is exactly the base case of oddr, so P(0) is true. Inductive Hypothesis: Suppose P(k) is true for an arbitrary integer $k \ge 0$. Inductive Step:

• **Case 1:** *k* + 1 is even.

If k + 1 is even, then there is an integer x s.t. k + 1 = 2x, so then k = 2x - 1 = 2(x - 1) + 1, so therefore <u>k is odd</u>. We know that since k+1 > 0, oddr(k+1) should return <u>loddr(k)</u>. By the Inductive Hypothesis, we know that since k is odd, oddr(k) returns True, so oddr(k+1) returns <u>loddr(k)</u> = False, and k + 1 is even, therefore P(k+1) is true.

• **Case 2:** k + 1 is odd.

If k + 1 is odd, then there is an integer x s.t. k + 1 = 2x + 1, so then k = 2x and therefore <u>k is even</u>. We know that since k + 1 > 0, oddr(k+1) should return <u>loddr(k)</u>. By the Inductive Hypothesis, we know that since k is even, oddr(k) returns False, so oddr(k+1) returns <u>loddr(k)</u> = True, and k + 1 is odd, therefore P(k+1) is true.

Then P(k+1) is true for all cases. Thus, we have shown P(n) is true for all integers $n \ge 0$ by induction.

4. Strong Induction: Stamp Collection

A store sells 3 cent and 5 cent stamps. Use strong induction to prove that you can make exactly n cents worth of stamps for all $n \ge 10$.

Hint: you'll need multiple base cases for this - think about how many steps back you need to go for your inductive step.

Solution:

Let P(n) be defined as "You can buy exactly n cents of stamps". We will prove P(n) is true for all integers $n \ge 10$ by strong induction.

Base Cases: (n = 10, 11, 12):

- n = 10: 10 cents of stamps can be made from two 5 cent stamps.
- n = 11: 11 cents of stamps can be made from one 5 cent and two 3 cent stamps.
- n = 12: 12 cents of stamps can be made from four 3 cent stamps.

Inductive Hypothesis: Suppose for some arbitrary integer $k \ge 12$, $P(10) \land P(11) \land \dots \land P(k)$ holds.

Inductive Step:

Goal: Show P(k+1), i.e. show that we can make k+1 cents in stamps.

We want to buy k + 1 cents in stamps. By the I.H., we can buy exactly (k + 1) - 3 = k - 2 cents in stamps. Then, we can add another 3 cent stamp in order to buy k + 1 cents in stamps, so P(k + 1) is true.

Note: How did we decide how many base cases to have? Well, we wanted to be able to assume P(k-2), and add 3 to achieve P(k+1). Therefore we needed to be able to assume that $k-2 \ge 10$. Adding 2 to both sides, we needed to be able to assume that $k \ge 12$. So, we have to prove the base cases up to 12, that is: 10, 11, 12.

Another way to think about this is that we had to use a fact from 3 steps back from k + 1 to k - 2 in the IS, so we needed 3 base cases.

Conclusion: So by strong induction, P(n) is true for all integers $n \ge 10$.

5. Strong Induction: Recursively Defined Functions

Consider the function f(n) defined for integers $n \ge 1$ as follows: f(1) = 1 for n = 1 f(2) = 4 for n = 2 f(3) = 9 for n = 3f(n) = f(n-1) - f(n-2) + f(n-3) + 2(2n-3) for $n \ge 4$

Prove by strong induction that for all $n \ge 1$, $f(n) = n^2$.

Complete the induction proof below.

Solution:

- 1 Let P(n) be defined as " $f(n) = n^2$ ". We will prove P(n) is true for all integers $n \ge 1$ by strong induction.
- 2 Base Cases (n = 1, 2, 3):
 - n = 1: $f(1) = 1 = 1^2$.
 - n = 2: $f(2) = 4 = 2^2$.
 - n = 3: $f(3) = 9 = 3^2$

So the base cases hold.

- 3 Inductive Hypothesis: Suppose for some arbitrary integer $k \ge 3$, P(j) is true for $1 \le j \le k$.
- 4 Inductive Step:

Goal: Show
$$P(k + 1)$$
, i.e. show that $f(k + 1) = (k + 1)^2$.

$$\begin{split} f(k+1) &= f(k+1-1) - f(k+1-2) + f(k+1-3) + 2(2(k+1)-3) & \text{Definition of f} \\ &= f(k) - f(k-1) + f(k-2) + 2(2k-1) \\ &= k^2 - (k-1)^2 + (k-2)^2 + 2(2k-1) & \text{By IH} \\ &= k^2 - (k^2 - 2k + 1) + (k^2 - 4k + 4) + 4k - 2 \\ &= (k^2 - k^2 + k^2) + (2k - 4k + 4k) + (-1 + 4 - 2) \\ &= k^2 + 2k + 1 \\ &= (k+1)^2 \end{split}$$

So P(k+1) holds.

5 **Conclusion:** So by strong induction, P(n) is true for all integers $n \ge 1$.

6. Strong Induction: A Variation of the Stamp Problem

A store sells candy in packs of 4 and packs of 7. Let P(n) be defined as "You are able to buy n packs of candy". For example, P(3) is not true, because you cannot buy exactly 3 packs of candy from the store. However, it turns out that P(n) is true for any $n \ge 18$. Use strong induction on n to prove this.

Hint: you'll need multiple base cases for this - think about how many steps back you need to go for your inductive step.

Solution:

Let P(n) be defined as "You are able to buy n packs of candy". We will prove P(n) is true for all integers $n \ge 18$ by strong induction.

Base Cases: (n = 18, 19, 20, 21):

- n = 18: 18 packs of candy can be made up of 2 packs of 7 and 1 pack of 4 (18 = 2 * 7 + 1 * 4).
- n = 19: 19 packs of candy can be made up of 1 pack of 7 and 3 packs of 4 (19 = 1 * 7 + 3 * 4).
- n = 20: 20 packs of candy can be made up of 5 packs of 4 (20 = 5 * 4).
- n = 21: 21 packs of candy can be made up of 3 packs of 7 (21 = 3 * 7).

Inductive Hypothesis: Suppose for some arbitrary integer $k \ge 21$, $P(18) \land \dots \land P(k)$ hold.

Inductive Step:

Goal: Show P(k+1), i.e. show that we can buy k+1 packs of candy.

We want to buy k+1 packs of candy. By the I.H., we can buy exactly k-3 packs, so we can add another pack of 4 packs in order to buy k+1 packs of candy, so P(k+1) is true.

Note: How did we decide how many base cases to have? Well, we wanted to be able to assume P(k-3), and add 4 to achieve P(k+1). Therefore we needed to be able to assume that $k-3 \ge 18$. Adding 3 to both sides, we needed to be able to assume that $k \ge 21$. So, we have to prove the base cases up to 21, that is: 18, 19, 20, 21.

Another way to think about this is that we had to use a fact from 4 steps back from k + 1 to k - 3 in the IS, so we needed 4 base cases.

Conclusion: So by strong induction, P(n) is true for all integers $n \ge 18$.