# **CSE 390Z:** Mathematics for Computation Workshop

# **Practice 311 Midterm Solutions**

Name: \_\_\_\_\_

UW ID: \_\_\_\_\_

#### Instructions:

- This is a **simulated practice midterm**. You will **not** be graded on your performance on this exam.
- Nevertheless, please treat this as if it is a real exam. That means that you may not discuss with your neighbors, reference outside material, or use your devices during the next 50 minute period.
- If you get stuck on a problem, consider moving on and coming back later. In the actual exam, there will likely be opportunity for partial credit.
- There are 4 problems on this exam, totaling 70 points.

#### 1. Predicate Translation [20 points]

Let the domain of discourse be animals. Translate the following statements to predicate logic, using the following predicates:

Bird(x) := x is a bird CanFly(x) := x can fly Loves(x, y) := x loves y

(a) (5 points) There is a bird that loves itself.

#### Solution:

 $\exists x (\mathsf{Bird}(x) \land \mathsf{Loves}(x, x))$ 

(b) (5 points) All flightless birds love a bird that can fly (the bird being loved can be different for each flightless bird).

### Solution:

 $\forall x (\neg \mathsf{CanFly}(x) \land \mathsf{Bird}(x) \to \exists y (\mathsf{Bird}(y) \land \mathsf{CanFly}(y) \land \mathsf{Loves}(x,y))$ 

(c) (5 points) Every bird loves exactly one non bird.

## Solution:

$$\begin{split} &\forall x(\mathsf{Bird}(x) \to \exists y(\neg \mathsf{Bird}(y) \land \mathsf{Loves}(x, y) \land \forall z((z \neq y \land \neg \mathsf{Bird}(z)) \to \neg \mathsf{Loves}(x, z)))) \\ &\mathsf{OR} \\ &\forall x(\mathsf{Bird}(x) \to \exists y(\neg \mathsf{Bird}(y) \land \mathsf{Loves}(x, y) \land \forall z(\mathsf{Loves}(x, z) \to (\mathsf{Bird}(z) \lor z = y)))) \\ &\mathsf{OR} \\ &\forall x(\mathsf{Bird}(x) \to \exists y(\neg \mathsf{Bird}(y) \land \mathsf{Loves}(x, y) \land \forall z((\mathsf{Loves}(x, z) \land \neg \mathsf{Bird}(z)) \to z = y))) \end{split}$$

\*Note that a great exercise is to show that the above 3 solutions are all logically equivalent :)

(d) (5 points) Translate the following into English:  $\exists x \forall y (\mathsf{Bird}(x) \land (\mathsf{CanFly}(y) \to \mathsf{Loves}(x, y))$ 

#### Solution:

There is a bird that loves all flying animals.

#### 2. Number Theory Proof [20 points]

Prove the following statement using an English proof:

The square of every odd integer can be written in the form 8k + 1 for some integer k.

Hint: You may want to use proof by cases for this problem.

#### Solution:

Let n be an arbitrary odd integer.

We want to show that  $n^2$  can be written in the form 8k + 1 for some integer k. In other words, we want to show  $n^2 = 8k + 1$ . Rearranging this equation, we get  $n^2 - 1 = 8k$ .

By definition of odd, n = 2j + 1 for some integer j. Then,  $n^2 - 1 = 4j^2 + 4j + 1 - 1 = 4(j^2 + j)$ . At this point, consider two cases:

• Case 1: *j* is even.

By definition of even, j = 2x for some integer x. Then

$$j^{2} + j = (2x)^{2} + (2x)$$
$$= 4x^{2} + 2x$$
$$= 2(2x^{2} + x)$$

Plugging this into the above equation, we get

$$n^{2} - 1 = 4(2(2x^{2} + x)) = 8(2x^{2} + x)$$

Let  $k = 2x^2 + x$ . Then

$$n^2 = 8k + 1$$

Since x is an integer,  $2x^2 + x$  must be an integer. So, there exists an integer k such that we can write  $n^2 = 8k + 1$ .

Case 2: j is odd.
 By definition of odd, j = 2x + 1 for some integer x. Then

$$j^{2} + j = (2x + 1)^{2} + (2x + 1)$$
  
= 4x<sup>2</sup> + 4x + 1 + 2x + 1  
= 4x<sup>2</sup> + 6x + 2  
= 2(2x<sup>2</sup> + 3x + 1)

Plugging this into the above equation, we get

$$n^{2} - 1 = 4(2(2x^{2} + 3x + 1)) = 8(2x^{2} + 3x + 1)$$

Let  $k = 2x^2 + 3x + 1$ . Then

$$n^2 = 8k + 1$$

Since x is an integer,  $2x^2 + 3x + 1$  must be an integer. So, there exists an integer k such that we can write  $n^2 = 8k + 1$ .

In all cases, we were able to write  $n^2$  in the form 8k + 1. Since n was an arbitrary odd integer, this proves the original claim.

#### 3. Induction [20 points]

Recall the Fibonacci sequence:  $f_0 = 0$ ,  $f_1 = 1$ , and  $f_n = f_{n-1} + f_{n-2}$  for all  $n \ge 2$ . Prove that  $f_{n+1} < \left(\frac{7}{4}\right)^n$  for all  $n \ge 1$ .

Hint: To prove a < inequality, start with the left side of the inequality and use algebra + your inductive hypothesis to show that it is less than the right side.

#### Solution:

1. Let P(n) be  $f_{n+1} < \left(\frac{7}{4}\right)^n$ . We will prove P(n) true for all integers  $n \ge 1$  by strong induction.

- 2. Base Cases
  - n = 1:  $f_{1+1} = f_2 = f_1 + f_0 = 1 + 0 = 1 < \frac{7}{4} = \left(\frac{7}{4}\right)^1$ , so P(1) holds.
  - n = 2:  $f_{2+1} = f_3 = f_2 + f_1 = 1 + 1 = 2 < \frac{49}{16} = \left(\frac{7}{4}\right)^2$ , using the fact that  $f_2 = 1$  from above. So, P(2) holds.
- 3. Inductive Hypothesis: Assume that for some arbitrary integer  $k \ge 2$ , P(j) holds for all  $2 \le j \le k$ .
- 4. Inductive Step:

Goal: Show 
$$P(k+1)$$
, i.e. show  $f_{(k+1)+1} < \left(\frac{7}{4}\right)^{k+1}$ 

$$f_{(k+1)+1} = f_{k+1} + f_k \qquad \text{Def Fibonacci}$$

$$< \left(\frac{7}{4}\right)^k + \left(\frac{7}{4}\right)^{k-1} \qquad \text{I.H.}$$

$$= \left(\frac{7}{4}\right)^k \cdot \left[1 + \left(\frac{7}{4}\right)^{-1}\right] \qquad \text{Factor out } \left(\frac{7}{4}\right)^k \text{ term}$$

$$= \left(\frac{7}{4}\right)^k \left(\frac{11}{7}\right) \qquad \text{Simplify fractions}$$

$$= \left(\frac{7}{4}\right)^{k+1} \left(\frac{7}{4}\right)^{-1} \left(\frac{11}{7}\right) \qquad \text{Split up } k \text{ exponent}$$

$$= \left(\frac{7}{4}\right)^{k+1} \cdot \frac{44}{49} \qquad \text{Simplify fractions}$$

$$< \left(\frac{7}{4}\right)^{k+1}$$

Thus P(k+1) holds.

5. Therefore, by the principles of strong induction, the claim P(n) holds for all integers  $n \ge 1$ .

#### 4. Miscellaneous [10 points]

For parts a-c, write if the statement is true or false. If it is false, provide a counterexample. You do not need to provide any proofs/reasoning if the statement is true.

(a) (2 points)  $4 \mid a^4$  for all integers a.

### Solution:

False. Counterexample is a = 1.  $4 \nmid 1^4$ 

(b) (2 points) If a|b and b|c, then a|c.

#### Solution:

True.

Given that a|b and b|c, we know that b = ax and c = by for some integers x, y. Substituting the first equation into the second, we get c = (ax)y = a(xy) which means a|c.

(c) (2 points) If  $a \equiv b \pmod{5}$  and  $a \equiv 1 \pmod{5}$  then 5|b.

### Solution:

False. Counterexample: a = 1 and b = 6.  $1 \equiv 6 \pmod{5}$  and  $1 \equiv 1 \pmod{5}$  but  $5 \nmid 6$ .

(d) (4 points) Which of the following are equivalent to  $\neg p \lor (q \to r)$ ? Select all that apply.

$$\Box \neg p \lor (q \lor r)$$
$$\Box \neg (p \land q) \lor r$$
$$\Box (p \to r) \lor \neg q$$
$$\Box \neg p \lor (q \land r)$$

#### Solution:

The second and third options are correct.