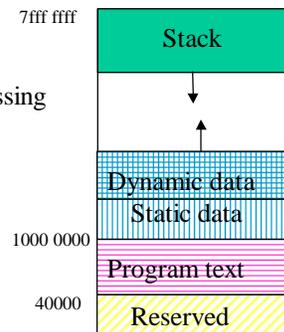


## Program and memory layout

- By convention the layout is:

- Note that only half of the addressing space is taken by user  
Other half is O.S.



## Procedures

- Procedures/functions are the major program structuring mechanism
- Calling and returning from a procedure requires a protocol between *caller* and *callee*
- Protocol is based on conventions

## Procedures/Functions -- Protocol

- Each machine (compiler?) has its own set of protocol(s)
- Protocol: combination of hardware/software
  - e.g., “jal” is hardware; use of register \$29 as \$sp is software
- Protocol: sequence of steps to be followed at each call and each return
  - controlled by hardware and/or software
- In RISC machines
  - hardware performs simple instructions
  - software (compiler/assembler) controls sequence of instructions

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## Program stack

- Each executing program (process) has a *stack*
- Stack = dynamic data structure accessed in a LIFO manner
- Program stack automatically allocated by O.S.
- At the start of the program, register \$sp (\$29 in Mips) is automatically loaded to point to the first empty slot on top of stack
  - After that it will be your responsibility to manage \$sp
- By convention, stack grows towards lower addresses
  - to allocate new space (i.e., when you *push*), decrement \$sp
  - to free space on top of stack (*pop*), increment \$sp

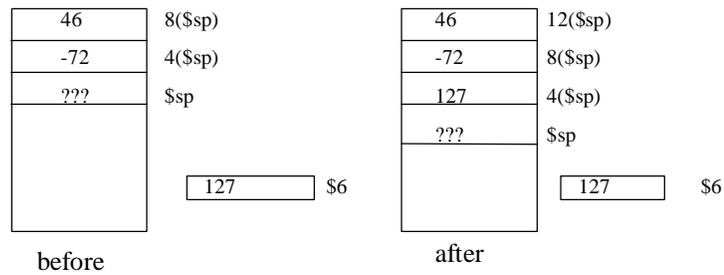
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## Push operation

- *push* adds an item on top of stack
  - one instruction to manipulate the data, e.g. “sw \$6,0(\$sp)”
  - one instruction to adjust the stack pointer e.g., “subu \$sp,\$sp,4”



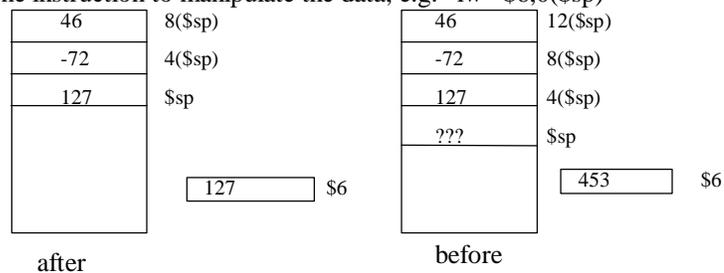
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## Pop operation

- *pop* removes the item on top of stack and stores it in a register
  - one instruction to adjust the stack pointer e.g., “addu \$sp,\$sp,4”
  - one instruction to manipulate the data, e.g. “lw \$6,0(\$sp)”



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## Procedure call requirements (caller/callee)

- Caller must pass the return address to the callee
- Caller must pass the parameters to the callee
- Caller must save what is *volatile* (registers) and could be used by the callee
- Callee must save the return address (in case it becomes a caller)
- Callee must provide (stack) storage for its own use
- Caller/callee should support recursive calls

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## Mechanism

- Registers are used for
  - passing return address in \$ra
    - jal target
  - passing a small number of parameters (up to 4 in \$a0 to \$a3)
  - keeping track of the stack (\$sp)
  - returning function values (in \$v0 and \$v1)
- Stack is used for
  - saving registers to be used by callee
  - saving info about the caller (return address)
  - passing parameters if needed
  - allocating local data for the called procedure

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## Procedure calls and register conventions

Register	Name	Function	Comment
\$0	Zero	Always 0	No-op on write
\$1	\$at	Reserved for assembler	Don't use it
\$2-3	\$v0-v1	Expr. Eval/funct. Return	
\$4-7	\$a0-a3	Proc./func. Call parameters	
\$8-15	\$t0-t7	Temporaries; volatile	Not saved on proc. Calls
\$16-23	\$s0-s7	Temporaries	Should be saved on calls
\$24-25	\$t8-t9	Temporaries; volatile	Not saved on proc. Calls
\$26-27	\$k0-k1	Reserved for O.S.	Don't use them
\$28	\$gp	Pointer to global static memory	
\$29	\$sp	Stack pointer	
\$30	\$fp	Frame pointer	
\$31	\$ra	Proc./func return address	

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## Who does what on a call (one sample protocol)

- Caller
  - Saves any volatile register (\$t0-\$t9) that has contents that need to be kept
  - Puts up to 4 arguments in \$a0-\$a3
  - If more than 4 arguments, pushes the rest on the stack
  - calls with jal instruction
- Callee
  - saves \$ra on stack
  - saves any non-volatile register (\$s0-s7) that it will use

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## Who does what on return

- Callee
  - restores any non-volatile register (\$s0-\$s7) it has used
  - restores \$ra
  - puts function results in \$v0-\$v1
  - adjusts \$sp
  - returns to caller with “jr \$ra”
- Caller
  - restores any volatile register it had saved
  - examines \$v0-\$v1 if needed

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## Example of a call sequence

- Assume 2 arguments in \$t0 and \$t3 and we want to save the contents of \$t6 and \$t7

```
move    $a0,$t0    #1st argument in $a0
move    $a1,$t3    #2nd argument in $a1
subu    $sp,$sp,8  #room for 2 temps on stack
sw      $t6,8($sp) #save $t6 on stack
sw      $t7,4($sp) #save $t7 on stack
jal     target
```

- Assume the callee does not need to save registers

```
target: sw  $ra,0($sp) #save return address
        subu $sp,$sp,4 # on stack
```

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## Return from the previous sequence

- The callee will have put the function results in \$v0-\$v1

```
addu    $sp,$sp,4           #pop
lw      $ra,0($sp)         #return address in $ra
jr      $ra                 #to caller
```

- The caller will restore \$t6 and \$t7 and adjust stack

```
lw      $t6,8($sp)
lw      $t7,4($sp)
addu    $sp,$sp,8
```