Writing Into The Cache

- Write-through ... every write to the cache goes “through” to the memory, keeping it consistent
- Copy-back ... writes go only as “far” as the cache, and then they are written back to memory when the block is flushed
- Other concepts
  - Write buffer
  - Write buffer stall
  - Fill on write

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Virtual Memory

Virtual Memory is the process of mapping a logical address space numbered from 0 to the physical address space of the computer. The mapping is performed on a page basis.

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Virtual Memory

Virtual memory solves several critical problems:

- Simplifying RAM memory management by separating the address of memory from its physical location
- Providing protection for users by giving them their own address spaces
- Simplifying compilation and usage of libraries by allowing all programs to begin addressing memory from zero

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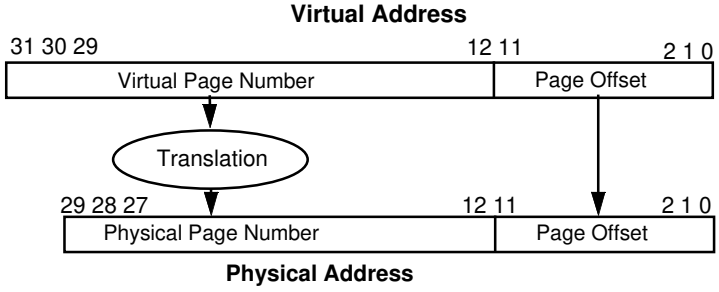
Virtual to Physical Translation

Physical memory (RAM) is divided into pages -- contiguous sequences of memory typically in the 4KB -- 16KB range

The pages are blocks of a fully associative cache for the memory of the program which is stored in the next lower level of the memory hierarchy (disk)



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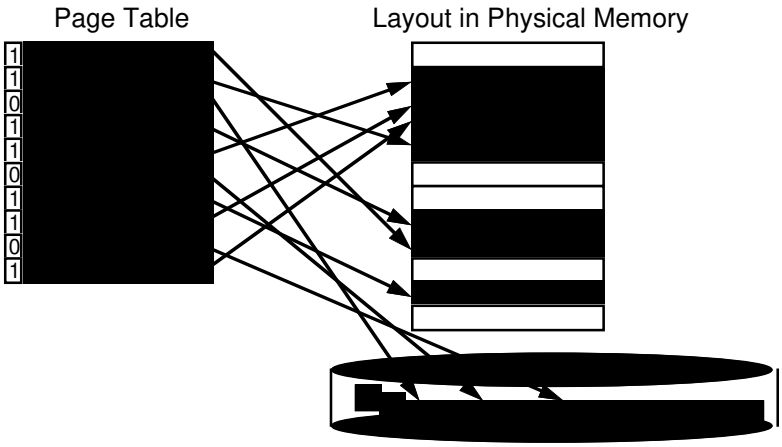
Page Table



Page table and disk addresses may be kept in separate tables

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VM Considerations

Large page sizes allow the huge miss penalties to be amortized over many references

Avoiding premature page replacment, e.g. by address collisions, extends page life and reduces fault rate ... thus full associativity

Page faults can be handled in software, which can use cleverness to reduce fault rate

Write back is needed because of the high cost of writes

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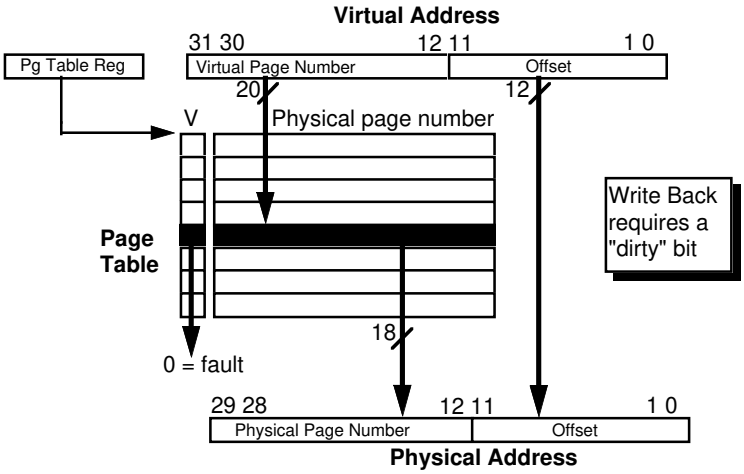
Implementing Address Translation



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Page Replacement Strategy

When the Valid = 0, a page fault is signaled.

Some page must be replaced -- pick the page that will be used furthest into the future: Opt

Replace the least recently used (LRU) page.

- LRU strategies are effective, but expensive.
- "Use" bits can be a decent approximation.

Space required for page tables can be substantial:

- 4K pages imply 20 bits of virtual address.
- At x bytes per entry implies xMB for full page table.

Represent only the prefix of the table using base + extent.

Grow low addresses and high addresses separately by using the msb to indicate which part of the VM space is being used.

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