

Flow of Control -- Conditional branch instructions

- You can compare directly
 - Equality or inequality of two registers
 - One register with 0 (>, <, ≥, ≤)
- and branch to a target specified as
 - a signed displacement expressed in *number of instructions* (not number of bytes) from the instruction *following* the branch
 - in assembly language, it is **highly** recommended to use labels and branch to labeled target addresses because:
 - the computation above is too complicated
 - some pseudo-instructions are translated into two real instructions

10/8/2004

CSE378 Instr. encoding. (c't'd)

1

Examples of branch instructions

```
Beq    rs,rt,target    #go to target if rs = rt
Beqz   rs, target     #go to target if rs = 0
Bne    rs,rt,target    #go to target if rs != rt
Bltz   rs, target     #go to target if rs < 0
```

etc.

but note that you cannot compare directly 2 registers for <, > ...

10/8/2004

CSE378 Instr. encoding. (c't'd)

2

Comparisons between two registers

- Use an instruction to set a third register


```
slt    rd,rs,rt    #rd = 1 if rs < rt else rd = 0
sltu   rd,rs,rt    #same but rs and rt are considered unsigned
```
 - Example: Branch to Lab1 if \$5 < \$6


```
slt    $10,$5,$6    # $10 = 1 if $5 < $6 otherwise $10 = 0
bnez   $10,Lab1     # branch if $10 = 1, i.e., $5 < $6
```
 - There exist pseudo instructions to help you!


```
blt    $5,$6,Lab1   # pseudo instruction translated into
                    # slt $1,$5,$6
                    # bne $1,$0,Lab1
```
- Note the use of register 1 by the assembler and the fact that computing the address of Lab1 requires knowledge of how pseudo-instructions are expanded

10/8/2004

CSE378 Instr. encoding. (c't'd)

3

Unconditional transfer of control

- Can use "beqz \$0, target"
 - Very useful but limited range ($\pm 32K$ instructions)
- Use of Jump instructions


```
j    target    #special format for target byte address (26 bits)
jr   $rs       #jump to address stored in rs (good for switch
                #statements and transfer tables)
```
- Call/return functions and procedures


```
jal target    #jump to target address; save PC of
                #following instruction in $31 (aka $ra)
jr   $31      #jump to address stored in $31 (or $ra)
```

Also possible to use

```
jalr rs,rd    #jump to address stored in rs; rd = PC of
                #following instruction in rd with default rd = $31
```

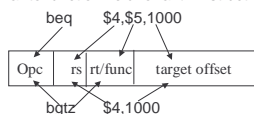
10/8/2004

CSE378 Instr. encoding. (c't'd)

4

Branch addressing format

- Need Opcode, one or two registers, and an offset
 - No base register since offset added to PC
- When using one register (i.e., compare to 0), can use the second register field to expand the opcode
 - similar to function field for arith instructions



10/8/2004

CSE378 Instr. encoding. (c't'd)

5

How to address operands

- The ISA specifies *addressing modes*
- MIPS, as a RISC machine has very few addressing modes
 - register mode*. Operand is in a register
 - base or displacement or indexed mode*
 - Operand is at address "register + 16-bit signed offset"
 - immediate mode*. Operand is a constant encoded in the instruction
 - PC-relative mode*. As base but the register is the PC

10/8/2004

CSE378 Instr. encoding. (c't'd)

6

Some interesting instructions. Multiply

- Multiplying 2 32-bit numbers yields a 64-bit result
 - Use of HI and LO registers
- ```
Mult rs,rt #HI/LO = rs*rt
Multu rs,rt
```
- Then need to move the HI or LO or both to regular registers
- ```
mflo  rd      #rd = LO
mfhi  rd      #rd = HI
```
- Once more the assembler can come to the rescue with a pseudo inst
- ```
mul rd,rs,rt #generates mult and mflo
 #and mfhi if necessary
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

7

## Some interesting instructions. Divide

- Similarly, divide needs two registers
  - LO gets the quotient
  - HI gets the remainder
- If an operand is negative, the remainder is not specified by the MIPS ISA.

10/8/2004

CSE378 Instr. encoding. (c't'd)

8

## Logic instructions

- Used to manipulate bits within words, set-up masks etc.
  - A sample of instructions
- ```
and   rd,rs,rt      #rd=AND(rs,rt)
andi  rd,rs,immed
or    rd,rs,rt
xor   rd,rs,rt
```
- Immediate constant limited to 16 bits (zero-extended). If longer mask needed, use Lui.
 - There is a pseudo-instruction NOT
- ```
not rt,rs #does 1's complement (bit by bit
 #complement of rs in rt)
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

9

## Example of use of logic instructions

- Create a *mask* of all 1's for the low-order byte of \$6. Don't care about the other bits.
- ```
ori   $6,$6,0x00ff  # $6[7:0] set to 1's
```
- Clear high-order byte of register 7 but leave the 3 other bytes unchanged
- ```
lui $5,0x00ff # $5 = 0x00ff0000
ori $5,$5,0xffff # $5 = 0x00ffff
and $7,$7,$5 # $7 = 0x00..... (...whatever was
 #there before)
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

10

## Shift instructions

- Logical shifts -- Zeroes are inserted
- ```
sll  rd,rt,shamt    #left shift of shamt bits; inserting 0's on
                        #the right
srl  rd,rt,shamt    #right shift of shamt bits; inserting 0's
                        #on the left
```
- Arithmetic shifts (useful only on the right)
- ```
– sra rd,rt,shamt # Sign bit is inserted on the left
```
- Example let \$5 = 0xff00 0000
- ```
sll  $6,$5,3        # $6 = 0xf800 0000
srl  $6,$5,3        # $6 = 0x1fe0 0000
sra  $6,$5,3        # $6 = 0xffe0 0000
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

11

Example -- High-level language

```
int a[100];
int i;

for (i=0; i<100; i++){
    a[i] = 5;
}
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

12

Assembly language version

Assume: start address of array a in r15.
We use r8 to store the value of i and r9 for the value 5

```
add    $8,$0,$0    #initialize i
li     $9,5        #r9 has the constant 5
Loop: mul   $10,$8,4    #r10 has i in bytes
        #could use a shift left by 2
addu   $14,$10,$15   #address of a[i]
sw     $9,0($14)    #store 5 in a[i]
addiu  $8,$8,1      #increment i
blt    $8,100,Loop  #branch if loop not finished
        #taking lots of liberty here!
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

13

Machine language version (generated by SPIM)

```
[0x00400020] 0x00004020 add $8, $0, $0 ; 1: add $8,$0,$0
[0x00400024] 0x34090005 ori $9, $0, 5 ; 2: li $9,5
[0x00400028] 0x34010004 ori $1, $0, 4 ; 3: mul $10,$8,4
[0x0040002c] 0x01010018 mult $8, $1
[0x00400030] 0x00005012 mflw $10
[0x00400034] 0x0147021 addu $14, $10, $15 ; 4: addu $14,$10,$15
[0x00400038] 0xad90000 sw $9, 0($14) ; 5: sw $9,0($14)
[0x0040003c] 0x25080001 addiu $8, $8, 1 ; 6: addiu $8,$8,1
[0x00400040] 0x29010064 slti $1, $8, 100 ; 7: blt $8,100,Loop
[0x00400044] 0x1420ff9 bne $1, $0, -28 ;[Loop-0x00400044]
```

10/8/2004

CSE378 Instr. encoding. (c't'd)

14