Question 1 (8 points). Preprocessor. Consider the following program:

```c
#include <stdio.h>
#define aardvark 5
#define sum(x,y) x+y
#define times(x,y) x*y

#ifndef aardvark
int magic = 2;
#define DEBUG(s) printf("%s\n", s)
#else
int magic = 3;
#define DEBUG(s)
#endif

void foo(int* a, int* b, int* c) {
    *a = 1 * *c;
    *b = 2 * times(*b, *a);
    *c = 3 * sum(*a, *b);
    DEBUG("in foo");
}

int main() {
    int w = aardvark;
    int x = 1;
    int y = magic;
    int z = 2;
    foo(&x, &y, &w);
    printf("%d %d %d %d\n", w, x, y, z);
    foo(&w, &w, &z);
    printf("%d %d %d %d\n", w, x, y, z);
    return 0;
}
```

Write the output of this program below.

```
45 5 30 2
8 5 30 32
```
Question 2 (10 points). make. Suppose we have a project with multiple source files that have the \#include dependencies shown below.

```
***************
* vehicle.h *
***************
#ifndef VEHICLE_H
#define VEHICLE_H
...
#endif

***************
* vehicle.c *
***************
#include "vehicle.h"
...

***************
* wheel.h *
***************
#ifndef WHEEL_H
#define WHEEL_H
...
#endif

***************
* car.h *
***************
#endif

***************
* car.c *
***************
#include "car.h"
...

***************
* main.h *
***************
#include "main.h"
int main() {
...
}
```

If we’re being lazy, we could create an executable program from these files by executing the command `gcc -Wall -g -std=c11 -o main *.c`.

On the next page, create a Makefile that builds an executable program named `main` from these files as done by the `gcc` command above, but only recompiles and relinks the minimum number of files needed after any changes are made to the source files. Your answer should be done in two steps:

- First, construct the dependency graph that shows the dependencies between the source files, the compiled .o files, and the final executable file `main` that is created by linking the .o files.
- After drawing the dependency graph, write the final Makefile.

Write your answers on the next page.
(2a) (5 points). Draw the dependency graph (diagram) showing dependencies between the files on the previous page, the .o files created by compiling the .c files, and the final executable program main.

![Dependency Diagram]

(2b) (5 points). Give the contents of a Makefile that will build the program as described by the dependency diagram above. The default target that is built if we just type make with no arguments should be the executable program main.

```
main: vehicle.o car.o main.o
    gcc -Wall -g -std=c11 -o main vehicle.o car.o main.o

vehicle.o: vehicle.c vehicle.h
    gcc -Wall -g -std=c11 -c vehicle.c

car.o: car.c car.h vehicle.h wheel.h
    gcc -Wall -g -std=c11 -c car.c

main.o: main.c vehicle.h car.h
    gcc -Wall -g -std=c11 -c main.c
```
Question 3 (12 points). Debugging + scripting. While working on the memory manager for HW6, your partner has discovered that the bench program crashes with a segmentation fault when it is run with no arguments.

    $ ./bench
    Segmentation fault

You suspect that the problem is because you are accessing the command-line arguments (argv) without properly checking that they are set first.

(3a) (5 points). Your partner doesn't know how to use gdb, but you do! Below, describe to your partner the exact commands that they can execute to run gdb on the bench executable and see if your hypothesis is true. This includes the commands that you will give to gdb itself.

    # Compile the program (make sure the -g flag is set!)
    $ make
    $ gdb ./bench

    Within gdb:
    # run until it hits the segfault
    run
    # see what function segfaulted and what the callstack was
    # hypothesis: it's in main where we access the argv - is the
    #             crash in main
    bt
    # list code around where we are
    list
    # if the hypothesis was correct, and we're in main, we can print
    # our variables
    print argv[1]
    print argc
    # if we're not in main, but one of our arguments is derived from
    # argv, we can go up in the callstack to main
    up

After the issues that your partner has been having, you decide to write a simple bash script to do some basic validation. On the next page, write a script that does the following three things:

- Runs clint.py on all source files (.c and .h). Assume clint.py is in the same directory.
- Builds the executable with make.
- Runs the bench program with no arguments (bench target in the Makefile).

If any of the three commands exits with a status code that is not 0, then you should print an error message to stderr and exit with status code 1. Otherwise, if all three exit successfully with status code 0, exit with status code 0 without printing anything (redirect output for the other commands). Write the script on the following page – some bash hints are provided.
(3b) (5 points). Write a script to perform the three validation steps described on the last page.

Some of the tests that can appear in a [ ] or [[ ]] test command in a bash script:
- string comparisons: =, !=
- numeric comparisons: !=eq, -ne, -gt, -ge, -lt, -le
Shell variables: $# (# arguments), $? (last command result), $0, $1, … (specific arguments), shift (discard first argument)
Blackhole file: /dev/null (used for suppressing output)

#!/bin/bash

# clint.py
# First method for running a command testing its exit code
if ! ./clint.py *.h *.c > /dev/null ; then
    echo "clint.py failed" >&2 ; exit 1
fi

# Basic make
# Second method for testing an exit code
make > /dev/null
if [ $? -ne 0 ]
then
    echo "make failed" >&2
    exit 1
fi

# Run bench
# Third method for testing an exit code
make bench > /dev/null || { echo "bench failed" >&2 ; exit 1 ; }
exit 0

(3c) (2 points). Give the shell commands to put the validate.sh script from the previous question into your git repository and share it with your partner.

# Your side
git add validate.sh
git commit -m "script to run before commits to validate repo"
git push

// Partner's side
git pull
Question 4 (14 points). Tries. In HW5 we used a trie to store the words in a dictionary based on the T9 digit sequences that encoded the words. For this problem, assume that a node is defined as follows:

```c
typedef struct Node {
    char* word; // C-string if this node has a word attached, otherwise NULL
        // are subtrees for digits 2-9;
        // next[0] is the synonym ('#') link.
} Node; // For 0<=i<10, next[i]==NULL if
        // next[i] is an empty subtree.
```

For this problem, write a function called `isAlphabetized` that takes a `struct Node*` as an argument and determines whether the words in the trie are alphabetized within each T9 digit sequence. The function returns 1 if the trie is alphabetized, 0 if not.

You ONLY need to determine whether the words in each T9 group (those sharing the same T9 digit sequence) are in alphabetical order. You do NOT need to compare ordering across different T9 digit sequences. For instance, you should verify that all words with the sequence 227 are alphabetized relative to each other, but you do not need to check that they are in any particular order relative to words with the sequence 2273.

You should assume that all necessary header files have already been `#included` and you do not need to add any `#includes`. An empty trie or a NULL trie passed to your function is by definition alphabetized. Duplicate words do not break alphabetization.

Hints: recursion really, really, really is your friend. Remember that all words with the same T9 digit sequence are represented as a node with a non-NUL word and all nodes that follow in the 0th (#) child position. Alphabetically later words are considered "greater".

A bit of (maybe) useful reference information about strings:

- `char* strcpy(dest, src)`
- `char* strncpy(dest, src, n)`
- `char* strcat(dest, src)`
- `char* strncat(dest, src, n)`
- `int strcmp(string1, string2)`
- `int strncmp(string1, string2, n)`
- `char* strstr(string, search_string)`
- `int strlen(s)`
- `int strnlen(s, max_length)`
- `int strxlen(s)`
- `Character tests: isupper(c), islower(c), isdigit(c), isspace(c)`
- `Character conversions: toupper(c), tolower(c)`

Write your answer on the next page.
Question 4 (cont). Write your implementation of the trie `isAlphabetized` function below.

// Returns 1 if words in the trie are alphabetized within each T9 digit sequence, 0 if they are not. Words are not compared across different T9 digit sequences.

```c
int isAlphabetized(struct Node* root) {
    if (root == NULL) {
        return 1;
    }
    for (int i = 0; i < 10; i++) {
        if (!isAlphabetized(root->next[i])) {
            return 0;
        }
    }
    return 
        root->word == NULL ||
        root->next[0] == NULL ||
        root->next[0]->word = NULL ||
        strcmp(root->word, root->next[0]->word) <= 0;
}
```

// Alternative approach, explicit next[0] iteration

```c
int isAlphabetized(struct Node* root) {
    if (root == NULL) {
        return 1;
    }
    for (int i = 2; i < 10; i++) {
        if (!isAlphabetized(root->next[i])) {
            return 0;
        }
    }
    if (root->word != NULL) {
        while (root->next[0] != NULL &&
               root->next[0]->word != NULL) {
            if (strcmp(root->word, root->next[0]->word) > 0) {
                return 0;
            }
            root = root->next[0];
        }
        return 1;
    }
```
Question 5 (12 points). Testing. Describe three black-box tests for the `isAlphabetized` function from the previous problem. For full credit, the three tests must verify different things about the implementation, and must describe the specific input or setup for the test and the expected result. In your description of test inputs/outputs you should specify the characteristics of the trie passed as input to the function. A drawing is not necessary, but we do expect specific examples of word(s) and a reasonable explanation of how they are placed in the trie.

Some possible examples (definitely not complete):

INPUT:
   (empty) `isAlphabetized(NULL)`
OUTPUT:
   Returns 1

INPUT:
   (one node) `root(word=ban)` , rest of children are NULL
OUTPUT:
   Returns 1

INPUT:
   (basic alphabetical) `root(word=ban) [0]-> cam [0]-> can`
OUTPUT:
   Returns 1

INPUT:
   (basic non-alphabetical) `root(word=ban) [0]-> can [0]-> cam`
OUTPUT:
   Returns 0

INPUT:
   (basic duplicates) `root(word=ban) [0]-> ban [0]-> ban`
OUTPUT:
   Returns 1

INPUT:
   (children of root) `root[2] -> ban [0]-> can`
   `root[3] -> dad [0]-> fad`
OUTPUT:
   Returns 1

INPUT:
   (children not alphabetized) `root[2] -> ban [0]-> can`
   `root[3] -> fad [0]-> dad`
OUTPUT:
   Returns 0

INPUT:
   (no modification) `root[2] -> ban [0]-> can`
   `root[3] -> fad [0]-> dad`
OUTPUT:
   Double check that tree structure is exactly the same after
Question 6 (14 points). Memory management. Recall from HW6 that we can represent the free list for the getmem/freemem storage allocator as a linked list of blocks. The beginning of each block is described by the following C struct:

```c
typedef struct FreeNode { // node on free list:
    uintptr_t size; // number of bytes in this block, not including the
                    // size of this header
    struct FreeNode* next; // next block on free list or
                            // NULL if none
} FreeNode;
```

In HW6, we used malloc to request large chunks from the system. Clients then used getmem/freemem to reserve smaller pieces of those chunks. However, we did not require you to ever call free to release the chunks that we reserved with malloc. We will do that now.

(6a) (6 points). Write a function `freeAll()` that frees all memory that was allocated by your memory manager with malloc. You should assume that all blocks that the client reserved with getmem have been released with freemem, which means that the free list will store only chunks that were allocated with malloc.

You should also assume that there is a global pointer variable that stores the address of the first block on the free list; your code can access that variable:

```c
FreeNode* freeList; // free list blocks; NULL if freeList is empty
```

Write your function below.

```c
void freeAll() {

    while (freeList != NULL) {
        // We must save the next pointer before calling free.
        FreeNode* next = freeList->next;
        free(freeList);
        freeList = next;
    }
}
```
(6b) (2 points). It's actually not a good idea to assume that all blocks that the client reserved with getmem have been released with freemem, since a client program can behave incorrectly. Describe in detail what might go wrong with the calls to free in freeAll in this case – why would free fail? Hint: malloc/free are similar to getmem/freemem from HW6; you can assume that free acts like freemem.

Two ways this could go wrong:
1) If some blocks are still reserved, then they won't be in the free list and we might not call free() on all of the blocks that we malloc'ed.
2) Since we split big malloc'ed chunks into smaller chunks for getmem, we may have blocks on the free list that are not directly allocated with malloc. If we try to free those chunks before they are merged back with their original partners, we would call free() with a pointer that was not allocated with malloc, which has undefined behavior (ie free won't be able to find the block information for that malloc'ed block).

(6c) (6 points). One way to solve this problem is to do "reference counting." We maintain a global variable blocksAllocated, which starts at 0. Every time getmem is called and returns a valid pointer (not NULL), we increase blocksAllocated by 1. Every time freemem is called with a valid pointer (not NULL), we decrease blocksAllocated by 1. In freeAll, we can then assert that blocksAllocated == 0 before freeing any blocks.

Define an integer blocksAllocated to accomplish this, then implement getmemWrapper and freememWrapper to perform the reference counting (incrementing or decrementing of blocksAllocated). These functions should call getmem and freemem to actually accomplish the allocation/release of the blocks.

// Define blocksAllocated here:

uintptr_t blocksAllocated = 0;

void* getmemWrapper(uintptr_t size) {
    void* result = getmem(size);
    if (result != NULL) {
        blocksAllocated++;
    }
    return result;
}

void freememWrapper(void* ptr) {
    if (ptr != NULL) {
        blocksAllocated--;
    }
    freemem(ptr);
}
Question 7 (12 points). C++.

To explore a bit more C++, we’ve designed a simple class to represent a tuple. A tuple is an ordered sequence of values — like a C array, except that we'll do bounds checking to make sure that we don't access elements that aren't valid. We refer to the number of elements in the tuple as the number of "dimensions".

The class declaration in file Tuple.h looks like this:

```cpp
#ifndef TUPLE_H
#define TUPLE_H

class Tuple {
   public:
      // Construct a tuple with "dim" dimensions.
      explicit Tuple(size_t dim);

      // Destructor
      virtual ~Tuple();

      // accessor - returns value at the index
      int get(size_t index) const;

      // setter - sets the value at the index.
      void set(size_t index, int value);

   private:
      // Array storing the tuple's value in each dimension
      int* point_;
      size_t dimensions_;  
};
#endif   // TUPLE_H
```

On the next page, give implementations for the four functions (constructor, destructor, get accessor and set setter) that are declared but not implemented above.

Reminder: you'll need to declare an array on the heap since we don't know how big it is until the constructor. You can use either new or malloc to achieve this. If the requested index is not valid for get or set (it is larger than the largest valid dimension), then get should return 0 and set should do nothing.

Write your answer on the next page.
**Question 7 (cont).** Provide your implementation of the Tuple class declared on the previous page as it would be written in the implementation file Tuple.cpp.

```cpp
#include "Tuple.h"

// write your implementation of class Tuple below.

Tuple::Tuple(size_t dim) {
    point_ = new int[dim](); // the () sets all elements to 0
    dimensions_ = dim;
}

Tuple::~Tuple() {
    delete [] point_;
}

int Tuple::get(size_t index) const {
    if (index < dimensions_) {
        return point_[index];
    }
    return 0;
}

void Tuple::set(size_t index, int value) {
    if (index < dimensions_) {
        point_[index] = value;
    }
}
```
**Question 8 (10 points). C++ virtual mystery.** Consider the following program, which compiles and executes without error.

```cpp
class B {
    public:
        B() { cout << "B()" << endl; }
        void foo() { cout << "B::foo" << endl; }
        virtual void bar() { foo(); cout << "B::bar" << endl; }
        virtual void baz() { bar(); cout << "B::baz" << endl; }
};

class Bob : public B {
    public:
        Bob() { cout << "Bob()" << endl; }
        void foo() { cout << "Bob::foo" << endl; }
        virtual void baz() { cout << "Bob::baz" << endl; }
};

int main() {
    B* m = new B();
    m->baz();
    Bob* n = new Bob();
    n->baz();
    B* p = (B*) n;
    p->bar();
    p->foo();

    return 0;
}
```

In the box to the right, write the output that is produced when this program is executed.

```
Output:

B()
B::foo
B::bar
B::baz
B()
Bob()
Bob::baz
B::foo
B::bar
B::foo
```

In the box to the right, write the output that is produced when this program is executed.
Question 9 (8 points). Concurrency. The following code in C++ manages a chef who is cooking in a shared kitchen. The function will boil a pot for a given number of minutes. Suppose we use this code in a program with two threads (T1/T2), both of which boil many pots.

(9a) (4 points). This code is not thread-safe. Give a brief explanation of what can go wrong if the two threads are using these functions concurrently. You can write out a bad interleaving to the right of the function above, circle sections of the code that have a data race, or write a descriptive paragraph to the right of the function.

```cpp
int burnersAvailable = 4;

bool boilPot(Pot* pot,
             int minutes) {

    m.lock();
    if (burnersAvailable > 0) {
        burnersAvailable--;
        m.unlock();
        placeOnBurner(pot);
        sleep(minutes);
        removeFromBurner(pot);
        m.lock();
        burnersAvailable++;
        m.unlock();
        return true;
    }
    m.unlock();
    return false;
}
```

(9a)

// Bad interleaving:
if (burnersAvailable > 0) {
    burnersAvailable--;
    ...
    // burnersAvailable can go < 0
    // if original burnersAvailable
    // was 1.
    // Data race: if one thread does
    // burnersAvailable++ while the
    // other does burnersAvailable-- or
    // burnersAvailable++, or does the
    // test for burnersAvailable > 0,
    // then there may be corruption of
    // the variable and/or the return.
}

m.unlock();
return false;
```

(9b) (4 points). We can solve this issue with locks (std::mutex). Assuming that there is a preexisting mutex called m, add lock and unlock calls (i.e., m.lock() and m.unlock()) to the function above to make the function completely thread-safe. Do not lock a larger section of code than necessary, so as to allow parallelism when possible.

You cannot lock while the sleep() is happening otherwise only one pot can boil at a time! Also must unlock before all returns.