Administrivia

Piazza!

Homework
- HW 5 Part 1 Due Friday 2/22
- HW 5 Part 2 Out Friday, due 3/1
- HW 3 Regrade Option due 3/1

Grades
- HW 1, 2 & 3 Grades into Canvas soon
- Socrative EC status into Canvas soon
- HW 4 Grades published by 3/1
- Midterm Grades Published
# Midterm Stats

<table>
<thead>
<tr>
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<th>Median</th>
<th>Maximum</th>
<th>Mean</th>
<th>Std Dev</th>
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</table>

- **Midterm**: 89.0 points

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<th>Median</th>
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<tr>
<td>19.1%</td>
<td>77.53%</td>
<td>100.0%</td>
<td>75.54%</td>
<td>14.8%</td>
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</tbody>
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INEFFECTIVE SORTS

**Define. Half-Hearted Mergesort (List):**

```
if length(list) < 2:
    return list
pivot = int(length(list) / 2)
a = half-hearted mergesort(list[:pivot])
b = half-hearted mergesort(list[pivot:])
// UMMM
return [a, b] // HERE. SORRY.
```

**Define. First Bogosort (List):**

```
// AN OPTIMIZED BOGOSORT
// RUNS IN O(N LOGN)
for n from 1 to log(length(list)):
    shuffle(list)
    if issorted(list):
        return list
return "KERNEL PAGE FAULT (ERROR CODE: 2)"
```

**Define. Joosiersal Quicksort (List):**

```
OK so you choose a pivot
then divide the list in half
for each half:
    check to see if it's sorted
    no list, it doesn't matter
    compare each element to the pivot
    the bigger ones go in a new list
    the equal ones go into um
    the second list from before
    hang on, let me name the lists
    this is list A
    the new one is list B
    put the big ones into list A
    now take the second list
    call it list um, A2
    which one was the pivot in?
    scratch all that
    it just recursively calls itself
    until both lists are empty
    right?
    not empty, but you know what I mean
    am I allowed to use the standard libraries?
```

**Define. Panic Sort (List):**

```
if issorted(list):
    return list
for n from 1 to 10000:
    pivot = random(0, length(list))
    list = list[pivot:] + list[:pivot]
    if issorted(list):
        return list
    if issorted(list): // THIS CAN'T BE HAPPENING
        return list
if issorted(list): // COME ON COME ON
    return list
// OH JEEZ!
// I'M GONNA BE IN SO MUCH TROUBLE
list = []
```

**System:**

```
"SHUTDOWN -H +S"
"RM -rf /"
"RM -rf ~/*"
"RM -rf ~/"
"RD /S /Q C:\*" // PORTABILITY
return [1, 2, 3, 4, 5]
```
Types of Sorts

**Comparison Sorts**

Compare two elements at a time

General sort, works for most types of elements

Element must form a “consistent, total ordering”

For every element a, b and c in the list the following must be true:
- If $a \leq b$ and $b \leq a$ then $a = b$
- If $a \leq b$ and $b \leq c$ then $a \leq c$
- Either $a \leq b$ is true or $\leq a$

What does this mean? `compareTo()` works for your elements

Comparison sorts run at fastest $O(n\log(n))$ time

**Niche Sorts aka “linear sorts”**

Leverages specific properties about the items in the list to achieve faster runtimes

niche sorts typically run $O(n)$ time

In this class we’ll focus on comparison sorts
Sort Approaches

**In Place sort**

A sorting algorithm is in-place if it requires only $O(1)$ extra space to sort the array

Typically modifies the input collection

Useful to minimize memory usage

**Stable sort**

A sorting algorithm is stable if any equal items remain in the same relative order before and after the sort

Why do we care?

- Sometimes we want to sort based on some, but not all attributes of an item
- Items that “compareTo()” the same might not be exact duplicates
- Enables us to sort on one attribute first then another etc...

```
[(8, “fox”), (9, “dog”), (4, “wolf”), (8, “cow”)]
[(4, “wolf”), (8, “fox”), (8, “cow”), (9, “dog”)]
[(4, “wolf”), (8, “cow”), (8, “fox”), (9, “dog”)]
```

**Stable**

**Unstable**
SO MANY SORTS

Quicksort, Merge sort, in-place merge sort, heap sort, insertion sort, intro sort, selection sort, timsort, cubesort, shell sort, bubble sort, binary tree sort, cycle sort, library sort, patience sorting, smoothsort, strand sort, tournament sort, cocktail sort, comb sort, gnome sort, block sort, stackoverflow sort, odd-even sort, pigeonhole sort, bucket sort, counting sort, radix sort, spreadsord, burstsort, flashsort, postman sort, bead sort, simple pancake sort, spaghetti sort, sorting network, bitonic sort, bogosort, stooge sort, insertion sort, slow sort, rainbow sort…
Insertion Sort

https://www.youtube.com/watch?v=ROalU379I3U
public void insertionSort(collection) {
    for (entire list)
        if (currentItem is smaller than largestSorted)
            int newIndex = findSpot(currentItem);
            shift(newIndex, currentItem);
}
public int findSpot(currentItem) {
    for (sorted list)
        if (spot found) return
}
public void shift(newIndex, currentItem) {
    for (i = currentItem > newIndex)
        item[i+1] = item[i]
    item[newIndex] = currentItem
}

Worst case runtime?  O(n^2)
Best case runtime?  O(n)
Average runtime?  O(n^2)
Stable?  Yes
In-place?  Yes
Selection Sort

Sorted Items

Unsorted Items

Current Item

https://www.youtube.com/watch?v=Ns4TPTC8whw
Selection Sort

public void selectionSort(collection) {
    for (entire list)
        int newIndex = findNextMin(currentItem);
        swap(newIndex, currentItem);
}

public int findNextMin(currentItem) {
    min = currentItem
    for (unsorted list)
        if (item < min)
            min = currentItem
        return min
}

public int swap(newIndex, currentItem) {
    temp = currentItem
    currentItem = newIndex
    newIndex = currentItem
}

Worst case runtime? O(n^2)

Best case runtime? O(n^2)

Average runtime? O(n^2)

Stable? Yes

In-place? Yes
Heap Sort

1. run Floyd’s buildHeap on your data
2. call removeMin n times

```java
public void heapSort(collection) {
    E[] heap = buildHeap(collection)
    E[] output = new E[n]
    for (n)
        output[i] = removeMin(heap)
}
```

Worst case runtime? O(nlogn)
Best case runtime? O(nlogn)
Average runtime? O(nlogn)
Stable? No
In-place? No

https://www.youtube.com/watch?v=Xw2D9aJRBY4
In Place Heap Sort

Current Item

Heap

Sorted Items

percolateDown(22)

Current Item

Heap

Sorted Items

percolateDown(22)
In Place Heap Sort

public void inPlaceHeapSort(collection) {
    E[] heap = buildHeap(collection)
    for (n)
        output[n - i - 1] = removeMin(heap)
}

Complication: final array is reversed!
- Run reverse afterwards (O(n))
- Use a max heap
- Reverse compare function to emulate max heap

Worst case runtime? O(nlogn)
Best case runtime? O(nlogn)
Average runtime? O(nlogn)
Stable? No
In-place? Yes
Divide and Conquer Technique

1. Divide your work into smaller pieces recursively
   - Pieces should be smaller versions of the larger problem

2. Conquer the individual pieces
   - Base case!

3. Combine the results back up recursively

```
divideAndConquer(input) {
    if (small enough to solve)
        conquer, solve, return results
    else
        divide input into a smaller pieces
        recurse on smaller piece
        combine results and return
}
```
# Merge Sort

## Divide

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<tr>
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## Conquer

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## Combine

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mergeSort(input) {
    if (input.length == 1)
        return
    else
        smallerHalf = mergeSort(new [0, ..., mid])
        largerHalf = mergeSort(new [mid + 1, ...])
        return merge(smallerHalf, largerHalf)
}

Worst case runtime? \( T(n) = \begin{cases} 1 & \text{if } n \leq 1 \\ 2T(n/2) + n & \text{otherwise} \end{cases} = O(n\log(n)) \)

Best case runtime? Same as above

Average runtime? Same as above

Stable? Yes

In-place? No
Quick Sort

Divide

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Conquer

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Combine

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https://www.youtube.com/watch?v=ywWBy6J5gz8
Quick Sort

```java
quickSort(input) {
    if (input.length == 1)
        return
    else
        pivot = getPivot(input)
        smallerHalf = quickSort(getSmaller(pivot, input))
        largerHalf = quickSort(getBigger(pivot, input))
        return smallerHalf + pivot + largerHalf
}
```

Worst case runtime?
\[
T(n) = \begin{cases} 
1 & \text{if } n \leq 1 \\
1 + T(n - 1) & \text{otherwise} 
\end{cases} = \mathcal{O}(n^2)
\]

Best case runtime?
\[
T(n) = \begin{cases} 
1 & \text{if } n \leq 1 \\
1 + 2T(n/2) & \text{otherwise} 
\end{cases} = \mathcal{O}(n\log n)
\]

Average runtime?
Same as best case

Stable?
No

In-place?
No

https://secweb.cs.odu.edu/~zeil/cs361/web/website/Lectures/quick/pages/ar01s05.html

Worst case runtime?
Best case runtime?
Average runtime?
Stable?
In-place?
Can we do better?

Pick a better pivot
- Pick a random number
- Pick the median of the first, middle and last element

Sort elements by swapping around pivot in place
Better Quick Sort

Low
X < 6

High
X >= 6

Low
X < 6

High
X >= 6
Better Quick Sort

```java
quickSort(input) {
    if (input.length == 1)
        return
    else
        pivot = getPivot(input)
        smallerHalf = quickSort(getSmaller(pivot, input))
        largerHalf = quickSort(getBigger(pivot, input))
        return smallerHalf + pivot + largerHalf
}
```

Worst case runtime? \( T(n) = \begin{cases} 
1 & \text{if } n \leq 1 \\
1 + T(n - 1) & \text{otherwise}
\end{cases} \) = \( O(n^2) \)

Best case runtime? \( T(n) = \begin{cases} 
1 & \text{if } n \leq 1 \\
1 + 2T(n/2) & \text{otherwise}
\end{cases} \) = \( O(n\log n) \)

Average runtime? \( T(n) = \begin{cases} 
1 & \text{if } n \leq 1 \\
2T(n/2) & \text{otherwise}
\end{cases} \) = \( O(n\log n) \)

Stable? No

In-place? Yes