1. True or False

For each statement below, state that its true or false. Explain.

- (a) Binary search in a sorted array is O(n).
- (b) Master Theorem can always be used to find the big-Theta of any recurrence.
- (c) Any Java object can be hashed.
- (d) The number of collisions in a hash table is dependent on the table capacity and the hash function.
- (e) Heaps can be implemented as an array with no gaps at each index.
- (f) Binary Search Trees can be implemented as an array with no gaps at each index.
- (g) Binary Search Trees have a better big-O runtime comepared to AVL trees because AVL trees are bushier.
- (h) AVL trees and BSTs have the same big-O runtime for inserting elements.
- (i) 3-Heaps have better tight big-O runtime for insert than 2-heaps or 4 heaps.
- (j) AVL trees, BSTs, and Heaps are all DAGs.
- (k) For a graph with negative edge weights, adding a constant to make all edge weights positive and then running Dijkstra's will give the same shortest path from a node s to g as the one with negative edge weights.
- (l) MSTs cannot be used to find the shortest path between two nodes.
- (m) Graphs can be implemented with a Dictionary.

2. ADTs and Implementations

For each of the following ADTs, list the implementations that we covered in class and an advantage of each other the others. (Hint: in what situations would you use each one?)

(a) List

(b) Dictionary/Map

3. Sorting Design Question

For each situation below, choose an appropriate sorting algorithm from the list below. Explain your answer.

Selection Sort, Insertion Sort, Merge Sort, Quick Sort, Heap Sort

(a) At class pictures the photographers usually sort the children by height before assigning places. The photographer doesn't care who was originally where in line but he/she does need to get them sorted as fast as possible and there is a lot of extra space in the room to arrange them.

(b) A company has lists of numbers that each need to be sorted. Because the lists can be long, short, reversed, in-order or a mix of all these situations, they need a sort that will always have the same tight big-O runtime regardless of the condition of the list.

(c) When you were younger, you most likely had a box of Crayola crayons. When you first buy them however, all the colors are not sorted in order. Which sort can you use to sort the crayons by color in the box such that you only take one or two of the crayons out of the box at any given time?

(d) A customer requires that you use a sort that has recursion. Which one can you choose? Explain why this request is not a good idea.

```
public class Friend implements Comparable<Friend> {
2
3
           private String name;
 4
           public Friend(String name) {
               this.name = name;
6
           }
           // assume this method runs in O(name.length()) runtime
9
           public boolean equals(Object o) {
               if (!(o instanceof Friend)) {
                   return false;
               }
13
14
               Friend other = (Friend) o;
16
               return this.name.equals(other.name);
           }
17
18
           public int hashCode() {
19
               boolean hasLowercase = false;
20
               for (int i = 0; i < name.length(); i++) {</pre>
21
                   if (name.charAt(i) >= 'a') {
22
                       return 0;
23
                   }
24
               }
25
               return 1;
26
           }
28
29
           public int compareTo(Friend other) {
30
               int sizeThis = this.name.length();
31
               int sizeOther = other.name.length();
               int index = 0;
               while (index < sizeThis && index < sizeOther) {</pre>
34
                   if (this.name.charAt(index) < other.name.charAt(index)) {</pre>
35
                       return -1;
36
                   } else if (this.name.charAt(index) > other.name.charAt(index)) {
37
                       return 1;
38
                   } else {
                       index++;
40
                   }
41
               }
42
               return sizeThis - sizeOther;
43
44
           }
45
46
           public String toString() {
47
               return this.name;
48
           }
49
       }
```

Warm-up:

4.

- (a) Give the best case runtime and the worst case runtime for the hashCode method for the Friend class as a simplified Tight O bound, where m is the length of the Friend's name. The .length() and .charAt method run in constant time.
- (b) Give the best case runtime and the worst case runtime for the compareTo method for the Friend class as a simplified Tight O bound, where m1 is the length of this Friend's name and m2 is the length of the other Friend's name. The .length() and .charAt method run in constant time.

```
public static IDictionary<Friend, IList<Friend>> simulateFriendshipGraph(DoubleLinkedList<String> names,
            int numFriends) {
           IDictionary<Friend, IList<Friend>> graph = new ChainedHashDictionary<>();
2
           IDictionary<String, Friend> stringToFriend = new ChainedHashDictionary<>();
3
 4
           for (int i = 0; i < names.size(); i++) {</pre>
               String name = names.get(i);
6
               Friend friend = new Friend(name);
 7
               stringToFriend.put(name, friend);
           }
 9
10
           Iterator<String> itr = names.iterator();
           for (int i = 0; i < names.size(); i++) {</pre>
               Friend friend = stringToFriend.get(itr.next());
13
               IList<Friend> neighbors = new DoubleLinkedList<>();
14
               graph.put(friend, neighbors);
           }
16
           for (KVPair<String, Friend> pair : stringToFriend) {
18
               String originalName = pair.getKey();
19
               Friend friend = pair.getValue();
20
               names.delete(names.indexOf(originalName));
21
               IList<Friend> neighbors = graph.get(friend);
22
               IList<String> removed = new DoubleLinkedList<>();
23
               for (int j = 0; j < numFriends; j++) {</pre>
24
                   // generates a random number from 0 to names.size() - 1 inclusive on both ends in constant time
25
                   int randomIndex = new Random(10).nextInt(names.size());
26
27
                   String name = names.get(randomIndex);
28
                   neighbors.add(stringToFriend.get(name));
29
                   removed.add(name);
                   names.delete(randomIndex);
31
               }
               while (!removed.isEmpty()) {
34
                   names.add(removed.remove());
35
               }
36
               names.add(originalName);
37
38
39
           }
40
41
           return graph;
       }
42
```

You can assume that line 19, 20, and 26 all run in constant time. Assume that each String in the names parameter is of length m. You can assume that numFriends will be less than names.size() so the code works. Also, the hashCode method for String objects will run in m time, so you will want to account for that whenever hashCode is called. All the runtimes should be in terms of constants, m, numFriends, and n (the .size() of names).

- (a) For the for loop from lines 5 to 9
 - i. list any method calls that are n runtime or a bigger complexity class
 - ii. what is the worst case runtime (consider CHD to use separate chaining and appropriate resizing strategies) for the loop from lines 5 to 9? Give your answer as a simplified, tight big Oh bound.
- (b) For the for loop from lines 12 to 16
 - i. What is maximum number of —V—, the number of vertices in the graph?
 - ii. list any method calls that are n runtime or a bigger complexity class
 - iii. what is the worst case runtime (consider CHD to use separate chaining and appropriate resizing strategies) for the loop from lines 12 to 16? Give your answer as a simplified, tight big Oh bound.
- (c) For the loop from lines 18 to 32
 - i. list any method calls that are n runtime or a bigger complexity class
 - ii. how many times (maximum) does the loop from lines 18 to 32 run?
 - iii. how many times (maximum) is the line 31 executed?
- (d) What is the overall worst case runtime of this method? Give your answer as a simplified, tight O bound.
- (e) About the behavior of the method:
 - i. What is this code doing? (in 1 sentence)
 - ii. Is the graph directed or undirected?
 - iii. Is the graph weighted or unweighted?
 - iv. Is the names parameter mutated or not mutated after the method is over?

```
public static void printRandomReachingFriendship(IDictionary<Friend, IList<Friend>> graph) {
2
           IList<Friend> list = new DoubleLinkedList<>();
 3
           IPriorityQueue<Friend> heap = new ArrayHeap<>();
 4
           for (KVPair<Friend, IList<Friend>> pair : graph) {
               heap.add(pair.getKey());
           }
 9
           Friend start = heap.removeMin();
           ISet<Friend> included = new ChainedHashSet<>();
13
           DoubleLinkedList<Friend> toProcess = new DoubleLinkedList<>();
14
           toProcess.add(start);
16
           while (!toProcess.isEmpty()) {
               Friend next = toProcess.get(0);
               toProcess.delete(0);
18
               System.out.println("next person: " + next);
19
               for (Friend neighbor : graph.get(next)) {
20
                   if (!included.contains(neighbor)) {
21
                       included.add(neighbor);
22
                       toProcess.add(neighbor);
                   }
^{24}
                   mysterySort(toProcess, 0, toProcess.size() - 1);
25
               }
26
           }
       }
28
29
       public static void mysterySort(DoubleLinkedList<Friend> list, int low, int high) {
30
           if (low < high) {</pre>
31
              /* pi is partitioning index, list[pi] is
                now at right place */
               int pi = partition(list, low, high);
34
               // Recursively sort elements before
35
               // partition and after partition
36
               mysterySort(list, low, pi-1);
37
               mysterySort(list, pi+1, high);
38
           }
       }
40
41
       public static int partition(DoubleLinkedList<Friend> list, int low, int high) {
42
           Friend pivot = list.get(low);
43
           int i = low - 1;
44
           int j = high + 1;
45
           while (true) {
46
               do {
47
                   i++;
48
               } while (list.get(i).compareTo(pivot) < 0);</pre>
49
               do {
51
                   j--;
               } while (list.get(j).compareTo(pivot) > 0);
53
54
               if (i >= j) {
                   return j;
               }
58
               // swap list.get(i) and list.get(j)
               Friend temp = list.get(i);
60
               list.set(i, list.get(j));
61
               list.set(j, temp);
62
           }
63
       }
64
```

You can assume that System.out.println runs in constant time. Also assume this method is passed a graph that was constructed by the method on the previous page. For simplicity you can assume that compareTo and hashCode take constant time, but if you want to challenge yourself try to compute the runtimes considering what you gave earlier for the worst/base case runtimes. All your answers for this part should be defined in terms of |V|, |E|, where they represent the total number of vertices and total number of edges in the graph parameter, and numFriends.

- (a) For the loop on lines 5-7 give a simplified tight O bound for:
 - i. the best case runtime
 - ii. the worst case runtime
- (b) For the loop from lines 16 to 27:
 - i. how many times are lines 17-19 executed in the worst case? Bonus question: best case?
 - ii. how many times max will mysterySort get called in the worst case?
 - iii. how many times are lines 22 and 23 executed in the worst case?
 - iv. list any method calls that are V or E runtime or bigger complexity class
 - v. What is the worst case runtime for the loop from 16 to 28? Give your answer as a simplified tight Oh. For now, say the runtime of mysterySort is S runtime.
- (c) for mysterySort:
 - i. What sort does mysterySort look the most like?
 - ii. Imagine the first time partition is called, and the difference between low and high is the size of the list parameter, called p. What is the worst case runtime for partition in terms of p?
 - iii. What is the worst case (describe it, not the runtime) for mysterySort? Hint: model the code as a recurrence / consider what sort you think this is and the worst case for this sort?
 - iv. What is the best case (describe it, not the runtime) for mysterySort?
 - v. How many times will we recurse in the worst case of mysterySort (aka how many times do we call partition)? Assume that the size of the list parameter is called p.
 - vi. Putting that together, what's the worst case runtime for mysterySort? Assume that the size of the list parameter is called p.
- (d) Is the way mysterySort used above more like the best case or the worst case runtime situation?
- (e) What is the complexity class of the max size of toProcess?
- (f) Combine your answers above to substitute in for the S runtime you used for mysterySort earlier, and provide a new simplified, tight O bound for the runtime of printRandomReachingFriendship.

5. Topo Sort

Consider the graph below:



- (a) What are the characteristics of this graph?
- (b) Give two possible orderings produced using Topological sort.

(c) Which graph algorithm can you modify to run Topo Sort?