

CSE 373: Graph traversal

Michael Lee

Friday, Feb 16, 2018

Goal: How do we traverse graphs?

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- ▶ Traverse a graph to find if there's a connection from one node to another?

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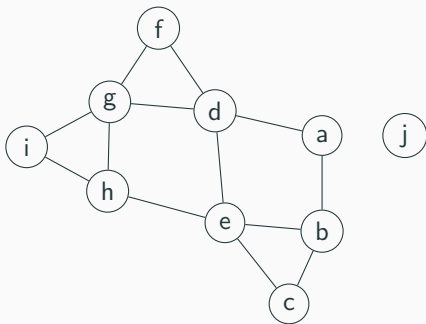
For example, can we...

- ▶ Traverse a graph to find if there's a connection from one node to another?
- ▶ Determine if we can start from our node and touch every other node?
- ▶ Find the shortest path between two nodes?

Solution: Use graph traversal algorithms like breadth-first search and depth-first search

Breadth-first search (BFS) example

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search(v):  
    visited = empty set  
  
    queue.enqueue(v)  
    visited.add(v)  
  
    while (queue is not empty):  
        curr = queue.dequeue()  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
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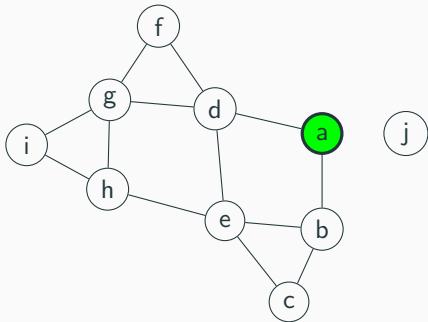
Current node:

Queue: a,

Visited: a,

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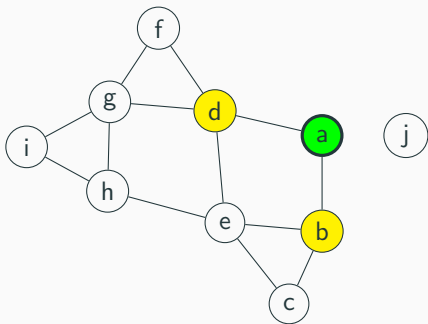
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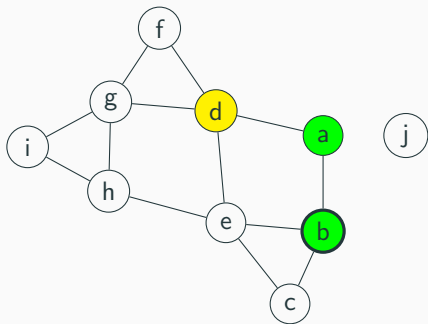
Current node: a

Queue: b, d,

Visited: a, b, d,

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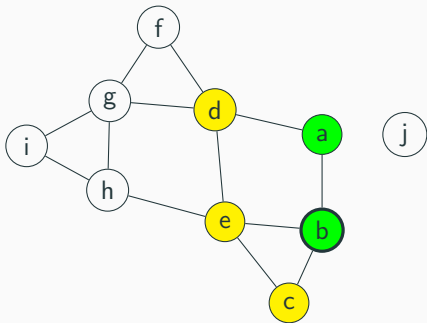
Current node: b

Queue: d,

Visited: a, b, d,

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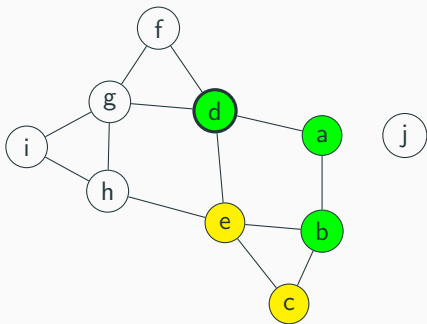
Current node: b

Queue: d, c, e,

Visited: a, b, d, c, e,

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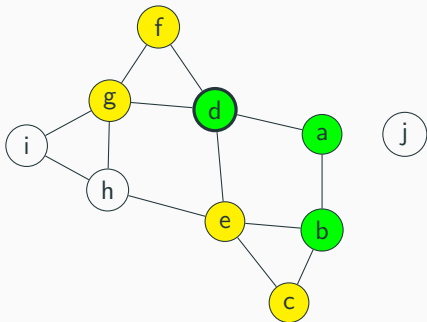
Current node: d

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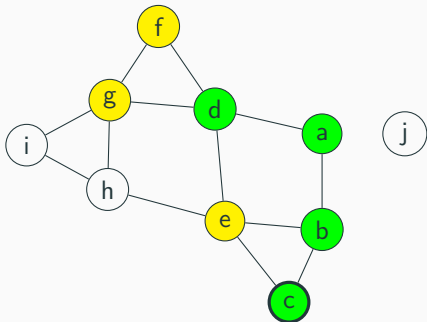
Current node: d

Queue: c, e, f, g,

Visited: a, b, d, c, e, f, g,

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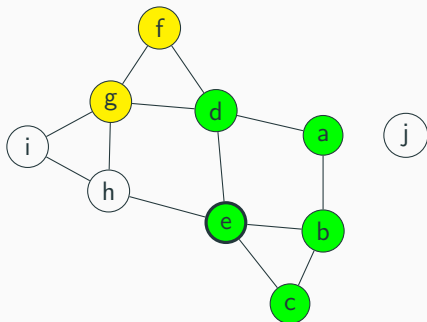
Current node: c

Queue: e, f, g,

Visited: a, b, d, c, e, f, g,

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Current node: e

Queue: f, g,

Visited: a, b, d, c, e, f, g,

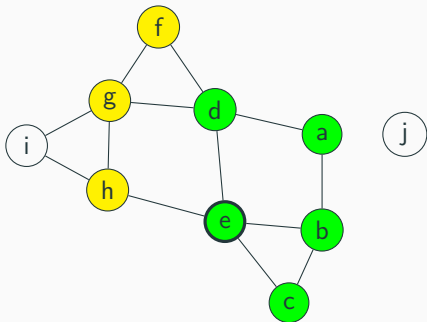
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```

Current node: e

Queue: f, g, h,

Visited: a, b, d, c, e, f, g, h,



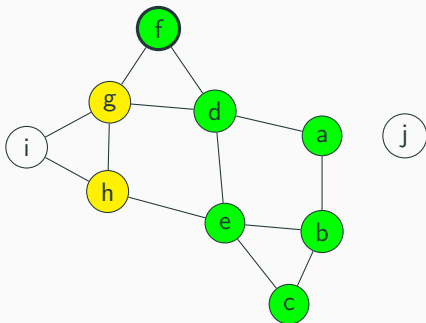
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```

Current node: f

Queue: g, h,

Visited: a, b, d, c, e, f, g, h,



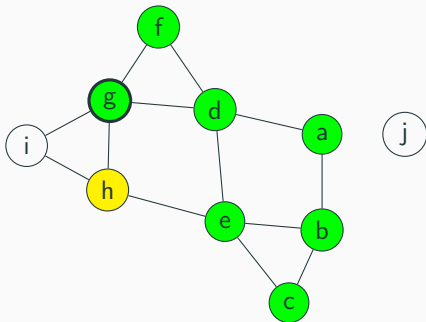
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Current node: g

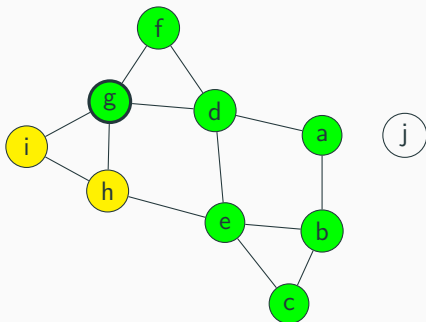
Queue: h,

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Current node: g

Queue: h, i,

Visited: a, b, d, c, e, f, g, h, i,

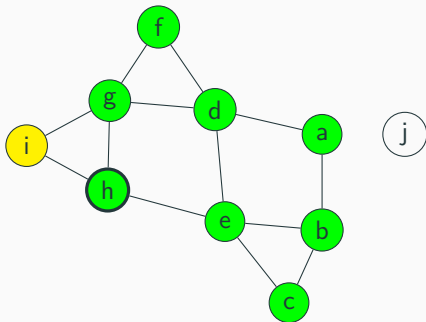
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Current node: h

Queue: i,

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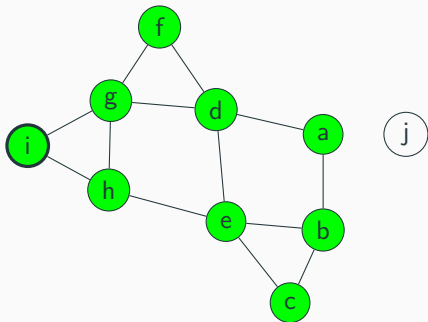
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So, $\mathcal{O}(|V| + 2|E|)$, which simplifies to $\mathcal{O}(|V| + |E|)$.

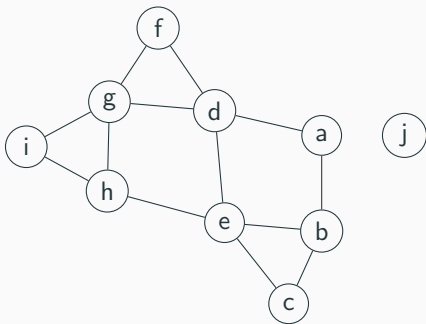
Breadth-first search (BFS)

Pseudocode:

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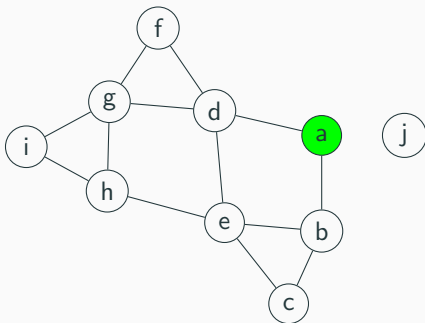
An interesting property...

Note: We visited the nodes in “rings” – maintained a gradually growing “frontier” of nodes.



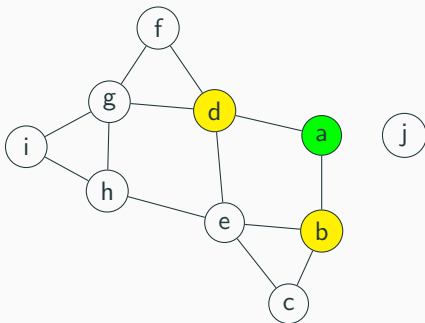
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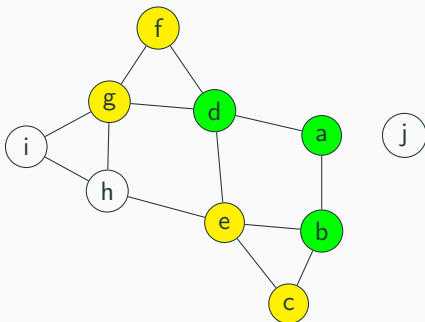
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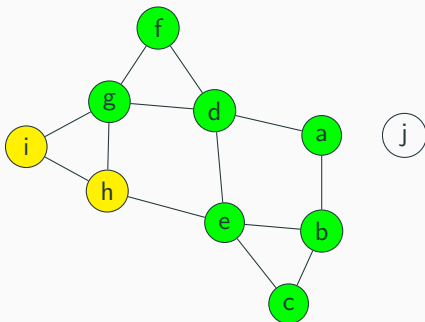
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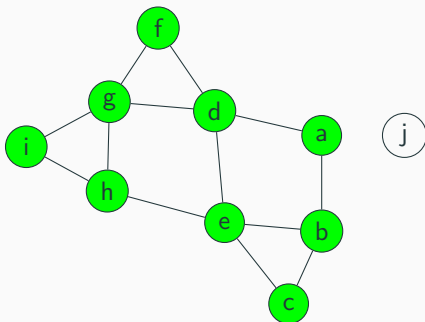
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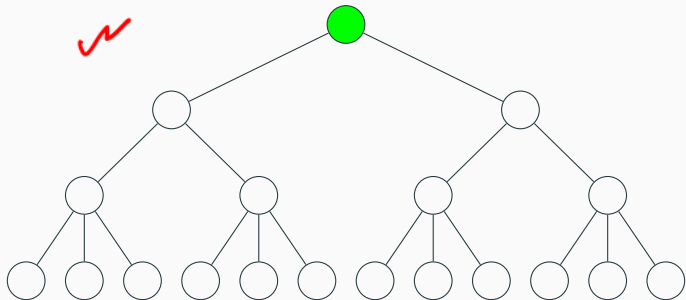
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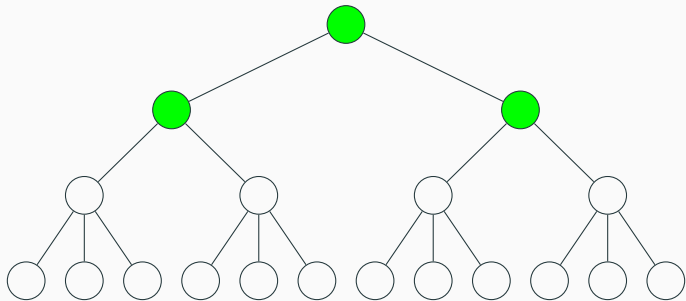
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What does this look like for trees?



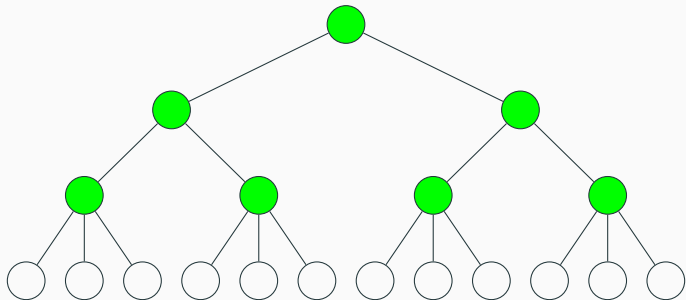
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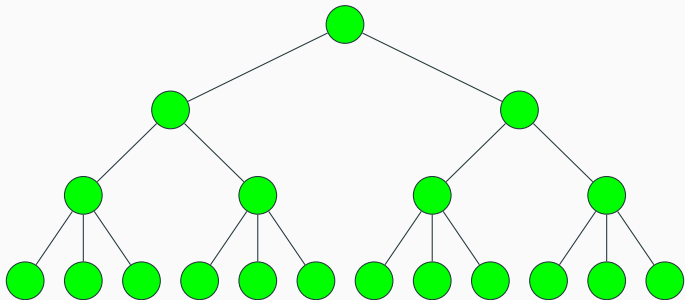
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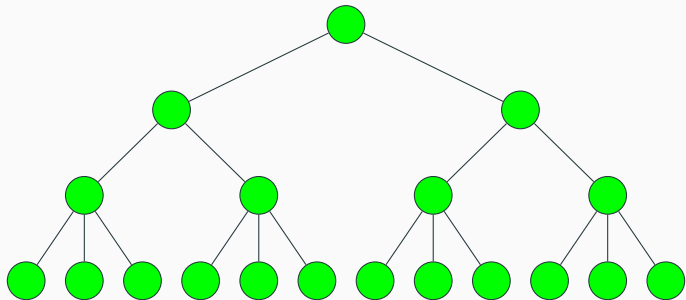
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The algorithm traverses the width, or “breadth” of the tree

Depth-first search (DFS)

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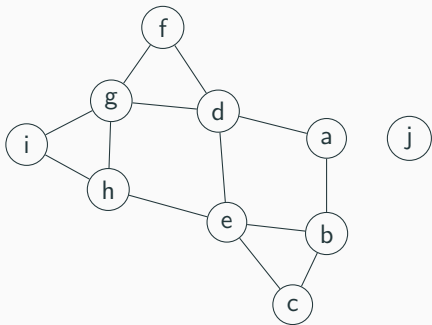
```
search(v):
    visited = empty set
    stack.push(v)
    visited.add(v)

    while (stack is not empty):
        curr = stack.pop()
        visited.add(curr)

        for (w : v.neighbors()):
            if (w not in visited):
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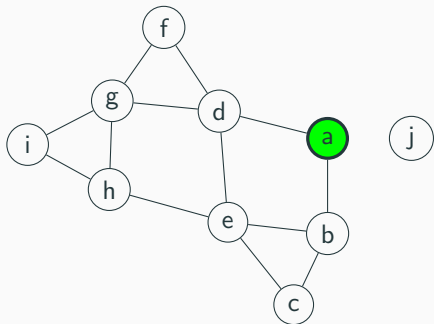
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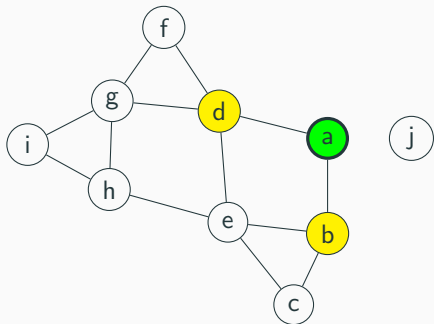
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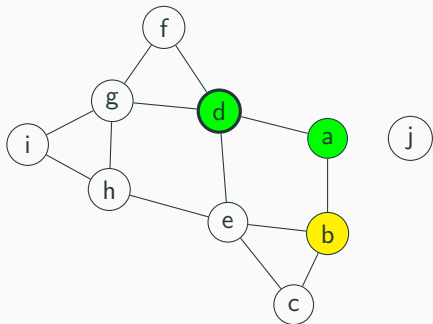
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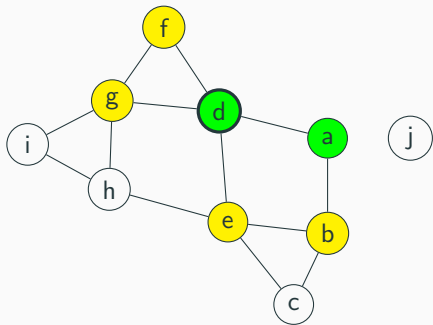
Current node: d

Stack: b,

Visited: a, b, d,

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                stack.push(w)
```



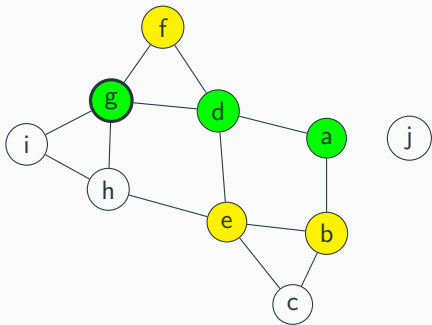
Current node: d

Stack: b, e, f, g,

Visited: a, b, d, e, f, g,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



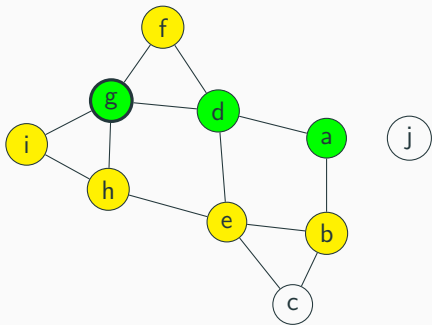
Current node: g

Stack: b, e, f,

Visited: a, b, d, e, f, g,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



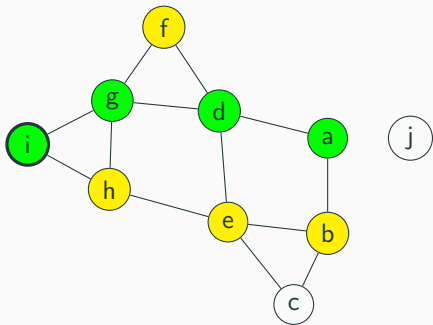
Current node: g

Stack: b, e, f, h, i,

Visited: a, b, d, e, f, g, h, i,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



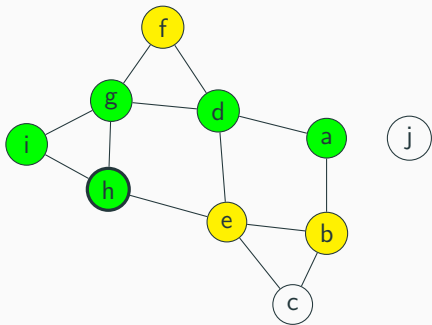
Current node: i

Stack: b, e, f, h,

Visited: a, b, d, e, f, g, h, i,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
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    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



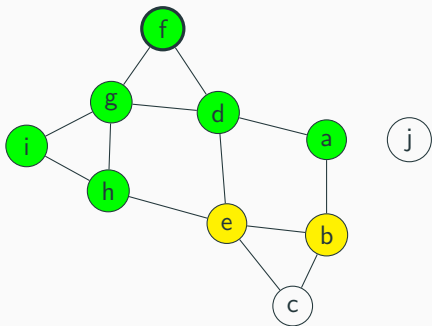
Current node: h

Stack: b, e, f,

Visited: a, b, d, e, f, g, h, i,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



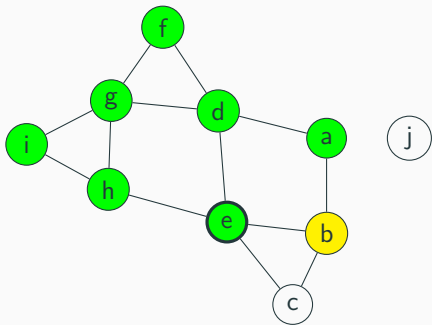
Current node: f

Stack: b, e,

Visited: a, b, d, e, f, g, h, i, e,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



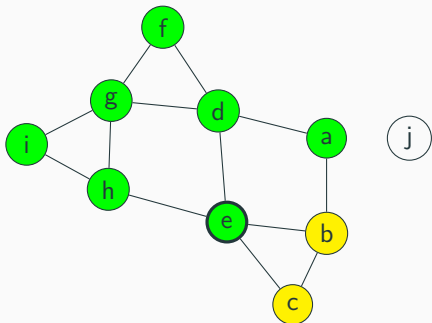
Current node: e

Stack: b, e,

Visited: a, b, d, e, f, g, h, i, e,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
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    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



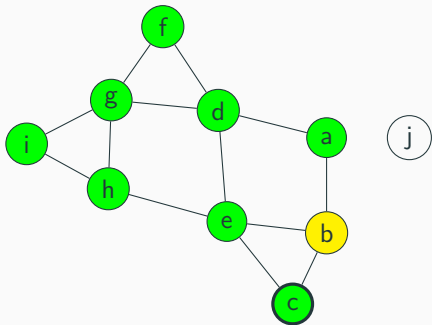
Current node: e

Stack: b, e, c,

Visited: a, b, d, e, f, g, h, i, e, c,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
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    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



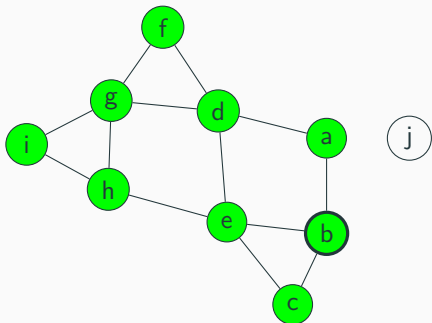
Current node: c

Stack: b,

Visited: a, b, d, e, f, g, h, i, e, c,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
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    while (stack is not empty):  
        curr = stack.pop()  
        visited.add(curr)  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



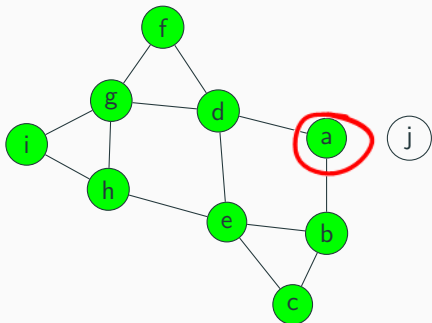
Current node: b

Stack:

Visited: a, b, d, e, f, g, h, i, e, c,

Depth-first search (DFS) example

```
search(v):  
    visited = empty set  
    stack.push(v)  
  
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        curr = stack.pop()  
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        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)
```



Current node:

Stack:

Visited: a, b, d, e, f, g, h, i, e, c,

Depth-first search (DFS)

Depth-first traversal, core idea:

1. Instead of using a queue, use a stack. Otherwise, keep everything the same.

Depth-first search (DFS)

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Depth-first search (DFS)

Depth-first traversal, core idea:

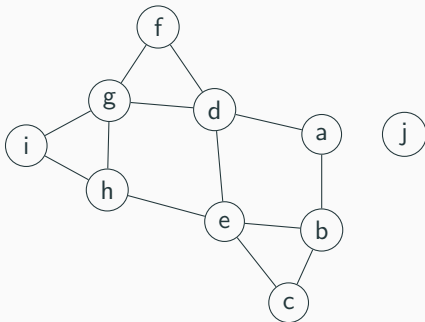
1. Instead of using a queue, use a stack. Otherwise, keep everything the same.
2. Runtime: also $\mathcal{O}(|V| + |E|)$ for same reasons as BFS

Pseudocode:

```
search(v):  
    visited = empty set  
  
    stack.push(v)  
    visited.add(v)  
  
    while (stack is not empty):  
        curr = stack.pop()  
  
        for (w : v.neighbors()):  
            if (w not in visited):  
                stack.push(w)  
                visited.add(curr)
```

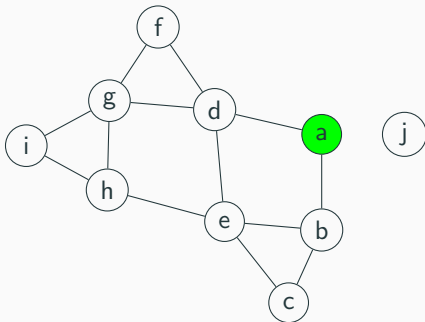

An interesting property...

Note: Rather than growing the node in “rings”, we randomly wandered through the graph until we got stuck, then “backtracked”.



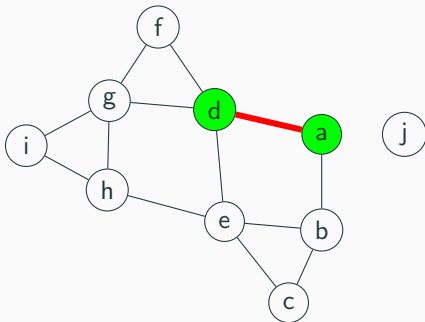
An interesting property...

Note: Rather than growing the node in “rings”, we randomly wandered through the graph until we got stuck, then “backtracked”.



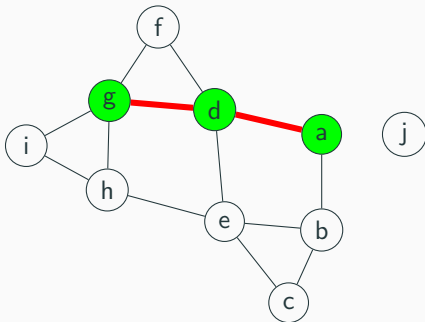
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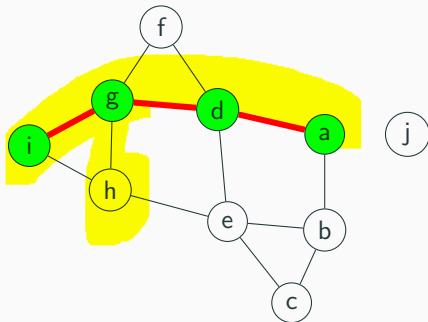
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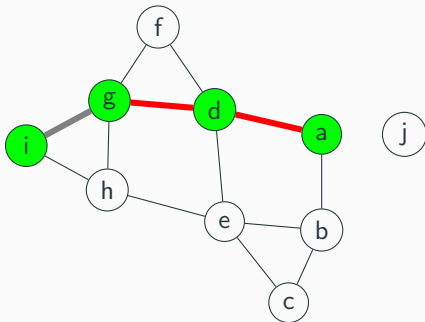
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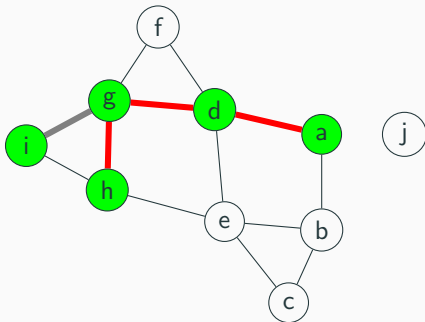
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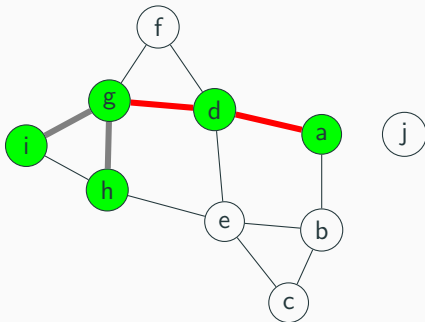
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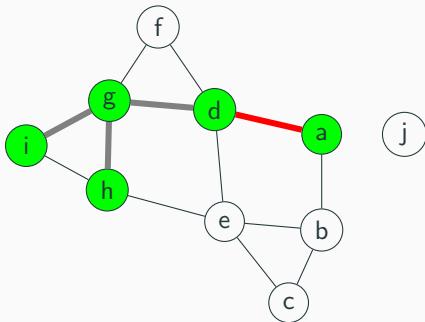
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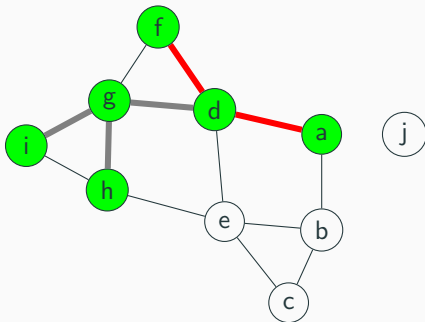
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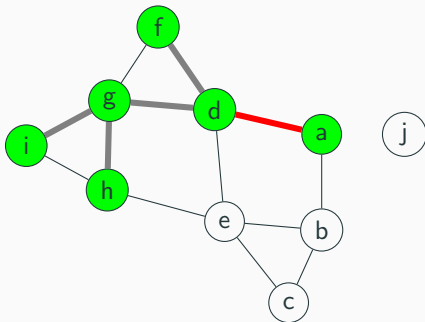
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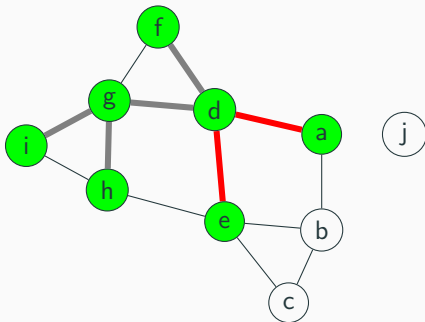
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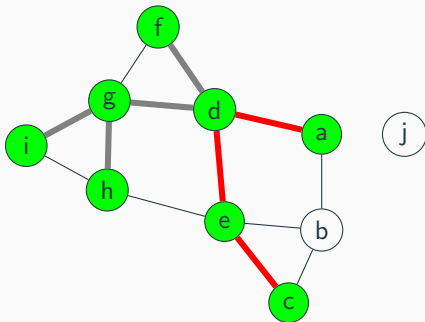
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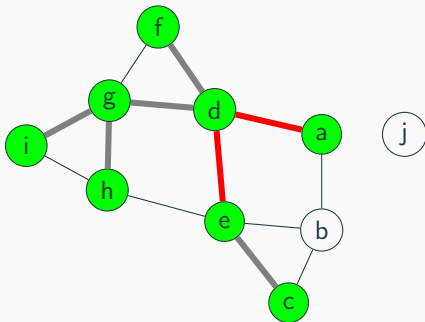
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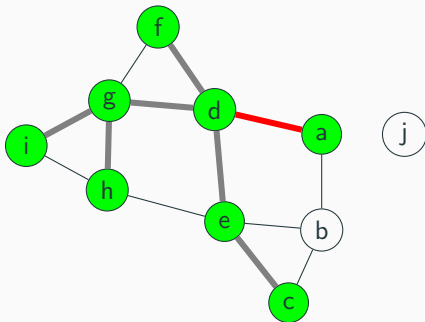
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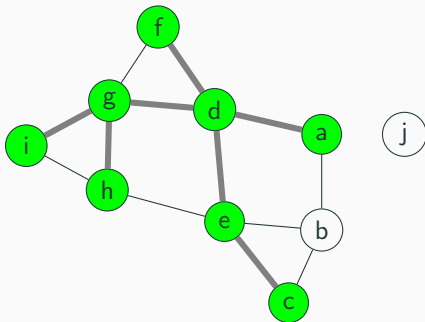
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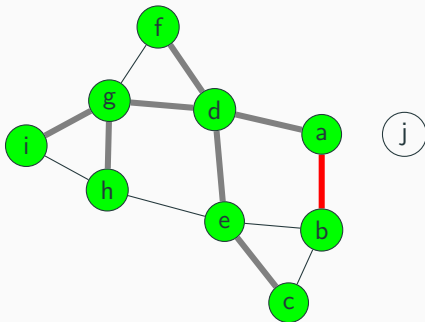
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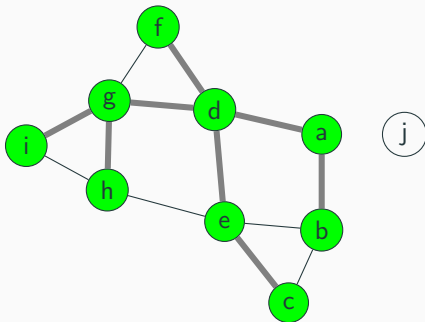
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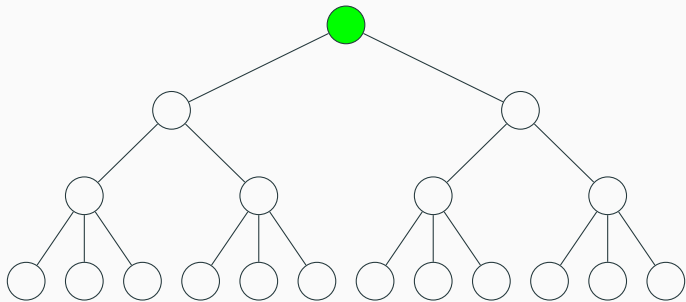
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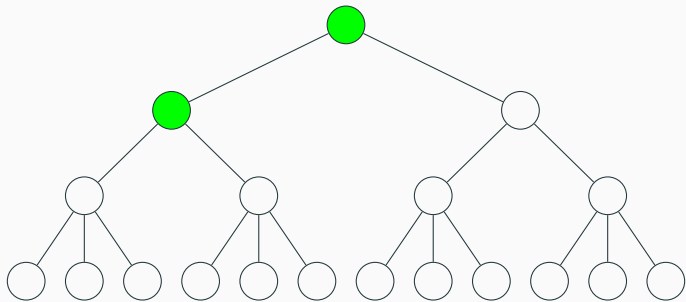
An interesting property...

What does this look like for trees?



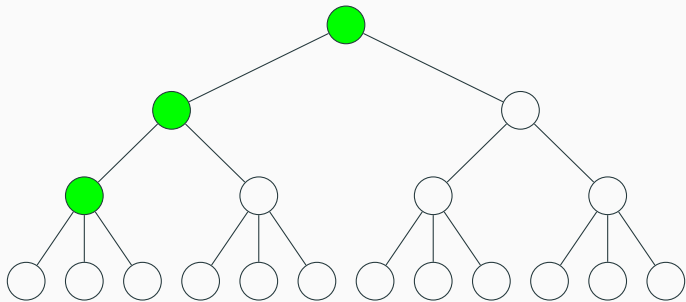
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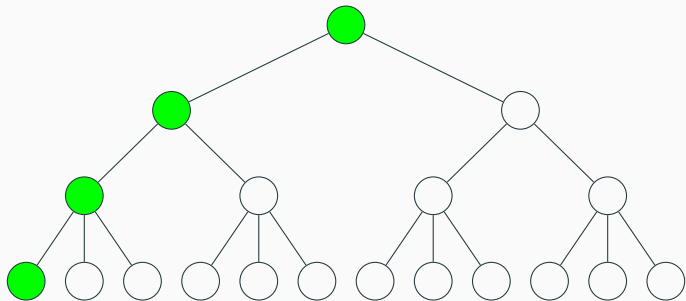
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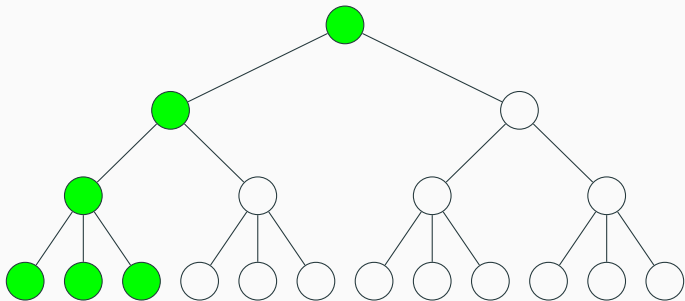
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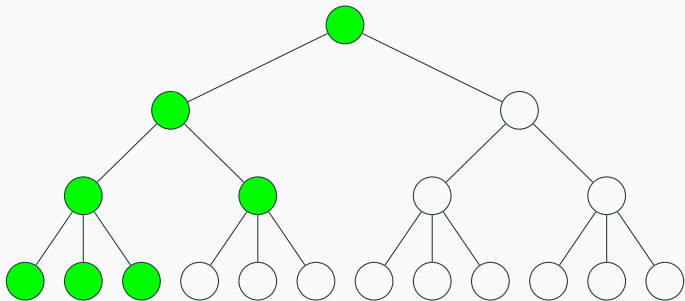
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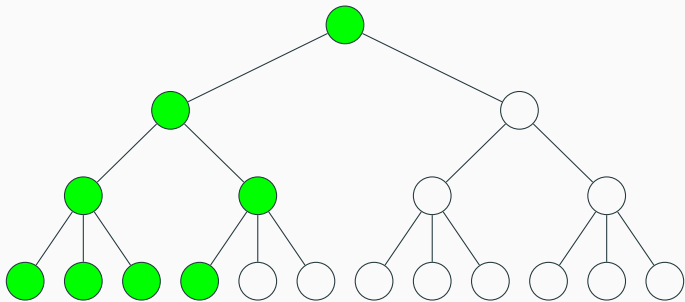
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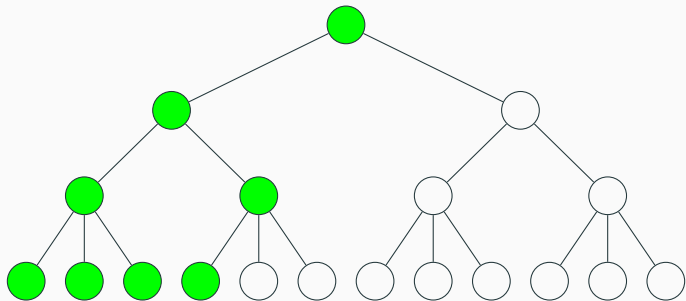
An interesting property...

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An interesting property...

What does this look like for trees?



The algorithm traverses to the bottom first: it prioritizes the “depth” of the tree

Note: rest of algorithm omitted

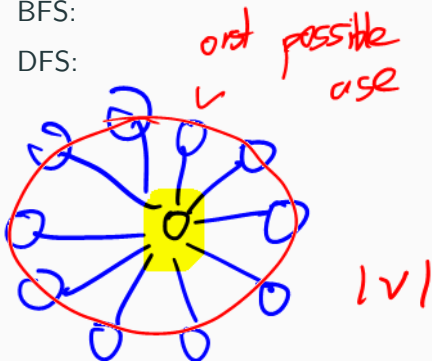
Question: When do we use BFS vs DFS?

Compare and contrast

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Related question: How much memory does BFS and DFS use in the worst case?

- ▶ BFS:
- ▶ DFS:



Compare and contrast

Question: When do we use BFS vs DFS?

Related question: How much memory does BFS and DFS use in the worst case?

- ▶ BFS: $\mathcal{O}(|V|)$ – what if every node is connected to the start?
- ▶ DFS: $\mathcal{O}(|V|)$

So, in the worst case, BFS and DFS both have the same worst-case runtime and memory usage.

They only differ in what order they visit the nodes.

Compare and contrast

How much memory does BFS and DFS use in the **average** case?

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Related question: how much memory do they use when we want to traverse a tree?

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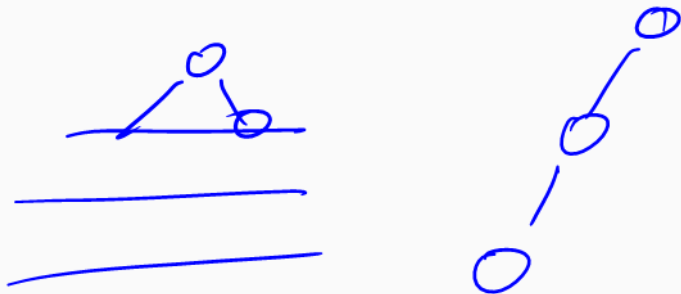
- ▶ BFS:
- ▶ DFS:

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For graphs:

- ▶ Use BFS if graph is “narrow”, or if solution is “near” start
- ▶ Use DFS if graph is “wide”

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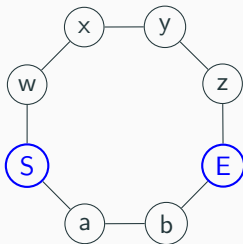
For graphs:

- ▶ Use BFS if graph is “narrow”, or if solution is “near” start
- ▶ Use DFS if graph is “wide”

In practice, graphs are often large/very wide, so DFS is often a good default choice. (It’s also possible to implement DFS recursively!)

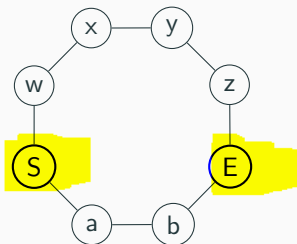
Design challenge

Question: How would you modify BFS to find the shortest path between every node?



Design challenge

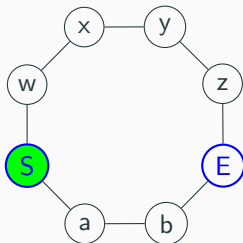
Question: How would you modify BFS to find the shortest path between every node?



Observation: Since BFS moves out in rings, we will reach the end node via the path of length 3 first.

Design challenge

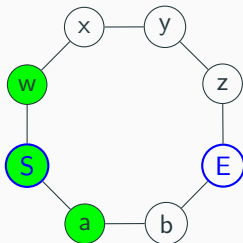
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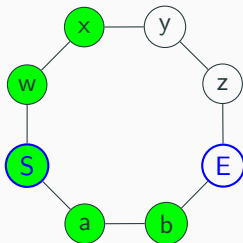
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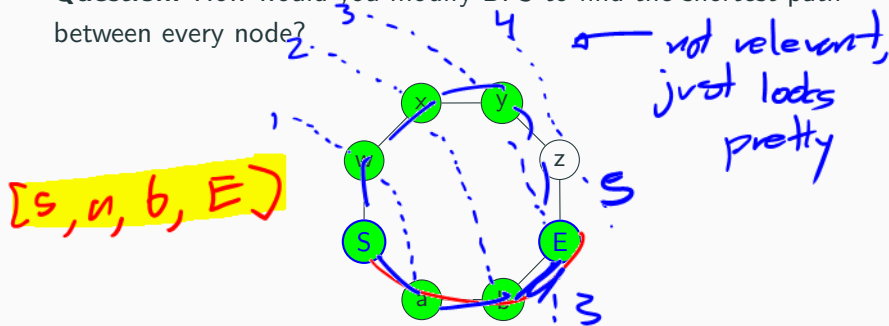
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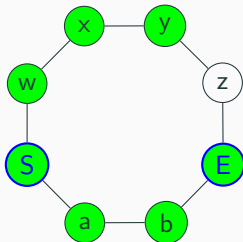
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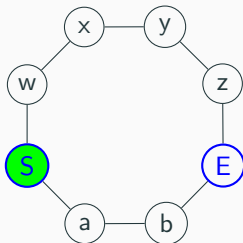
Observation: Since BFS moves out in rings, we will reach the end node via the path of length 3 first.

Idea: when we enqueue, store where we came from in some way. (e.g. mark node, use a dictionary...)

After BFS is done, *backtrack*.

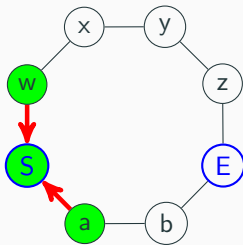
Design challenge: pathfinding

Question: How would you modify BFS to find the shortest path between every node?



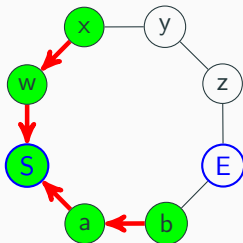
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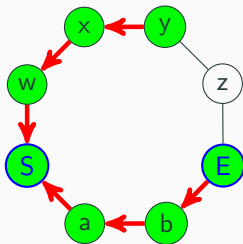
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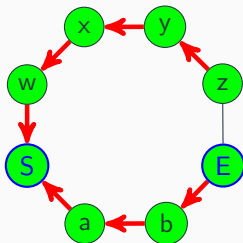
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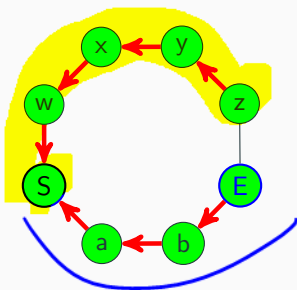
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Design challenge: pathfinding

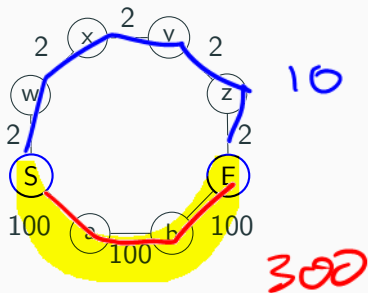
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Now, start from any node, follow arrows, then reverse to get path.

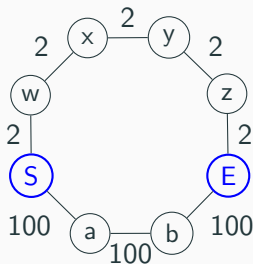
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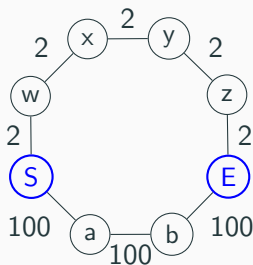


Weighted graph

A **weighted graph** is a kind of graph where each edge has a numerical “weight” associated with it.

Design challenge: pathfinding

Question: What if the edges have weights?



Weighted graph

A **weighted graph** is a kind of graph where each edge has a numerical “weight” associated with it.

This number can represent anything, but is often (but not always!) used to indicate the “cost” of traveling down that edge.

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We need a better algorithm.

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Today: **Dijkstra's algorithm**

Dijkstra's algorithm

Core idea:

1. Assign each node an initial cost of ∞

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Metaphor: Treat edges as canals and edge weights as distance. Imagine opening a dam at the starting node. How long does it take for the water to reach each vertex?

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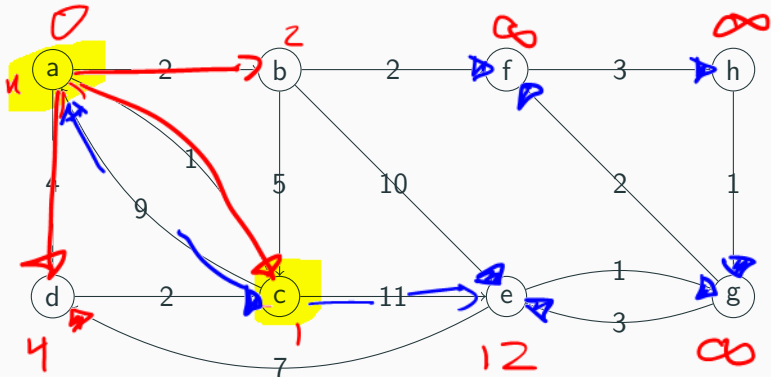
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Pronunciation: DYKE-struh ("dijk" rhymes with "bike")

Dijkstra's algorithm

Suppose we start at vertex "a":



Dijkstra's algorithm

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 - ▶ Could use a heap!
- ▶ If we're using a heap, how do we update node costs?
 - ▶ Could add a `changeKeyPriority(...)` method to heap
 - ▶ Alternatively, add the node and the cost to the heap again (and ignore duplicates)