

# CSE 373: Tradeoffs and Abstractions

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## Warmup questions:

Instructions:

- ▶ Recall: What's an ADT? What's a data structure? An implementation of a data structure?
- ▶ Skim the Queue ADT on your handout.
- ▶ Discuss: How would you implement a queue?

## Possible queue implementations

# Announcements

Course overload link: [Given in lecture only]

Other announcements:

- ▶ Overloading + looking for a partner? Talk to me after class.
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Setup tips and tricks:

- ▶ Suspect the spec is out-of-date? Shift-refresh in your browser
- ▶ Use Java 8, not 9
- ▶ When running into weird Eclipse issues, try restarting it

## Places to get practice

- ▶ Section 1 handouts
- ▶ Practice-it: <https://practiceit.cs.washington.edu>
- ▶ CSE 143 class website (17au or older)
- ▶ Project 1

Need help? Visit office hours!

ADTs are just a tool for *communicating* with other programmers

This course focuses on *implementing* ADTs: implementing data structures

# Why?

Why?

Why can't we just use `java.util.*`?



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The reality: nothing's perfect

But we can work around many *tradeoffs* by carefully *adapting* data structures and *abstracting* algorithms!

There are (often highly *non-obvious*) ways to organize information to enable *efficient* computations over data.

However, no method is perfect: there exists unavoidable **tradeoffs**.

# Tradeoffs

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Core questions:

- ▶ What operations do I really need?
- ▶ What assumptions am I making about how my software will be used? (e.g. more lookups or inserts)

## Case study: The List ADT

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- ▶ **size**: return the number of elements in the stack

# Tradeoffs

Goal: implement the List ADT

Compare and contrast: array list vs linked list

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- ▶ Amount of space used overall:
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  - ▶ Array list: No wasted space
  - ▶ Linked list: Slightly more space per element

## A question:

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One idea:

```
for (int i = 0; i < myList.size(); i++) {  
    System.out.println(myList.get(i));  
}
```

How efficient is this if myList is an array list? A linked list?

## A problem:

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We want to make linked list iteration fast. How?

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Idea!

- ▶ **Adapt** the list ADT
- ▶ **Abstract** the idea of iteration

## A solution?

```
Iterator<String> iter = myList.iterator();  
while (iter.hasNext()) {  
    String item = iter.next();  
    System.out.println(item);  
}
```



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- ▶ **size**: return the number of elements in the stack
- ▶ **iterator**: returns an iterator over the list

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Supported operations:

- ▶ **hasNext**: returns ‘true’ if there’s another element left to yield and false otherwise
- ▶ **next**: returns the next element (if there is one)

## Implementing an iterator: A plan of attack

What is this 'efficiency' thing anyways?



## Parting thoughts

Reminder: Overloading/partner concerns, talk to me after class

Supplemental resources: see resources page on class website for...

- ▶ Strategies on effectively testing code
- ▶ Info on JUnit
- ▶ Math review (logs, exponents, summations)

Have suggestions for more resources docs we should write?

Use feedback form.