1. Minimum Spanning Trees

Draw the minimum spanning tree that results when running Kruskal's algorithm on the graph below. Please draw the vertices in the same order as in the graph below. Write final state of the array disjoint sets data structure from the run that results from using union-by-$\text{SIZE}$ with NO path compression.

If there is a tie between edges in Kruskal's algorithm, choose the edge for which the endpoints, when written in alphabetical order, are alphabetically first. For example, if (D, A) and (B, C) were tied, you would chose (D,A) first because “AD” is alphabetically before “BC”.

If two trees to be unioned have the same size, make the root of the unioned tree the representative element that comes alphabetically first. For example, if a tree rooted at A and a tree rooted at F are to be unioned but have the same size, then A would be used as the root of the combined tree.

When drawing the final array, assume that the mapping from vertex letter to array index is alphabetical: $A \rightarrow 0$, $B \rightarrow 1$, etc.
2. Dynamic Programming

This problem will walk you through the steps of designing a dynamic program. The problem we are solving is the longest palindrome problem: given a string, \( S \), what is the length of the longest palindrome that is a substring of \( S \)? A palindrome is a string that reads the same forwards as backwards. For example, “racecar”, “eve”, and “I”, are all palindromes. We also consider the empty string to be a palindrome (of length 0).

(a) First we need to figure out what our subproblems are. Since we are working with strings, a natural subproblem to use is substrings. Let \( \text{OPT}(i, n) \) denote the length of the longest palindrome in the substring of length \( n \) starting at index \( i \). Write an expression for the recursive case of \( \text{OPT}(i, n) \). (Hint: All palindromes above a certain size have palindromes as substrings).

(b) Next we need a base case for our \( \text{OPT} \) recurrence. Write an expression for the base case(s) of this recurrence. (Hint: Which size strings are always palindromes?)

(c) Now that we have a complete recurrence, we need to figure out which order to solve the subproblems in. Which subproblems does the recursive case \( \text{OPT}(i, n) \) require to be calculated before it can be solved?

(d) Given these dependencies, what order should we loop over the subproblem in?

(e) We have all of the pieces required to put together a dynamic program now. Write pseudocode for the dynamic program that computes the length of the longest palindromic substring of \( S \).