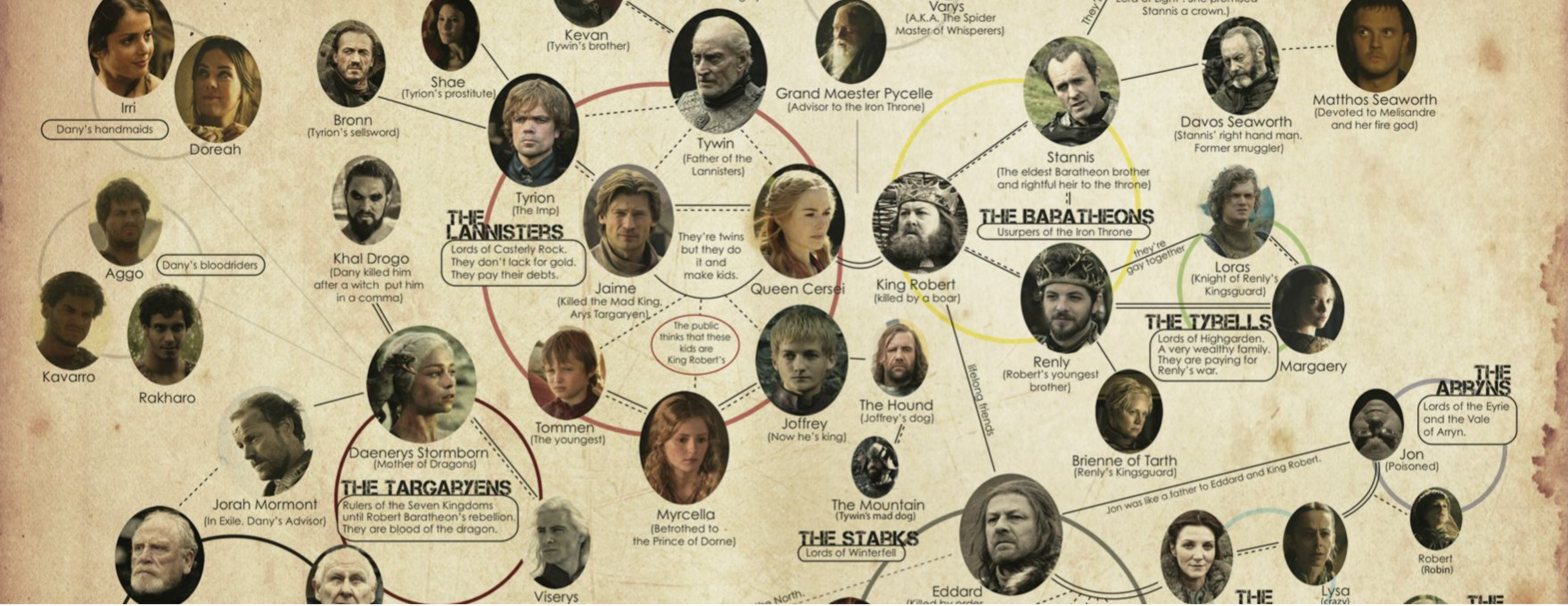


Implementing Graphs

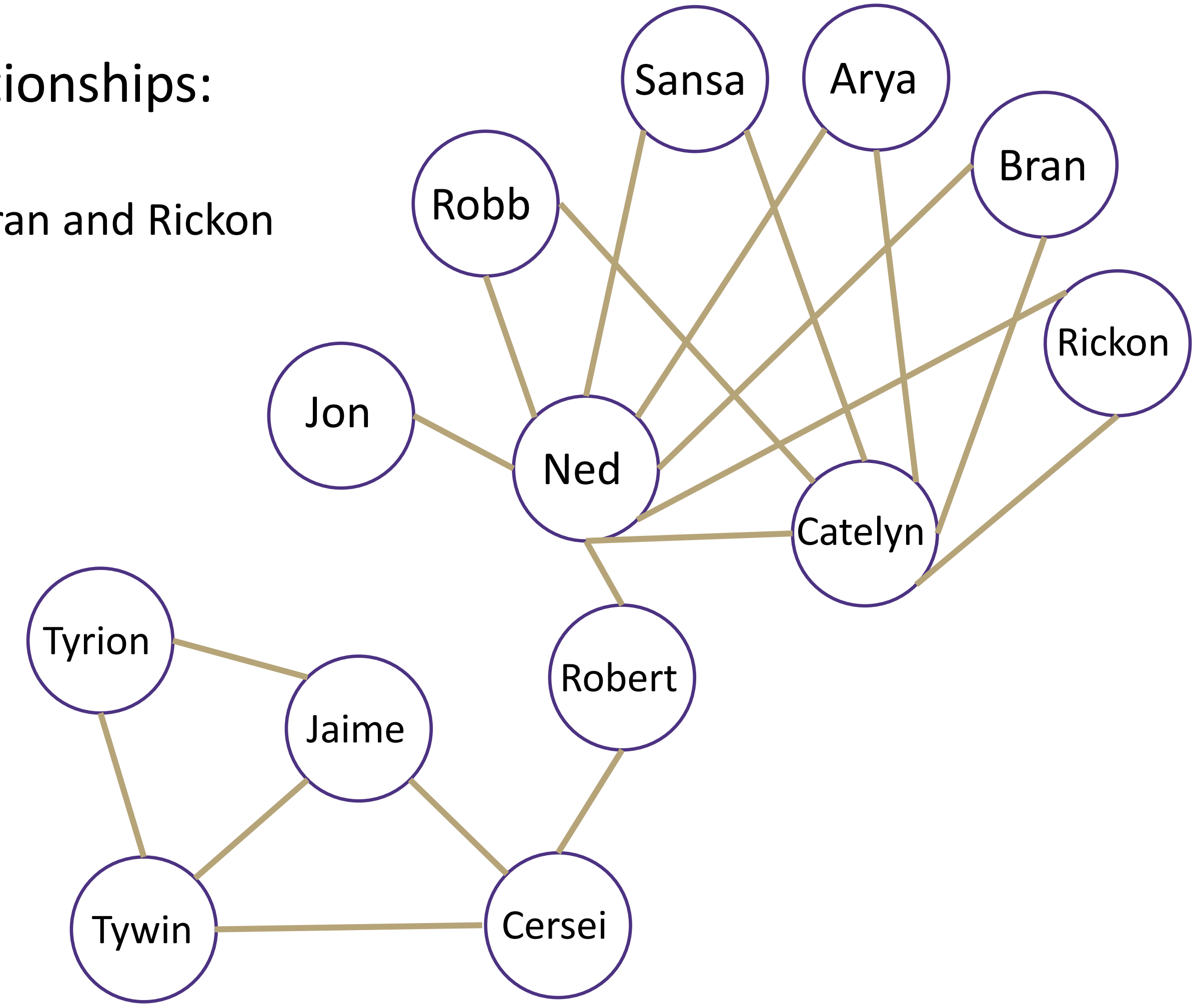
Data Structures and Algorithms

https://visualdatahub.wordpress.com/2013/04/02/a-web-of-thrones-character-relationships/

Warm Up

Draw a graph that represents the following relationships:

- Ned and Catelyn are married
- Ned and Catelyn are the parents of Robb, Sansa, Arya, Bran and Rickon
- Ned is the “father” of Jon
- Ned and Robert are in allies
- Robert and Cersei are married
- Cersei and Jaime are in love
- Tywin is the father of Cersei, Jamie and Tyrion



V = 13

E = 19

Announcements

I’m not here next week

Email me with grade questions

Graph Vocabulary

Graph Direction

- **Undirected graph** – edges have no direction and are two-way

  V = { Dany, Drogo, Jon }

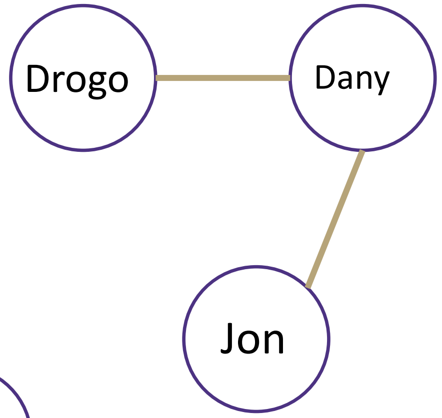
  E = { (Dany, Drogo), (Dany, Jon) } *inferred (Drogo, Dany) and (Jon, Dany)*

- **Directed graphs** – edges have direction and are thus one-way

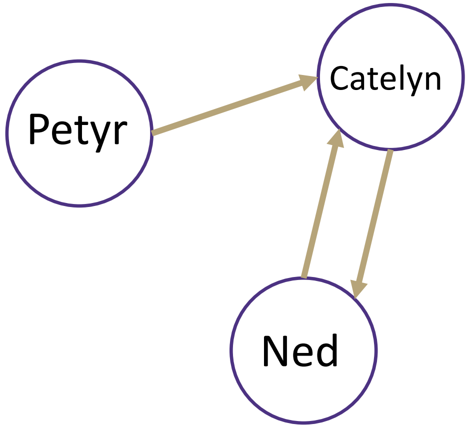
  V = { Petyr, Catelyn, Ned }

  E = { (Petyr, Catelyn), (Catelyn, Ned), (Ned, Catelyn) }

Undirected Graph:



Directed Graph:



Degree of a Vertex

- **Degree** – the number of edges containing that vertex

  Drogo : 1, Danny : 1, Jon : 1

- **In-degree** – the number of directed edges that point to a vertex

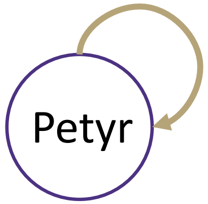
  Petyr : 0, Catelyn : 2, Ned : 1

- **Out-degree** – the number of directed edges that start at a vertex

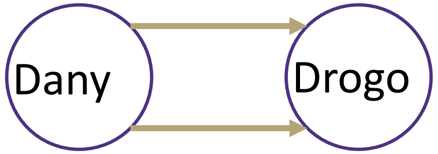
  Petyr : 1, Catelyn : 1, Ned : 1

Graph Vocabulary

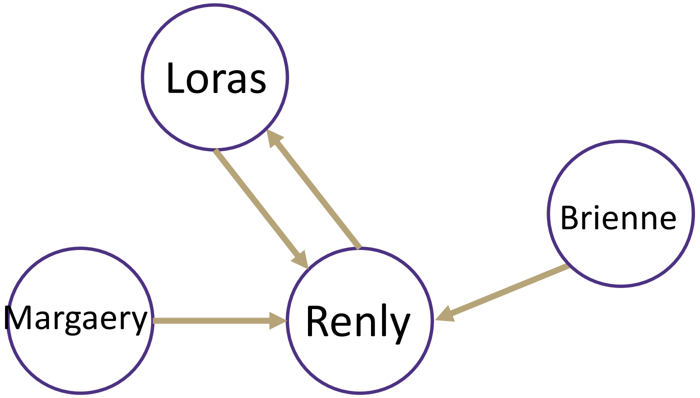
Self loop – an edge that starts and ends at the same vertex



Parallel edges – two edges with the same start and end vertices

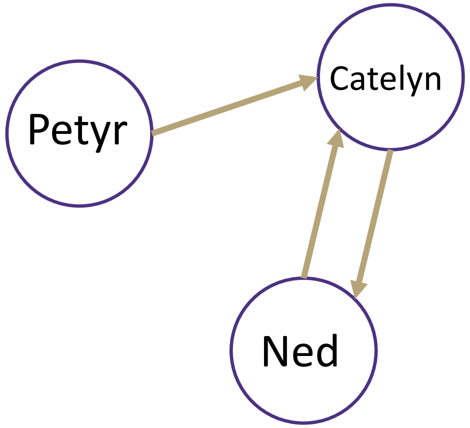


Simple graph – a graph with no self-loops and no parallel edges

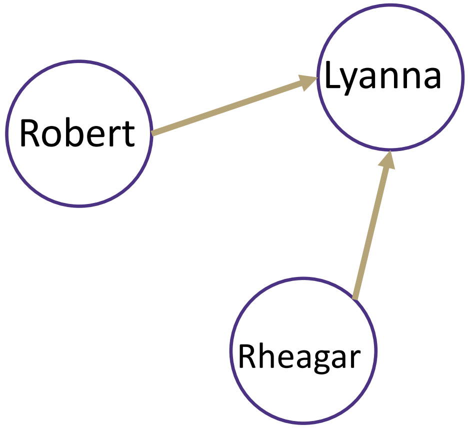


Food for thought

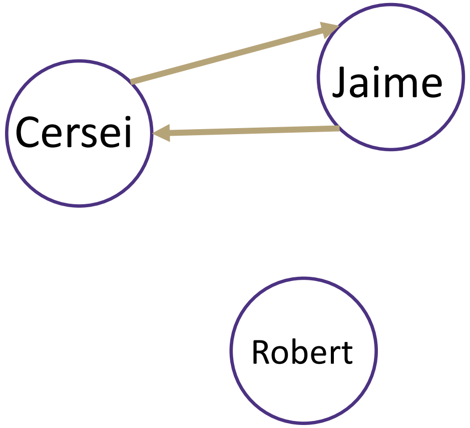
Is a graph valid if there exists a vertex with a degree of 0? Yes



Petyr has an “in degree” of 0

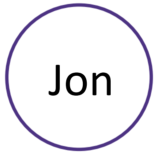


Lyanna has an “out degree” of 0



Robert has both an “in degree” and an “out degree” of 0

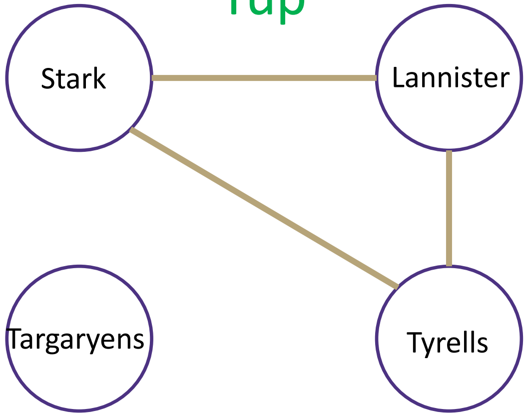
Is this a valid graph?

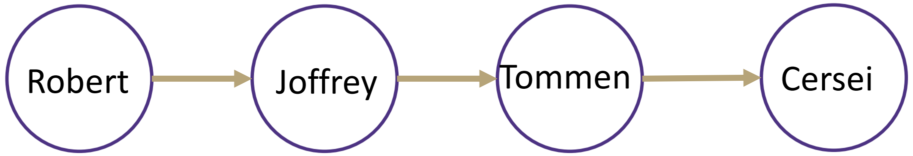


Yes!

Are these valid?

Yup





Sure

Implementing a Graph

Implement with nodes…

Implementation gets super messy

What if you wanted a vertex without an edge?

How can we implement without requiring edges to access nodes?

Implement using some of our existing data structures!

Adjacency Matrix

Assign each vertex a number from 0 to V – 1

Create a V x V array of Booleans

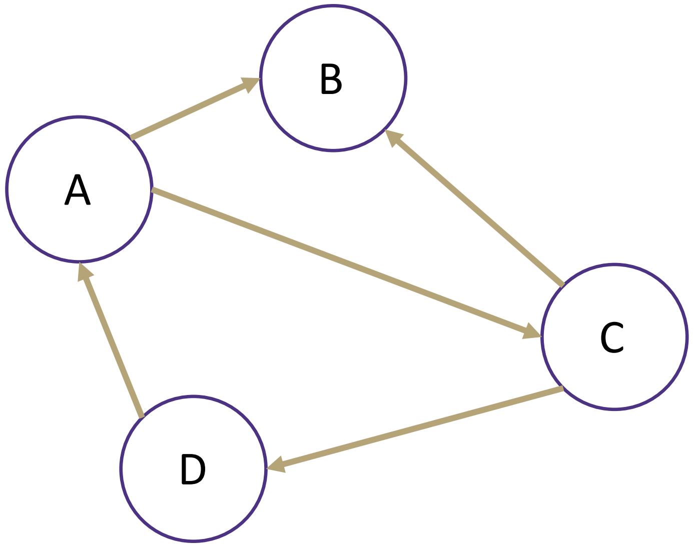
If (x,y) ∈ E then arr[x][y] = true

	A	B	C	D
A		T	T	
B				
C		T		T
D	T			

Runtime (in terms of V and E)

- get out - edges for a vertex O(v)
- get in – edges for a vertex O(v)
- decide if an edge exists O(1)
- insert an edge O(1)
- delete an edge O(1)
- delete a vertex
- add a vertex

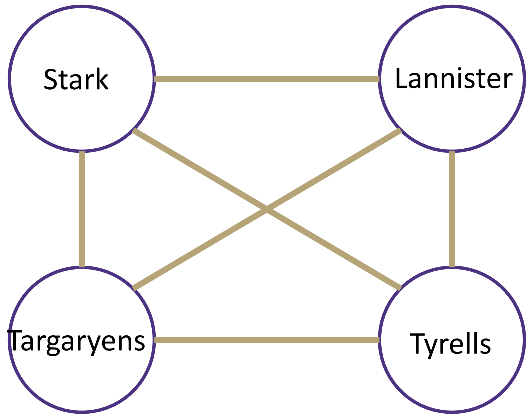
How much space is used?
V^2



Graph Vocabulary

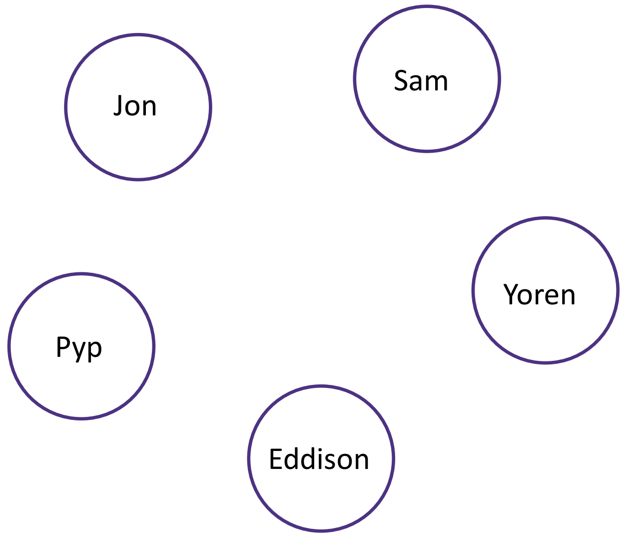
Dense Graph – a graph with a lot of edges

$E \in \Theta(V^2)$



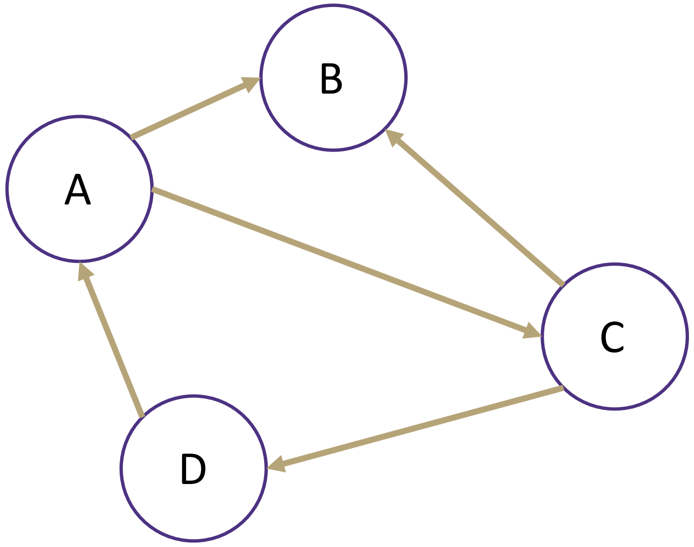
Sparse Graph – a graph with “few” edges

$E \in \Theta(V)$



An Adjacency Matrix seems a waste for a sparse graph…

Adjacency List



Create a Dictionary of size V from type V to Collection of E

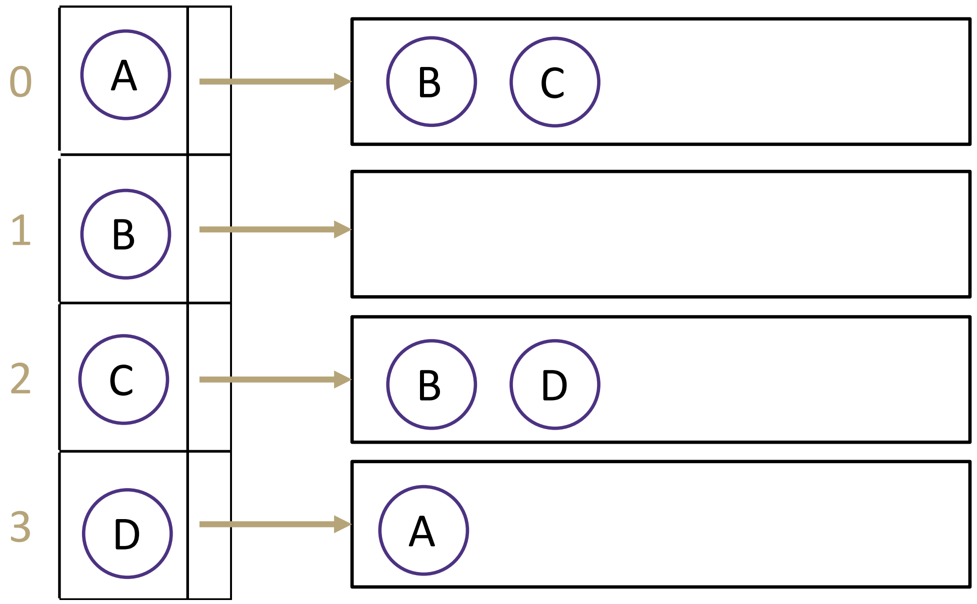
If (x,y) ∈ E then add y to the set associated with the key x

Runtime (in terms of V and E)

- get out - edges for a vertex O(1)
- get in - edges for a vertex O(V + E)
- decide if an edge exists O(1)
- insert an edge O(1)
- delete an edge O(1)
- delete a vertex
- add a vertex

How much space is used?

V + E

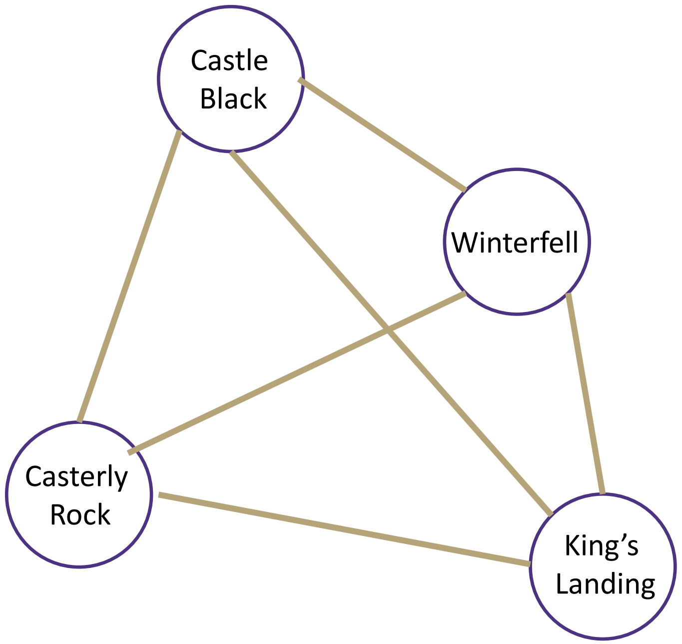


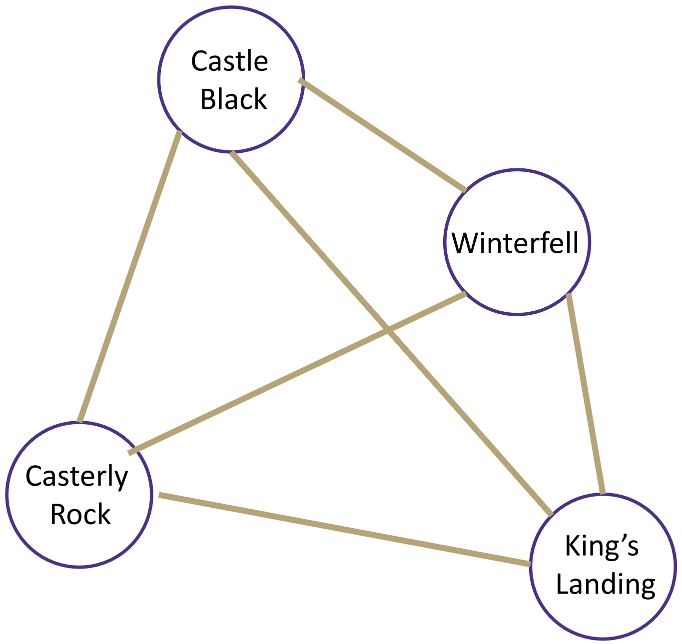
Walks and Paths

Walk – continuous set of edges leading from vertex to vertex

A list of vertices where if I is some int where 0 < 1 < Vn every pair (Vi, Vi+1) in E is true

Path – a walk that never visits the same vertex twice

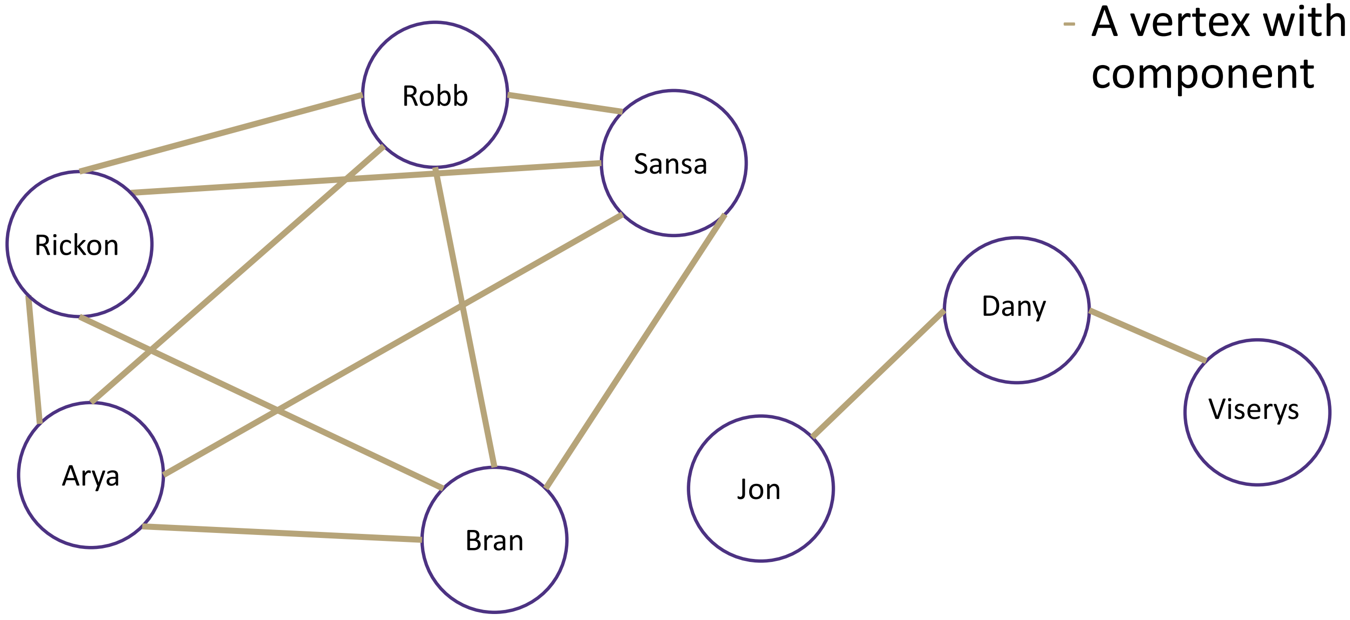




Connected Graphs

Connected graph – a graph where every vertex is connected to every other vertex via some path. It is not required for every vertex to have an edge to every other vertex

There exists some way to get from each vertex to every other vertex



Connected Component – a *subgraph* in which any two vertices are connected via some path, but is connected to no additional vertices in the *supergraph*

- There exists some way to get from each vertex within the connected component to every other vertex in the connected component
- A vertex with no edges is itself a connected component

Graph Algorithms

Traversing a Graph

In all previous data structures:

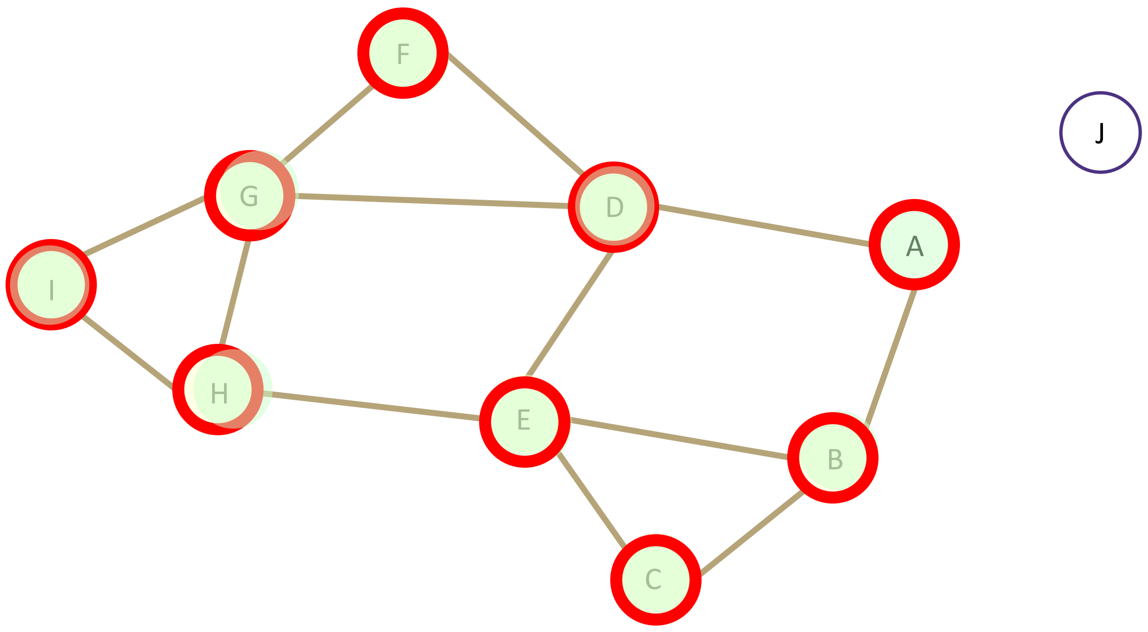
1. Start at first element
2. Move to next element
3. Repeat until end of elements

For graphs – Where do we start? How do we decide where to go next? When do we end?

1. Pick any vertex to start, mark it "visited"
2. Put all neighbors of first vertex in a "to be visited" collection
3. Move onto next vertex in "to be visited" collection
4. Mark vertex "visited"
5. Put all unvisited neighbors in "to be visited"
6. Move onto next vertex in "to be visited" collection
7. Repeat…

Breadth First Search

```
search(graph)
   toVisit.enqueue(first vertex)
   while(toVisit is not empty)
      current = toVisit.dequeue()
      for (V : current.neighbors())
         if (V is not in queue)
            toVisit.enqueue(v)
      visited.add(current)
```



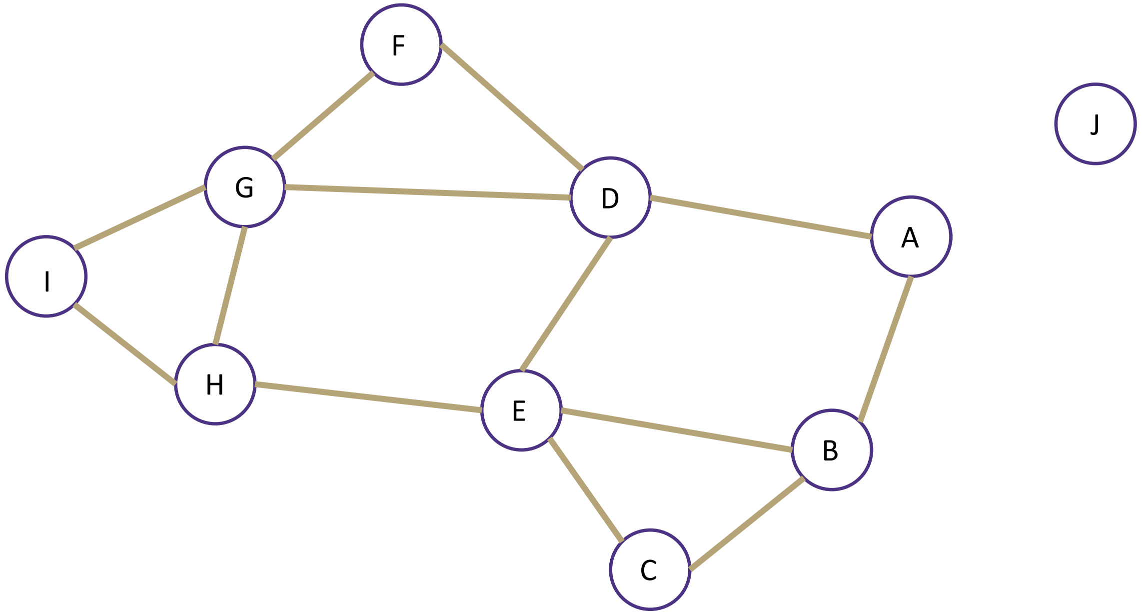
Current node: I

Queue: B D E C F G H I

Visited: A B D E C F G H I

Breadth First Search Analysis

```
search(graph)
   toVisit.enqueue(first vertex)
   while(toVisit is not empty)
      current = toVisit.dequeue()
      for (V : current.neighbors())
         toVisit.enqueue(v)
      visited.add(current)
```



Visited: A B D E C F G H I

How many times do you visit each node? 1 time each

How many times do you traverse each edge? Max 2 times each

- Putting them into toVisit
- Checking if they're in toVisit

Runtime? O(V + 2E) = O(V + E) "graph linear"

Depth First Search (DFS)

BFS uses a queue to order which vertex we move to next

Gives us a growing “frontier” movement across graph

Can you move in a different pattern? Can you use a different data structure?

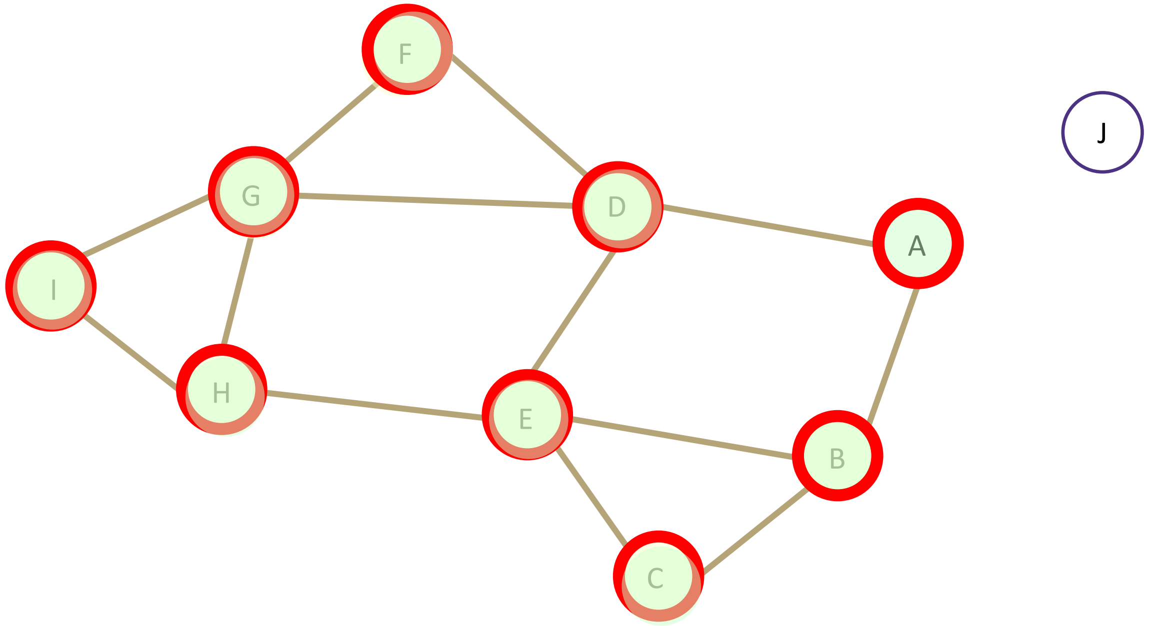
What if you used a stack instead?

```
bfs(graph)
   toVisit.enqueue(first vertex)
   while(toVisit is not empty)
      current = toVisit.dequeue()
      for (V : current.neighbors())
         if (V is not in queue)
            toVisit.enqueue(v)
      visited.add(current)
```

```
dfs(graph)
   toVisit.push(first vertex)
   while(toVisit is not empty)
      current = toVisit.pop()
      for (V : current.neighbors())
         if (V is not in stack)
            toVisit.push(v)
      visited.add(current)
```

Depth First Search

```
dfs(graph)
   toVisit.push(first vertex)
   while(toVisit is not empty)
      current = toVisit.pop()
      for (V : current.neighbors())
         if (V is not in stack)
            toVisit.push(v)
      visited.add(current)
```



Current node: D

Stack: D CB EI HGF

Visited: A B E H G F I C D

How many times do you visit each node? 1 time each

How many times do you traverse each edge? Max 2 times each

- Putting them into toVisit
- Checking if they're in toVisit

Runtime? O(V + 2E) = O(V + E) "graph linear"