



CSE373: Data Structures and Algorithms
Algorithm Design Paradigms

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 Winter 2016

This lecture material represents the work of multiple instructors at the University of Washington. Thank you to all who have contributed!

Algorithm Design Techniques

- Greedy
 - Shortest path, minimum spanning tree, ...
- Divide and Conquer
 - Divide the problem into smaller subproblems, solve them, and combine into the overall solution
 - Often done recursively
 - Quick sort, merge sort are great examples
- Dynamic Programming
 - Brute force through all possible solutions, storing solutions to subproblems to avoid repeat computation
- Backtracking
 - A clever form of exhaustive search

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Dynamic Programming: Idea

- Divide a bigger problem into many smaller subproblems
- If the number of subproblems grows exponentially, a recursive solution may have an exponential running time ☹
- Dynamic programming to the rescue! ☺
- Often an individual subproblem may occur many times!
 - Store the results of subproblems in a table and re-use them instead of recomputing them
 - Technique called **memoization**

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Fibonacci Sequence: Recursive

- The fibonacci sequence is a very famous number sequence
- 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...
- The next number is found by adding up the two numbers before it.
- Recursive solution:

```
fib(int n) {
  if (n == 1 || n == 2) {
    return 1
  }
  return fib(n - 2) + fib(n - 1)
}
```

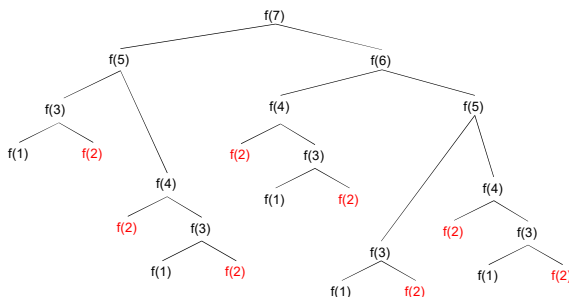
- Exponential running time!
 - A lot of repeated computation

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Repeated computation



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Fibonacci Sequence: memoized

```
fib(int n):
  results = Map() # Empty mapping container.
  results.put(1, 1)
  results.put(2, 1)
  return fibHelper(n, results)

fibHelper(int n, Map results):
  if (!results.contains(n)):
    results.put(n, fibHelper(n-2)+fibHelper(n-1))
  return results.get(n)
```

Now each call of **fib(x)** only gets computed once for each x!

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Comments

- Dynamic programming relies on working “from the bottom up” and saving the results of solving simpler problems
 - These solutions to simpler problems are then used to compute the solution to more complex problems
- Dynamic programming solutions can often be quite complex and tricky
- Dynamic programming is used for **optimization** problems, especially ones that would otherwise take exponential time
 - Only problems that satisfy the **principle of optimality** are suitable for dynamic programming solutions
 - i.e. the subsolutions of an optimal solution of the problem are themselves optimal solutions for their subproblems
- Since exponential time is unacceptable for all but the smallest problems, dynamic programming is sometimes essential

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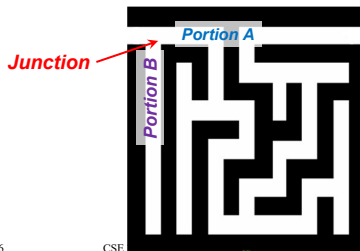
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Backtracking: Idea

- Backtracking is a technique used to solve problems with a large search space, by systematically trying and eliminating possibilities.
- A standard example of backtracking would be going through a maze.
 - At some point, you might have two options of which direction to go:



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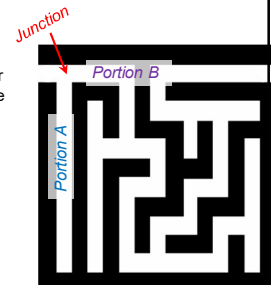
Backtracking

One strategy would be to try going through **Portion A** of the maze.

If you get stuck before you find your way out, then you “backtrack” to the junction.

At this point in time you know that **Portion A** will **NOT** lead you out of the maze,

so you then start searching in **Portion B**



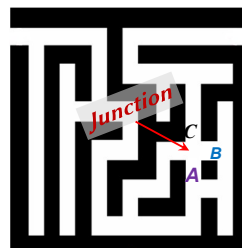
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Backtracking

- Clearly, at a single junction you could have even more than 2 choices.
- The backtracking strategy says to try each choice, one after the other,
 - if you ever get stuck, “backtrack” to the junction and try the next choice.
- If you try all choices and never found a way out, then there IS no solution to the maze.

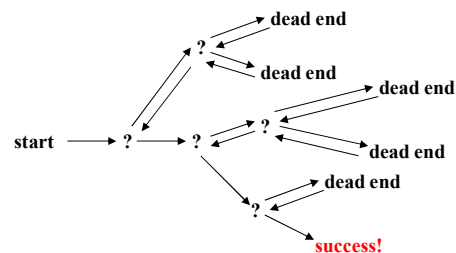


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Backtracking (animation)



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Backtracking

- Dealing with the maze:
 - From your start point, you will iterate through each possible starting move.
 - From there, you recursively move forward.
 - If you ever get stuck, the recursion takes you back to where you were, and you try the next possible move.
- Make sure you don't try too many possibilities,
 - Mark which locations in the maze have been visited already so that no location in the maze gets visited twice.
 - (If a place has already been visited, there is no point in trying to reach the end of the maze from there again.)

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Backtracking

The neat thing about coding up backtracking is that it can be done recursively, without having to do all the bookkeeping at once.

- Instead, the stack of recursive calls does most of the bookkeeping
- (i.e., keeps track of which locations we've tried so far.)

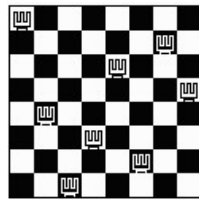
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Backtracking: The 8 queens problem

- Find an arrangement of 8 queens on a single chess board such that no two queens are attacking one another.
- In chess, queens can move all the way down any row, column or diagonal (so long as no pieces are in the way).
 - Due to the first two restrictions, it's clear that each row and column of the board will have exactly one queen.



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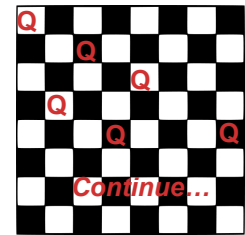
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Backtracking

The backtracking strategy is as follows:

- Place a queen on the first available square in row 1.
- Move onto the next row, placing a queen on the first available square there (that doesn't conflict with the previously placed queens).
- Continue in this fashion until either:
 - You have solved the problem, or
 - You get stuck.



When you get stuck, remove the queens that got you there, until you get to a row where there is another valid square to try.

Animated Example:

<http://www.hbmeyer.de/backtracking/achtdamen/eight.htm#u>

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Backtracking – 8 queens Analysis

- Another possible brute-force algorithm is generate all possible permutations of the numbers 1 through 8 (there are $8! = 40,320$),
 - Use the elements of each permutation as possible positions in which to place a queen on each row.
 - Reject those boards with diagonal attacking positions.
- The backtracking algorithm does a bit better
 - constructs the search tree by considering one row of the board at a time, eliminating most non-solution board positions at a very early stage in their construction.
 - because it rejects row and diagonal attacks even on incomplete boards, it examines only 15,720 possible queen placements.
- 15,720 is still a lot of possibilities to consider
 - Sometimes we have no other choice but to do the best we can ☺

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