



CSE373: Data Structures and Algorithms

Lecture 1: Introduction; ADTs; Stacks/Queues

Lauren Milne
Summer 2015

Welcome!

We will learn *fundamental data structures* and *algorithms* for organizing and processing information

- “Classic” data structures / algorithms
- Analyze efficiency
- When to use them

Today in class:

- Introductions and course mechanics
- What this course is about
- Start *abstract data types* (ADTs), *stacks*, and *queues*

Course staff



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Works with Richard Ladner in Accessibility

Does triathlons and skijoring in free time

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Junior majoring in Electrical Engineering
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Originally from Tabasco, Mexico!

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Course Information

- <http://courses.cs.washington.edu/courses/cse373/15su/>
- Textbook: Weiss 3rd Edition in Java
- Course email list: **cse373a_15sp@u.washington.edu**
- College of Arts & Sciences Instructional Computing Lab
 - <http://depts.washington.edu/aslab/>
 - Or your own machine
- Will use Java for the programming assignments
- Eclipse is recommended programming environment

Office Hours

- M 1:30-2:30 (Lauren)
- Tu 12-1 pm (Jiechen)
- W 12-1 pm (Nao)
- W 4-5 pm (Mert)
- Th 9:30-10:30 am (Yanling)
- F 12-1 pm (Mauricio)

Do these work for everyone?

Collaboration and Academic Integrity

DON'T CHEAT!

Seriously, read the policy online, follow Gilligan's Island rule.

What this course will cover

- Introduction to Algorithm Analysis
- Lists, Stacks, Queues
- Trees, Hashing, Dictionaries
- Heaps, Priority Queues
- Sorting
- Disjoint Sets
- Graph Algorithms
- Introduction to Parallelism and Concurrency

Goals

- Be able to **make good design choices** as a developer
- Be able to **justify** and **communicate** your design decisions
- This is not a course about Java!

To-do list

In next 24-48 hours:

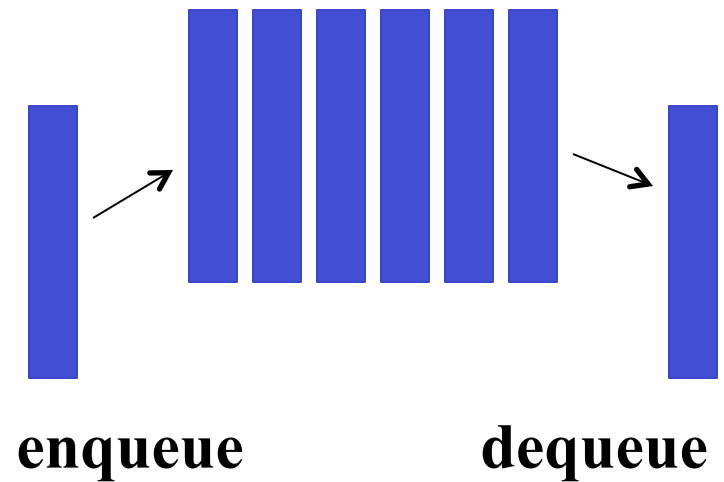
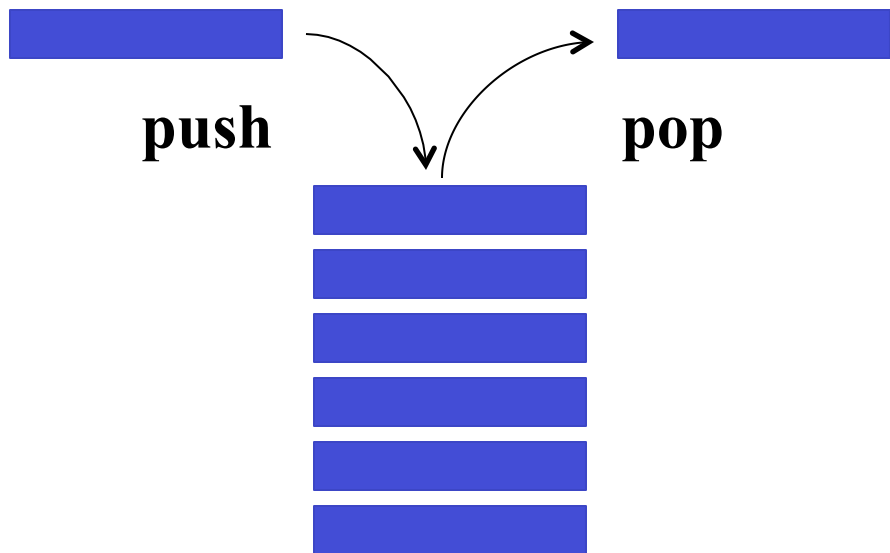
- Read the web page
- Read all course policies
- Read Chapters 3.1 (lists), 3.6 (stacks) and 3.7 (queues) of the Weiss book
 - Relevant to Homework 1, **due next week!**
- Set up your Java environment for Homework 1

<http://courses.cs.washington.edu/courses/cse373/15su/>

Data structures

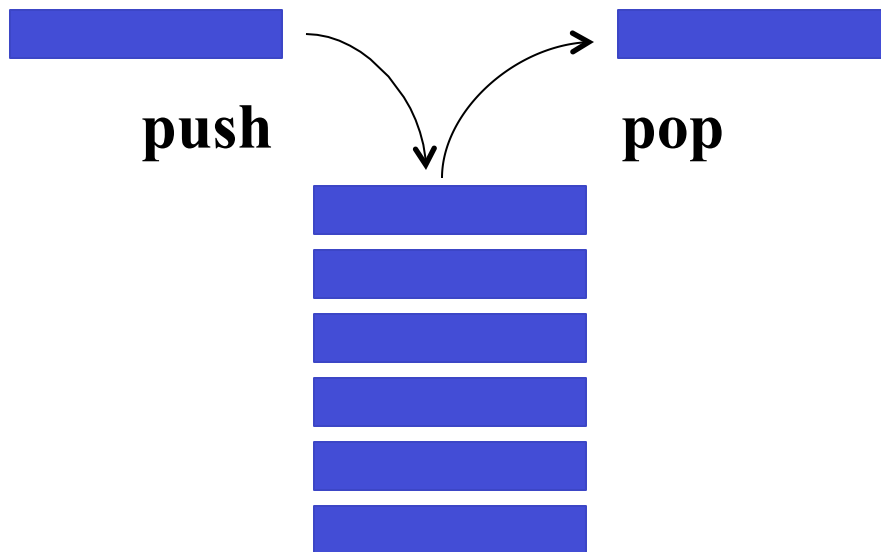
A data structure is a **way to organize information** to enable *efficient* computation over that information

What are some examples?



Data structures

A data structure is a **way to organize information** to enable *efficient* computation over that information



It supports certain operations, each with a:

- Meaning
- Performance

Trade-offs

A data structure strives to provide many useful, efficient operations.

But there are unavoidable trade-offs:

- Time vs space
- One operation's efficiency vs another
- Generality vs simplicity vs performance

So, when should I use different data structures?

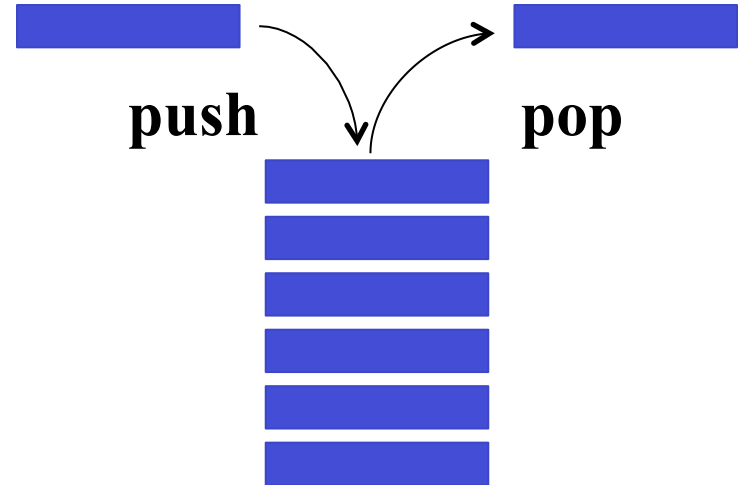
Terminology

- Abstract Data Type (ADT)
 - Mathematical description of a “thing” with set of operations
 - Not concerned with implementation details
- Data structure
 - A specific organization of data and family of algorithms for implementing an ADT
- Algorithm
 - A high level, language-independent description of a step-by-step process
- Implementation of a data structure
 - A specific implementation in a specific language

Example: Stacks

- The **Stack ADT** supports operations:

- **isEmpty**
- **push**
- **pop**
- What else?



- A Stack **data structure** could use a linked-list or an array and associated **algorithms** for the operations
- One **implementation** is in the library `java.util.Stack`

Why Useful

Stack ADT arises all the time in programming

- Recursive function calls
- Balancing symbols in programming (parenthesis)
- Evaluating postfix notation: $3\ 4\ +\ 5\ *$
- Conversion from infix $((3+4)*5)$ to postfix notation

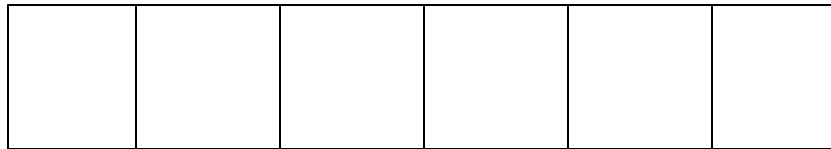
We can code up a **reusable library**

We can **communicate** in high-level terms

Stack Implementations

- stack as an array

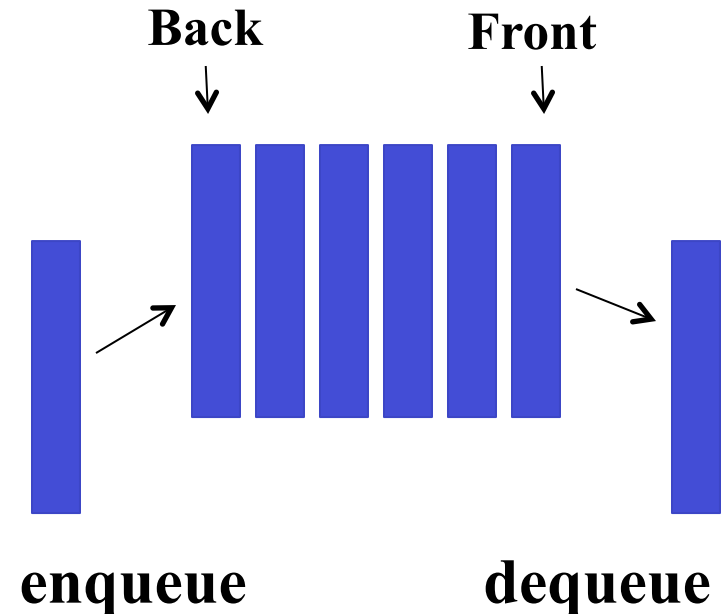
1. Initially empty
2. Push('a')
3. Push('b')
4. Pop()



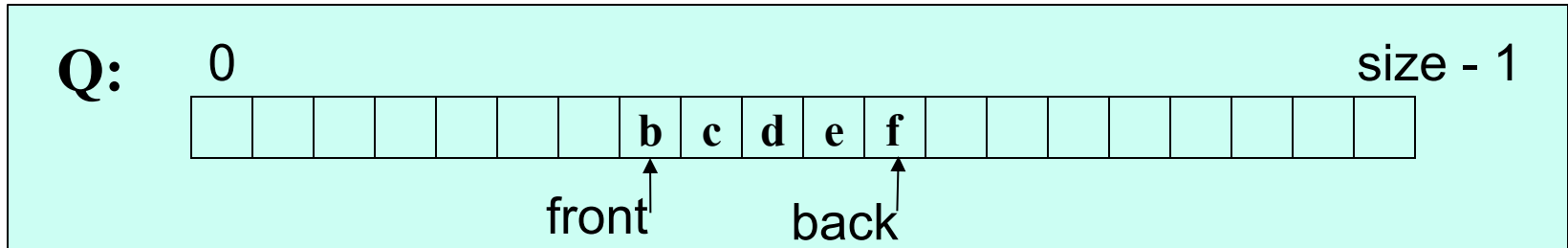
- stack as a linked list

The Queue ADT

- Operations
 - **enqueue**
 - **dequeue**
 - **is_empty**
 - **What else?**
- Just like a stack except:
 - Stack: LIFO (last-in-first-out)
 - Queue: FIFO (first-in-first-out)



Circular Array Queue Data Structure



// Basic idea only!

```
enqueue(x) {  
    next = (back + 1) % size  
    Q[next] = x;  
    back = next  
}
```

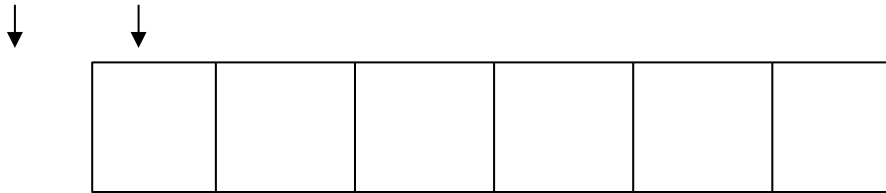
// Basic idea only!

```
dequeue() {  
    x = Q[front];  
    front = (front + 1) % size;  
    return x;  
}
```

- What if **queue** is empty?
- What if **array** is full?
- How to *test* for empty?
- What is the *complexity* of the operations?
- Can you find the k^{th} element in the queue?
- Why do we use a circular array for a queue vs a standard array for a stack?

In Class Practice

back front



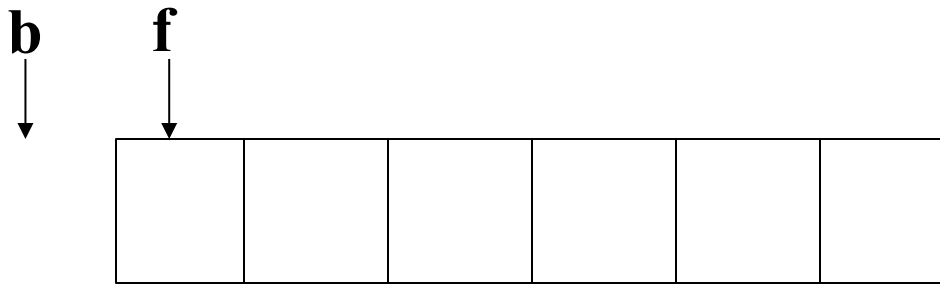
```
// Basic idea only!
```

```
enqueue(x) {  
    next = (back + 1) % size  
    Q[next] = x;  
    back = next  
}
```

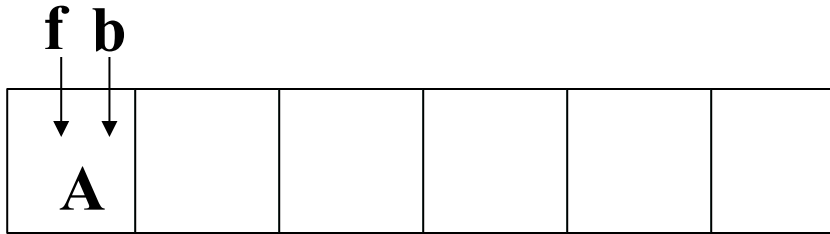
```
// Basic idea only!
```

```
dequeue() {  
    x = Q[front];  
    front = (front + 1) % size;  
    return x;  
}
```

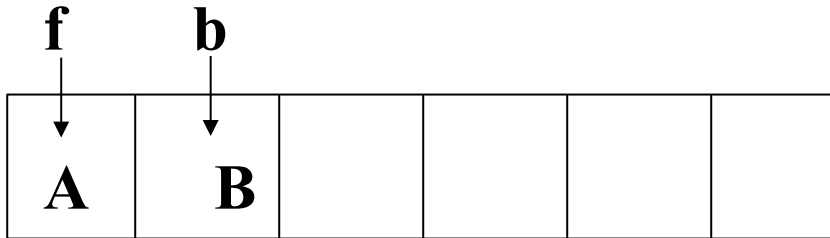
```
enqueue('A')  
enqueue('B')  
enqueue('C')  
o = dequeue()  
o = dequeue()  
enqueue('D')  
enqueue('E')  
enqueue('F')  
enqueue('G')  
enqueue('H')  
enqueue('I')
```



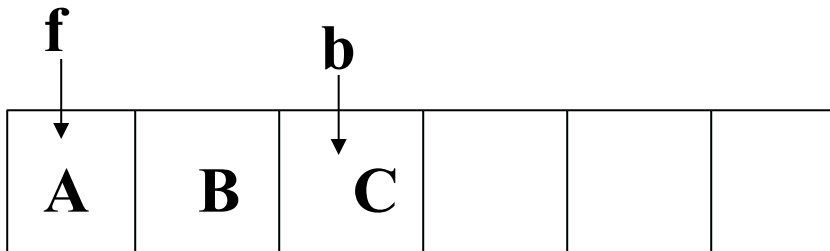
enqueue('A')



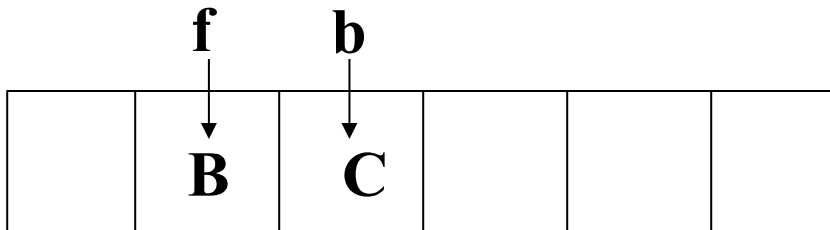
enqueue('B')

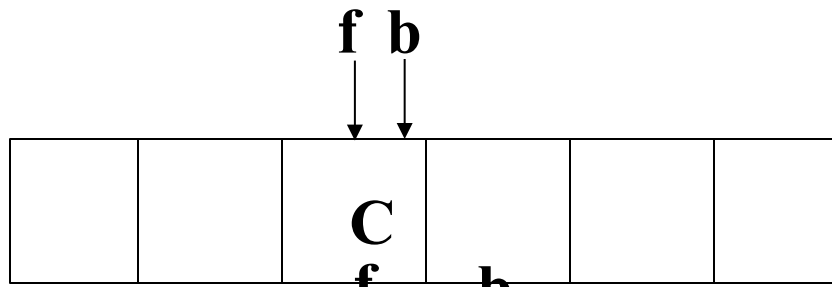


enqueue('C')

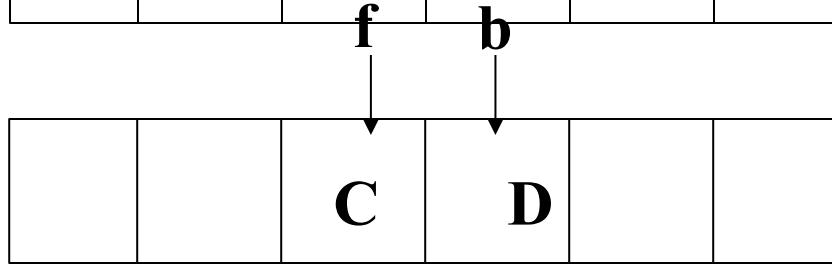


o = dequeue()

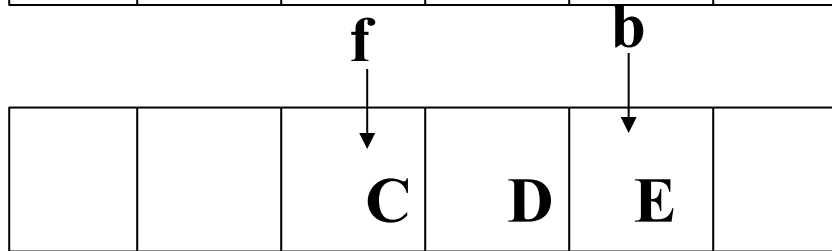




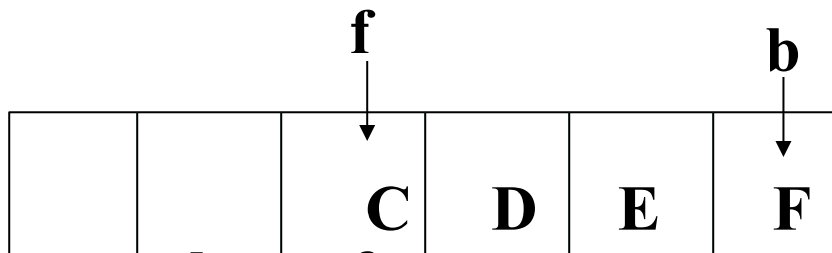
o = dequeue



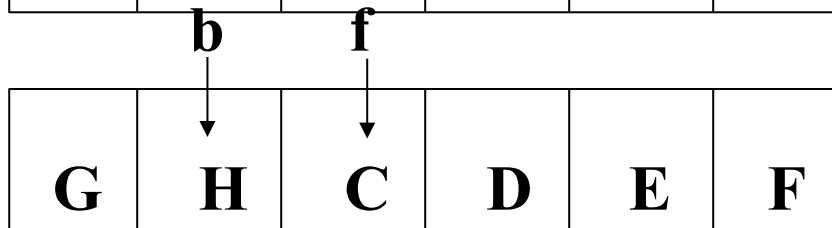
enqueue('D')



enqueue('E')



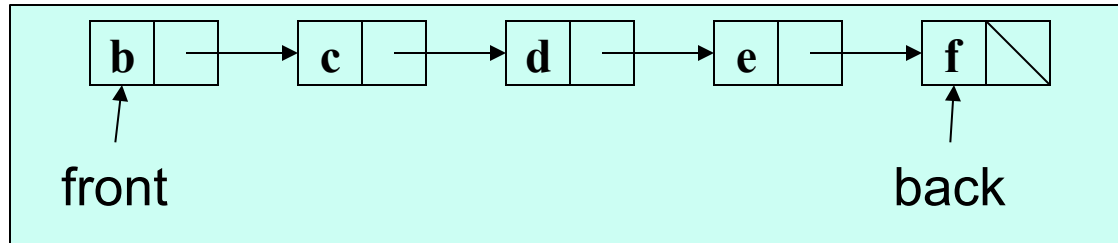
enqueue('F')



enqueue('G'), enqueue('H')

enqueue('i')

Linked List Queue Data Structure



// Basic idea only!

```
enqueue(x) {  
    back.next = new Node(x);  
    back = back.next;  
}
```

// Basic idea only!

```
dequeue() {  
    x = front.item;  
    front = front.next;  
    return x;  
}
```

- What if **queue** is empty?
 - Enqueue?
 - Dequeue?
- Can **list** be full?
- How to *test* for empty?
- What is the *complexity* of the operations?
- Can you find the k^{th} element in the queue?

Circular Array vs. Linked List

Array:

- May waste unneeded space or run out of space
- Space per element excellent
- Operations very simple / fast
- Constant-time access to k^{th} element
- For operation insertAtPosition, must shift all later elements
 - Not in Queue ADT

List:

- Always just enough space
- But more space per element
- Operations very simple / fast
- No constant-time access to k^{th} element
- For operation insertAtPosition must traverse all earlier elements
 - Not in Queue ADT

Conclusion

- Abstract data structures allow us to define a new data type and its operations.
- Each abstraction will have one or more implementations.
- Which implementation to use depends on the application, the expected operations, the memory and time requirements.
- Both stacks and queues have array and linked implementations.
- We'll look at other ordered-queue implementations later.