

# CSE 373 Fall 2015, Homework 2

Due Friday, October 16<sup>th</sup> in homework Dropbox

*Please show work where applicable*

Name: \_\_\_\_\_

uwid (not your student number): \_\_\_\_\_

**1) For each of the following, show that  $f \in O(g)$ . That is, you will need to find values for  $c$  and  $n_0$  such that the definition of big-O holds true as we did with the examples in lecture.**

a)  $f(n) = 12n$

$$g(n) = \frac{n}{5}$$

b)  $f(n) = 6n^2 + 1000$

$$g(n) = n^4$$

c)  $f(n) = 6\log(n)$

$$g(n) = .5n$$

2) For each of the following program fragments, determine the asymptotic runtime in terms of n

a)

```
public void mysteryOne(int n) {
    int x = 0;
    for (int i = n; i >= 0; i--) {
        if ((i % 5) == 0) {
            break;
        } else {
            for (int j = 1; j < n; j *= 2) {
                x++;
            }
        }
    }
}
```

b)

```
public void mysteryTwo(int n) {
    int x = 0;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < ((n * n - 1) / 3); j++) {
            x += j;
        }
    }
}
```

c)

```
public void mysteryThree(int n) {
    for (int i = 0; i < n; i++) {
        methodTwo(i);
    }
}

private void methodTwo(int x) {
    if (x > 0) {
        methodTwo(x - 1);
    }
}
```

3) For each of the following, determine if  $f \in O(g)$ ,  $f \in \Omega(g)$ ,  $f \in \Theta(g)$ , several of these, or none of these.

a)  $f(n) = \log n$        $g(n) = \log \log n$

b)  $f(n) = 2^{2n}$        $g(n) = 2^n$

c)  $f(n) = 25n^3$        $g(n) = n^3 + 25n$

#### 4) Pseudocode and recurrence relations

**a)** Write pseudocode for a function that calculates the largest difference between any two numbers in an array of positive integers with a runtime in  $\Theta(n^2)$ .

*For example, the largest difference between any two numbers in the following array would be 19.*

*$a = [4, 6, 3, 9, 2, 1, 20]$*

**b)** Can this function be written with a runtime in  $\Theta(n)$ ? If yes, write pseudocode below. If no, why? What would have to be different about the input in order to do so?

**c)** Can this function be written with a runtime in  $\Theta(1)$ ? If yes, write pseudocode below. If no, why? What would have to be different about the input in order to do so?

### 5) Recurrence Relations

**a)** Find the tightest Big-Oh bound for the following recurrence relation  $T(n) = n + T(n/2)$ . Justify your answer.

**b)** Find a Big-Oh bound for the following recurrence relation  $T(n) = n + 2T(n/2)$ . Justify your answer.

## 6) Growth Rates

a) Order the following functions from slowest to fastest growth rate

- $2^{72}$
- $n^2 \log n$
- $2^{n/2}$
- $\log n$
- $n \log n^2$
- $n^6$
- $n \log \log n$
- $n \log^2 n$
- $n$
- $n^2$
- $n \log n$
- $2^n$
- $\log^2 n$
- $2/n$
- $n^{1/2}$

### 7) Big-Oh Definition

Suppose  $T1(n)$  is  $O(f(n))$  and  $T2(n)$  is  $O(f(n))$ . Which of the following are always true (for all  $T1$ ,  $f$ , and  $T2$ )? You do not need to prove an item is true (just saying true is enough for full credit), but if an item is false, you need to give a counterexample to demonstrate it is false. To give a counterexample, give values for  $T1(n)$ ,  $T2(n)$ , and  $f(n)$  for which the statement is false (for example, you could write, "The statement is false if  $T1(n) = 100n$ ,  $T2(n) = 2n^2$ , and  $f(n) = n^3$ "). Hints: Think about the definitions of big-O, big- $\Omega$ , and big- $\Theta$ .

a)  $T1(n)/T2(n)$  is  $O(1)$ .

b)  $T1(n) + T2(n)$  is  $\Omega(f(n))$ .

c)  $T1(n) - T2(n)$  is  $O(f(n))$ .

d)  $T1(n)$  is  $O(T2(n))$ .

e)  $T2(n)$  is  $\Theta(T1(n))$ .