Lists (a first look)

CSE 373 Data Structures

Readings

- Reading
 - > Chapter 3
 - > You can start peeking at Chapter 6

We will cover

- List ADT (a first look)
- List implementation
 - > Array
 - > Linked list
 - > Doubly linked list
- An example application (long integers)
- Circular list

List ADT

- What is a List?
 - > Ordered sequence of elements $A_1, A_2, ..., A_N$
- Elements may be of arbitrary type, but all are of the same type
- Elements have values
- Elements have positions (first, kth, last etc..)

Common operations on lists

- Constructor for an empty list
- Queries: size(); isEmpty();
- Insert and delete
 - Must indicate where: first, last, kth, after some element etc...
- Find, set, replace
 - > With a given value, find previous etc...
- Will look at a "list interface" in the Java sense later

Simple Examples of List Use

- Polynomials
 - > $25 + 4x^2 + 75x^{85}$
 - An element is a term whose value must indicate the power and the coefficient
- Unbounded Integers
 - > 4576809099383658390187457649494578
 - > Do not fit within a single computer word
 - > An element has for value a single digit

List Implementations

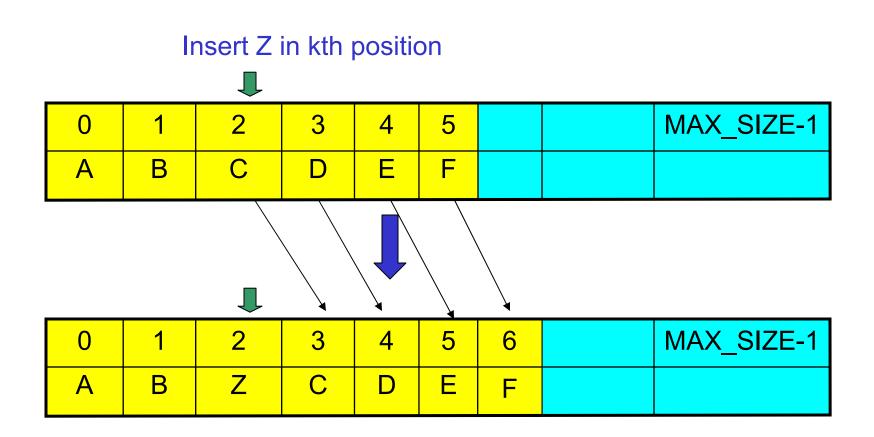
- Two types of implementation:
 - > Array-Based
 - > Linked list (pointer based)

List: Array Implementation

- Basic Idea:
 - > Pre-allocate a big array of size MAX_SIZE
 - > Keep track of current size using a variable count
 - Shift elements when you have to insert or delete (except of course for insertlast and deletelast)

0	1	2	3	 count-1	MAX_SIZE-1
A ₁	A ₂	A ₃	A ₄	 A _N	

List: Array Implementation

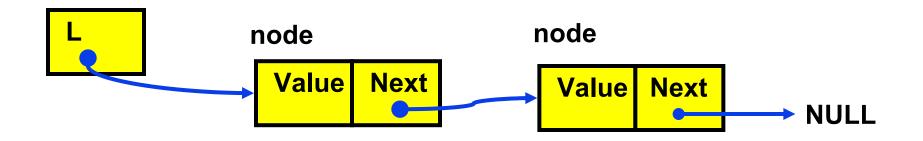


Array Insert_kth Running Time

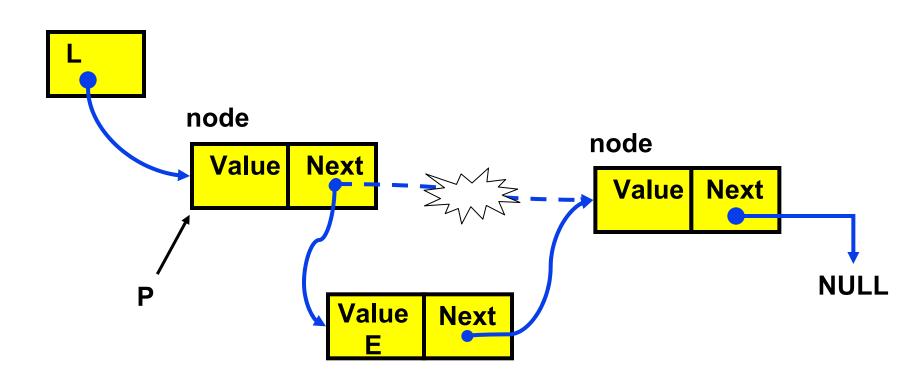
- Running time for N elements?
- On average, must move half the elements to make room – assuming insertions at positions are equally likely
- This is O(N) running time.
- Worst case is insert at position 0. Must move all N items one position before the insert. Still O(N)

Linked Implementation

- Basic Idea:
 - Allocate little blocks of memory (nodes) as elements are added to the list
 - > Keep track of list by linking the nodes together
 - Change links (pointers) when you want to insert or delete



Linked list: Insert_after

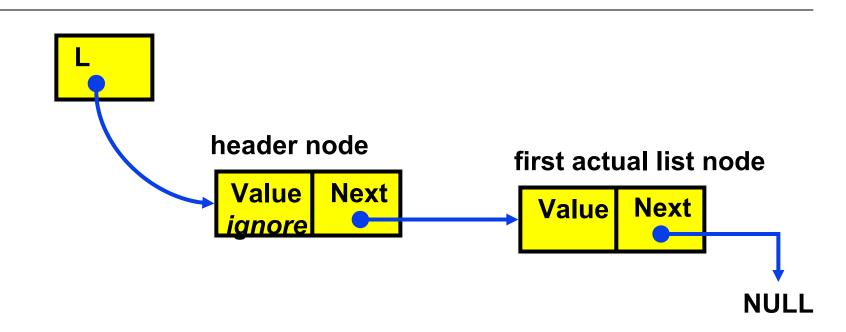


Insert the value E after P

Insertion After

```
InsertAfter(p : node, e : thing): {
x : node; //declares the type of x
x := new node;
x.value := v;
x.next := p.next;//be sure to do in right order
p.next := x;
}
```

Linked List with Header Node



Advantage: "insert after" and "delete after" can be easily done at the beginning of the list (insert_first and delete_first)

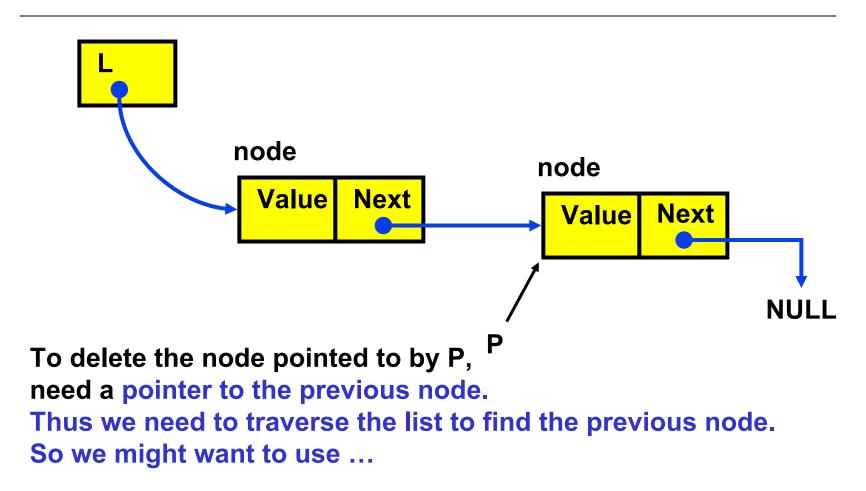
Linked list Implementation Caveats

- Whenever you break a list, your code should fix the list up as soon as possible
 - Draw pictures of the list to visualize what needs to be done
- Pay special attention to boundary conditions:
 - > Empty list
 - > Single item same item is both first and last
 - > Two items first, last, but no middle items

Linked List Insert Running Time

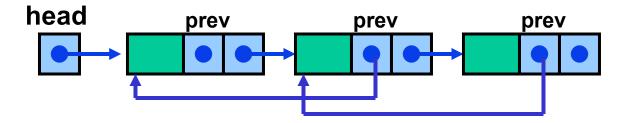
- Running time for N elements?
- "Insert_after" takes constant time (O(1))
- Does not depend on input size
- Compare to array based list Insert_kth which is O(N)
- However, how about Insert_last?

Linked List Delete



Doubly Linked Lists

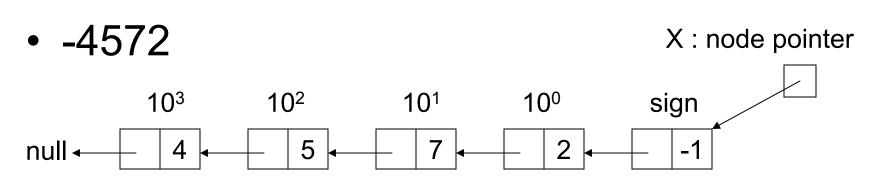
- In singly linked lists, findPrevious (and hence Delete) is slow [O(N)] because we cannot go directly to previous node
- Solution: Keep a "previous" pointer at each node



Double Link Pros and Cons

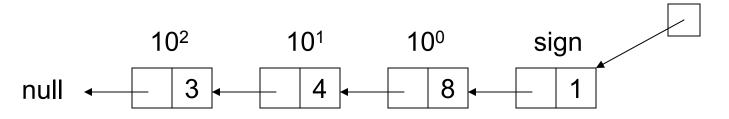
- Advantage
 - > Delete (not DeleteAfter) and FindPrev are faster
- Disadvantages:
 - More space used up (double the number of pointers at each node)
 - More book-keeping for updating the two pointers at each node (pretty negligible overhead)

Unbounded Integers Base 10

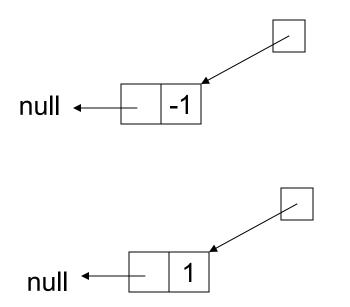


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Y : node pointer

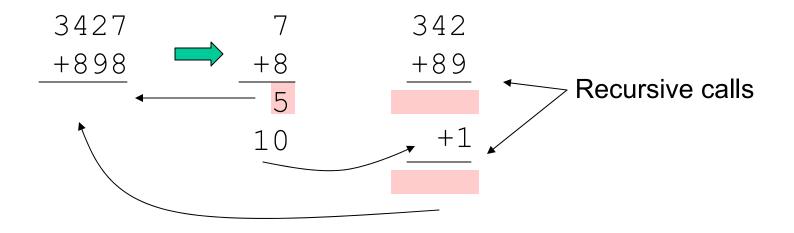


Zero



Recursive Addition

• Positive numbers (or negative numbers)



Recursive Addition

• Mixed numbers

