

CSE 373 – Data Structures and Algorithms
Autumn 2003

Dry assignment #1.

Due date: 10/10/03 (see submission instructions in course web-page).

1. (12 points) Explain in at most two sentences what is wrong, if anything, with the following induction proof.
- Claim: All horses are the same color.
 - Proof: We prove this by showing that any set of horses contains only horses of a single color; in particular, this is true for the set of all horses. Let H be an arbitrary set of horses. We show by induction on n , the number of horses in H , that all horses in H are the same color.
 - Basis: The cases $n = 0$ and $n = 1$ are immediately seen to be true.
 - Induction step: Consider any number n of horses in H . Call these horses $h_1, h_2, h_3, \dots, h_n$. By the induction hypothesis, any set of $n-1$ horses contains only horses of a single color. Consider the set H_1 obtained by removing horse h_1 from H , and the set H_2 obtained by removing horse h_2 from H . There are $n-1$ horses in each of these sets, hence the induction hypothesis applies to each: H_1 has all horses of a single color (say, c_1), and H_2 has all horses of a single color (say, c_2). But, since horse h_n is common to *both* sets H_1 and H_2 , the two colors c_1 and c_2 must be the same. This completes the induction step.

2. (24 points) True or False? Give a brief explanation.

a. $\sum_{k=1}^n k = O(n)$

b. $\sum_{k=1}^n k = \Omega(n)$

c. $2^n = \Theta(3^n)$

d. $3n^2 + n + n \cdot \log(n) = \Omega(n^2)$

e. $3n^2 + n + n \cdot \log(n) = \Omega(n \cdot \log(n))$

f. $\frac{n^2}{2^n} = O(1)$

3. (20 points) For each of the following questions, briefly explain your answer.

- a. If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes $O(n)$ on some inputs?
- b. If I prove that an algorithm takes $O(n^2)$ worst-case time, is it possible that it takes $O(n)$ on all inputs?
- c. If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes $O(n)$ on some inputs?
- d. If I prove that an algorithm takes $\Theta(n^2)$ worst-case time, is it possible that it takes $O(n)$ on all inputs?

4. (24 points) The sequence 0,1,1,3,5,11,21,43 is given by

$$S_0=0, S_1=1.$$

$$S_k=S_{k-1}+2*S_{k-2} \quad (k>1)$$

Write (in pseudocode) a function that gets as input a **sorted** array $a[]$ and its length n and returns the maximal index k such that all the numbers S_0, S_1, \dots, S_k appear in $a[]$ (or -1 if such an index does not exist).

The time complexity of your function should be $O(n)$. The space complexity should be $O(1)$.

Examples: For $a=\{0, 1, 1, 2, 3, 4, 5, 11, 15, 17, 21, 56, 67\}$ the returned value should be 6 (S_0, S_1, \dots, S_6 are in the array).

For $a=\{-5, -2, 0, 1, 1, 1, 2, 3, 3, 4, 5, 5\}$ the returned value should be 4

For $a=\{-8, 1, 2, 3, 4, 5\}$ the returned value should be -1

5. (20 points) Write (in pseudocode) a recursive function 'MaxPair' that gets an array $a[]$ of integers and its size n (it is known that $n>1$), and returns the maximal sum of two consecutive elements in $a[]$ (that is $\text{Max}(a[j-1]+a[j]) : 1 \leq j \leq n-1$). You are not allowed to use loops in your solution.

What is the time and space complexity?