Fundamentals

CSE 373
Data Structures
Lecture 5

Mathematical Background

- Today, we will review:
  - Logs and exponents
  - Series
  - Recursion
  - Motivation for Algorithm Analysis

Powers of 2

- Many of the numbers we use will be powers of 2
- Binary numbers (base 2) are easily represented in digital computers
  - each "bit" is a 0 or a 1
  - $2^0=1$, $2^1=2$, $2^2=4$, $2^3=8$, $2^4=16$, $2^8=256$, …
  - an n-bit wide field can hold 2^n positive integers:
    - $0 \leq k \leq 2^n-1$

Unsigned binary numbers

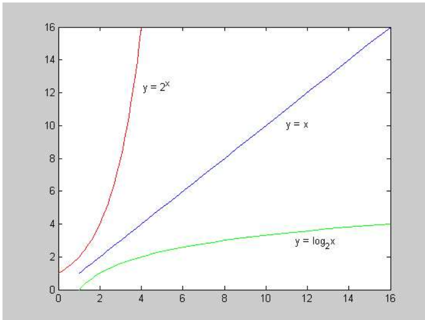
- Each bit position represents a power of 2
- For unsigned numbers in a fixed width field
  - the minimum value is 0
  - the maximum value is 2^n-1, where n is the number of bits in the field
- Fixed field widths determine many limits
  - 5 bits = 32 possible values ($2^5 = 32$)
  - 10 bits = 1024 possible values ($2^{10} = 1024$)

Binary and Decimal

$2^8=256$	$2^7=128$	$2^6=64$	$2^5=32$	$2^4=16$	$2^3=8$	$2^2=4$	$2^1=2$	$2^0=1$	$Decimal_{10}$
							1	1	3
					1	0	0	1	9
					1	0	1	0	10
					1	1	1	1	15
				1	0	0	0	0	16
				1	1	1	1	1	31
		1	1	1	1	1	1	1	127
	1	1	1	1	1	1	1	1	255

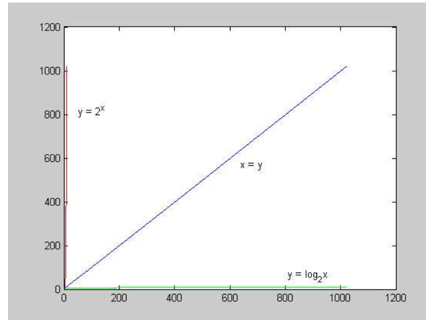
Logs and exponents

- Definition: $\log_2 x = y$ means $x = 2^y$
  - the log of x, base 2, is the value y that gives $x = 2^y$
  - $8 = 2^3$, so $\log_2 8 = 3$
  - $65536= 2^{16}$, so $\log_2 65536 = 16$
- Notice that $\log_2 x$ tells you how many bits are needed to hold x values
  - 8 bits holds 256 numbers: 0 to 2^8-1 = 0 to 255
  - $\log_2 256 = 8$




```
x = 0:.1:4
y = 2.^x
plot(x,y,'r')
hold on
plot(y,x,'g')
plot(y,y,'b')
```

2^x and $\log_2 x$




```
x = 0:10
y = 2.^x
plot(x,y,'r')
hold on
plot(y,x,'g')
plot(y,y,'b')
```

2^x and $\log_2 x$

Floor and Ceiling

$\lfloor X \rfloor$ Floor function: the largest integer $\leq$ X

$\lfloor 2.7 \rfloor = 2$ $\lfloor -2.7 \rfloor = -3$ $\lfloor 2 \rfloor = 2$

$\lceil X \rceil$ Ceiling function: the smallest integer $\geq$ X

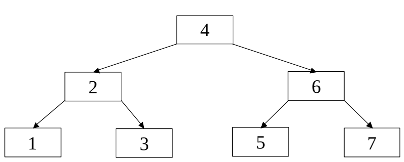
$\lceil 2.3 \rceil = 3$ $\lceil -2.3 \rceil = -2$ $\lceil 2 \rceil = 2$

Facts about Floor and Ceiling

1. $X - 1 < \lfloor X \rfloor \leq X$
2. $X \leq \lceil X \rceil < X + 1$
3. $\lfloor n/2 \rfloor + \lceil n/2 \rceil = n$ if n is an integer

Example: $\log_2 x$ and tree depth

- 7 items in a binary tree, $3 = \lfloor \log_2 7 \rfloor + 1$ levels



Properties of logs (of the mathematical kind)

- We will assume logs to base 2 unless specified otherwise
- log AB = log A + log B
  - › $A = 2^{\log_2 A}$ and $B = 2^{\log_2 B}$
  - › $AB = 2^{\log_2 A} \bullet 2^{\log_2 B} = 2^{\log_2 A + \log_2 B}$
  - › so $\log_2 AB = \log_2 A + \log_2 B$
  - › note: $\log AB \neq \log A \bullet \log B$

Other log properties

- log A/B = log A – log B
- log (A^B) = B log A
- log log X < log X < X for all X > 0
  - log log X = Y means $2^{2^Y} = X$
  - log X grows slower than X
    - called a "sub-linear" function

A log is a log is a log

- Any base x log is equivalent to base 2 log within a constant factor

$$\log_x B = \log_x B$$
$$B = 2^{\log_2 B} \qquad x^{\log_x B} = B$$
$$x = 2^{\log_2 x} \qquad (2^{\log_2 x})^{\log_x B} = 2^{\log_2 B}$$
$$2^{\log_2 x \log_x B} = 2^{\log_2 B}$$
$$\log_2 x \log_x B = \log_2 B$$
$$\log_x B = \frac{\log_2 B}{\log_2 x}$$

Arithmetic Series

- $S(N) = 1 + 2 + \ldots + N = \sum_{i=1}^{N} i$
- The sum is
  - S(1) = 1
  - S(2) = 1+2 = 3
  - S(3) = 1+2+3 = 6
- $\sum_{i=1}^{N} i = \frac{N(N+1)}{2}$

Why is this formula useful?

Algorithm Analysis

- Consider the following program segment:

```
x:= 0;
for i = 1 to N do
  for j = 1 to i do
    x := x + 1;
```

- What is the value of x at the end?

Analyzing the Loop

- Total number of times x is incremented is executed =

$$1 + 2 + 3 + \ldots = \sum_{i=1}^{N} i = \frac{N(N+1)}{2}$$

- Congratulations - You've just analyzed your first program!
  - Running time of the program is proportional to N(N+1)/2 for all N
  - O(N^2)

Analyzing Mergesort

```
Mergesort(p : node pointer) : node pointer {
Case {
  p = null : return p; //no elements
  p.next = null : return p; //one element
  else
    d : duo pointer; // duo has two fields first,second
    d := Split(p);
    return Merge(Mergesort(d.first),Mergesort(d.second));
}
}
```

T(n) is the time to sort n items.

$T(0), T(1) \le c$

$T(n) \le T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + dn$

Mergesort Analysis Upper Bound

$$
\begin{aligned}
T(n) &\le 2T(n/2) + dn \qquad \text{Assuming } n \text{ is a power of 2}\\
&\le 2(2T(n/4) + dn/2) + dn\\
&= 4T(n/4) + 2dn\\
&\le 4(2T(n/8) + dn/4) + 2dn\\
&= 8T(n/8) + 3dn\\
&\vdots\\
&\le 2^k T(n/2^k) + kdn\\
&= nT(1) + kdn \qquad \text{if } n = 2^k\\
&\le cn + dn\log_2 n\\
&= O(n \log n)
\end{aligned}
$$

Recursion Used Badly

- Classic example: Fibonacci numbers F_n

0,1, 1, 2, 3, 5, 8, 13, 21, …

Leonardo Pisano Fibonacci (1170-1250)

› $F_0 = 0$, $F_1 = 1$ (Base Cases)

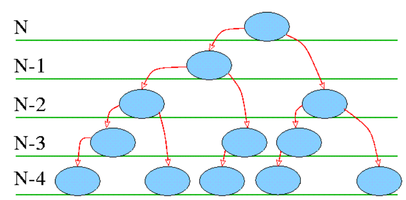
› Rest are sum of preceding two $F_n = F_{n-1} + F_{n-2}$ $(n > 1)$

Recursive Procedure for Fibonacci Numbers

```
fib(n : integer): integer {
  Case {
  n ≤ 0 : return 0;
  n = 1 : return 1;
  else : return fib(n-1) + fib(n-2);
  }
  }
```

- Easy to write: looks like the definition of F_n
- But, can you spot the big problem?

Recursive Calls of Fibonacci Procedure



- Re-computes fib(N-i) multiple times!

Fibonacci Analysis Lower Bound

T(n) is the time to compute fib(n).

$T(0), T(1) \ge 1$

$T(n) \ge T(n-1) + T(n-2)$

It can be shown by induction that $T(n) \geq \phi^{n-2}$ where

$$\phi = \frac{1+\sqrt{5}}{2} \approx 1.62$$

Iterative Algorithm for Fibonacci Numbers

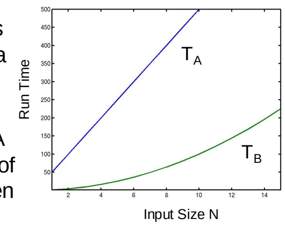
```
fib_iter(n : integer): integer {
fib0, fib1, fibresult, i : integer;
fib0 := 0; fib1 := 1;
case {_
  n < 0 : fibresult := 0;
  n = 1 : fibresult := 1;
  else :
    for i = 2 to n do {
      fibresult := fib0 + fib1;
      fib0 := fib1;
      fib1 := fibresult;
    }
  }
return fibresult;
}
```

Recursion Summary

- Recursion may simplify programming, but beware of generating large numbers of calls
  - › Function calls can be expensive in terms of time and space
- Be sure to get the base case(s) correct!
- Each step must get you closer to the base case

Motivation for Algorithm Analysis

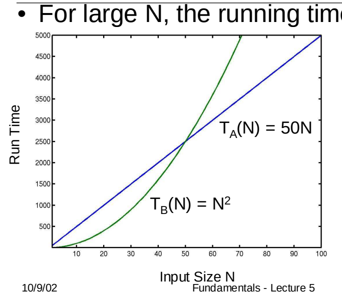
- Suppose you are given two algorithms A and B for solving a problem
- The running times $T_A(N)$ and $T_B(N)$ of A and B as a function of input size N are given



Which is better?

More Motivation

- For large N, the running time of A and B



Now which algorithm would you choose?

Asymptotic Behavior

- The "asymptotic" performance as $N \to \infty$, regardless of what happens for small input sizes N, is generally most important
- Performance for small input sizes may matter in practice, if you are <u>sure</u> that <u>small</u> N will be common <u>forever</u>
- We will compare algorithms based on how they scale for large values of N

Order Notation

- Mainly used to express upper bounds on time of algorithms. "n" is the size of the input.
- $T(n) = O(f(n))$ if there are constants c and n_0 such that $T(n) \leq c\, f(n)$ for all $n \geq n_0$.
  - › $10000n + 10\ n \log_2 n = O(n \log n)$
  - › $.00001\ n^2 \neq O(n \log n)$
- Order notation ignores constant factors and low order terms.

Why Order Notation

- Program performance may vary by a constant factor depending on the compiler and the computer used.
- In asymptotic performance ($n \to \infty$) the low order terms are negligible.

Some Basic Time Bounds

- Logarithmic time is O(log n)
- Linear time is O(n)
- Quadratic time is $0(n^2)$
- Cubic time is $O(n^3)$
- Polynomial time is $O(n^k)$ for some k.
- Exponential time is $O(c^n)$ for some c > 1.

Kinds of Analysis

- Asymptotic – uses order notation, ignores constant factors and low order terms.
- Upper bound vs. lower bound
- Worst case – time bound valid for all inputs of length n.
- Average case – time bound valid on average – requires a distribution of inputs.
- Amortized – worst case time averaged over a sequence of operations.
- Others – best case, common case, cache miss