

## More Lists

CSE 373  
Data Structures  
Lecture 4

## Recall Unbounded Integers

- -4572
  - 10<sup>3</sup>
  - 10<sup>2</sup>
  - 10<sup>1</sup>
  - 10<sup>0</sup>X : node pointer  
null ← [ ] 4 ← [ ] 5 ← [ ] 7 ← [ ] 2 ← [ ] -1 ← sign
- 348
  - 10<sup>2</sup>
  - 10<sup>1</sup>
  - 10<sup>0</sup>Y : node pointer  
null ← [ ] 3 ← [ ] 4 ← [ ] 8 ← [ ] 1 ← sign
- Zero
  - 10<sup>3</sup>
  - 10<sup>2</sup>
  - 10<sup>1</sup>
  - 10<sup>0</sup>null ← [ ] 1 ← null ← [ ] -1 ← sign

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## Alternative Addition

- Use an auxiliary function
  - › AddAux(p,q : node pointer, cb : integer)  
which returns the result of adding p and q and the carry/borrow cb.
  - › Add(p,q) := Add(p,q,0)
  - › Advantage: more like what we learned in grade school.

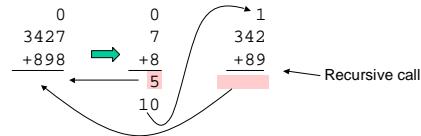
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## Auxiliary Addition

- Positive numbers (or negative numbers)



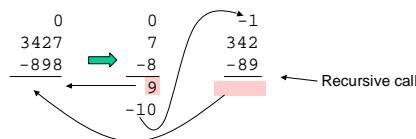
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## Auxiliary Addition

- Mixed numbers



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## Copy

- Class participation.
- Design a recursive algorithm to make a copy of an integer.

```
Copy(p : node pointer) : node pointer {  
    ???  
}
```

next value  
node [ ] [ ]

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## Comparing Integers

```

IsZero(p : node pointer) : boolean {
    return p.next = null;
}
IsPositive(p : node pointer) : boolean {
    return not IsZero(p) and p.value = 1;
}
Negate(p : node pointer) : node pointer { //destructive
    if p.value = 1 then p.value := -1
    else p.value := 1;
    return q;
}
LessThan(p,q :node pointer) : boolean {
    p1,q1 : node pointer;
    p1 := Copy(p); q1 := Copy(q);
    return IsPositive(Add(q1,Negate(p1))); // x < y iff 0 < y - x
        //We assume Add and Negate are destructive
}

```

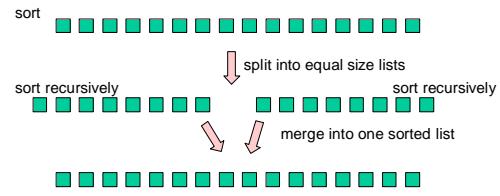
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## List Mergesort

- Overall sorting plan



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## Mergesort pseudocode

```

Mergesort(p : node pointer) : node pointer {
    Case {
        p = null : return p; //no elements
        p.next = null : return p; //one element
        else
            d : duo pointer; // duo has two fields first,second
            d := Split(p);
            return Merge(Mergesort(d.first),Mergesort(d.second));
    }
}

```

Note: Mergesort is destructive.



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## Split

```

Split(p : node pointer) : duo pointer {
    d : duo pointer;
    Case {
        p = null : d := new duo; return d
        p.next = null : d := new duo; d.first := p ; return d
        else :
            d := Split(p.next.next);
            p.next.next := d.first;
            d.first := p.next;
            p.next := d.second;
            d.second := p;
            return d;
    }
}

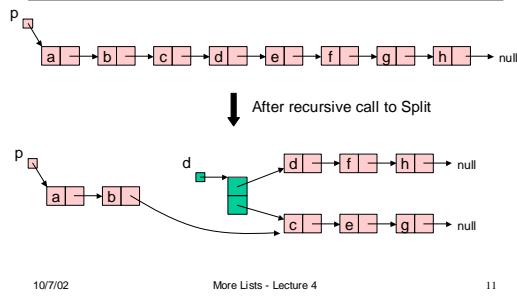
```

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## Split Example

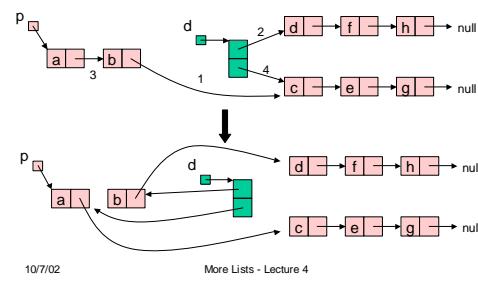


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## Split Example



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## Merge

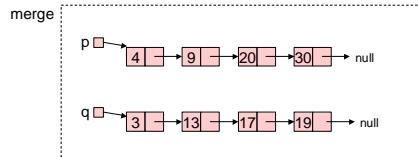
```
Merge(p,q : node pointer): node pointer{
    case {
        p = null : return q;
        q = null : return p;
        LessThan(p.value,q.value) :
            p.next := Merge(p.next,q);
            return p;
        else :
            q.next := Merge(p,q.next);
            return q;
    }
}
```

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## Merge Example

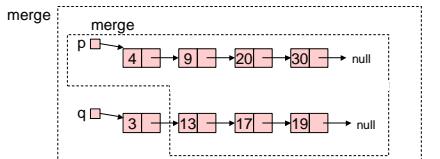


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## Merge Example

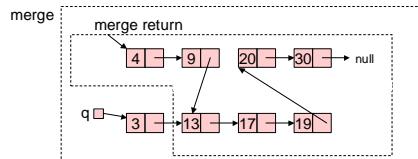


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## Merge Example



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## Implementing Pointers in Arrays – “Cursor Implementation”

- This is needed in languages like Fortran, Basic, and assembly language
- Easiest when number of records is known ahead of time.
- Each record field of a basic type is associated with an array.
- A pointer field is an unsigned integer indicating an array index.

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## Idea

Pointer World	Nonpointer World	
n nodes	D N	
data next	1 2 3 4 5 . . n	• D[ ] : basic type array
		• N[ ] : integer array
data : basic type		• Pointer is an integer
next : node pointer		• null is 0
		• p.data is D[p]
		• p.next is N[p]
		• Free list needed for node allocation

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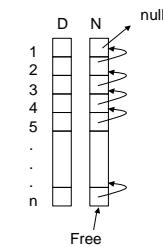
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## Initialization

Free = n	D	N
1		0
2		1
3		2
4		3
5		4
.		
n		n-1

means  
↔



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## Example of Use

L → a → b → c → null

n = 8	D	N
L = 4		
Free = 7		
1		3
2	c	0
3		1
4	a	6
5		8
6	b	2
7		5
8		1

```
InsertFront(L : integer, x : basic type) {
    q : integer;
    if not(Free = 0) then q := Free
    else return "overflow";
    Free := N[Free];
    D[q] := x;
    N[q] := L;
    L := q;
}
```

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## Try DeleteFront

- Class Participation
  - Define the cursor implementation of DeleteFront which removes the first member of the list when there is one.
- › Remember to add garbage to free list.

```
DeleteFront(L : integer) {
    ???
}
```

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## Copy Solution

```
Copy(p : node pointer) : node pointer {
    if p = null then return null
    else {
        q : node pointer;
        q := new node;
        q.value := p.value;
        q.next := Copy(p.next);
        return q;
    }
}
```

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## DeleteFront Solution

```
DeleteFront(L : integer) {
    q : integer;
    if L = 0 then return "underflow"
    else {
        q := L;
        L := N[L];
        N[q] := Free;
        Free := q;
    }
}
```

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