

DESIGN OF DIGITAL CIRCUITS AND SYSTEMS

Finite State Machine Review

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Relevant Course Information

- ❖ HW1 due on Monday (4/6)
 - Homework can be completed in groups of up to **4**
- ❖ Lab 1 report due Friday (4/10)
 - Labs can be completed in groups of up to **2**
- ❖ Lab demos:
 - Lab demo sign up sheet sent out soon (check with partner)
 - 15 minutes for demos, early labs will be quicker
 - Make sure LabsLand is set up and synthesized *beforehand*
- ❖ Quiz 1 is Thursday, April 9 in last 25 min of lecture
 - Draw FSM state diagram & make design decisions

Lecture 1 Review

❖ Useful operators:

- Ternary operator: `<cond> ? <then> : <else>`
- Concatenation: `{sig, ..., sig}`
- Replication: `{n{m}}`

❖ A **parameter** is a named constant

```
parameter N = 8;           // bus width
parameter period = 100;   // timing constant
```

❖ A parameterized module:

- `module <name> #(<parameter list>) (<port list>);`
- Parameters can be given default values
 - *e.g.*, `#(parameter N = 8)`

Review Question

- ❖ There are two forms of bit extensions: zero-extension (add 0s) and sign-extension (copy MSB)

■ **4-bit to 8-bit examples:**

0b <u>0</u> 100	<u>zero</u> →	0b <u>0000</u> 0100
	<u>sign</u> →	0b <u>0000</u> 0100
0b <u>1</u> 100	<u>zero</u> →	0b <u>0000</u> 1100
	<u>sign</u> →	0b <u>1111</u> 1100

- ❖ Write out SystemVerilog pseudocode for a parameterized *extender* module

- Inputs `sign` (1 bit), `in` (M bits); output `out` (N bits $> M$)
- `out` should either be the sign-extended version of `in` (`sign = 1`) or the zero-extended version of `in` (`sign = 0`)

Review Question (Possible) Solution

```

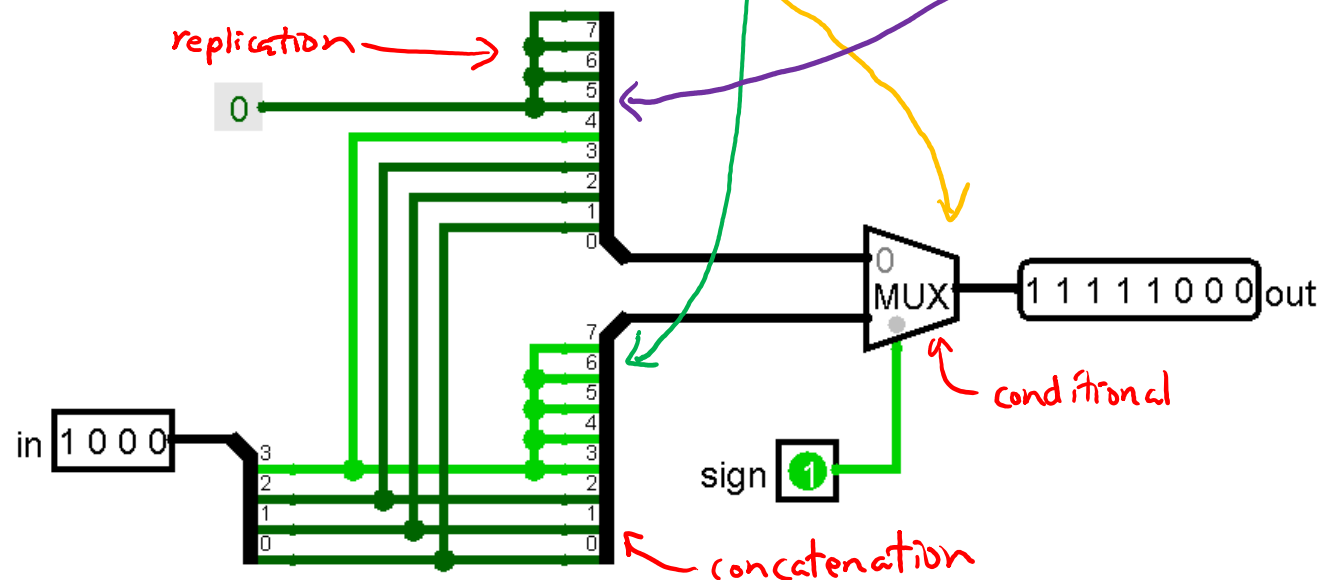
module extender #(parameter M = 4, N = 8)
    (output [N-1:0] out,
     input [M-1:0] in,
     input sign);
    assign out = sign ? {{(N-M){in[M-1]}},in} : {{(N-M){1'b0}},in};
endmodule // extender
    
```

conditional/ternary

replication

concatenation

❖ Hardware if $M = 4$ and $N = 8$:



Lecture Outline (1/3)

- ❖ **SystemVerilog Review & Tips (Continued)**
- ❖ FSMs
- ❖ Test Benches

Structural vs. Behavioral Revisited

- ❖ Not a strict definition of these terms, so exact classification is not that important
- ❖ Structural:
 - Instantiating modules (library and user-defined) and defining port connections
 - `assign`: continuous assignment
 - Used with nets

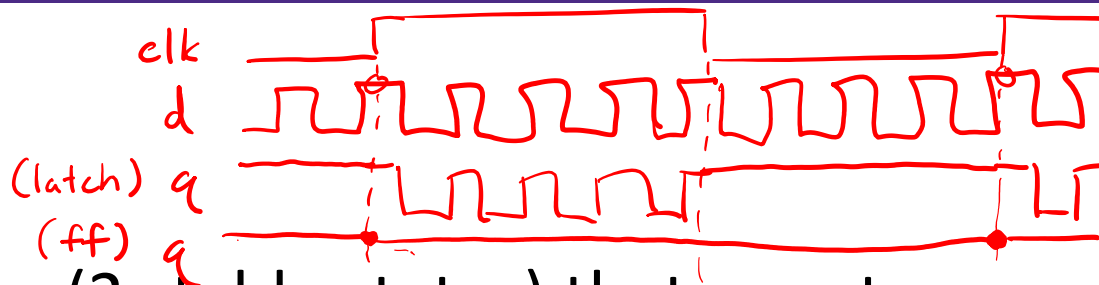
Verilog Procedural Blocks

- ❖ A *procedural block* is made up of behavioral code in the form of procedural statements whose effects are interpreted sequentially
 - The block itself is awakened/triggered in a non-sequential manner
- ❖ **initial**: Block triggered once at time zero
 - Non-synthesizable (*i.e.*, for simulation/testbenches only)
- ❖ **always**: Block triggered by a *sensitivity list* *always @ (_____)*
 - Any object that is assigned a value in an **always** statement must be declared as a variable (*e.g.*, **logic** or **reg**).

SystemVerilog Procedural Blocks

- ❖ SystemVerilog introduced variants on always that are generally more robust and more specialized
- ❖ `always_comb`: Block intended for combinational logic ~~@(*)~~
 - Sensitivity list is automatically built
- ❖ `always_latch`: Block intended for latch-based logic
 - Sensitivity list is automatically built
- ❖ `always_ff`: Block intended for sequential logic (*i.e.*, synchronous/clocked) `@(posedge clk)`
 - Sensitivity list must be specified

Latch vs. Flip-Flop



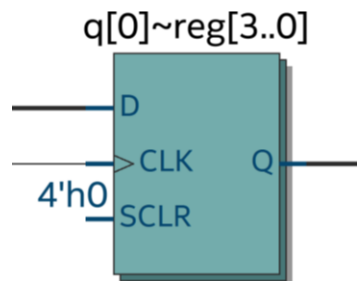
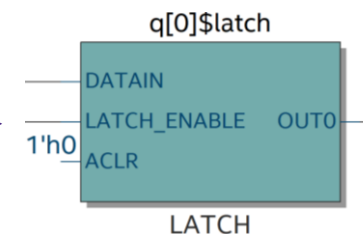
- ❖ Both are bistable multivibrators (2 stable states) that can store information
- ❖ A latch is *asynchronous*; a flip-flop is *edge-triggered*

```

module my_latch(input logic clk,
               input logic [3:0] d,
               output logic [3:0] q);

    always_latch
        if (clk) q <= d;
endmodule
    
```

q will update to d as long as clk = 1
← "latched" when clk = 0



```

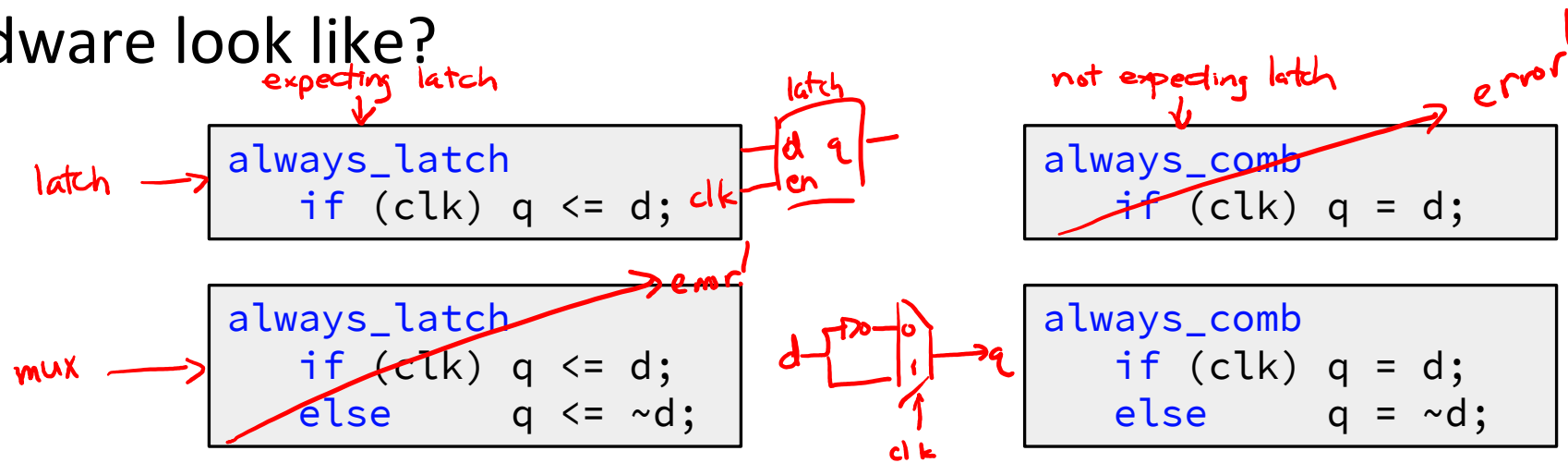
module my_ff(input logic clk,
            input logic [3:0] d,
            output logic [3:0] q);

    always_ff @(posedge clk)
        q <= d;
endmodule
    
```

only updates at rising edge of clock

Inferred Latches

- ❖ **Warning:** Easy to write code with inadvertent latches
 - Check your synthesis output for “Inferred latch”
 - Usually from incomplete assignments – unspecified branch infers latch behavior
- ❖ **Question:** Which of the following will synthesize and, if so, what will the hardware look like?



- Demo: Tools → “Netlist Viewers” → “RTL Viewer”

case Statement (1/2)

- ❖ Create combinational logic and is easier to read than lots of `if/else` statements
 - Must always be inside an `always` block
 - Each case has an implied C-style break

```
module seven_seg(bcd, segs);  
    input logic [3:0] bcd;  
    output logic [6:0] segs;  
  
    always_comb  
        case (bcd)  
            //                abc_defg  
            0: segs = 7'b011_1111;  
            1: segs = 7'b000_0110;  
            2: segs = 7'b101_1011;  
            3: segs = 7'b100_1111;  
            4: segs = 7'b110_0110;  
            5: segs = 7'b110_1101;  
            6: segs = 7'b111_1101;  
            7: segs = 7'b000_0111;  
            8: segs = 7'b111_1111;  
            9: segs = 7'b110_1111;  
  
        endcase  
  
    endmodule
```

case Statement (2/2)

- ❖ Create combinational logic and is easier to read than lots of `if/else` statements
 - Must always be inside an `always` block
 - Each case has an implied C-style break
 - Remember to use `default` to avoid incomplete assignments!

```
module seven_seg(bcd, segs);  
    input logic [3:0] bcd;  
    output logic [6:0] segs;  
  
    always_comb  
        case (bcd)  
            //                abc_defg  
            0: segs = 7'b011_1111;  
            1: segs = 7'b000_0110;  
            2: segs = 7'b101_1011;  
            3: segs = 7'b100_1111;  
            4: segs = 7'b110_0110;  
            5: segs = 7'b110_1101;  
            6: segs = 7'b111_1101;  
            7: segs = 7'b000_0111;  
            8: segs = 7'b111_1111;  
            9: segs = 7'b110_1111;  
            → default: segs = 7'bX;  
        endcase  
  
endmodule
```

Other SystemVerilog Resources

- ❖ SystemVerilog Language Reference Manual
 - On website, Verilog → Reference Manual (586 pages...)
 - 371 digest, Verilog → Tutorial (16 pages!)
- ❖ SystemVerilog articles
 - <https://www.systemverilog.io/>
 - <http://www.verilogpro.com/>
 - <https://www.chipverify.com/systemverilog/systemverilog-tutorial>
- ❖ One style guide for SystemVerilog
 - <https://www.systemverilog.io/styleguide>
 - We won't enforce, but good guidelines

TECHNOLOGY

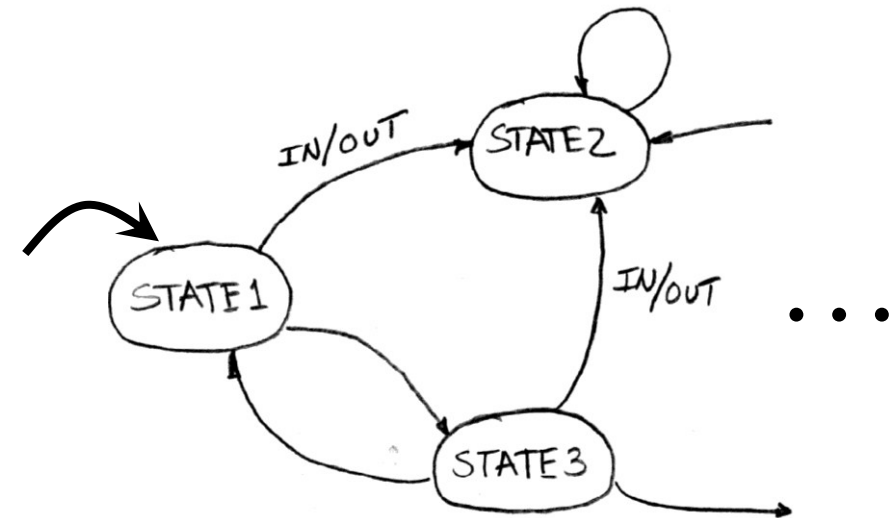
BREAK

Lecture Outline (2/3)

- ❖ SystemVerilog Review & Tips (Continued)
- ❖ **Finite State Machine Design**
- ❖ Test Benches

Finite State Machines (FSMs)

- ❖ A convenient way to conceptualize computation over time using a *state transition diagram*
 - Consists of a *set of states*, an *initial state*, and a *transition function*
- ❖ FSM implementations come in 3 blocks:
 - State register (SL)
 - Next state logic (CL)
 - Output logic (CL)



FSM Implementation Notes

- ❖ States must be assigned a binary encoding
 - More readable by using parameters or an enum
 - Encoding choices can affect logic simplification
- ❖ Reset signal can be synchronous (responds to `clk`) or asynchronous (responds to `reset`)
 - Determined by whether or not `reset` is in sensitivity list
- ❖ State logic (next state logic + state update) can be written as 1 combined block or 2 separate blocks
- ❖ If input is asynchronous, may want to add a two-flip-flop *synchronizer* to deal with metastability



FSM SystemVerilog Design Pattern

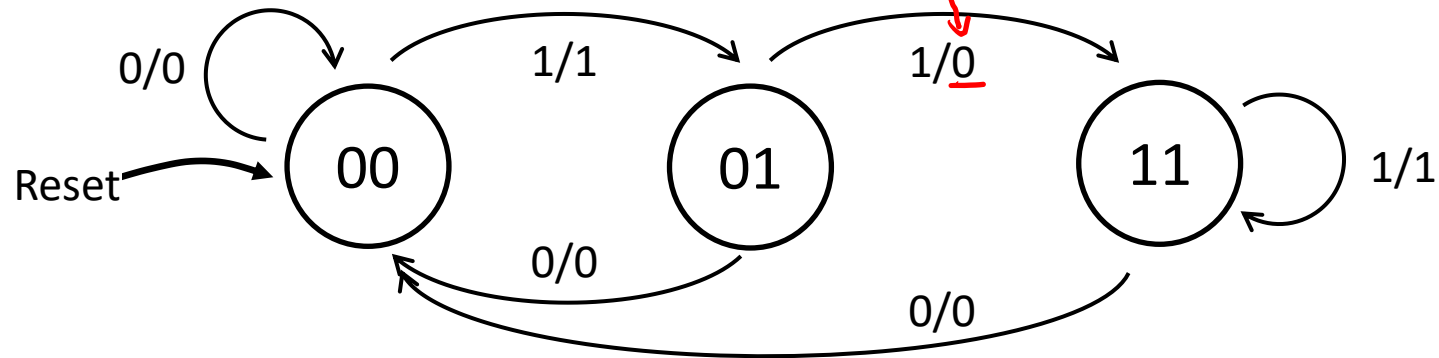
❖ Which, if any, construct(s) would you expect to use for each of the following basic sections of a module that implements an FSM?

alternatives exist!

<i>not logic!</i>	▪ // define states and state variables	<i>enum logic [N-1:0] {A, B, C} ps, ns;</i>	initial	assign	always_comb	always_ff	None
<i>combinational logic</i>	▪ // next state logic (ns)		initial	assign	always_comb	always_ff	None
<i>combinational logic</i>	▪ // output logic		initial	assign	always_comb	always_ff	None
<i>sequential logic</i>	▪ // state update logic (ps)		initial	assign	always_comb	always_ff	None

FSM Example: String Manipulator

- ❖ Takes in a stream of inputs and removes the second 1 from every consecutive string of 1's.



- Example inputs: 0 1 0 1 1 0 1 1 1 0 1 1 1 1 ...
 outputs: 0 1 0 1 0 0 1 0 1 0 1 0 1 1 ...

FSM Example: String Manipulator Implementation

```
module fsm (input logic clk, reset, in,
           output logic out);
```

```
① // present and next state
  enum logic [1:0] {S0, S1, S3} ps, ns;
```

```
② // next state logic
  always_comb
```

```
  case (ps)
```

```
    S0: if (in) ns = S1;
        else ns = S0;
```

```
    S1: if (in) ns = S3;
        else ns = S0;
```

```
    S3: if (in) ns = S3;
        else ns = S0;
```

```
  endcase
```

```
③ // output logic
```

```
  assign out = in & (ps[1] | ~ps[0]);
```

```
  ...
```

could also be: ns = in ? S1: S0;

ps != S1

```
  ...
```

```
④ // sequential logic (DFFs)
  // synchronous reset
```

```
  always_ff @(posedge clk)
```

```
    if (reset)
```

```
      ps <= S0; // reset state
```

```
    else
```

```
      ps <= ns;
```

```
endmodule // fsm
```

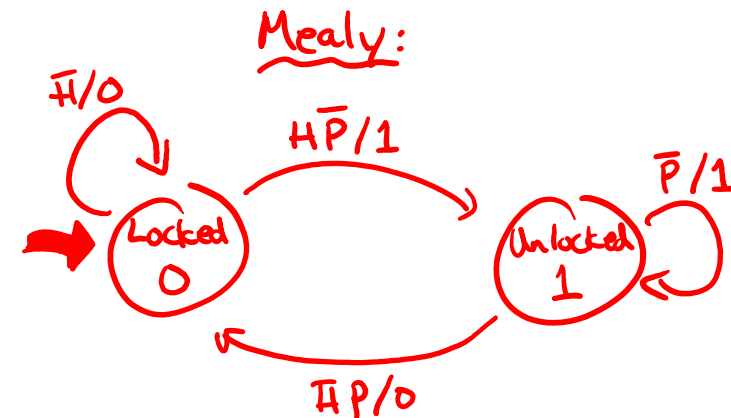
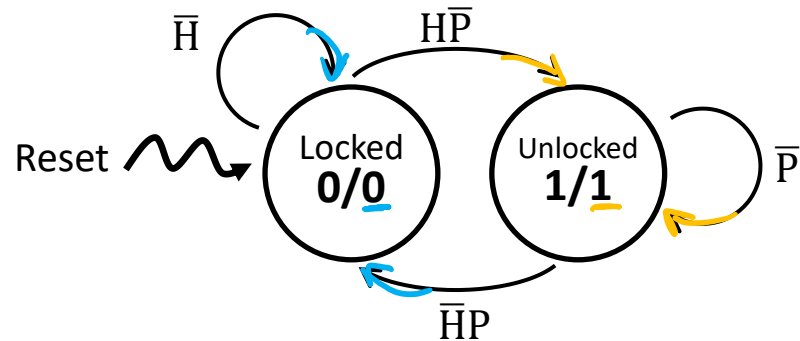
Moore vs. Mealy

- ❖ **Moore** machines define their outputs based on states ($\textcircled{00/1}$) and **Mealy** machines define outputs based on transitions ($\xrightarrow{0/1}$)
 - Mealy machines are more *flexible*
 - Moore outputs are function of state; Mealy outputs are function of state *and inputs*
 - All FSMs can be expressed in either form, but some systems are more naturally expressed one way versus the other
 - Feel free to use either in this class if not specified
- ★ However, there *are* implementation differences!

Mealy \leftrightarrow Moore Conversions (1/2)

Not testable material

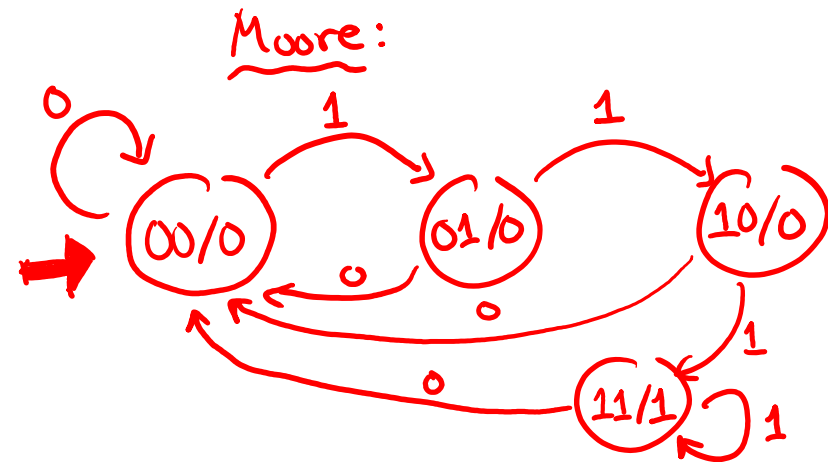
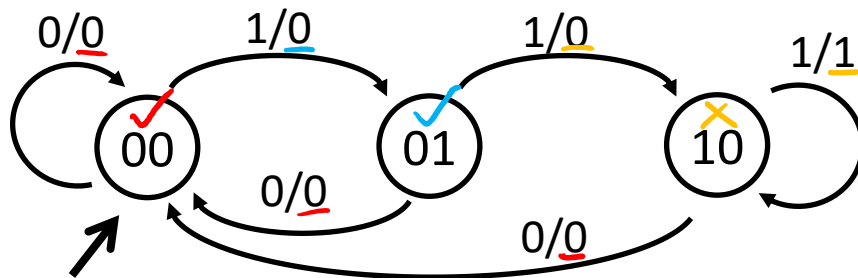
- ❖ **Moore** \rightarrow **Mealy**: Copy the state output to every transition *entering* the state
- ❖ Example: FSM for a *turnstile*, which is locked until someone swipes their Husky ID (input H) and then locks once you push through (input P) the unlocked gate. Outputs a light that glows red (0) or green (1).



Mealy \leftrightarrow Moore Conversions (2/2)

Not testable material

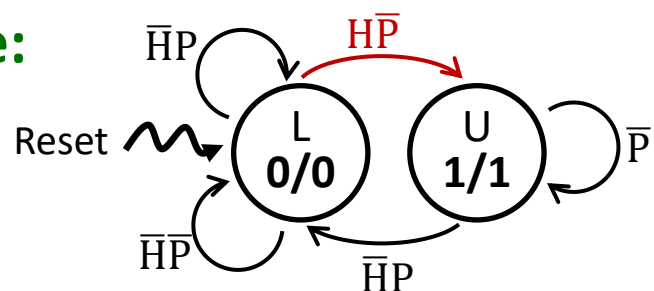
- ❖ **Mealy** \rightarrow **Moore**: More complicated process; if incoming transitions differ in output, may need to “split” the state
- ❖ Example: The threeOnes FSM from Lecture 1



Moore vs. Mealy Outputs (1/2)

- ❖ Compare a Moore and Mealy FSM for the turnstile. Complete the statements and waveform below, assuming no delays:

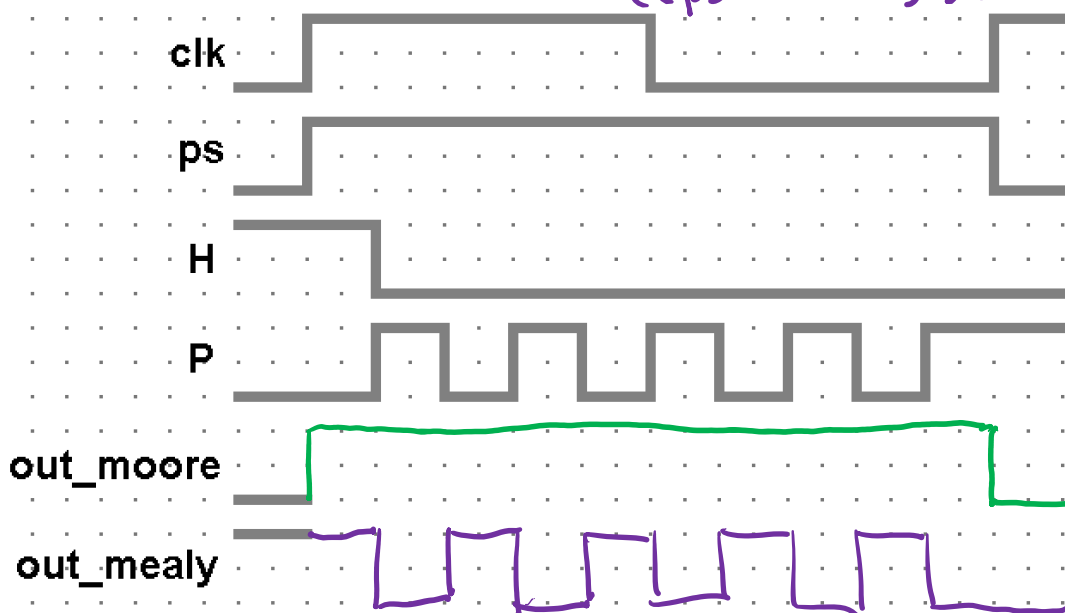
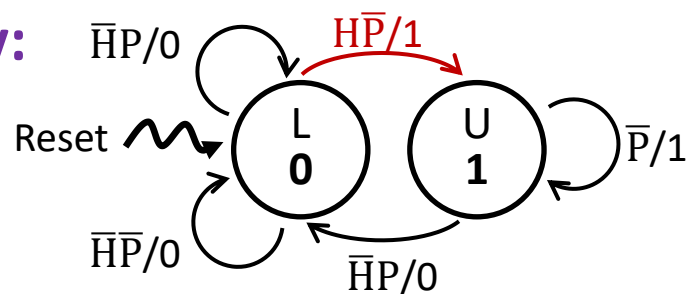
Moore:



```

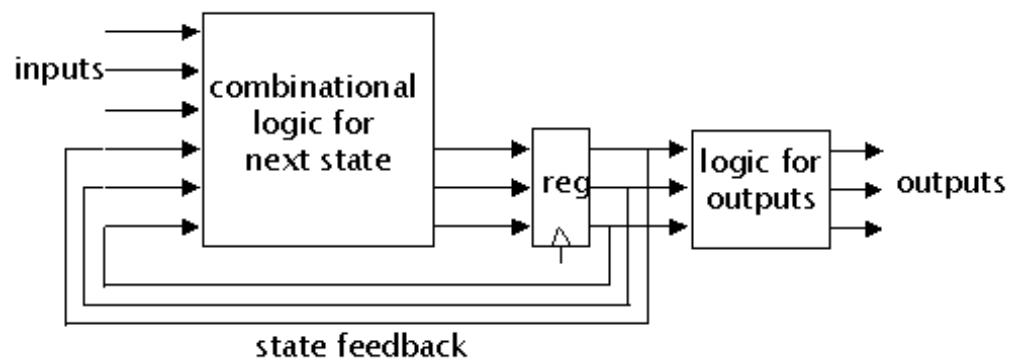
    assign out_moore = (ps == U);
    assign out_mealy = ((ps == U) & ~P) |
                       ((ps == L) & H & ~P);
  
```

Mealy:



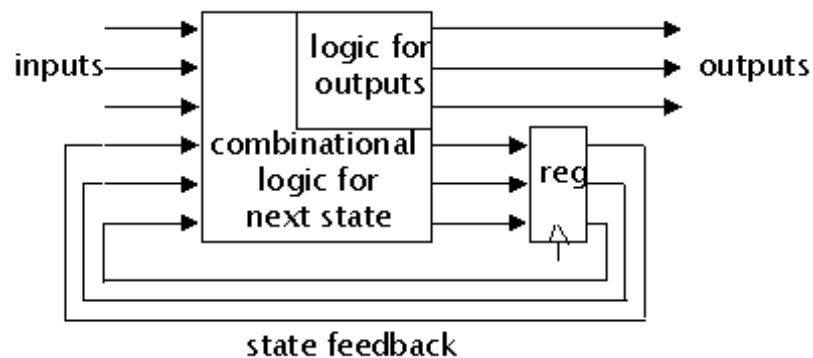
Moore vs. Mealy Outputs (2/2)

❖ Moore:



- Outputs change *synchronously* with state changes

❖ Mealy:



can add register to synchronize output changes

- Input changes can cause *immediate* output changes

(often change sooner than Moore outputs)

Lecture Outline (3/3)

- ❖ SystemVerilog Review & Tips (Continued)
- ❖ Finite State Machine Design
- ❖ **Test Benches**

Test Benches

- ❖ Special modules *needed for simulation only!*
 - Software constraint to mimic hardware
- ❖ ModelSim runs entirely on your computer
 - Tries to simulate your FPGA environment without actually using hardware – no physical signals available
 - Must create fake inputs for FPGA's physical connections
 - *e.g.*, LEDR, HEX, KEY, SW, CLOCK_50
 - Unnecessary when code is loaded onto FPGA
- ❖ Need to define both input signal combinations as well as their timing

Test Bench Timing Controls

- ❖ Delay: `#<time>` `#10` `#{period/2}`
 - Delays by a specific amount of simulation time
 - ❖ Edge-sensitive: `@(<pos/neg>edge <signal>)` `@(posedge clk)`
 - Delays next statement until specified transition on signal
 - ❖ Level-sensitive Event: `wait(<expression>)` `wait(clk)`
 - Waits until `<expression>` evaluates to TRUE
 - ❖ Stop simulation: `$stop;`
 - ❖ Timescale: ``timescale <time unit> / <precision>`
 - e.g., ``timescale 1 ns / 1 ps`
- Handwritten notes:*
- Red arrow pointing to `<time unit>`: each arbitrary simulation time unit (#1)
- Red arrow pointing to `<precision>`: rounding

Extender Test Bench

```

`timescale 1 ns / 1 ns
module extender_tb();

    parameter M = 4, N = 8;
    logic [M-1:0] in;
    logic [N-1:0] out;
    logic sign;

    extender #(M, N) dut (.*);

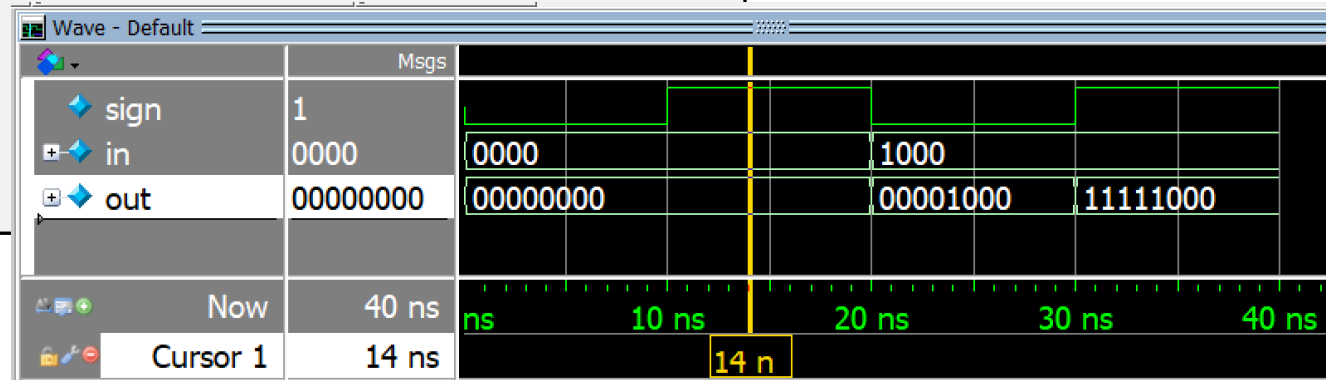
    int i;
    initial begin
        for (i = 0; i < 2**2; i++) begin
            sign = i[0]; in = {i[1], {(M-1){1'b0}}}; #10;
        end // for
        $stop;
    end // initial
endmodule // extender_tb

```

MSB of in
sign

$i = 0b0\dots000$
 $0b0\dots001$
 $0b0\dots010$
 $0b0\dots011$

exponentiation (2^2)



FSM Test Bench Notes

- ❖ Your main goal is to test every transition that we care about – may take extra clock cycles
- ❖ For simulation, you need to generate a clock signal
 - Assume we have `parameter clock_period;`

Explicit Edges:

```
initial
  clk = 0;

always_comb begin
  #(clock_period/2) clk <= 1;
  #(clock_period/2) clk <= 0;
end
```

Toggle:

```
initial begin
  clk <= 0;
  forever #(clock_period/2) clk <= ~clk;
end
```

String Manipulator Test Bench

```

module fsm_tb();

    logic clk, reset, in, out;

    fsm dut (.*);

    // simulated clock
    parameter period = 100;
    initial begin
        clk <= 0;
        forever
            #(period/2)
            clk <= ~clk;
    end // initial clock

    ...

```

```

...
initial begin
    reset <= 1; in <= 0; @(posedge clk);
    reset <= 0; in <= 0; @(posedge clk);
    in <= 0; @(posedge clk);
    in <= 1; @(posedge clk);
    in <= 0; @(posedge clk);
    in <= 1; @(posedge clk);
    in <= 1; @(posedge clk);
    in <= 0; @(posedge clk);
    in <= 1; @(posedge clk);
    in <= 1; @(posedge clk);
    in <= 1; @(posedge clk);
    @ (posedge clk);

    $stop; // end simulation
end // initial signals

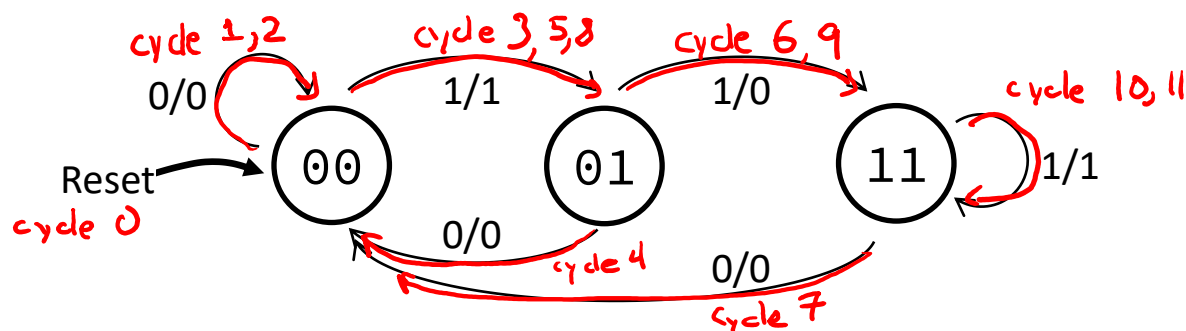
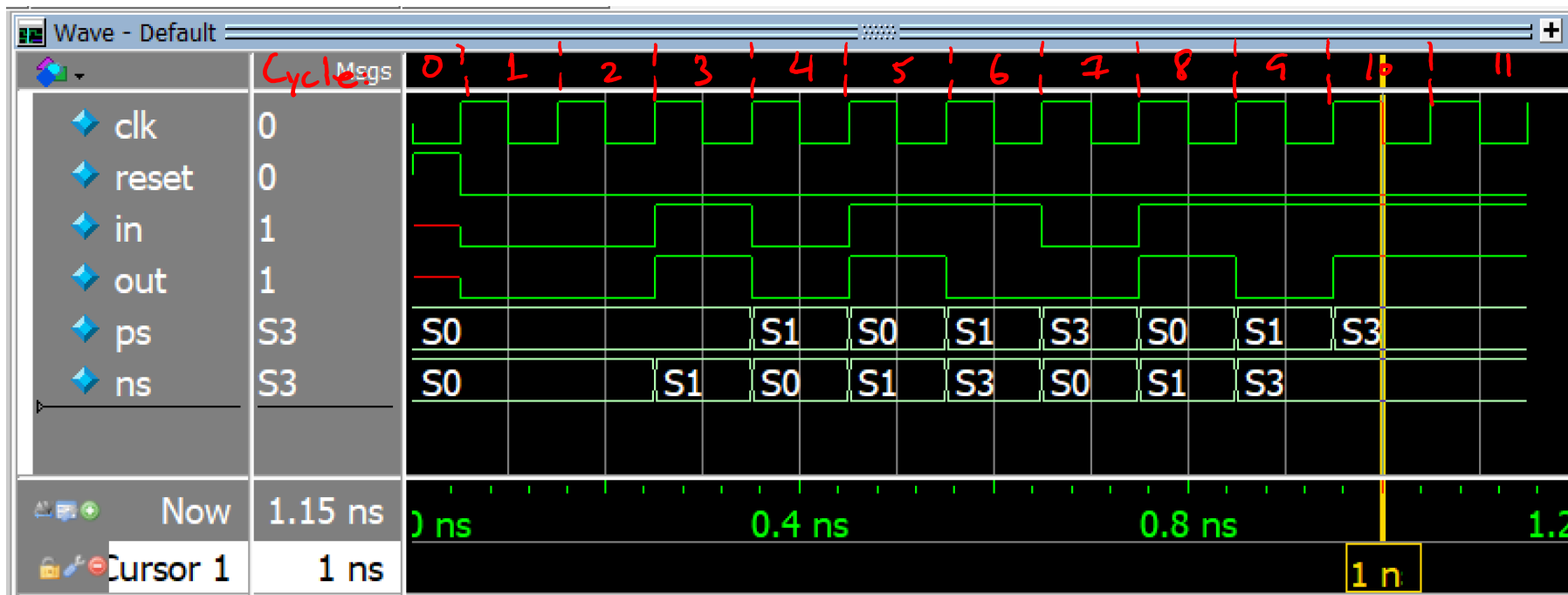
endmodule // fsm_tb

```


① signal changes

② the clock edge that reads the new signal values

String Manipulator Waveforms



Checking Responses

- ❖ Visually checking simulated waveforms quickly becomes impractical for large designs simulated over thousands of clock cycles
 - Displaying and explaining your waveforms for labs can be tedious
- ❖ There are simulator-independent system tasks to write messages to the user/tester!
 - Look similar to `printf()` in C
 - `$<system_task>(<format_string>, <sig_1>, <sig_2>, ...)`
 - Will look at `$display` today and others later on

Checking Responses: \$display

- ❖ Triggers once when encountered, prints the given format string and adds a new line:

```
// define test inputs
int i;
initial begin
  for (i = 0; i < 2**2; i++) begin
    sign = i[0]; in = {i[1], {(M-1){1'b0}}}; #10;
    $display("t = %0t, %b %s %b",
             (current simulation time) $time, in, sign ? "-+->" : "-0->", out);
  end // for
  $stop;
end // initial
```

```
Transcript
VSIM U>
# t = 10, 0000 -0-> 00000000
# t = 20, 0000 -+-> 00000000
# t = 30, 1000 -0-> 00001000
# t = 40, 1000 -+-> 11111000
```

Format Specifiers

Table 5.7: Format Specifiers.

Specifier	Meaning
%h	Hexadecimal format
%d	Decimal format (signed)
%o	Octal format
%b	Binary format
%c	ASCII character format
%v	Net signalstrength
%m	Hierarchical name of current scope
%s	String
%t	Time
%e	Real in exponential format
%f	Real in decimal format
%g	Real in exponential or decimal format

Table 5.8: Special characters.

Symbol	Meaning
\n	New line
\t	Tab
\\	\ character
\"	" character
\xyz	Where xyz is are octal digits - the character given by that octal code
%%	% character

\ escape character

- **Warning:** these differ from the specifiers for C's `printf()`
- *Minimum* field width is specified by numbers between the '%' and specifier letter
 - e.g., %3d will pad out to 3 digits if necessary,
%0d will show just the minimum number of digits needed