

Design of Digital Circuits and Systems

Testing: Assertions, OOP

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Relevant Course Information

- ❖ Quiz 4 this Thursday @ 11:40 am
 - Algorithms to Hardware
- ❖ Lab 5 report due Friday (5/23)
- ❖ Lab 6 proposal due next week (5/28)
 - (1) Describe your major project behavior, features, components/modules, and user interaction in a few paragraphs
 - (2) Include at least a top-level block diagram (preferably with signals labeled on it; other diagrams welcome)
 - (3) Include images/sketches of VGA output
 - “Proposal Workshop” in lecture on 5/27

Testbenches

- ❖ HDL module that tests another module
 - Typically called the *device under test* (dut) or *unit under test* (uut)
 - No ports (*i.e.*, inputs or outputs)
 - Not synthesizable
- ★ Note: even if written in the same HDL, testbenches may give different simulation results on different simulators

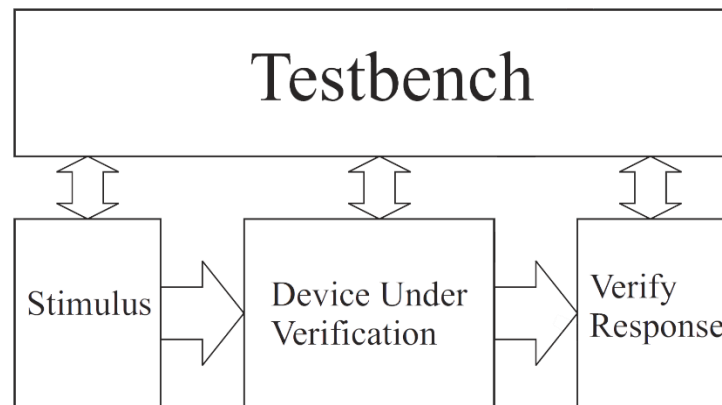


Figure 8.1: Modular testbench structure.

Test Vectors from a File

❖ Can be convenient to load test vectors from a file

- Use \$readmemb and \$readmemh
- Can also save you recompiling time! → can change values in file without recompiling

```
logic [W-1:0] test_vectors[0:15]; //array to hold test vectors
```

```
// define test inputs
```

```
integer i;
```

```
initial begin
```

```
    $readmemh("tests.txt", test_vectors);
```

```
    Reset = 1; Start = 0; @(posedge clk);
```

```
    Reset = 0; @(posedge clk);
```

```
    for (i = 0; i < 2**4; i++) begin
```

```
        Start = 1; Num = test_vectors[i]; @(posedge clk);
```

```
        Start = 0; @(posedge Ready);
```

```
    end
```

```
    @(posedge clk); // extra cycle of output
```

```
    $stop();
```

could combine signals into single vector!
{A,B,C} = test_vectors[i];

Dumping Responses

- ❖ The results of a simulation can be “dumped” to a file for later viewing in a waveform viewer or analysis
 - `$dumpfile` specifies the name of the file
 - "dump.vcd" by default (**V**alue **C**hange **D**ump)
 - Found in `<Project>\simulation\modelsim`
 - `$dumpvars` saves all of the variables from that point onward to that file
 - You can use arguments to specify which variables you want

```
// define test inputs
integer i;
initial begin

    $dumpfile("values.vcd");
    $dumpvars;

    Reset = 1; Start = 0;    @(posedge clk);
    Reset = 0;              @(posedge clk);
```

EDA Playground



- ❖ The advanced verification features we will discuss cannot be run in ModelSim so we will use EDA Playground instead
 - A web application that will let you use more powerful commercial simulators
 - Homework 6 will walk you through the registration process and a short tutorial
 - ★ ■ To use the waveform viewer in EDA playground, you *must* generate a .vcd file during your simulation! *\$dumpvars;*

Checking Responses (Review)

- ❖ Visually checking simulated waveforms quickly becomes impractical for large designs simulated over thousands of clock cycles
 - Even for isPrime, we are constantly scanning right for Done, then scanning up and down for P.
 - Displaying and explaining your waveforms for labs has been tedious for a while now
- ❖ There are simulator-independent system tasks to write messages to the user/tester!
 - Look similar to printf() in C
 - `$<system_task>(<format_string>, <sig_1>, <sig_2>, ...)`

Format Specifiers (Review)

Table 5.7: Format Specifiers.

Specifier	Meaning
%h	Hexadecimal format
%d	Decimal format (signed)
%o	Octal format
%b	Binary format
%c	ASCII character format
%v	Net signalstrength
%m	Hierarchical name of current scope
%s	String
%t	Time
%e	Real in exponential format
%f	Real in decimal format
%g	Real in exponential or decimal format

escape character
Table 5.8: Special characters.

Symbol	Meaning
\n	New line
\t	Tab
\\	\ character
\"	" character
\xyz	Where xyz is are octal digits - the character given by that octal code
%%	% character

- **Warning:** these differ from the specifiers for `printf`
- The *minimum* field width is specified by numbers between the '%' and specifier letter
 - e.g., %3d will pad out to 3 digits if necessary,
%0d will show just the minimum number of digits needed

Checking Responses: \$display (Review)

- ❖ Triggers once when encountered, prints the given format string and adds a new line:

```
// define test inputs
integer i;
initial begin
    Reset = 1; Start = 0; @(posedge clk);
    Reset = 0;          @(posedge clk);
    for (i = 0; i < 2**W; i++) begin
        Start = 1; Num = i; @(posedge clk);
        Start = 0;          @(posedge Ready);
        { $display("T = %4t, isPrime(%2d) = %s",
                  $time, Num, P ? "Yes" : "No");
          }
    end
    @(posedge clk); // extra cycle of output
    $stop();
end
```

current simulation time

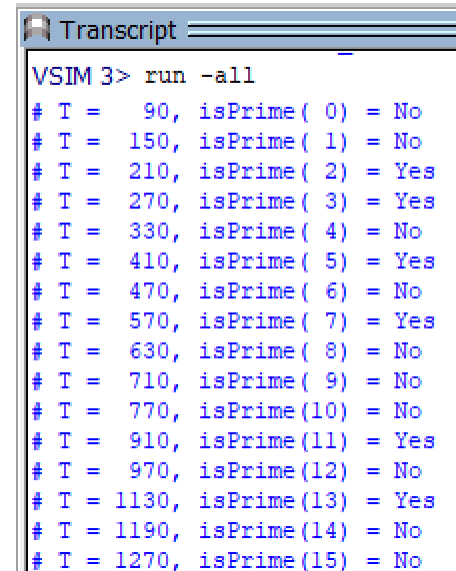
```
Transcript
VSIM 4> run -all
# T = 90, isPrime( 0) = No
# T = 150, isPrime( 1) = No
# T = 210, isPrime( 2) = Yes
# T = 270, isPrime( 3) = Yes
# T = 330, isPrime( 4) = No
# T = 410, isPrime( 5) = Yes
# T = 470, isPrime( 6) = No
# T = 570, isPrime( 7) = Yes
# T = 630, isPrime( 8) = No
# T = 710, isPrime( 9) = No
# T = 770, isPrime(10) = No
# T = 910, isPrime(11) = Yes
# T = 970, isPrime(12) = No
# T = 1130, isPrime(13) = Yes
# T = 1190, isPrime(14) = No
# T = 1270, isPrime(15) = No
```

Checking Responses: \$write

- ❖ Triggers once when encountered, prints the given format string without a new line:

```
// define test inputs
integer i;
initial begin
    Reset = 1; Start = 0; @(posedge clk);
    Reset = 0;          @(posedge clk);
    for (i = 0; i < 2**W; i++) begin
        Start = 1; Num = i; @(posedge clk);
        Start = 0;          @(posedge Ready);
        { $write("T = %4t, isPrime(%2d) = %s\n",
                $time, Num, P ? "Yes" : "No ");
        }
    end
    @(posedge clk); // extra cycle of output
    $stop();
end
```

Same messages? **Yes**



```
VSIM 3> run -all
# T = 90, isPrime( 0) = No
# T = 150, isPrime( 1) = No
# T = 210, isPrime( 2) = Yes
# T = 270, isPrime( 3) = Yes
# T = 330, isPrime( 4) = No
# T = 410, isPrime( 5) = Yes
# T = 470, isPrime( 6) = No
# T = 570, isPrime( 7) = Yes
# T = 630, isPrime( 8) = No
# T = 710, isPrime( 9) = No
# T = 770, isPrime(10) = No
# T = 910, isPrime(11) = Yes
# T = 970, isPrime(12) = No
# T = 1130, isPrime(13) = Yes
# T = 1190, isPrime(14) = No
# T = 1270, isPrime(15) = No
```

Checking Responses: \$monitor

- ❖ Triggers when encountered, then triggers anytime one of its signal changes (adds a new line):

```
// define test inputs
integer i;
initial begin
    { $monitor("T = %4t, isPrime(%2d) = %s\n",
              $time, Num, P ? "Yes" : "No ");

    Reset = 1; Start = 0; @(posedge clk);
    Reset = 0;           @(posedge clk);
    for (i = 0; i < 2**W; i++) begin
        Start = 1; Num = i; @(posedge clk);
        Start = 0;         @(posedge Ready);
    end
    @(posedge clk); // extra cycle of output
    $stop();
end
```

triggers when Num or P changes

Same messages? *No*

Transcript

```
VSIM 6> run -all
# T =    0, isPrime( x) = He
# T =   30, isPrime( 0) = He
# T =   70, isPrime( 0) = No
# T =   90, isPrime( 1) = No
# T =  150, isPrime( 2) = No
# T =  190, isPrime( 2) = Yes
# T =  210, isPrime( 3) = Yes
# T =  270, isPrime( 4) = Yes
# T =  310, isPrime( 4) = No
# T =  330, isPrime( 5) = No
# T =  390, isPrime( 5) = Yes
# T =  410, isPrime( 6) = Yes
# T =  450, isPrime( 6) = No
# T =  470, isPrime( 7) = No
# T =  550, isPrime( 7) = Yes
# T =  570, isPrime( 8) = Yes
# T =  610, isPrime( 8) = No
# T =  630, isPrime( 9) = No
# T =  710, isPrime(10) = No
# T =  770, isPrime(11) = No
# T =  890, isPrime(11) = Yes
# T =  910, isPrime(12) = Yes
# T =  950, isPrime(12) = No
# T =  970, isPrime(13) = No
```

garbage

Lecture Outline

- ❖ Testbenches (yet again)
- ❖ **Assertions**
- ❖ Object-Oriented Programming

Assertion-Based Verification

- ❖ `$display`, `$write`, `$monitor`
 - Can indicate the response of the circuit in textual form
 - Still must be *verified manually/visually*, even if you also print the expected response alongside it

- ❖ **Assertions** are SystemVerilog features that can print messages when an expected condition fails
 - `assert` – *immediate* assertion that follows simulation event semantics
 - `assert property` – *concurrent* assertion based on clock semantics

Immediate Assertions

- ❖ An **immediate assertion** is an if-else statement with a default-generated else:

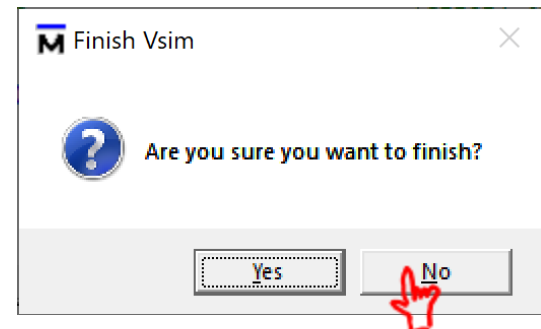
<pre>assert (P == 1);</pre>	↔	<pre>if (P == 1); // nothing if true else \$error("Assertion error.");</pre>
-----------------------------	---	--

- Must be contained inside of a procedural block
- ❖ Can also explicitly define *pass* and *fail* statements:

```
// defined pass, default fail  
assert (P == 1) $display("%2d is prime", Num);  
  
// default pass (nothing), defined fail  
assert (P == 1) else $error("%2d is not prime", Num);  
  
// defined pass, defined fail  
assert (P == 1) $display("%2d is prime", Num);  
else $error("%2d is not prime", Num);
```

Failure Messages

- ❖ Messaging: `$info`, `$warning`, `$error`
 - Ordered in increasing severity (less severe are suppressible)
 - Same argument format as `$display`, `$monitor`
 - All print additional debugging line (time, scope, file, line),
but simulation continues
- ❖ Break: `$fatal`
 - Takes an `error_code` as extra (1st) argument that is passed to `$finish`, which terminates the simulation
 - ModelSim produces this pop-up box:
 - Click “No”, otherwise ModelSim will exit



Short Tech Break

Concurrent Assertions

- ❖ **Concurrent assertions** run continuously throughout simulation based on a sampling clock and can test for much more complex behaviors
 - Do not need to be placed inside another procedural block
 - Assert that a specified property is true
 - Like immediate assertions, can specify pass/fail code
 - *Unfortunately, these do not work in ModelSim* (need a more fully-featured simulator)
- ❖ Example: assert that **Ready** and **Done** are never true at the same time

```
property ready_nand_done;  
    @(posedge clk) ~(Ready & Done);  
endproperty  
assert property (ready_nand_done);  
else // fail code
```

property definition {
assertion →
pass code (points to the assertion line)
else // fail code

Properties

❖ Defined between `property` and `endproperty`

- Includes the ability to define an argument list!

• e.g.,

```
property Nand(logic A, logic B);  
    @(posedge clk) ~(A & B);  
endproperty  
assert property (Nand(Ready, Done));
```

- Can be defined in-line, but this is stylistically discouraged

❖ Complex properties are typically active over (*i.e.*, they span) a period of time

- Specified using a combination of implications and sequences

• e.g.,

```
property handshake;  
    @(posedge clk) Req |-> ##[1:2] Ack;  
endproperty
```

← Ack should be asserted 1-2 clock cycles after Req

Implications (Mathematics)

❖ $p \Rightarrow q$ is read as “ p implies q ”

- A statement meaning: if p is true, then q must also be true
- The statement evaluates to true or false based on whether the actual values of p and q *support* the implication:

p	q	$p \Rightarrow q$
<u>false</u>	false	true
<u>false</u>	true	true
true	false	false
true	true	true

no way to contradict the implication

← contradicts the implication

- Logically equivalent to $\neg p \vee q$ or $p \rightarrow q$: 1

Implications (SystemVerilog)

- ❖ **Implications** are notated by $A \mid \rightarrow C$ and $A \mid \Rightarrow C$
 - A is the *antecedent* (LHS), C is the *consequent* (RHS)
 - The consequent is only evaluated if the antecedent is true
 - In the context of assertions and properties, evaluating to true is a pass and false is a fail
- ❖ **Implication timing:**
 - An overlapped implication ($\mid \rightarrow$) evaluates C in the *same* clock cycle that A was true
 - A non-overlapped implication ($\mid \Rightarrow$) evaluates C on the *next* clock cycle after A was true
- ❖ Practice: write an equivalent implication to $\sim(A \& B)$

evaluated simultaneously, so overlapped: $A \mid \rightarrow \sim B$

p, ~q

A	B	NAND
0	0	1
1	0	1
1	1	0

Sequences

- ❖ A **sequence** is a series of Boolean expressions with defined relationships *in time*
 - Any Boolean expression is, by itself, an implicit sequence
 - Sequences can be constructed from other sequences and *sequence operators*
 - You can name a sequence and give it arguments using **sequence** and **endsequence**
- ❖ Common sequence operators:
 - $##N$ – delays next sequence by N cycles
 - $[*N]$ – N consecutive repetitions of the LHS
 - $[=N]$ – N non-consecutive repetitions of the LHS
 - Any N can be replaced by the inclusive range $A:B$

Req \mapsto $##[1:2]$ Ack

A \mapsto ~~##~~3 B

*A \mapsto B $[*3]$*

A \mapsto B $[=3]$

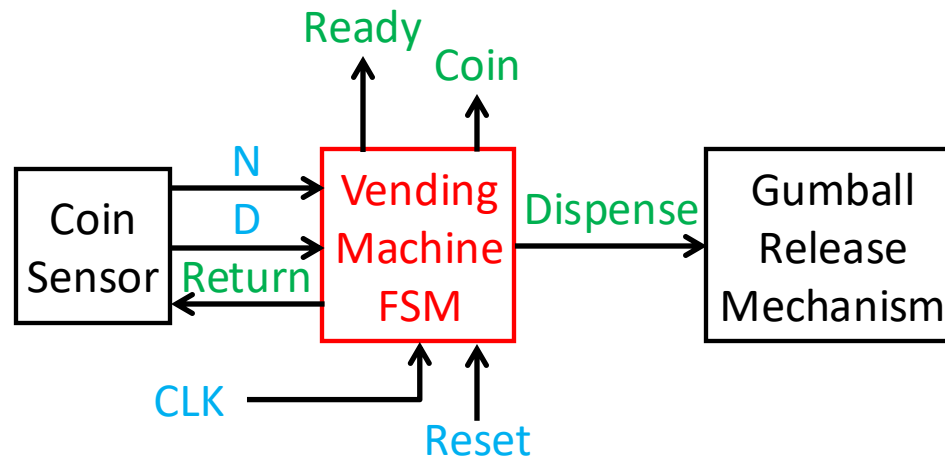
Sequences

❖ Example: rewritten handshake property

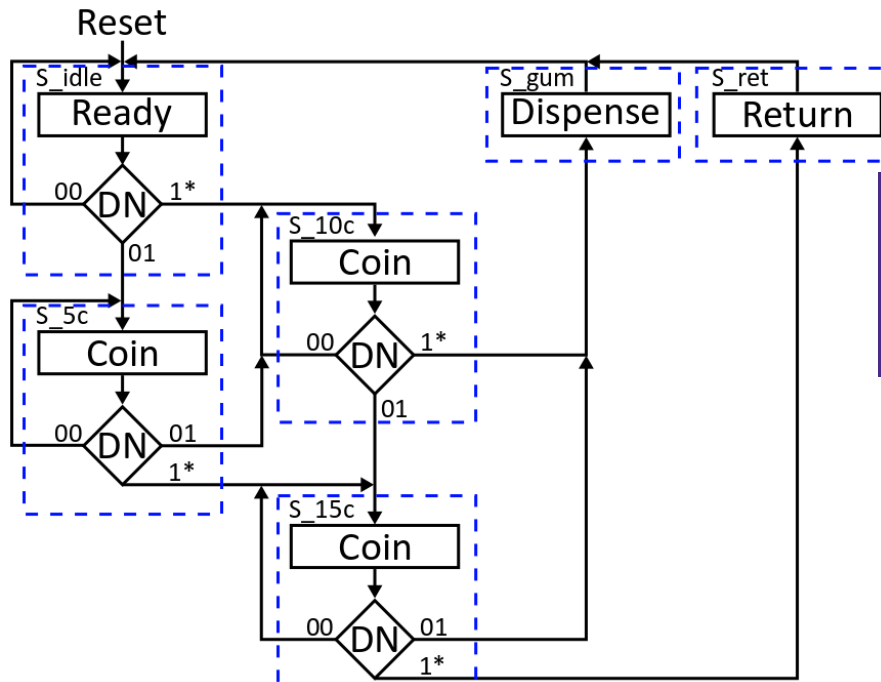
```
sequence request;  
    Req;  
endsequence  
  
sequence acknowledge;  
    ##[1:2] Ack;  
endsequence  
  
property handshake;  
    @(posedge clk) request |-> acknowledge;  
endproperty
```

Assertion Example

- ❖ Modified vending machine specs:
 - The machine only accepts dimes (D, 10¢) and nickels (N, 5¢)
 - Once 20¢ has been inserted, a gumball is dispensed; if more than 20¢ is inserted, all coins are returned
 - The machine has two lights
 - One to show that it is ready for the next transaction (Ready)
 - One to show that further coins need to be inserted (Coin)



Vending Machine ASM Chart & State Table



State	Next State			Ready	Dispense	Return	Coin
	D	$\bar{D}N$	$\bar{D}\bar{N}$				
S_idle	S_10c	S_5c	S_idle	1	0	0	0
S_5c	S_15c	S_10c	S_5c	0	0	0	1
S_10c	S_gum	S_15c	S_10c	0	0	0	1
S_15c	S_ret	S_gum	S_15c	0	0	0	1
S_gum	S_idle	S_idle	S_idle	0	1	0	0
S_ret	S_idle	S_idle	S_idle	0	0	1	0

Testing the Vending Machine

- ❖ **Dispense** and **Ready** should never be asserted at the same time
 - Write an *immediate* assertion to double-check this fact in an **always** block:

```
always @(*)  
    assert (~ (Dispense & Ready))  
    else $error("Dispense and Ready both asserted");
```

- Now write a *concurrent* assertion to double-check this fact on each clock edge:

```
property DispenseNandReady;  
    @(posedge clk) ~(Dispense & Ready);  
end property  
  
assert property (DispenseNandReady)  
    else $error("Dispense and Ready both asserted");
```

Testing the Vending Machine

- ❖ Write properties to double-check the following expected behaviors:

- From the idle state, inserting a coin should cause the **Coin** output to be asserted:

property idle_to_coin;
@(posedge clk) (ps == S_idle) & (D | N) ==> Coin;
endproperty

*↑D
↑N*

*can replace with:
Ready*

- ❖ Scope reminder:

- You may want to express an immediate assertion or property using states (**parameter**, **enum**)
- Make sure that the assertion or property is inside the appropriate module then (not the test bench)

Testing the Vending Machine

- ❖ Write properties to double-check the following expected behaviors:
 - In every clock cycle, exactly 1 of **Ready**, **Coin**, **Dispense**, and **Return** should be asserted:

↑ messy using regular logic : any of

0001
0010
0100
1000

many convenient system tasks exist to help with writing properties:

property outputs;
@(posedge clk)

endproperty

\$onehot ({Ready, Coin, Dispense, Return});

Aside: Default Clocking

- ❖ Instead of putting the clock edge in every property, it is possible to define a default clocking block:

```
default clocking clock_block;  
    @(posedge clk)  
endclocking
```

- Then you can omit the `@(posedge clk)` clause in properties and assertions!

Short Tech Break

Lecture Outline

- ❖ Testbenches (yet again)
- ❖ Assertions
- ❖ **Object-Oriented Programming**

Object-Oriented Programming

- ❖ SystemVerilog allows for OOP
 - Including inheritance and polymorphism
 - *For verification – not synthesizable (no good in ModelSim)*
- ❖ Encapsulates the data together with the code/routines that manipulates them (need a more fully-featured simulator)
 - Proper usage can yield gains in productivity, maintainability, and thoroughness
- ❖ Facilitates testing – testbench's goal is to apply stimuli and then check to see if the result is correct
 - We can model our testbenches as objects that perform a sequence of actions: create a **transaction**, transmit it, receive the result, check the result, report any issues

OOP Terminology

Blueprint for a house



Class

A complete house



Object

House Address
123 Elm Street



Handle

Turn on/off switches



Methods / actions

State of all the
Light switches



Properties / data members

Defining a Class

- ❖ A class is defined between `class` and `endclass`

```
class Transaction;  
    bit [31:0] addr; } define properties  
    function void display();  
        $display("Transaction: %h", addr);  
    endfunction } define methods  
endclass
```

The image shows a VHDL class definition for a `Transaction` class. The code is enclosed in a box. Handwritten red annotations explain the syntax: an arrow points from the word `Transaction` to the label *name*; a bracket on the right side of the `addr` property is labeled *define properties*; a bracket on the right side of the `display` function is labeled *define methods*; and an arrow points from the word `void` to the label *return type*.

- ❖ Can be defined at the top-level or within a `module` or `package`
 - Typically define each class in a separate file, or can group related classes in packages

Aside: Packages

- ❖ A **package** creates an explicitly named scope that contains declarations intended to be shared

- Can contain types, variables, tasks, functions, sequences, properties, classes, etc.
- Must be a top-level block

```
package pack;  
    class Trans;  
        // class body  
    endclass  
endpackage
```

- ❖ Package components can be accessed directly via the scope resolution operator (`::`) or imported

```
module use_trans();  
    initial begin  
        pack::Trans tr;  
        // test code  
    end  
endmodule
```

```
module use_trans();  
    import pack::*;  
    initial begin  
        Trans tr;  
        // test code  
    end  
endmodule
```

could also just import pack::Trans

everything

Constructing and Using Objects

- ❖ Create class handle, instantiate an object instance, use dot notation to access properties and methods:

```

module use_trans();
    initial begin
        // separate
        ① pack::Trans tr;
        ② tr = new();
    end
endmodule

```

```

module use_trans();
    initial begin
        // combined
        ① ② pack::Trans tr = new();
        ③ tr.display();
        ③ $write("%0d", tr.addr);
    end
endmodule

```

Handwritten annotations for the combined code:

- Red arrow from ③ `tr.display()` to `display`: *method name*
- Red arrow from ③ `tr.addr` to `addr`: *property name*
- Red arrow from `new()` to `new()`: *argument list*

- ❖ Can define/override the class constructor:

```

class Transaction;
    bit [31:0] addr;

    function new();
        addr = 371;
    endfunction

    // rest of class definition...

```

Handwritten annotations for the class constructor:

- Red box around `new()` in `function new();`: *special method*

Classes Exercise

- ❖ A MemTrans class to generate transactions for memory modules
- ❖ Create the class with the following:
 - data_in property of logic type (8 bits)
 - addr property of logic type (4 bits)
 - write property of logic type (1 bit)
 - void function that prints out the values of data_in and addr in hex and write in binary → \$ display
 - A reasonable constructor → set to all zeros?
- ❖ Create a mem_test module that instantiates a MemTrans object and invokes its function } put in an initial block

Classes Exercise Sample Solution

```
class MemTrans;
  logic [7:0] data_in;
  logic [3:0] addr;
  logic write;

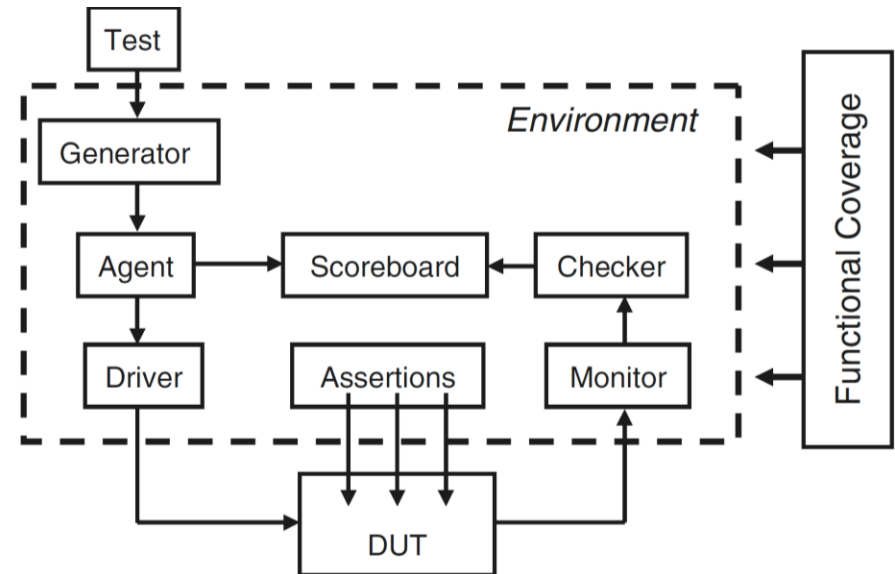
  function void print();
    $display("data_in = 0x%2h", data_in);
    $display("addr = 0x%h", addr);
    $display("write = %b", write);
  endfunction

  function new();
    {data_in, addr, write} = 13'd0;
  endfunction
endclass
```

```
module mem_test ();
  MemTrans tr;
  initial begin
    tr = new();
    tr.print();
  end
endmodule
```

Layered Testbenches

- ❖ Each block is an object and passes transaction objects
 - *Generator* creates transactions
 - *Driver* talks to design
 - *Monitor* receives response
 - *Scoreboard* compares response to expectations
- ❖ Transactions can be transferred and held in FIFO buffers for queuing



Looking Ahead

- ❖ Classes are required for SystemVerilog's constrained **randomization** features

- ❖ Randomized testing
 - Difficult to completely test large designs
 - Can be hard to anticipate all edge cases
 - Want to find unexpected errors
 - Designed tests only cover what you are anticipating